

Dennis Bathory Kitz

THE CUSTOM TRS-80 & OTHER MYSTERIES



The complete guide to customizing TRS-80 software and hardware

Dennis Bathory Kitsz
THE CUSTOM TRS-80
& OTHER MYSTERIES

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Editor's Note

ABOUT THE AUTHOR

Dennis Bathory Kitsz is a composer and active defender of the contemporary arts. However, this approach to life has resulted in pecuniary disaster which could only be ameliorated by working occasionally as a librarian, book editor, printing press operator, truck driver, newspaper editor, laborer, secretary, graphics designer, electronics technician, and typist. He is currently the director of the Dashuki Music Theatre, the president of Green Mountain Micro, and a regular columnist for several microcomputing magazines. He lives in Roxbury, Vermont with his wife Claire Manfredonia, as well as Aida, Mehitabel and Smokimoto (the cats), Fritz (the dog), and Fred and Ethel (the finches).

I'd like to say a word of thanks to Dennis Kitsz for helping us through the entire range of Murphy's laws that were proven, over and over again, in the preparation of this book. . .

I'd also like to thank:
Jim Perry, for teaching me to swim the old-fashioned way;
Jim Murphy of Bishop Graphics, for his typing ability;
Thomas Scott Nelson, for wandering through at a good time;
Nancy DeDiemar, for the kind of help you don't find anymore;
H.C. Pennington, for lending a hand with the impossible parts;
Bruce Stuart, for lending a hand with Harv;
and Charles Trapp for sticking with me to the end . . . which kept refusing to arrive.

And for those of you who have waited so long for this book, we believe you'll find the wait was worthwhile. There have been many updates of information that couldn't have been included a year ago, and several extras — including an additional chapter providing 111 cures for the common crash.

David E. Moore
February 1982

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Preface

A few years ago, when CB was king, a company riding high with those 10-4's hesitantly announced a very expensive new product. A few typed sheets with hand drawings were all store sales personnel had to explain this small 'home and business computer'. Neither Tandy nor we few first customers had any idea that the TRS-80 would become the first true home computer, the Model T of the microprocessor age.

There's no question that the 'Trash-80' carries the burden of some weak engineering and corner-cutting decisions made at the time of the CB decline, when this personal computer could have been Radio Shack's Waterloo instead of its new Wunderkind. But, like Rabelais's Gargantua, the new child soon outgrew its expectations – as well as its clothes.

And so everyone from the inner sanctum of the Tandy Corporation to the cluttered backrooms of a thousand hobbyists began the attempt to keep this new child content and amused. Thus was born the custom TRS-80. From the very beginning, and from seemingly nowhere, came forth educated hordes who would prod, poke, paint, primp, and prime this humble computer into becoming more than itself – a five-hundred-dollar supercomputer.

The machine couldn't always do it. Tandy took the rap, but also took the cash. Bruised and broken 80's littered the electronic landscape, yet also unrecognizably modified TRS-things took on tasks as diverse as business accounting, industrial control, and music making.

In this book, pathways to creating your own customized TRS-80 will be explored. The machine will be added to, opened and altered, its software patched, its uses expanded, and its breakdowns cured. If you expect that your personal computer can do more than a fancy game of video violence, then I hope you will join me in these explorations.

As this book was completed, Tandy announced that the TRS-80 – now called the Model I – is history. Production has ended, and that's a good

thing. The pressure is now off; there is no more Trash-80 to defend, but there still is a TRS-80 to put to use.

This book has been put together to please everyone. Of course it will not. Although I can't expect to help lead each of you through the intricacies of a TRS-80, I hope that, whether your wish is to jump right in and solder every wire or to learn slowly and deliberately the theory and practice of the machine, you can gain some insight from this volume.

Chapters have been arranged with basic theory and concepts toward the front. Special sections discussing the computer's software have been included to add some dimension to the concepts, and some of my own opinions, thoughts and tirades have been boxed throughout the text.

I have seen and used nearly all the commercial products presented in this book; comments, therefore, are based on first-hand knowledge, unless otherwise noted.

I have attempted in the appendices to present lists of those suppliers, publishers, terms, etc., that every TRS-80 user might want to know and not know how to find.

As the TRS-80 joins the computer museum along with the ENIAC, UNIVAC, IBM 370, and others, it still holds a more special place than many of those – with possibly half of million of its kin in use in the United States, Canada, Great Britain, Australia, Germany, and other countries. With that in mind, this book may need an occasional update.

I would appreciate receiving updated information, corrections, suggestions and criticism. Though I cannot promise a personal answer, those suggestions will be reflected in future printings of this book.

Dennis Bathory Kitsz

Roxbury, Vermont

February 1981

Acknowledgments

Someday the interdependence and cooperation among authors and programmers in this infant field will be chronicled. I can think of a hundred books or programs or articles or newsletters, each one of which gave me an essential part of the insight needed to create this volume, and without any single one of which that creation would have been impossible. Among the programmers, authors, and friends . . .

Philip K. Hooper the Codesmith, *for teaching me the simplicity and elegance of machine code*; Ron Gillen of Lab Service, Inc., *for keeping me awash in new information*; Nick "Spike" Maggio of the Philadelphia/Castor Avenue Radio Shack, *for risking life, limb and manager-ship on Tandy's red tape battlefield, and for teaching me ventriloquism*; Roger Fuller of Fuller Software and Bryan Mumford of Mumford Micro Systems, *for two reference books I couldn't do without*; Jim Perry, *for his entrepreneurial acumen and refreshing lack of humility*; Michael Comendul of 80 Microcomputing, *for pointing out that English is a difficult second language for most programmers*; Debra Marshall of the same publication, *for pointing out that it doesn't have to be that way*; Dave Moore and Thomas Scott Nelson of IJG, *for saving this book from disaster*; Charley Butler and Joni Kosloski of The Alternate Source, *for my author's carte blanche*; Dave Beetle of the pioneer-

ing On-Line, *for really starting it all*; Harv Pennington *for jollies*; and for many and varied favors: Bill Johnson (*Cleveland Users Group*), Vaughn Jupe, Ron Troxell (*Personal Micro Computers*), Don Stoner (*The Peripheral People*), Lee Perryman (*The Associated Press*), Wayne Green (*himself*), Eric Maloney (*Kilobaud*), Jack Decker, Mike Barton (*MSB Electronics*), Leo Waltz, Jerry Sabin, Fred Blechman, Don Stevens, Walt Auch III, Vince Schulz, Bill Archbold (*Archbold Electronics*), Al Abrahamson (*Norwalk Users Group*), Bill Barden, Don Valentine (*Tecmark Associates*), Harold J. Matts, Stan Ockers, Thomas Frederick (*ABS Suppliers*), Les Logan (*TCS Newsletter*), Brian Harron (*Ottawa Users Group*), John Bilotta, Gregg Shadel, Don C. Tatum (*Barre-Montpelier Radio Shack*), Andrew Law, *the many manufacturers represented herein, and the usual coterie of others forgotten and maligned. Special thanks to the anonymous author of the TRS-80 Technical Reference Manual for a job well done.*

This book is dedicated to my wife, Claire Manfredonia, because she'll do most anything for a good gag. One day she walked into a serious meeting at an engineering company and hit me in the face with a blueberry pie. I won't tell you about the marshmallow fluff.

d.b.k.

March 1982

Introduction

The Tools You Will Need

Your basic TRS-80, with some attachments and software, is a thousand-dollar item. So I'll not encourage you to use dime-store tools. Buy the best you can afford, keep them clean, and reserve them just for use on the '80. Don't double up tools with the family auto. You may not need them all, but here is my customizer's toolbox:

A medium-sized flat-blade screwdriver and Phillips blade screwdriver (a reversible combination is ideal). With these you open cases and remove cabinets.

A jeweler's set of flat and Phillips blade screwdrivers; hex nut drivers are optional. These drivers can be used to align tape heads, help make delicate wire bends, adjust trimmer controls and even repair watches.

One very thin screwdriver for lifting integrated circuits out of sockets. This will be its only purpose, but the first time you break the pins off a \$10 jumper cable, you'll wish you'd used it!

Small scissor-type cutters (manicuring types are excellent). These will be used for snipping leads in tight spots.

Small diagonal wire cutters and/or front-cutting 'nippers'. Your general purpose cutters. They are fast and easy to use, but not to be used for heavy wire around the house.

Needlenose pliers (two pairs, normal and 90-degree types). You'll need these for bending leads, also extracting bits and pieces you've dropped into a nest of wiring.

An X-acto type knife, with a strong blade and handle you feel comfortable with. Since this will be used to cut delicate solder traces, you should be able to handle it deftly. I use a single edged razor blade, but have leather fingers!

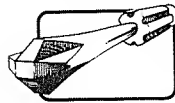
A scalpel, if you can get one. For very delicate trimming and scraping; **a dental pick** for pulling off solder balls or lifting parts off a board (get this item from an obliging dentist — they are often discarded when worn); **tweezers and needlepoint hooks**. The latter come in handy for tracing incorrect wire-wrapped connections.

Rat-tail, triangular, and flat files. These are only for sprucing up the cosmetics, so if you don't care how it looks, save a few bucks.

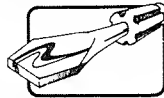
A wire-wrapping tool. The decision on this can be tough. If you can afford it, get one of the electrically operated slit-and-wrap types. Stay away from 'just wrap' tools, since they depend on the sharpness and quality of the sockets; also they are useless for wrapping capacitors or resistors. I use a simple double-ended tool sold by Radio Shack for about \$5. It wears out after a thousand or so connections, but it fits my hand well, and is not clumsy like some electric units.

A soldering iron. The decision is not easy. Should you spend top dollar and get an expensive one or buy a cheap unit that can be discarded when it wears out? I use a \$5

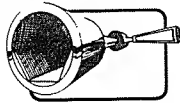
To help you prepare for each project, the following graphic symbols have been used as a key to the tools needed for each project:



Phillips screwdriver



Flat-blade screwdriver



Sockets



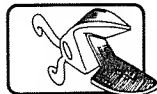
Thin lifting tool



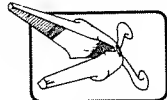
Scissors type cutters



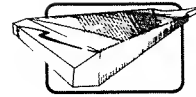
Solder



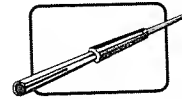
Wire cutters



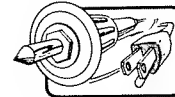
Needlenose pliers



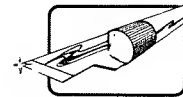
Various files



Wire wrapping tool



Soldering iron



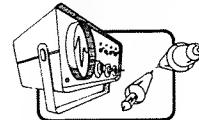
X-acto knife or blade



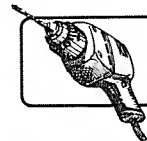
Solder wick



Multimeter



Oscilloscope



Drill

soldering iron which can be junked when it gets beat, but my editor uses the best he can get (a \$30 temperature-controlled one).

I file a set of \$1 tips to my satisfaction, and lubricate the threads with white heat sink grease. This way I have a few different tips at my disposal; with plated bits you **never** file the tips.

A Multimeter. The voltage regulators in your TRS-80 are very good, so any problems will usually show up as gross errors. This offers you a way out of buying an expensive multimeter; for most of these projects, the \$10 pocket variety will suffice.

However, for lots of repair work a better meter is in order; I use a \$40 type (not digital!) for my work.

An oscilloscope. For the projects, no. But for repairs, yes. Don't panic thinking of a thousand dollars for a digital scope, because an old color television scope will do perfectly well; they can be found in the bargain bins for \$50 to \$100. If it saves you

a \$100 repair bill, you've paid for it. Mine is an old RCA type WO-90Q, built for early color TV, and just fine for the bulk of TRS-80 work.

You will also need supplies in the tool box. Among these are:

Solder. Get the best you can afford. There's nothing so unpleasant as a great glob of the stuff between two traces on a board. Order the multicore rosin flux type, and stay away from most the off-the-shelf stuff. Remember, multicore rosin type only, and the finer the gauge the better. **Never use acid flux solder**, as used by plumbers and electricians.

Soldering wick. Marketed under the names *Spirig*, *Solder Up* and *Solder Wick*, it's a copper braid impregnated with soldering flux. When heated with the soldering iron it absorbs Solder off the board, thus freeing components. Don't do without this stuff unless you like fried circuit boards and burnt fingers.

Wirewrap wire. Also called by the trade name *Kynar*, this is 28- or 30-gauge single-strand wire used to interconnect the pins of wire wrap sockets. It comes in an assortment of colors; get them all, so you can keep data, address, power and ground lines separate.

Multiconductor cable. The more flexible wire is easier on the coordination, but also the most expensive. Best buy is *Spectra Twist*, and its kin, from surplus houses. If you need jumper cables, buy them; Making a two-ended, 40-pin jumper cable can be three hours of maddening work.

Bus wire. This is solid, uninsulated stuff. A small roll will do for a lifetime. I use it for wiring, securing bulky capacitors to circuit boards, holding bundles of things together and for making special tools.

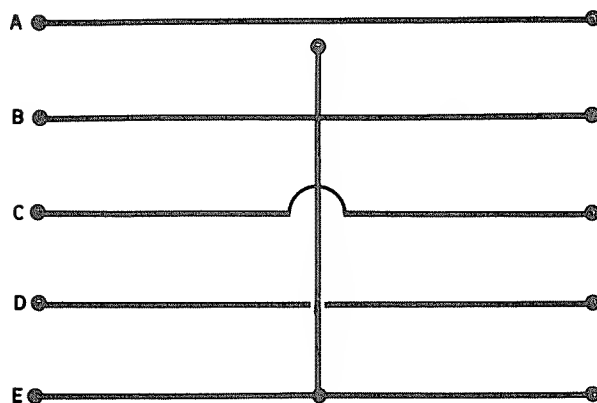
Miscellaneous. Sockets, perforated board, mounting hardware, and such will always be needed.

Details about supplies needed for each project in this book will be presented with the project. Except for integrated circuits, most of the items are available right off the shelf at a local Radio Shack or other electronics supply house.

Schematics

Schematic drawings of electronic circuits are identical to maps. They show routes, direction, junctions, relative importance and functions of locales, two-way and one-way streets, traffic flow and congestion and so forth. At first, the symbols may seem like the mysterious hieroglyphics of a secret society, but their symbolism can soon become as familiar as a roadmap. Even strange places can be assessed from afar.

First, the symbols. A line is a wire running from some point in the circuit to another. Consider the sketches below:

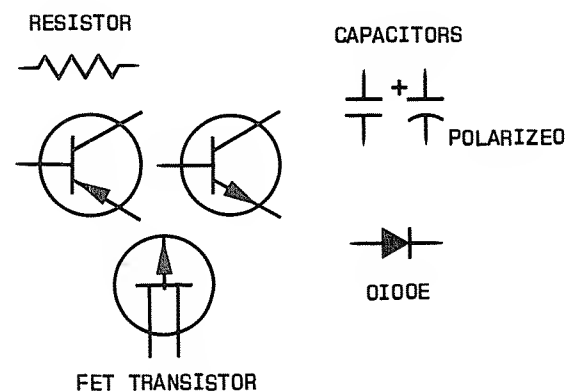


The first drawing is a simple wire. The electrical path moves from one point to another, in either direction. By following the path of a wire through a circuit, the pattern of connections can be discovered. When wires are forced to cross one another, but not to connect with each other, it must be made clear. On a roadmap, non-intersecting roads are shown either by a break in one of the intersecting lines, or in showing interstate highways, merely by crossing one 'below' the other in a different color.

Sketches b, c and d are the three ways of drawing wires which do not connect to each other. The first, simply crossing them, is the most common. The second method places a semicircular bump in the crossing path, and is used by Sams Publications in this country and commonly in Europe. Occasionally the broken path crossing shown in sketch d is used.

When wires connect, a dot is used to clarify that a connection is to be made. Occasionally, you may come across earlier schematics which use the 'bump' method of showing unconnected wires. On these schematics, the lack of a bump indicates wires are connected.

The wires (or patterns of copper etched on circuit boards) connect electronic components. Some of them are:

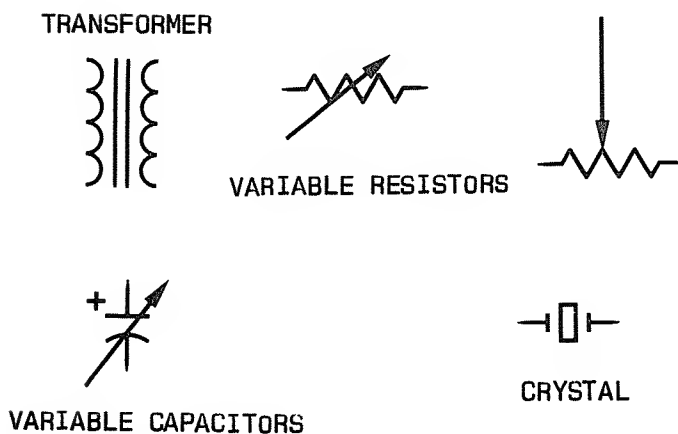


Since this is a lesson in reading schematics and not electronic theory, I recommend that you turn to an excellent book by Forrest Mims, 'Engineer's Notebook', sold by Radio Shack, for an introduction to what each of these parts does. Briefly, the symbol for a resistor has the flavor of a long wire being compressed, meaning the electrical flow is somehow being resisted. The innards of a capacitor generally consist of metal foil separated by a non-conducting paper or plastic, and the capacitor's schematic symbol is fairly representative, with two plates facing each other but not joining.

Some capacitors are designed to fit into a circuit in only one direction; the positive (+) sign identifies that direction. These capacitors are identified on their bodies by a positive or negative sign. Another one direction (polarized) device is the diode. It consists of an arrowhead striking a barrier, implying that current may flow in the direction of the arrowhead, but not back across the plate. The body of a diode may have the diode symbol imprinted on it, or a band to indicate the 'barrier' end.

The transistor usually has three connections (such connections are called 'leads' on small parts such as these). These leads are identified as collector, base and emitter or source, gate and drain, depending on the transistor type. This will be shown on the diagram, and the transistor will be imprinted with the information, or it will be provided on the package in which the transistor is sold.

A few other symbols are:



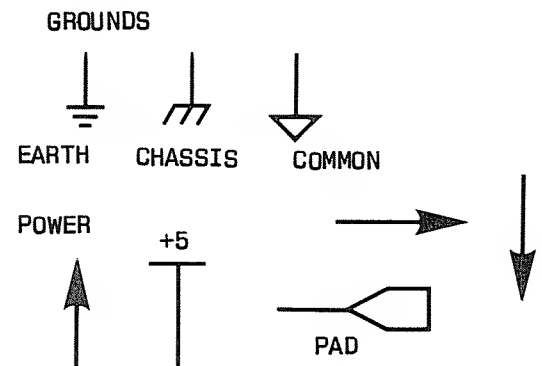
The first is a transformer, whose job it is to take current fed into one coil and induce that current into a second coil. An iron or ferrite center (the parallel lines in the symbol) aids in efficient transfer of that current.

The next three symbols look like resistors and capacitors, which they are. The added arrows show that their values may be varied; hence, they are called variable resistors and variable capacitors. The variable resistor is best known as the volume control on a television, and the variable capacitor is found as the tuning control on a table radio.

The last symbol is a crystal, a piece of cut quartz capable of vibrating (resonating) under certain electrical conditions. Because a crystal is a very accurate, fixed, molecular device, it is capable of resonating (also called oscillating) at

precise intervals. It is used for the master control of all pulses in the TRS-80.

A few directional symbols are now in order:



The first are known as grounds, and they are used to indicate a potential of zero or neutral voltage. The first of the trio is an earth ground, commonly used in radio, television and hi-fi schematics, but purists use it only to describe an actual connection to a ground spike or cold water pipe. The second is a chassis ground, indicating an electrical connection to the metal case which encloses the circuit. It is often (though incorrectly) interchanged with the earth ground.

The last of the three grounds is a 'common' or neutral ground, and the one which is used to indicate the zero voltage line in the computer. All other voltages within the computer system are described in terms of their relation to this ground.

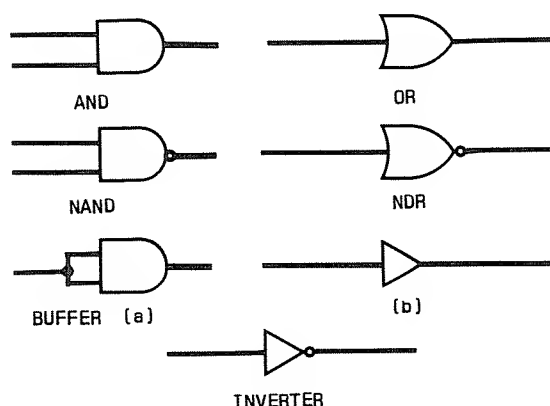
The next quartet of symbols indicate power. The up arrow generally points to an actual voltage value (such as +5 or +12). The horizontal line indicates merely a 'high' level, that is, a connection is made to the normal positive power supply for the circuits in the system (+5 volts in the TRS-80).

Non-positive voltages have no standard symbols. Negative (or below ground) voltages can have either a horizontal arrow or a down arrow, pointing to the voltage desired at that point. The schematics tells you that a connection is made to the voltage level shown.

Another use of a horizontal arrow is to point to important connections to be made elsewhere on the schematic or on other sheets of the schematic. In the former case, the arrow is used because actually drawing the wire may clutter the schematic, making it illegible. When you see an arrow, be sure to find the other end of the connection described (indicating words such as 'clock', 'mem' or 'port FF' may be used as guides to where the connection is made).

Another useful symbol is the last of the group above, the pad. It indicates a significant connection, usually to another device or circuit board. Using this symbol makes it clear that the connection is to be made somewhere off the board on which you are working. In this book, I have not used these symbols where indicating a connection to the TRS-80; instead, the cable to the TRS-80 is shown with the connecting wires striking a wide vertical band marked 'TRS-80 Edge Connector'. Other types of off-board connections, however, are shown with the pads.

The most common families of parts found in computer circuits, however, are shown below:

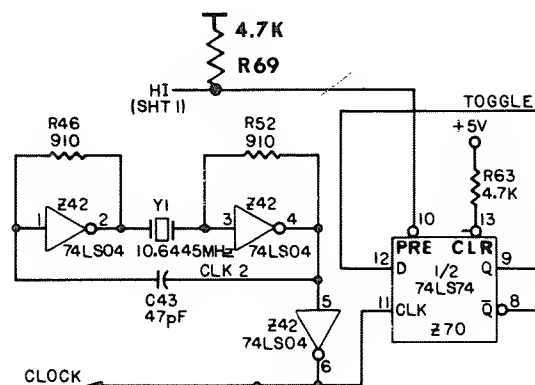


These symbols represent integrated circuits, those multiple lead, buglike packages that handle the bulk of the work in the computer. Briefly, these are logical building blocks. Sometimes there are several blocks in one integrated circuit, and these various blocks may be scattered throughout the circuit diagram. This can be confusing when actually building a circuit, but since pin (lead) numbers are given, you only have to remember where you put the part.

You should know that the TRS-80 *Technical Reference Handbook* uses what are called 'functional' schematic elements, meaning identical parts are not necessarily drawn the same throughout the schematic. I have chosen not to use this method, which, although it makes circuit operation less evident, is clearer when doing actual wiring.

Complete logical and physical diagrams of every circuit used in this book are given in an appendix. Those diagrams will help give you an idea of how these logical blocks are packaged inside 8-, 14-, 16-, 18-, 20-, 24-, 28- and 40-pin cases.

Basically, that covers reading a schematic roadmap. Below is a section of circuit. See how the logic elements are connected to each other as well as to two resistors, a capacitor and a crystal. Notice also that the logic elements are all marked 'Z19', since they are separate blocks within a single component. An arrowhead indicates a wire leading off the board, and power and ground connections are shown. The numbers on the logic elements are the pin numbers for the component connections:



Be Tolerant

Every electronic component is manufactured to work within specific limits, whether they be accuracy, temperature, speed, power use or other limit. These are the components parameters or *tolerances*. The circuits in this book have been designed to use the most commonly available parts, so the matter of tolerances is rarely important. However, sometimes those tolerances are important, such as when talking about memory speed or power supply voltages.

Power supply should be within five percent of the voltage specified; a supply indicated at five volts may vary only from 4.5 volts to 5.5 volts. By using the power supply regulators shown in the schematics, these voltages should not be of concern. Unless you are familiar with power supply design, do not attempt to use other methods of regulation.

Very few of the resistors have tolerances noted on the schematics. The rule of thumb is one quarter watt at five percent, but if you can only obtain half watt units, or 10 or 20 percent resistors, don't be concerned. The quarter watt

resistors are a bit less costly and are a bit more aesthetically appealing. Consider also that if a resistor is specified as 1,000 ohms, a 20 percent deviation gives a range of 800 ohms to 1,200 ohms. Thus, the standard values of 910 ohms or 1,200 ohms should do as well.

Capacitors are notoriously sloppy in their tolerances, especially electrolytic types (those whose polarity is marked on the schematics). These normally vary from 20 percent low to more than 100 percent high – thus, when a 500 microfarad capacitor is noted, it can range from 400 to 1,000 microfarads. Also, there is some revision in the standard numbering method used for parts values: 470 microfarads is now being called 500 microfarads, for example. So when you try to obtain a capacitor value marked in the parts list, remember that a nearby higher value is fine.

Voltage parameters for polarized (electrolytic) capacitors are important. Never get an electrolytic capacitor with a value less than that specified, but do not hesitate to take one with a higher voltage parameter. That is, a capacitor specified at 47 microfarads, 16 volts, can be replaced with one specified at 50 microfarads, 35 volts. It may be physically larger, but it will work equally well.

If you walk into a store and hand the sales clerk a parts list, don't be surprised if you are asked a few more questions. You might be faced with choosing between parts which are identical as far as the parts list in this book is concerned, but which include other parameters.

Resistors can be carbon composition, carbon film, glass or wire wound. These days, carbon film is common and cheap, and that's your first choice. Carbon composition is the next choice at a lower quality, and glass is excellent but at a higher cost. Forget wire wound, because they can contribute unwanted side effects.

Ordinary capacitors are manufactured in many ways: ceramic, polystyrene, polyester, silver mica, polycarbonate and paper. For the bypass capacitors necessary for all the circuits in this book, ceramic types are your choice. Cheap. If you get silver mica, so much the better, but you'll pay a price. Watch out for polystyrenes or polyesters if you plan to solder, because they are delicate and you can damage them with too much heat. Otherwise they are excellent, but quality overkill. Polycarbonates are slick types, and you might consider using these if you build the 8-track mass storage system. Run the other way if you see paper capacitors.

Electrolytic capacitors come in two basic types – metal cans (covered with plastic), and those manufactured using tantalum (an expensive metal of great strength and purity). For most digital projects, choose the ordinary cans. Tantalums of the same value, although smaller, high quality, and very pretty looking, are costly and not required here.

Digital integrated circuit part numbers are generic, which means that a 74LS00 circuit might be sold as an SN74LS00 or an NEC-74LS00. The prefix characters refer to manufacturers. On the other hand, those parts whose numbers contain 'LS' may not be substituted by parts marked 'S' or 'C' or by those with no markings. 74LS00 may not be replaced by 7400, 74S00, or 74C00, nor may they be exchanged for each other. When integrated circuits are specified, try not to substitute with other circuit 'families'.

This section will not make you a master schematic reader; only practice will do that. Pick up copies of the Engineer's Notebook mentioned above, as well as various of the project books sold by Radio Shack and others.

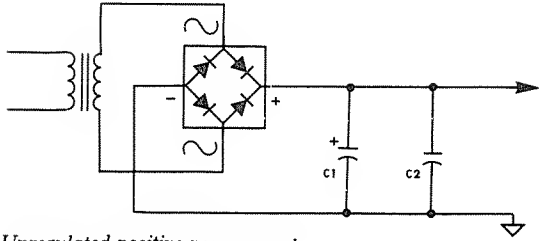
Building a Power Supply

All the projects presented in this book will either be modifications to the TRS-80, in which case they will draw power from the computer itself, or outboard devices which will need power. There are several ways to get this power:

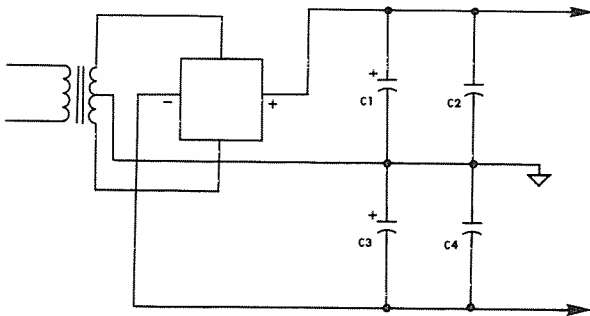
1. Build a power supply from scratch for each device, including transformer, rectifiers, capacitors and regulators.
2. Build a main power supply providing substantial current at between 7.5 and 15 volts, and put a voltage regulator and a capacitor on board each project.
3. Use an assembled power supply providing a 7.5 to 9 volt output for each project, and use a voltage regulator and a capacitor on each board.
4. Use an assembled power supply providing 5 volts and feed all boards from it.

My recommendation? Probably an assembled power supply of 7.5 to 9 volts feeding each project. Power supplies are important to projects, and they should be well regulated, free of residual 60 Hz (Hertz equals 'cycles per second') ripple, and should not tend to transmit signals between projects. Some things to consider include:

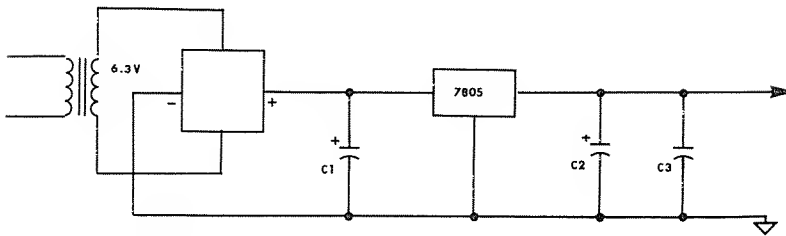
Building a Power Supply



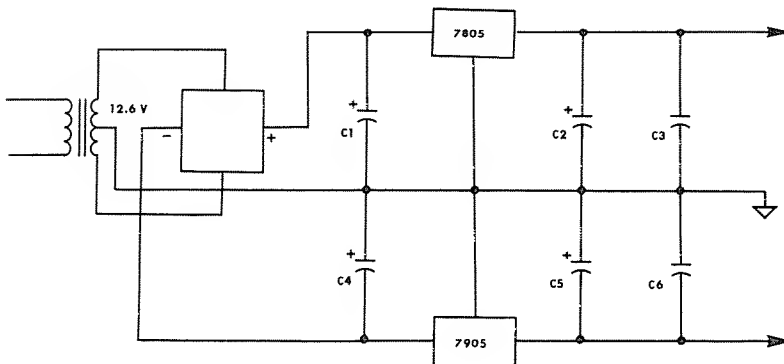
Unregulated positive power supply.



Unregulated bipolar power supply.



+5 volt regulated power supply.



+5/-5 volt regulated power supply.

1. Building a power supply is not very time-consuming, but it is probably going to be bulkier than the project being constructed. Then there's the question of where do you put it? If it's in the project, you have to be very careful to shield the AC primary and secondary leads. Also, you have to be sure it's working properly before you can begin to test your project.

2. Small, assembled power supplies are inexpensive. They are normally sold as 'battery eliminators', with their current capabilities specified. The AC leads are encapsulated in plastic with the rest of the supply. Although you need a separate house-current outlet for each of these supplies, the work you do (both building and testing) is lessened and the safety to you and your project is increased.

3. Larger power supplies are expensive to buy and complicated to build. Unless they and each project being fed contain plenty of transient suppression (in other words, lots of extra capacitors), the actions of one device may affect another. But they do tend to be more immune to house current fluctuations than small homemade or purchased power supplies. With regulators on each project, moreover, you can provide more immunity to spikes and fluctuations. They are truly 'brute force' circuits.

4. Large regulated power supplies are highly stable, but expensive. They are capable of feeding a whole range of boards, less house current outlets are needed, and the level of regulation is usually substantial enough to prevent fluctuations at the board level. However, onboard filtering is still necessary to prevent interaction among external devices.

The circuit diagrams that follow present a simple, unregulated, positive power supply; a simple, unregulated, bipolar power supply; a regulated five-volt power supply; a regulated dual-voltage (+5 and +12 volts) power supply; a regulated bipolar (+5 and -5 volts) power supply; a regulated four-voltage (+12, +5, -5 and -12 volts) supply, and last (but by no means least) the design of the power supply (+12, +5, and -5 volts) used in the TRS-80. All these supplies can use the transformer / rectifier 'power supply' sold for the TRS-80 as their source.

Power Supplies

Digital electronic circuits constructed by hand can perform as well as neatly-laid-out commercial circuits. In some ways, though, these circuits have to be designed better than commercial ones, because professionally etched boards are usually designed with careful consideration given to signal paths. Wire-wrapping or soldering, on the other hand, can look like a rat's nest and sometimes act that way.

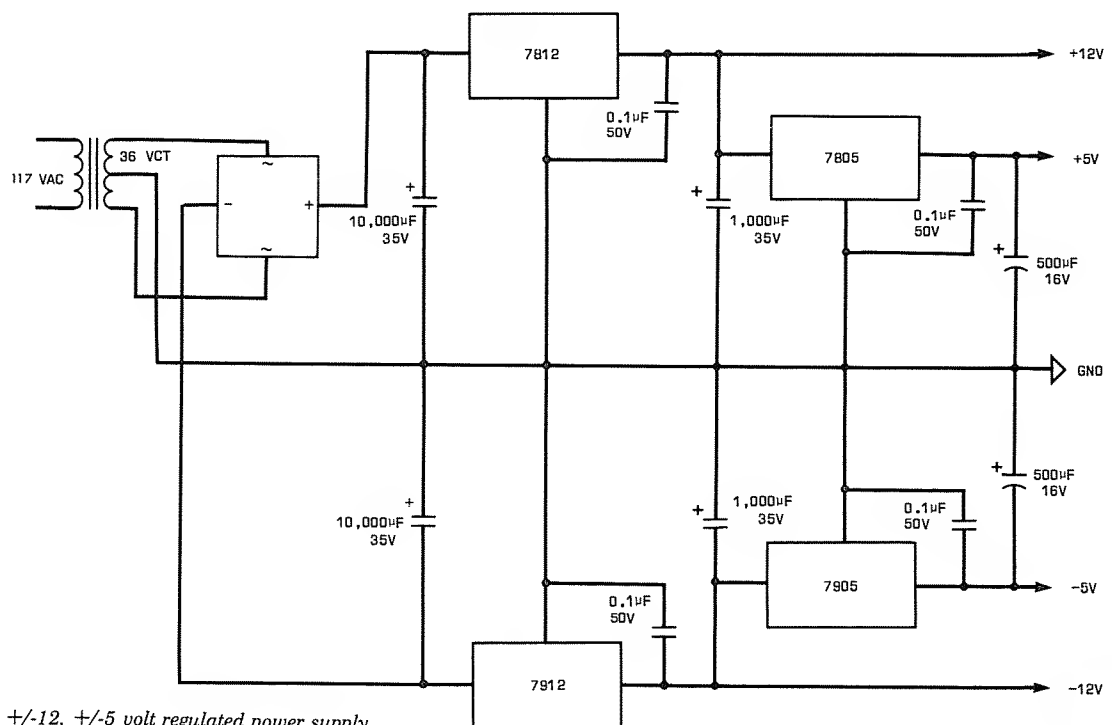
Most important: always use bypass capacitors in your projects. These are capacitors with a value of about 100 nanofarads (0.1 microfarads), ceramic discs or dipped mylar, attached between +5 volts and ground. These are *physically* placed near their respective integrated circuits. When I finish wire wrapping a circuit, I give it a quick test to make sure the wires are connected properly and the circuit behaves normally (barring occasional crashes). Then I attach a 100-nanofarad bypass capacitor between the positive power supply and ground, directly between the wire-wrap pins carrying the power.

Secondly: add termination resistors to the data lines whenever more than one project is connected to the computer. Termination resistors act as electronic clamps, holding the

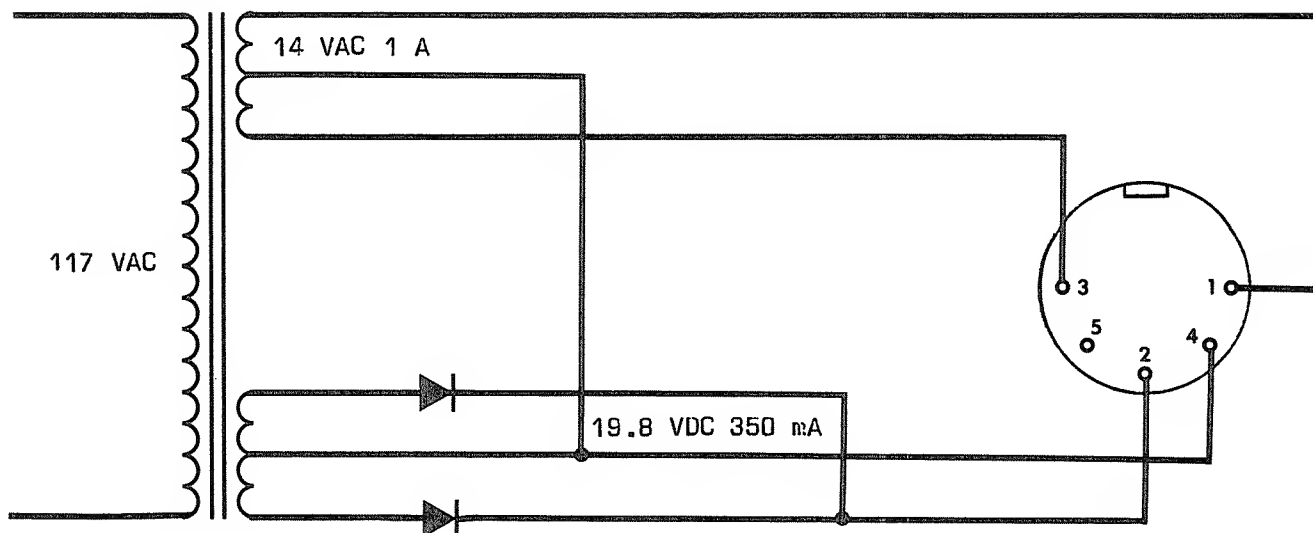
lines steady as the signals sweep through. These eliminate random noise, as well as signal 'overshoot', caused by capacitance introduced by the wires themselves. If you have an expansion interface of the new breed (no buffered cable), termination resistors are already in place. They are also standard with the *LNW* and *Microtek* interfaces (though I recommend doubling the value suggested by *LNW*, to at least 470 ohms and 1,000 ohms).

The simplest addition of termination resistors involves connecting a 1,000 ohm resistor from each data line to ground. To ensure even better signal clarity (and also to maintain the 'high' level the computer normally sees on its lines), also attach a 470 ohm resistor from each data line to +5 volts. In all, then, eight resistors will be needed for each project you add.

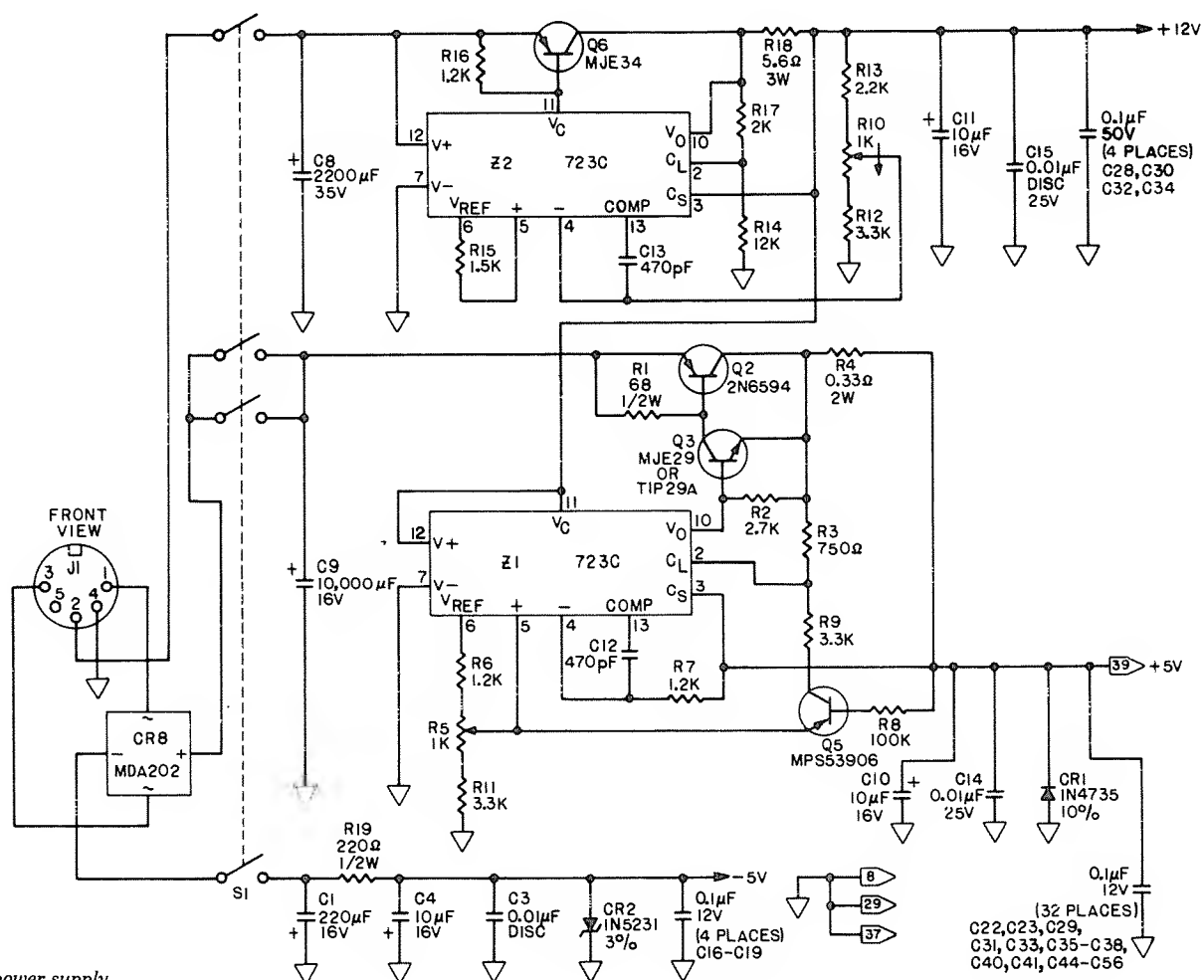
Finally, do not skimp on the power supply. Generally you will see a 2,200-microfarad capacitor followed by a regulator, a 470-microfarad capacitor and a 100-nanofarad capacitor. This is the absolute minimum configuration for a working power supply. If you have the room, increase the value of the first capacitor to 4,700 or 10,000 microfarads, and place a 100-nanofarad capacitor between input and ground physically near the voltage regulator.



+/-12V, +/-5 volt regulated power supply.



TRS-80 transformer adaptor and supply.



TRS-80 power supply.

Those Colors: What They Mean and How to Read Them

The color codes used for resistors, capacitors, and other parts are brought to you by the same folks that brought you phrases like 10W-40 and RS-232C: the standards-setting powers of the engineering industry. It becomes an international shorthand.

The colors are black, brown, red, orange, yellow, green, blue, purple, grey and white. If you can't immediately remember that, then pick up a piece of multi-conductor 'rainbow' cable. The colors are all there in the same order. The table below presents the color codes and how they can be read on the bodies of resistors, capacitors, and diodes.

FIRST AND SECOND COLOR BANDS		THIRD COLOR BAND	
BLACK	0	BLACK	0
BROWN	1	BROWN	X 10
RED	2	RED	X 100
ORANGE	3	ORANGE	X 1000
YELLOW	4	YELLOW	X 10,000
GREEN	5	GREEN	X 100,000
BLUE	6	BLUE	X 1,000,000
VIOLET	7	SILVER	100
GRAY	8	GOLD	10
WHITE	9		

FOURTH COLOR BAND IS THE TOLERANCE		
GOLD = 5%	SILVER = 10%	NONE = 20%



What do these values mean? Resistance is a kind of objection to electron flow, measured in ohms (pronounced with a long O). The abbreviation is a Greek omega (Ω). Thousands of ohms are kilo-ohms, or just kilohms, and abbreviated K (k in Europe). Millions of ohms are megohms, abbreviated simple M. The ability of a resistor to withstand electrical current is measured in Watts (W). Most

computer work is done with 1/4 Watt resistors.

For resistors without color bands, the values are stamped on using R (instead of omega) for ohms, K and M.

Capacitance is the inclination of a non-conducting object to store an electrical charge, measured in Farads. The abbreviation is a capital F. Since this is a very large amount of capacitance, real work is generally done in millionths of Farads, or microfarads (mF), and millionths of millionths of Farads, called picofarads (pF). Since many of the more popular capacitance ranges for computer work fall between these two figures, the abbreviation for thousandths of millionths of Farads, or nanofarads (nF) is common in Europe. The ability of a capacitor to withstand voltage is measured in voltage tolerance (V).

Capacitance is usually printed on the capacitor in mF; color bands are rare. Picofarads are marked 'p'; the absence of an abbreviation indicates microfarads. Note that these capacitor 'base values' are equivalent: 18=20, 27=30, 39=40, 47=50.

 Abbreviations and Conventions Used
 in this Book

AD-A15	Computer address lines.
C	A capacitor, specified in the parts list; positive and negative sides are marked.
CLK	Clock, usually a flip-flop input.
CLR	Clear, usually to flip-flop or latch.
CR	A diode, Radio Shack schematic reference.
D	A diode, specified in the parts list.
DD-D7	Computer data lines.
F	Ferrite; used only as part of mF, uF, pF.
GND	Common computer ground.
K	A relay, specified in the parts list.
-K	Resistance value suffix = x 1000.
LED	A light-emitting diode, specified in the parts list.
m-	Capacitance prefix for x 0.000 001.
M	A motor, specified in the parts list.
-M	Resistance value suffix = x 1,000,000.
MHz	Megahertz (million cycles per second).
p-	Capacitance prefix = x 0.000 000 001
P	Piggybacked integrated circuit.
PRE	Preset, usually to flip-flop or latch.
Q	Transistor, specified in the parts list.
Q	Flip-flop or latch output.
R	A resistor, specified in the parts list.
S	A switch, specified in the parts list.
T	Transformer, specified in the parts list.
u-	Capacitance prefix for x 0.000 001.
U	Integrated circuit (LNU references only).
Vcc	Collector-supply voltage, that is, the 'five-volt supply'.
W	Watt; resistor power reference.
X	DIP shunt or jumper, R/Shack reference.
X	Crystal oscillator.
Y	Demultiplexer output signal.
Z	An integrated circuit, specified in the parts list.
*	Used in text references for 'not', in place of a horizontal bar across the top of the abbreviation.
DV	Ground, zero-level, or center-ground.
+5V	Five-volt positive regulated power supply, the 'Five volt supply'.
-5V	5-volt negative regulated power supply.
+12V	12-volt positive regulated power supply.
117V	House current, 105-120 volts AC.

BASIC commands; Z-80 mnemonics and opcodes; CPU and other schematic signals are printed in UPPER CASE letters. Hexadecimal numbering is printed in **BOLD** letters.

1

Getting Inside

A TRS-80 is the proverbial black box. It's the first mass-marketed personal computer that was intended for simple home and business use. As far as retail sales are concerned – whether it's a car, a hi fi, or instant pudding – thinking of it as a black box is just fine. “Yes, it's a complete computer with a keyboard unit, video display, and cassette storage.”

To have this small computer work to even a fraction of its potential, though, you've got to have the power to control it. As with driving a fine automobile, that power comes with understanding, comes with being able to look through and into the box, comes with keeping it carefully tuned and customized so that it can respond to your wishes.

In this Chapter we'll open the electronic black box a little at a time. If this is all new to you, prepare for some exciting finds inside this box; if not, then follow along and discover just how much you've learned about microprocessors in their own age.

What You See

The main computer unit is in a small case looking something like a typewriter that has lost its carriage. There are 53 keys in a typewriter style layout, and perhaps on your unit a numeric keypad to the right. Connections on the back are marked *Power*, *Video*, and *Tape*, along with an on/off switch.

The unit is entirely silent. It is electronic in conception and execution.

A power supply lowers and isolates the 120-volt household supply to approximately one-seventh that value, and that feeds the

keyboard unit. An ordinary cassette player is cabled in, and a partly disemboweled television set plugs in place, serving as a crude video monitor.

The designers of the TRS-80 struggled with, and won, the battle of familiarity. Televisions, cassette recorders, and typewriters are among the most ordinary of home or office appliances. But by winning this battle of familiarity, the designers also clearly set themselves up to lose the battle of reliability (more on that later).

Let's first take a look at some of the things the computer does, and then very generally try to discover how and with what the machine does them.

When you power up your computer, you expect to be able to communicate with it. Unlike a television, it does not entertain, but rather evaluates and responds with an electronic psyche. If it were not capable of using a human-like language, we would be forced to use the machine's language. But since it will be asked to accomplish human tasks, we will demand that it speak a human sort of language – BASIC (Beginners' All-purpose Symbolic Instruction Code). BASIC has grown since its humble but inspired beginnings at Dartmouth College into a formidable tool capable of rocking other standard languages off their computational pedestals. If you have been using your TRS-80's BASIC, you know its fluidity. But it does not work alone.

The BASIC that is in the TRS-80 works hand-in-hand with the electronics to produce a video display for us to read, and to examine a keyboard for us to type on. The keyboard gives

input to the computer, the video display shows output from the computer. Input and output are grouped together in computer terms, and are collectively called *I/O*. The tape recorder, which saves and loads computer programs, would also be called *I/O*.

Forget the claims that the TRS-80 computer can do what this-or-that state-of-the-art computer could do five years ago. Maybe so, but it really can't. Because it isn't built like a piece of office furniture. Because it doesn't act, work, or 'think' the same. No matter how you paint it or pad it, a four cylinder sportscar will never be a luxury V-8 sedan. And Burger King is not a candidate for the four-star list. But they can do *what* they can do as well or better because of the simple, direct, streamlined nature of their operation and conception. Turn away the question of value judgment, and you discover that – considering portability, cost, ease of maintenance, accessibility and vastness of the software domain – the TRS-80 is probably a dimension *better* than that recently demised state-of-the-art dinosaur.

That about covers what we can say about the minimal TRS-80 setup – a keyboard computer unit with *I/O*. Before popping off the cover, let's name some of the rest of the *I/O* devices that can be hooked up to the machine. Those might include:

An Expansion Box.

This attachment expands not only the internal capacity of the computer, but also forms an electronic saddle, permitting other devices to ride on the back of the keyboard unit.

A Printer.

Obviously, an attachment to provide a permanent record of the machinations of the TRS-80.

A Diskette Drive.

A place to store information and programs when the computer is turned off; it is very like the cassette player, except that it is speedier and can be accessed differently. More on that later.

Voice Input/Output.

An ability to speak and be spoken to on the part of the computer. This is one of the experimental options.

Telephone Communications.

Communications with other computers,

similar or different, nearby or across the world.

Control Centers.

The power to change, activate or extinguish electrically-powered equipment throughout a home or office.

There are more devices which might be considered, including clocks to tell the time of day, little circuits to create sound effects and music, even other computers used as 'slaves'. In fact, where electrically controlled equipment is involved, almost anything can be attached directly or indirectly to the TRS-80.

Hesitation

When you open the computer's case, you'll see an enormous amount of electronic circuitry. Before you open it though, you should have an idea why there are so many *integrated circuits* and circuit board *traces* connecting them.

We live in an analog world. We judge size or volume or loudness not by how big or full or loud it *is*, but by how big or full or loud it seems in relation to something else, even if that 'something else' is merely what we are used to hearing in our normal world.

In other words, all our evaluations are made by analogy. "How big is it?" "As big as a basketball!" "Is she pretty?" "Pretty as a picture!" "Is it far?" "About a stone's throw from the corner." Our cliches are built on comparison or analogy. Ideally, then, we might like to build a machine that work for us in our own terms . . .

"Machine!"

"Duh, yessir, sir."

"Add fourteen and thirty-seven, machine!"

"Yup, yup. Lessee. Hmmm, fourteen is **this big**. And thirty-seven is **this big**. That makes a number *this big*. . . . Sir?"

"Yes, machine?"

"That's as much as thirty-seven and fourteen, sir! Looks just like fifty-one, sir."

"Thank you, machine."

This computer has a rather limited voice capacity, but what it did inside itself was electronically quite sophisticated. It took in a value and transformed it into an electrical voltage of fourteen units, stored it, accepted a

second value and transformed it into a voltage of thirty-seven units. Then it added the voltages together; the resulting voltage worked out to be equivalent to a value the size of fifty-one voltage units.

The manufacturers of the real world might be able to create electronic parts with this kind of accuracy, but chances are that these parts wouldn't be cheap. A TRS-80 made with them would be worth more than the previously mentioned V-8 auto.

So in the stone age of computer activity, it was decided that the simplest level of evaluations would have to be made. A low voltage or 'off' condition would be made equivalent to zero; a high voltage or 'on' condition would be set equal to one. That was it. All calculations would have to be done with ones and zeros – the binary system.

You've probably heard of, and perhaps used, binary numbers and there will be more on this system later in the book. But the point is that you can well imagine that doing work with real, human numbers means quite a basketful of the individual ones and zeros. Which means, consequently, many separate signal lines to carry those groups of numbers.

So Open it Already!

Turn the power off to your TRS-80 keyboard unit and flip over on its front. Carefully undo the cabinet, noting the different sizes of screws used to fasten it together. Holding the entire computer firmly together, again flip it on its back. Remove the top cover. Gently lift the keyboard forward and remove the five or six plastic spacers underneath it; in later 80's, one of these will be solid and the rest flexible. Note their positions. Now set the keyboard back, lift the electronics out of the case and place it on a spacious work surface.

Time out. As you make changes to your TRS-80, add memory and the like, you will be opening the case many times. Two things might happen. First, the keyboard grommets will get lost. Their purpose is twofold: to prevent the keyboard from shorting against the main circuit card, and to cushion delicate parts against a constant onslaught of none too gentle typing. If you lose them, cut new ones out of bottle corks. They'll work just fine.

Another more difficult problem is the keyboard cable itself. This is a band of springy copper leads covered with white plastic located

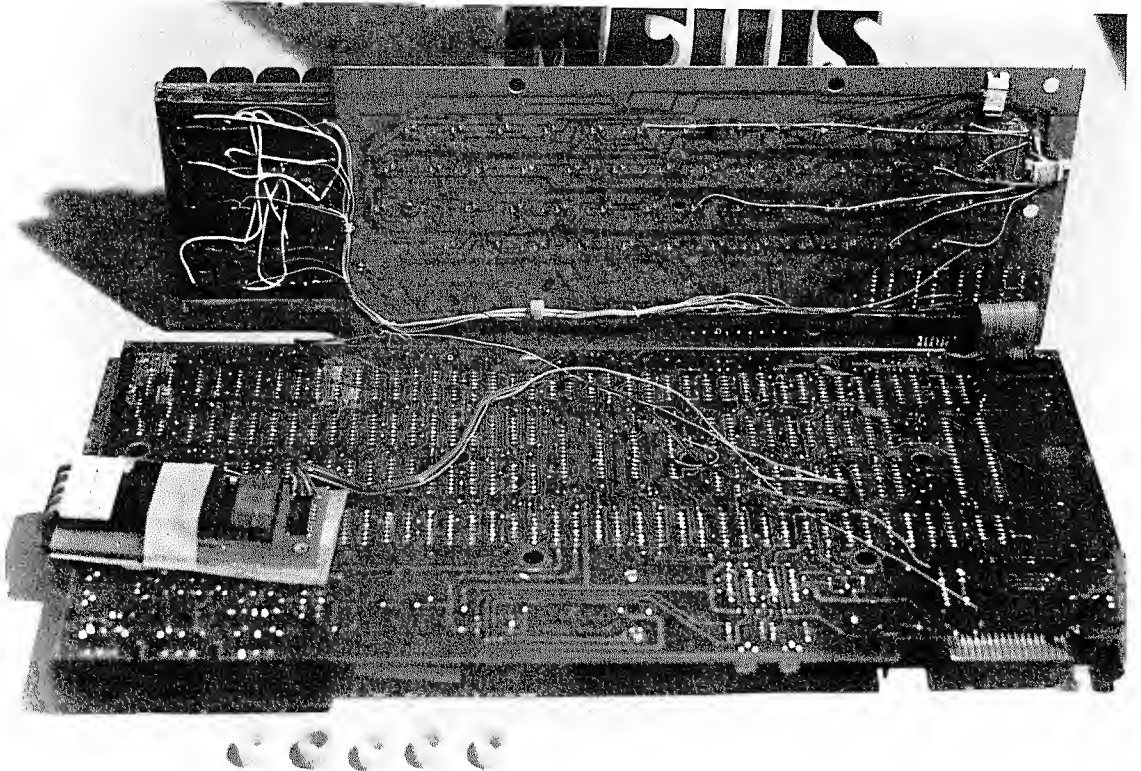


Photo 1-1. The TRS-80 opened and spread out.

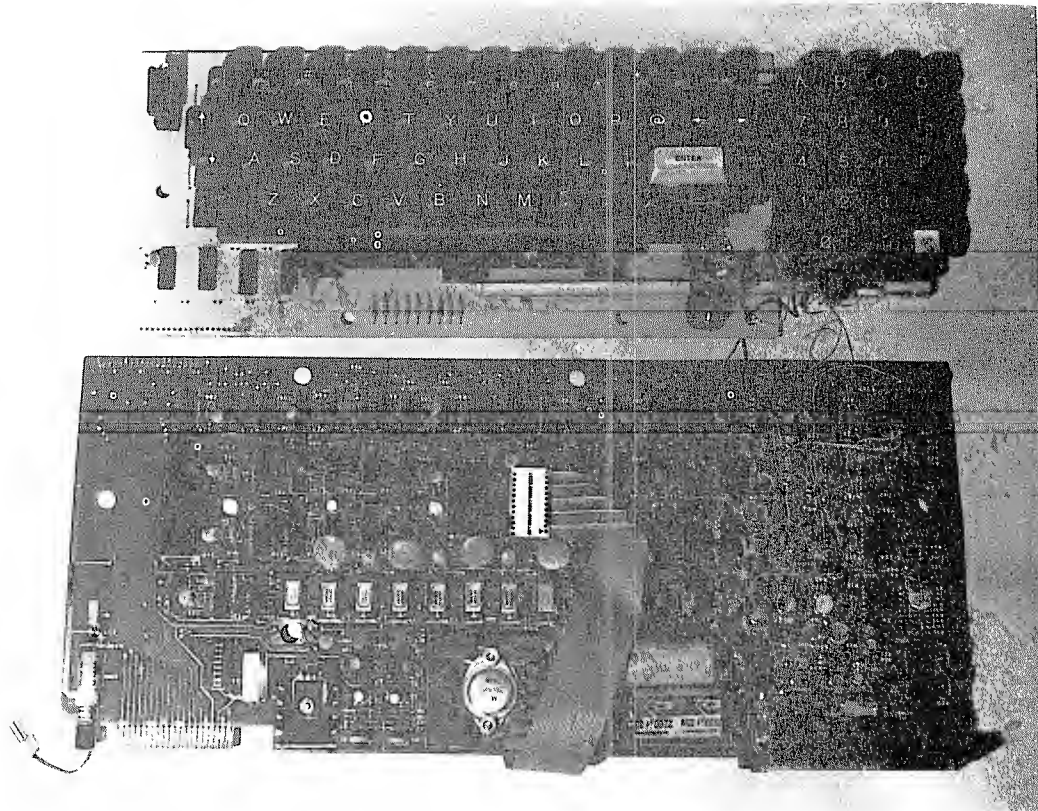


Photo 1-2. Closeup of keyboard cable area.

Complete TRS-80 spread out for cable replacement. Other visible changes include Level I/II switch on keyboard (bottom right of keyboard), new Level II interconnect cable (center), additional keyboard socket (top left), and external reset switch connector (bottom left).

at the bottom left of the keyboard. It electrically, and physically, attaches the keyboard to the main circuit board. Very subtle cracks can occur in the copper after about a half-dozen flexes. Symptoms can be lost letters or odd combinations of letters, a system constantly crashing or otherwise acting up, constantly repeating letters, or complete lockup of the keyboard. *Avoid flexing this cable quickly or bending it sharply.* Later we'll replace this cable with something better, but for the moment be gentle with it.

Now back to the machine. Spread the two sections connected by the keyboard cable out in front of you. The keyboard and the top of the circuit card will now be visible. The Central Processing Unit is a single integrated circuit, the Z-80; it is to the far left in the photograph on this page. The power supply is the block of 'heavy equipment' at the back; the two small potentiometers regulate the voltage to within five percent, so keep clear of these controls. The eight memory circuits are in a row near the power section, and the language memory (Level I and

newer Level II BASIC) is located in the center of the circuit card. On most Level II units, this language required three integrated circuits, and so these were placed on a separate two-by-three-inch card taped to the main board and connected with a 24-conductor cable. *Don't be tempted to remove that cable, yet!*

A group of circuits to the bottom left in the photograph hold the video image. Two important parts are socketed on the board; these are marked Z3 and Z71, and they are programmable shorting jumpers. What makes them programmable is the fact that with a gouging tool you can break the connections; not very subtle, but it works. Their purposes are different – one selects the amount of memory available in the keyboard unit, and the other selects which of the standard languages (Level I or Level II BASIC) is in use.

At the top left, the video and cassette output is controlled. Two small potentiometers on that side of the board position the video image on the screen, so if you received a TRS with its image off center, twiddling these will straighten it out. The

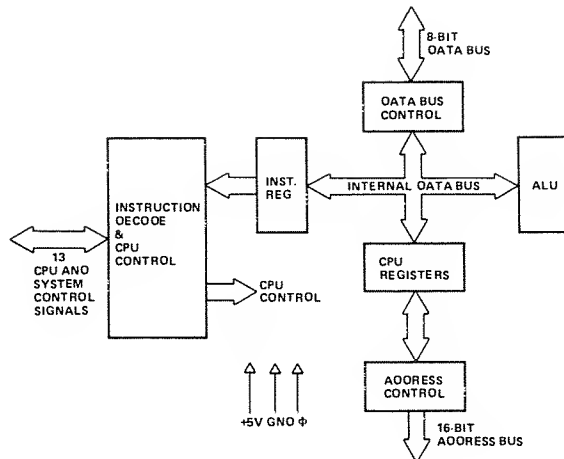


Figure 1-1. Z-80 block diagram from Zilog.

remainder of the board is dedicated to controlling the complex electronic traffic.

One last item is the edge card connector, which can be viewed at the top right hand side of the photograph. Through this connector, the TRS-80 may speak to the outside world. Much more on that connector later in this book.

On the keyboard itself there are a few circuits. Their job is to assist the Central Processing Unit (CPU) in determining which key is pressed. These circuits are sometimes remarkably sensitive; there will be more about these later.

Put it Back

Become familiar with the layout of this board, as it will assist you in locating work or repair areas. Functions of the computer parts are generally grouped together, so modifications will usually be restricted to a compact area of the board. Once you have a good idea of where everything is, put the keyboard unit back together carefully.

To really get to understand your machine you will want to obtain, either off the shelf at Radio Shack, from a Repair Center, or 'National Parts', the following manuals:

TRS-80 Micro Computer Technical Reference Handbook. Catalog No. 26-2103.

TRSDOS & Disk BASIC Reference Manual. Catalog No. 26-2104.

Printer Cable Service/Installation Manual (Order).

Expansion Interface Service Manual (Order). Make sure this includes the FD1771 disk controller appendix from Western Digital.

16K RAM Expansion Service Manual and Addenda (Order).

TTL Databook (National Semiconductor). Catalog No. 62-1370. This may not be stocked anymore; if so you can order it from National Semiconductor direct.

And finally, a *Z80-CPU / Z80A-CPU Technical Manual* should be ordered from Zilog, Inc., 10340 Bubb Road, Cupertino, CA 95014. It costs \$7.50.

These references will not only help you to make the modifications suggested in this book, but also to understand the operation of the computer, bend it to your wishes, and repair it if it fails. Additional references which you will find valuable are listed in Appendix II.

The Hidden Insides

It's not much of a mystery to TRS-80 users that all that hardware is controlled by software. That's one of the first things you learn. But it's also as simplistic as saying that the driver makes the car go, and just as misleading. Complete computers are called 'turnkey' systems because they imply simple, appliance-level setup and use. But a customized TRS-80 suggests something more, and with that 'something more' comes a requirement to have a better handle on hardware and software.

Let's define some terms. The TRS-80 is a personal computer, which is a popular way of saying a small, microprocessor-controlled, turnkey computer. The microprocessor, a Z-80, is a complicated, general-purpose, electronic switching center capable of accepting, changing, sending, and re-routing a complex array of electronic signals.

Both the TRS-80 as a whole, and the Z-80 in microcosm, have something called *architecture*. Architecture is the overall dimension by which these devices define their electronic space. Put simply, it is 'how they work'. As with all human work, this involves defining tasks and their order, executing those tasks, and producing a result.

In the Z-80, the architecture involves accepting electronic signal groups called *instructions*, decoding them into internal activities, and executing those activities. An arithmetic logic unit (ALU) performs simple mathematical functions, internal memory cells called *registers* hold signal information to be acted upon, and an internal *bus* controls the flow of electronic traffic. The order of entering signals is identified by means of an *address*, which

identifies a fixed numbered slot in the Z-80's electronic universe.

In the block diagram of this activity, shown opposite, note the terms '8-bit data bus' and '16-bit address bus'. The Z-80 is an integrated circuit with 40 external connections. The number 40 is arbitrary, chosen because manufacturing precision is currently limited to a physical 'package' of that size. That precision is also central to why the binary system is used, as mentioned earlier.

From the viewpoint of ten-fingered humans, it would be simpler to do our computing in familiar decimal form. As with the uninspired conversation between human and computer presented earlier, different levels of voltage could be used. But such levels of precision are difficult to produce commercially and impossible to diagnose when they fail.

That is the practical origin of the simple one/zero, on/off, true-false system of numbering used by the computer. To discover how these signals are arranged, we now turn back to the Z-80 itself and the 40 external pin connections. Information put into and called for out of the processor is called *data*, with a simple small value of zero. Any reasonable number of pins could have been dedicated to accepting data, but, based on the amount of work the processor had to do, eight of the 40 pins were assigned. This is the 8-bit (binary digit) data bus.

Likewise, there is the need to determine the order in which the processor is to do its work. In order to move from command to command, it identifies the order of those commands by their position, or *address*. Sixteen of the Z-80's pins are used to specify a number from 0000 0000 0000 0000 to 1111 1111 1111 1111.

When the power is turned on, this *address bus* switches to all zeros. The computer fetches its first command along the *data bus* from that all-zero location. It executes the instruction, fetches the next, and the next, and the next.

There's certainly more to this, but very basically that's the architecture of a microprocessor. Address lines for locations and order, and data lines for information. Now we'll turn to the TRS-80 and its architecture.

The TRS-80 contains the microprocessor, which is the Central Processing Unit (CPU) of the computer. Along its address bus is found the computer's internal information and instruction storage block – its *memory*. Also along this bus will be found video memory, a block which is reserved to display information to us on the screen; a keyboard, which is given its own set of addresses; the BASIC language; some unused areas; and single memory slots in which are housed windows to the cassette recorder, disk drives, printer and so forth.

This hardware is already familiar, so let's look at the hidden insides, the software. The BASIC system is found in permanent, Read-Only Memory (ROM). It consists of several major sections:

- A keyboard scanning routine to discover and interpret activities at the keyboard.
- A video processing routine which presents and updates the monitor information.
- Input/output controls for saving and loading program material and operating a printer.
- An interpreter capable of transforming the 'English' words which make up the BASIC we know, and determining what computational actions should be taken.
- Memory-management systems which apportion the computer's available memory into blocks which will not conflict.
- Arithmetic- and text-processing subroutines which can perform calculations and operations on numbers and alphanumeric characters.

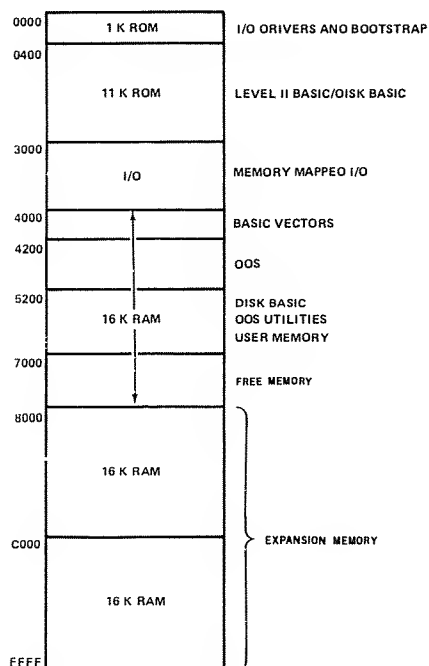


Figure 1-2. TRS-80 memory map.

Overall, the Level II BASIC language requires more than 12,000 separate 8-bit groups, known as *bytes*, to perform its work.

At the end of this Chapter is a detailed look at how the software operates from the time the machine is switched on to the time you read **MEMORY SIZE?** on the screen. In summary, that software disables any signals which might interrupt its operation, turns off the cassette relay and clears data from that output, restores the video screen to 64 characters per line, and sets up a block of memory for use by BASIC programs. A disk drive is searched for, and if one is found, a group of procedures are initiated in order that the disk program may take control of the TRS-80.

If no disk drive is found, the screen is cleared with blanks, the **MEMORY SIZE?** prompt is displayed, and a keyboard scanning process is begun. A valid response to that question is accepted, and, if necessary the entire bank of memory is tested. Error messages are updated on the video screen as needed. Finally, the memory size and available room for text strings is created, the **READY** prompt is displayed, and control of the TRS-80 is given to the user.

Power-Up Routines

The initialization routine of the TRS-80 is a complicated and very interesting aspect of the computer. It must, of course, set up all the parameters that will be used by BASIC programs, but it also conducts a series of tests and makes hardware adjustments to the device.

It double checks to assure the proper operation of memory, and to be certain that the parameters needed for proper operation of programs will be present. This section will take a look at the initialization process.

Here are the first few instructions:

0000		ORG	0000H
0000	F3	DI	
0001	AF	XOR	A
0002	C3 74 06	JP	0674H

At power-up, the Z-80 chip 'homes in' on address 0000, and begins its execution there. The first action is significant: DI (Disable Interrupt) keeps the clock 'heartbeat', generated by the expansion interface, from disturbing any actions of the computer – especially important, since the necessary software for handling that interrupt request is not in ROM, but rather a part of the disk BASIC, or what is also offered as 'Level III' BASIC.

So the interrupt is masked out. XOR A is the process of 'exclusive-ORing' the accumulator.

Exclusive OR is a logical operation which states: of two elements, either may be zero or one, but not both. Thus, whatever is present in the accumulator is XORed with itself. Since each bit is identical with its exclusive-OR partner, each bit will be set to zero. This effectively clears the accumulator, readying it for processes to come.

The final instruction of the group, JP 0674, gets out of the way of the Z-80's low memory, for it is in this area that the chip's restart codes – very frequently used subroutines – are found. Going on:

0674	03 FF	OUT	(OFFH),A
0676	21 02 06	LD	HL,0602H
0679	11 00 40	LD	DE,4000H
067C	01 36 00	LD	BC,0036H
067F	E0 80	LDIR	
0681	30	DEC	A
0682	30	DEC	A
0683	20 F1	JR	NZ,0676H

After the jump to 0674, the routine resets the output flip-flop at port FF (255 decimal). This flip-flop controls both cassette functions and the 32/64 character video, and by outputting the value in A (0, since it was exclusive-ORed earlier), the cassette will be off, no data will be present at its input, and video will come up normally.

Following this is an interesting (and encouraging) piece of code. Using the Z-80's powerful block move command LDIR, 54 bytes stored at address 06D2 are transferred to the RAM address area starting at 4000. These are the most important pieces of information the TRS-80 must have, so the writers of this program took great care to assure that this transfer would be certain. The LDIR instruction itself takes data stored at an address specified by register pair HL (in this case, 06D2), and moves a block whose length is specified by register pair BC (36 hex, or 54 decimal), to the location indicated by register pair DE (4000).

The interesting part is found just below. The value in A (0) is decremented twice (to FE), and the identical transfer instruction is repeated. This goes on until it reaches zero again, a total of 128 times! We may draw the conclusion that the Z-80 chip probably reaches full power and begins operating before memory gets to the point where it can reliably accept data . . . therefore, the instruction is repeated for a period of approximately 15 milliseconds.

Now a portion of RAM is cleared to zero with the following few commands:

06B5	06 27	LD	B,27H
06B7	12	LD	(DE),A
06B8	13	INC	DE
06B9	10 FC	DNZ	06B7H

Recall that after the previous setup process, the accumulator again contains zero. Here a block of RAM specified by the DE register (essentially where we left off) is loaded with zero. Using the fast DJNZ (decrement B, branch if not zero) instruction, 39 bytes are fixed at zero.

A few instructions follow that are very significant at power-up. Address 3840 contains the keyboard row where the BREAK key sits. It is connected to data line 4; thus the instruction AND 04 checks to see if it is held down. If it is not being held down, the result of the AND instruction will not be zero . . . and a jump to address 0075 will be made. This is why expansion interface owners without disks must press the break key at power-up:

0688	3A 40 38	LD	A, (3840H)
068E	E6 04	AND	04
0690	C2 75 00	JP	NZ, 0075H

How does the TRS-80 find out that a disk drive is in fact connected to the interface? The answer to that – and the reason that the computer ‘hangs up’ when an expansion interface is connected without a disk – is found in the next few bytes of code:

0693	31 70 40	LD	SP, 4070H
0696	3A EC 37	LD	A, (37ECH)
0699	3C	INC	A
069A	FE 02	CP	02
069C	0A 75 00	JP	C, 0075H

The stack pointer is set at 407D for use by potential future programs; it is out of the way of all the BASIC pointers set up in the first data transfer, an obvious but important action.

The accumulator is then loaded with the contents of memory location 37EC. There is no ‘memory’ per se at address 37EC; it is instead an instance of ‘memory mapping’. That is, when this memory cell is read, a signal is sent to the expansion interface. This signal strobes information from the Floppy Disk Controller (FDC) to the TRS-80. What will it find?

If no expansion interface is connected, there is no signal to strobe. Hence, the value will be floating, not pulled to ground (zero) on any bit. The computer sees all bits apparently ‘high’ at this location, and interprets it as binary 1111 1111, that is, hex FF.

The next instruction increments the accumulator, in this case resulting in FF plus 1, or 00. In the following instruction this value is compared to 2. A compare (in effect a subtraction, with no ‘result’) will cause the carry flag to be set, since 0 minus 2 is negative. (Note:

Why compare with 02? Why not just 01, as a carry would still be generated? My suspicion is that it is possible for those data lines to ‘float’ in the low state. In that case the CPU would ‘see’ 0000 0000, with the INC A instruction resulting in a value of 01 – which is still incorrect. So a compare with 02 guarantees the presence or absence of the disk controller.)

Once the carry flag has been set, the instruction JP C,0075 would be executed, sending the program to address 0075. For the moment, however, let’s assume that an expansion interface is connected to the TRS-80.

The FDC, when queried by the command LD A,(37ECH), will respond with 80. Incremented by one it becomes 81, and comparing it with 02 generates no carry. The JP C,0075H is thus ignored, and the program simply goes on to find:

069F	3E 01	LD	A, 1
06A1	32 E1 37	LD	(37E1H), A
06A4	21 EC 37	LD	HL, 37ECH
06A7	11 EF 37	LD	DE, 37EFH
06AA	36 03	LD	(HL), 3

The accumulator is set to 1, and address 37E1 is made to accept the contents of the accumulator. Again, 37E1 is a location memory mapped in the expansion interface. This code starts the disk drives rotating (or keeps them rotating), and selects drive zero.

That done, it loads HL with the disk controller address (37EC) and sends out a ‘restore’ command, which tells the controller to move to track zero. Thus, the data will come from drive zero and track zero. Register DE is prepared by loading it with the disk’s data address, 37EF. Now:

06AC	01 00 00	LD	BC, 0000
06AF	C0 60 00	CALL	0060H
0060		ORG	0060H
0060	08	DEC	BC
0061	78	LD	A, B
0062	B1	OR	C
0063	20 F8	JR	NZ, 0060H
0065	C9	RET	

This is a short but very useful subroutine. You may in fact want to call this yourself from time to time. Found at address 0060 is a simple delay loop; load the BC register pair (as is done just before the CALL instruction), and it is decremented and tested until it reaches zero. When it finally reaches zero, a return instruction sends the Z-80 back to the main program flow.

Why a delay? Merely to give the disk drive time to come up to speed, obvious but very

important. Moving ahead with this branch of the program:

0682		ORG	0682H
0682	C8 46	BIT	0,(HL)
0684	20 FC	JR	NZ,0682H

This is a loop which waits until the disk control chip says, "Okay, disk is up to speed and everything looks pretty good", and sends along a zero. The program loop tests this bit until it receives a zero. It is this loop which is maddening to you expansion interface owners who have no disk drive. Like all the previous memory-mapped addresses, 37EC will never have that zero. If a disk is present and all is well, the loop will have found the acknowledging zero sent by the FDC:

0686	AF	XOR	A
0687	32 EE 37	LD	(37EEH),A
068A	01 00 42	LD	8C,4200H
068D	3E 8C	LD	A,8CH
068F	77	LD	(HL),A
06C0	C8 4E	BIT	1,(HL)
06C2	28 FC	JR	NZ,06C0H
06C4	1A	LD	A,(0E)
06C5	02	LD	(8C),A
06C6	0C	INC	C
06C7	20 F7	JR	NZ,06C0H

The accumulator is cleared again, and the BC register is set to 4200. This will be an area of RAM set aside for disk use. 37EE is loaded with 0, and 37EC is loaded with byte 8C. This selects sector 0, and sets the read condition. Thus, having previously been set to drive zero, track zero, it will now read sector zero. The accumulator will look for the incoming bytes in the memory location specified by the DE register pair (37EF). This is the memory-mapped location through which the actual data will flow.

The accumulator picks up the data from DE, stores it in the RAM memory location indicated by BC (4200); the next instruction increments register C so that location 4201 is ready. The program loops back, waits for another message from the FDC, picks up another byte, and stores it. When register C finally reaches zero, the pointer will again contain 4200, and the loop terminates. Then:

06C9	C3 00 42	JP	4200H
------	----------	----	-------

Here the disk system takes over completely. As you recall, starting at 4200, data from the disk has been stored. By jumping to that location, the program direction is wrested from ROM and given to the first 256 bytes of the disk system bootstrap program.

Here, then, is a quick review of the activity so far: Interrupts are disabled, cassette is turned off

and data are cleared from that output, video is restored to normal, and significant pointers for BASIC program operation are set up. A disk drive is searched for and if one is found, a group of procedures is initiated in order to transfer control of the TRS-80 to these disk instructions.

A series of control signals and acknowledgments is exchanged between the floppy disk controller and the CPU, a page (256 bytes) of data is poured into a RAM buffer area, and program control is given over to this new series of commands.

If a disk drive is not found, or if the break key is held down during power-up, control is transferred to address 0075. At this point it should be noted that the 'reset' button on the TRS-80 is a non-maskable interrupt, that is, the only interrupt which the DI (Disable Interrupt) command first executed by the TRS-80 cannot mask out. When pressed, the reset button goes directly to address 0066, following a much shorter series of instructions reminiscent of the power-up routine.

Because it is likely most important RAM pointers are still intact, this sequence does not reset them:

0066		ORG	0066H
0066	31 00 06	LD	SP,0600H
0069	3A EC 37	LD	A,(37ECH)
006C	3C	INC	A
006D	FE 02	CP	02
006F	02 00 00	JP	NC,0000
0072	C2 CC 06	JP	06CCH

This group of instructions sets up the stack pointer, checks for the presence of a disk drive, and jumps to the complete initialization routine (reboot) if it finds one. If none is present, it goes to the READY sequence beginning at address 06CC.

Now let us return to the initialization program flow we have been following, which is found at 0075:

0075		ORG	0075H
0075	11 80 40	LD	0E,4080H
0078	21 F7 18	LD	HL,18F7H
007B	01 27 00	LD	8C,0027H
007E	E0 80	LDIR	

Using the LDIR instruction described earlier, a block of information located at 18F7 is transferred to RAM beginning at 4080. These bytes describe ports in use, error storage, INKEY\$ information, and so forth, as needed in the general operation of Level II BASIC (see Chapter 2).

A few specific addresses are delineated, and a large group of RAM bytes is then prepared. These jump to the familiar 'L3 ERROR' message because they are disk commands not available to Level II, yet patch points are prepared for them. The result of the following program statements is to fill addresses 4152 to 41A5 with the direction JP 012D.

008E	11 20 01	LD	DE,0120H
0091	06 1C	LD	B,1CH
0093	21 52 41	LD	HL,4152H
0096	36 C3	LD	[HL],0C3H
0098	23	INC	HL
0099	73	LD	[HL],E
009A	23	INC	HL
009B	72	LD	[HL],D
009C	23	INC	HL
009D	10 F7	DJNZ	0096H

Another group of ROM 'breakout' points follows; they all become returns to the main program flow. But notice something interesting about them – three bytes are set aside, but only one is filled with the return instruction (C9). This means that a jump command could be placed there. Let's first look at the series of instructions, then examine the possible benefits of changing them:

009F	06 15	LD	B,15H
00A1	36 C9	LD	[HL],0C9H
00A3	23	INC	HL
00A4	23	INC	HL
00A5	23	INC	HL
00A6	10 F9	DJNZ	00A1H

If we wanted to break into the BASIC operating system, this area of RAM is one place in which we could do it. Most of these are error codes of one kind or another. We could 'rescue' a program from displaying an error message, and halting, by patching in one of our own routines. If our routine were located at 5000, for example, the C9 instruction (followed by two unused bytes) could be replaced with a JP 5000 command, which needs all three byte positions: C3 00 50. Essentially, the authors of Level II BASIC provided many areas for expansion.

Now let's move on. BASIC programs begin at address 42E9. A pointer to that beginning is found as a zero at address 42E8. The next instruction sets that in place:

00AB	21 E8 42	LD	HL,42E8H
00AB	70	LD	[HL],B

The stack pointer is delineated, and a call is made to 1B8F, a subroutine to turn off or reset various devices, including the printer and cassette player. It is in part redundant, but a double-check is often worthwhile.

00AC	31 FB 41	LD	SP,41FBH
00AF	C0 8F 1B	CALL	1B8FH
00B2	C0 C9 01	CALL	01C9H

The call to 01C9 results in the screen being cleared and the cursor being placed at position 0. Finally, 'MEMORY SIZE?' appears:

00B5	21 05 01	LD	HL,0105H
00B8	C0 A7 2B	CALL	2BA7H
00BB	C0 B3 1B	CALL	1BB3H

At address 0105 is a block of ASCII bytes which spell out MEMORY SIZE. The subroutine starting at 2BA7 displays the string of data at the present location of the cursor, a byte at a time, until it finds a byte in the message whose value is 00. This terminates the display and advances the cursor. The call to 1BB3 is identical to the BASIC INPUT command, in that it displays the question mark and cursor, and halts for keyboard input.

If the keyboard input is the BREAK key, a carry is generated, and the program skips back to MEMORY SIZE and displays it again, waiting for keyboard input. The instruction RST 10 (ReStart at 0010) follows, which is the quickest way of calling a routine to locate the first character of an input. If one is found, the result of an OR instruction will not be zero. Here are the instructions that perform those functions:

00BE	3B F5	JR	C,00B5H
00C0	07	RST	10H
00C1	87	OR	A
00C2	20 12	JR	NZ,0006H

What if, on the other hand, there was no entry other than the ENTER key? You have no doubt noticed a slight pause in the action when you do not specifically set the memory size. Here's a look at that code:

00C4	21 4C 43	LD	HL,434CH
00C7	23	INC	HL
00C8	7C	LD	A,H
00C9	85	OR	L
00CA	2B 1B	JR	Z,00E7H
00CC	7E	LD	A,[HL]
00CD	47	LD	B,A
00CE	2F	GPL	
00CF	77	LD	[HL],A
00D0	8E	CP	[HL]
00D1	70	LD	[HL],B
00D2	2B F3	JR	Z,00C7H
00D4	1B 11	JR	00E7H

For the moment we will start at the instruction LD A,(HL). HL contains the address of a byte of RAM memory, the contents of which are placed in the accumulator. From the accumulator, they are also saved in the B register. The accumulator is complemented, which inverts all the ones to zeros and all the zeros to ones. This complemented value is then placed in the

memory location still specified by HL. The accumulator is compared with what has been placed in HL.

What, you ask? But this value was just placed in memory, why compare it? Because – and this is a very elegant piece of writing – if it does not compare:

1. The memory location is bad and only the block of memory below it should be used to be safe.
2. Or, this is the end of memory.

If this is good memory, then, the test for zero passes, the contents saved in register B are returned to memory, and the program loops back, incrementing HL to the next potential memory location.

We did skip a few instructions back there. They become important only after the first loop is complete. These commands OR the contents of H and L; when the result is zero, we are at address 0000 – full 48k memory has been found, and the test is complete.

Here's what we would find, alternatively, if we entered some value (or other characters) in response to MEMORY SIZE?:

0006	C0 5A 1E	CALL	1E5AH
0009	B7	OR	A
000A	C2 97 19	JP	NZ,1997H
000D	EB	EX	DE,HL
000E	2B	DEC	HL
000F	3E BF	LO	A,BFH
00E1	46	LO	B,(HL)
00E2	77	LO	(HL),A
00E3	8E	CP	(HL)
00E4	70	LO	(HL),B
00E5	20 CE	JR	NZ,00B5H

The call to 1E5A checks for numeric input, and jumps to 1997 ('?SN ERROR'), if it is not numeric. If the input is properly numeric, then registers DE and HL are exchanged; this action puts DE (left off at the lowest usable memory location above pre-set RAM needed by BASIC) in HL, where it can be manipulated conveniently.

Memory size minus one is usable; memory size and above is protected. So HL is decremented before being tested, then it is tested (in a manner similar, but not identical, to that done earlier). If the memory test fails, it's back to displaying MEMORY SIZE? again.

We're not quite there yet, however, as the figure entered for memory size may be too small. BASIC needs a bit of room to work with, so DE is set to 4414, and the subtraction subroutine at RST 18 is called. If a carry is generated, we're shipped off to the '?OM ERROR' message found at 197A. Here's what it all looks like:

00E7	2B	DEC	HL
00EB	11 14 44	LO	DE,4414H
00EB	0F	RST	18H
00EC	0A 7A 19	JP	C,197AH

A little more work is left to do. Recall that a value for available string space is set aside, and it is 50 bytes. Here is how it is done:

00EF	11 CE FF	LO	DE,0FFCEH
00F2	22 B1 40	LO	(40B1H),HL
00F5	19	ADD	HL,DE
00F6	22 A0 40	LO	(40A0H),HL

Register pair DE is set up with FFCE, which, if you are not yet weary of manipulation of hex numbers, is the two's complement of 50 decimal. That is, when FFCE is added to 0000, the result is FFCE hex, or 50 decimal less than the original figure. Try it to see that it works. This bit of code saves the value for top of available memory in 40B1, adds register DE to it, and saves that result (memory size minus 50 bytes for string space) in address 40A0.

There follows:

00F9	C0 40 1B	CALL	1B40H
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Here let me quote Roger Fuller, whose TRS-80 Supermap identifies this subroutine this way: Revelation 21:5 – "And behold . . . He shall make all things new."

This subroutine identifies and sets up all pointers necessary for the start of a BASIC program: Variables reset, previous program deleted, etc.

And now, the moment you've all been waiting for. Here it is:

00FC	21 11 01	LO	HL,0111H
00FF	C0 A7 28	CALL	28A7H
0102	C3 19 1A	JP	1A19H

The call to 28A7, you may recall, displays a string of ASCII characters. The string displayed in this case is . . .

RADIO SHACK LEVEL II BASIC

The final instruction is a jump to 1A19, the address of the 'READY' display.

To summarize this last portion of the initialization routine: all BASIC pointers, disk error codes, and ROM return codes are set up, the screen is cleared, and the MEMORY SIZE prompt is displayed. A valid response to that question is accepted, and, if necessary, the entire bank of memory is tested. Error messages are generated as needed. Finally, the memory size and available room for strings is recorded, the READY prompt is displayed, and control of the TRS-80 is given to the user.

NOTES

2

Copacetic Comprehension

There will doubtless be a day when books like this will be unnecessary. Personal computers will probably develop into the appliance area, with programmers, hobbyists, hardware designers and language specialists present only in the distant background of the market. But between now and then we are all faced with being either frustrated users or solderer-programmers, tailoring machines according to our personal demands.

To do this, certain skills are inevitably required. Among these are an understanding of non-decimal number systems, digital logic devices, machine-level languages, and a smattering of diagnostic sense. There are some fine books that cover all these topics (see Appendix II), so this chapter will only deal with them as far as needed to put this book to work. Among them are:

- Binary, decimal and hexadecimal number systems, how they arose, how and why they can be used, and where understanding them is essential.
- Common digital logic devices that appear in the TRS-80 and these projects, and how and where to use them.
- Some of the basic elements of machine language, and a few personal considerations on where it is best applied, and when BASIC is a better choice.

- A look inside the TRS-80, with an eye to diagnosing where troubles might lie and where changes might be in order.
- The basics of creating a workable power supply for the projects in this book.

Number Systems

Numbering is the single most overrated problem in computer programming. The answer (posed before the question) is this: numbers are merely *counting names*. That is, it makes no difference whether we think in tenths of a mile or eighths of an inch. Nor does it bother us that a day is made up of 24 hours, while an hour is 60 minutes. That a year is 365 days frightens us not, nor that months are a motley collection of sizes.

In parking lots, does it bother us that our vehicle may be parked in Row N as opposed to Row 14? There is no mystery when we mark off points with four scratches and a crosshatch. And does a dozen always conjure up 'twelve', or is a dozen something we have understood since youth?

Names are sizes are numbers; so it is with the number systems that we arbitrarily assign for the convenience of working with computers. When we are talking about electrical signals, it is clearest and easiest to think about ons and offs. Ons look pretty much like ones, and offs look like

zeros. It's a nice, clean concept, and one that illuminates the way we can refer to the machinery.

There's more convenience to naming a computer data condition 10110100 than to calling it an on off on on off on off off. Were data the only consideration, the binary one and zero method might have been satisfactory, without resorting to other means or stroking our memories.

Finding a location in a computer's memory is a much more difficult task. Although a memory location called . . .

11101000100110101

. . . might be easier to think about than . . .

onononoffonoffoffonoffonoffononoffonoffon

. . . it could use another step forward. In music, a long string of sixteenth notes like this –

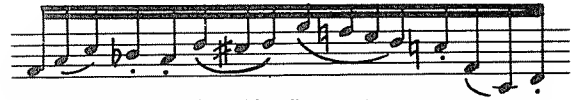


Figure 2-1. Illustration of illegible musical notation.

– is broken up to make it legible, so it looks instead like this –



Figure 2-2. Illustration of legible musical notation.

Likewise, that long binary string can be broken up from 11101000100110101 into convenient groups . . .

Converting Binary to Decimal

In the grade school years, students used to learn that a number like 5,163 contained a 3 in the ones place, a 6 in the tens place, a 1 in the hundreds place, and a 5 in the thousands place. It was to remind them that 5,163 was really 3 plus 60 (6 x 10) plus 100 (1 x 10 x 10) plus 5,000 (5 x 10 x 10 x 10).

The way other number systems are written follows this same pattern for their own bases. In base eight the number 5,163 would have a 3 in the ones place, a 6 in the *eights* place, a 1 in the *sixty-fours* place, and a 5 in the *five-hundred-twelves* place. That means that 5,163 is really 3 plus 48 (6 x 8) plus 64 (1 x 8 x 8) plus 2,560 (5 x 8 x 8 x 8). But notice how that's decimal thinking! Really in base eight there could be no '8'. . . it would have to be called '10'! 1, 2, 3, 4, 5, 6, 7, 10, 11, 12, 13, 14, 15, 16, 17, 20, and so on. So 5,163 in base eight is *still* 3 plus 60 plus 100 plus 5,000!

The binary system sneaks in the same way. A number like 1101 0001 0001 0011 turns into a 1 in the ones place, a 1 in the twos place, a 0 in the fours place, a 0 in the eights place, all the way up to a 1 in the thirty-two-thousand-seven-hundred-sixty-sevens place. In binary, the one on the far left is still a 1 in the quadrillions place, don't forget. But the message is how to convert all this to decimal. And here it is:

32768 16384 8192 4096 2048 1024 512 256 128 64 32 16 8 4 2 1
1 0 1 0 0 0 1 0 1 0 0 1 0 0 1 1

Just add the numbers: $1 \times 1 + 1 \times 2 + 0 \times 4 + 0 \times 8 + 1 \times 16 + \dots + 1 \text{ times } 32,768 = 41,619$. Voila. No matter how long the number is, and in whatever base:

1. Start at the left and produce a chart of the base number's powers, starting with 0 (X to the 0 power is always 1).
2. Lay the number to be converted underneath the base number chart.
3. Multiply each base number power by the digit in its place.
4. Sum the resulting numbers.

Does it work? Certainly. What is 163,341 in base 9? And in base 7? And in base 10?

Base 9 powers:	5	4	3	2	1	0
9 to that power:	59049	6561	729	81	9	1
Number to convert:	1	6	3	3	4	1
Multiplication:	1x59049	6x6561	3x729	3x81	4x9	1x1
Subtotals:	59049	+39366	+2187	+243	+36	+1
Converted result:	100882, base 10					

Base 7 powers:	5	4	3	2	1	0
7 to that power:	16807	2401	343	49	7	1
Number to convert:	1	6	3	3	4	1
Multiplication:	1x16807	6x2401	3x343	3x49	4x7	1x1
Subtotals:	16807	+14406	+1029	+147	+28	+1
Converted result:	32418, base 10					

Base 10 powers:	5	4	3	2	1	0
10 to that power:	100000	10000	1000	100	10	1
Number to convert:	1	6	3	3	4	1
Multiplication:	1x100000	6x10000	3x1000	3x100	4x10	1x1
Subtotals:	100000	+60000	+3000	+300	+40	+1
Converted result:	163341, base 10					

1101 0001 0011 0101

. . . although the legibility is improved, the human spark, the ability to look and recognize (that *aha!*) is not there. So the next step is to set about naming the sections. Since these on-off conditions can be written down as binary numbers, why not write them down in their decimal equivalents?

The question is rhetorical, of course, because not only can it be done, it is done. The only question is how to do it. Were a computer capable of swallowing all sixteen of those binary digits (bits) in one gulp, that question might be easily answered by calculating the conversion of 1101 0001 0011 0101 using a binary-to-decimal conversion table. The result, we find, is 53557.

But the computer, alas, cannot swallow all those bits in one bite . . . it can only swallow one byte full of bits (pardon). In other words, though a computer may need numbers sixteen bits long, only eight data lines exist to carry that data.

The component parts of the number 1101000100110101 are needed, eight bits at a time: 11010001 00110101.

There's the mathematical rub. 11010001 is 209 decimal, and 00110101 is 54 decimal. This seems hardly related to 53,557. Another solution is necessary, and it is a naming system as much as a

numbering system. It names each of the sixteen possible combinations of four binary digits:

Reading the Pins

Finding your way through digital circuits is much easier than finding your way through an ordinary table radio. Industry standards have made the process simple. Consumer integrated circuits are packaged in small, rectangular, plastic or ceramic cases with anywhere from 8 to 40 external connections known as pins.

Earlier integrated circuits – and many of the audio types currently being produced – were packaged in small metal cans and looked like transistors, with many wires protruding from the bottom. The wires were arranged around a keying tab on the edge of the can, and numbered like so:

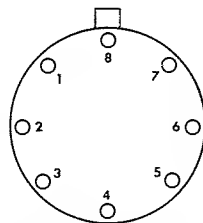


Figure 2-3. Can-type IC pin numbering.

As such circuits developed into more sophisticated and powerful devices, more pins were needed for input and output. A rectangular package was developed, but it was still numbered in a circle, starting (when looking down from the top) from the left of the notch, so:

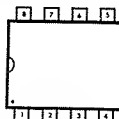


Figure 2-4. Dip-type IC pin number (8 pins).

All modern integrated circuits can be read from the top in this same way. 14- and 16-pin types start from the top left and read around:

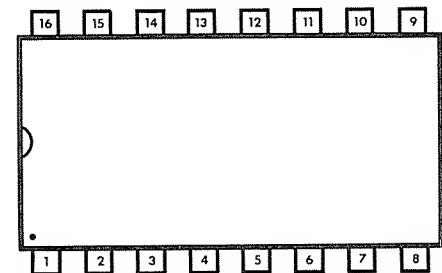
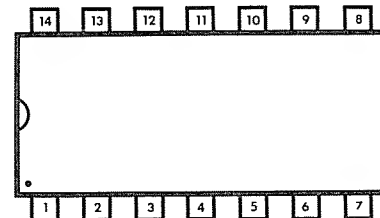


Figure 2-5. 14- and 16-pin Dip IC pin numbering.

You can read the pinouts of 18-, 20-, 24-, 28- and 40-pin circuits in the same manner. The highest numbered pin sits just opposite the lowest numbered pin. In the beginning this practice may seem confusing; it is. But after using the circuits – and counting their pins over and again – you will probably feel comfortable with the pin arrangement.

Just one thing: when you assemble TRS-80 add-ons, most of your work will be done from the bottom . . . which means reading backwards!

0000	is named	0	and is equal to decimal	0
0001	is named	1	and is equal to decimal	1
0010	is named	2	and is equal to decimal	2
0011	is named	3	and is equal to decimal	3
0100	is named	4	and is equal to decimal	4
0101	is named	5	and is equal to decimal	5
0110	is named	6	and is equal to decimal	6
0111	is named	7	and is equal to decimal	7
1000	is named	8	and is equal to decimal	8
1001	is named	9	and is equal to decimal	9
1010	is named	A	and is equal to decimal	10
1011	is named	B	and is equal to decimal	11
1100	is named	C	and is equal to decimal	12
1101	is named	D	and is equal to decimal	13
1110	is named	E	and is equal to decimal	14
1111	is named	F	and is equal to decimal	15

This may seem overdone; but A, B, C, D, E, and F are darn good names for binary values which exceed the number nine. If you don't have a name, make one up. For practical purposes, keep it within the symbols everyone has on the Royal typewriter.

Back to the number 1101000100110101. Crack it into those four legible pieces (1101 0001 0011 0101), and it can be named D135. To convert it to decimal, remember the old rule: the 5 is in the ones place, the 3 is this time in the sixteens place, the 1 is in the two-hundred-fifty-sixes place, and the D is in the four-thousand-ninety-sixes place. Thus, D135 is 5 plus 3 x 16 plus 1 x 256 plus (see the chart) 13 x 4,096, or . . . 53,557!

So, that long binary number can actually be digested by the computer as a byte of D1 and a byte of 35. After a while, the number system comes easily. My personal recommendation: work in it. Convert to decimal only when you absolutely must. Think in hexadecimal and binary. They are the tools with which you can speak to the computer.

Throughout this book, numbers in hexadecimal are printed in **BOLD**.

Digital Logic Devices

The binary number system and digital logic devices were developed together as a way of solving a practical dilemma: how to mass produce computers which could work quickly and accurately, and yet be inexpensive. As noted in Chapter 1, the problems of creating consistently accurate circuits, working with many different voltages levels, are formidable. Thus, simple yes-no, on-off logic was developed.

The intimidating term *Boolean algebra* is being used for the first, and last, time in this book – right in this sentence. You'll probably hear the phrase from time to time, but no matter – it's a professional's buzzword to keep the masses out. Forget it.

Back to digital logic devices. The essence of digital logic is to evaluate binary, on-off input; sometimes to determine a pattern of similarity or difference, sometimes to sense a change and sometimes to search for a signal. An appropriate result is produced as a result of the logical operation.

One of the logic building blocks is called a gate. A gate electronically evaluates its input to determine the pattern of similarity and difference of signals, and produces a specific output. A simple gate is shown below:



Figure 2-6. Simple AND gate.

Its job is to determine if the first AND second inputs are both at the one (high) level. Only under that condition will its output produce a high (one) signal. The table below shows how this AND gate works.

AND Gate		
If input #1 is	If input #2 is	The output result is
0	0	0
1	0	0
0	1	0
1	1	1

Table 2-1. AND gate action.

The table is called a *truth table*, and its purpose is to present every possible input and output condition for a given gate. Below is an OR gate. Stated in words, if either the first OR the second input is high, the output will be high. Examine the OR gate truth table; it really is quite logical.



Figure 2-7. Simple OR gate.

OR Gate		
Input 1	Input 2	Output
0	0	0
1	0	1
0	1	1
1	1	1

Table 2-2. OR gate action.

Given a huge set of interconnected gates and their known inputs, the final output of the group can be determined by using truth tables like

these. Gates may have more inputs than two (some have sixteen), and may produce the opposite results from the two described above (NOT-AND and NOT-OR gates, known as NAND and NOR gates). Truth tables reveal how the integrated circuit's design engineer specified the pattern of binary logic inside the circuit.

In this way, given a desired output and a known number of input signals, it is possible to determine what set of input values will trigger the desired output.

There are a number of other types of digital circuits. Most are created out of gates like those described above, but their features are unique enough to think about them separately. Among these other digital logic circuits are buffers, flip-flops, counters, latches, multiplexers and shift registers.

A buffer can be thought of as a two-input gate with both inputs tied together, like this:

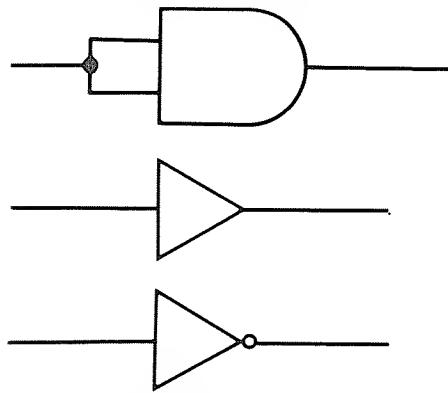


Figure 2-8. Buffer as (a) two-input gate and (b) buffer.

Its truth table is much simpler than that for two-input gates, because there are now only two input conditions. Either both inputs are high, or both inputs are low. Gates with 'true' outputs (AND, OR) will merely follow the input condition. When the inputs are high, the output is high; if the inputs go low, the output becomes low. Separate logic devices are manufactured that perform this 'follow-the-leader' function, and they are called *buffers*. They serve to isolate sections of a circuit, or rejuvenate a signal so it can feed many dozens of inputs in a large machine.

When a buffer reverses the condition of its input, (a high input is output low, and vice versa), the device is called an *inverter*. This kind of circuit can save the day in some cases, as when trying to locate a given binary number. Assume a circuit needs the binary number 1110 to turn on a pilot light. It is possible to choose four separate gates, each of which would provide an output

matching the desired number. These would be connected through more gates, and eventually the number could be discovered when the final signal was triggered properly. One way of detecting 1110 is shown below:

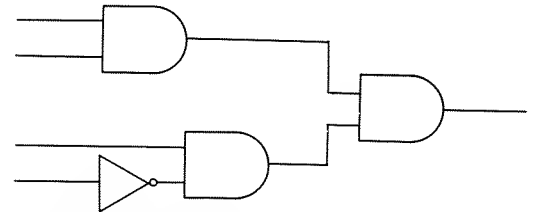


Figure 2-9. Bad decoding scheme for 1110.

But, although this circuit works, economy of cost and space and simple clarity dictate another solution. The last input could be inverted before it is evaluated, resulting in a pattern (1111) which could be quickly recognized by a multiple-input gate. The result is electronic simplicity and legibility; an improved decoding circuit is shown below. The ultimate result is the same.

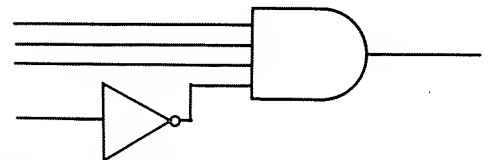


Figure 2-10. Good decoding scheme for 1110.

A flip-flop is a 'black box' which provides two outputs. When an input value is high (one), the first output will be high, and the second will be low. When the input value switches low (zero), the outputs will reverse. In other words, two opposite outputs for the price of one. But there is another significant use of the flip-flop.

Flip-flops also have an important input called a clock trigger, which is triggered only when its input returns to a given level. Only then will the outputs of the flip-flop reverse. That is, a given flip-flop clock may receive a 'zero' pulse. Its outputs will reverse. Then the zero pulse changes to a 'one' pulse. Nothing happens, but the trap is set to spring. When the one pulse changes back to a zero, the outputs reverse again. For every two changes at the clock, there will be but one change at the output. It takes four clock changes to produce two output changes.

Why is this useful? Because it is electronic, binary division. The truth table here shows how it works.

Binary Division with a Flip-Flop
Output of First Flip-Flop Connected to Clock of Second
Flip-Flops Change State Each Time Input Returns Low

Clock Input	Flip-Flop Output	Second Clock Input	Second Flip-Flop Output
0	0	0	0
1	0	0	0
0	1	1	0
1	1	1	0
0	0	0	1
1	0	0	1
0	1	1	1
1	1	1	1

{Input} {Input/2}

{Input/4}

Table 2-3. Binary division with a flip-flop.

Digital logic devices known as *counters* are combinations of gates and flip-flops that allow certain patterns to be counted: Binary, Binary Coded Decimal (BCD, where the highest number is decimal 10), Gray code and others.

Latches are very much like flip-flops, except that the input is 'captured' at the output by a trigger signal called an enable, a select, or a gating pulse. The input may change continuously, but the output only reflects the input when the enable is activated. Latches are very useful when hundreds of thousands of signals are flying around on one set of lines, and the computer must select only certain groups of signals. The cassette output of data is a latch; only the 500-baud (bits per second) pulses of data reach the cassette output, even though many different signals reach its input.

Multiplexers are sometimes misunderstood, but mostly because of their formidable name. A traffic light is a multiplexer – it allows several streams of traffic to meet at one intersection, but only one stream to proceed. The multiplexer is the electronic equivalent, having several inputs.

Gating signals select which of the inputs may reach the output. In a computer, this allows several devices to share a circuit (like the video, which must be sent to the screen, but also sends and receives characters from the rest of the computer).

Finally, *shift registers* treat bits of data like a bucket brigade sends up water: it goes in one end, and at each electronic 'go!', the bucket is sent along one position. The dots which make up the video display are produced by circuits which shift them out to the screen one bit at a time, in synchronization with the monitor's sweeping electron beam.

Into Machine Language

If we put faith in etymologists, then 'language' – which comes from the Latin meaning whole or part of a tongue – is an inappropriate word to put after 'machine'. Better to call it code, or blurms, or bingo chips even. Whatever it is, it becomes a tool for providing the user with all the power inside the TRS-80.

With a knowledge of binary and hexadecimal numbers, this language seems more fluid, and with a similar understanding of its electronic effects, it becomes the true 'lingua franca' of the computer.

In Chapter 1, I pointed out that BASIC is really just a disguised form of machine language. It is disguised because it presents itself in English-looking words, and has a large store of safety valves, error traps and messages – to prevent it from falling down an electronic rabbit hole.

For the moment I'll turn away from the metaphors, and present a practical simulation of that statement. The simulation can be familiarly written as follows:

CLS

When you command CLS (clear screen), BASIC enters a machine language subroutine to clear the screen. It automatically returns to a BASIC command-level 'READY' condition. Now I happen to know that CLS is located at memory address 01C9 (457 decimal), and that (non-disk system) 'READY' is found at 06CC (1740 decimal).

As an experiment, type . . .

SYSTEM (ENTER)
/1740 (ENTER)

. . . and you will be presented with the familiar 'READY', exactly as if you had pressed the Reset button. You've executed a machine-language GOTO, jumping directly to the 'READY' routine in ROM. As a second experiment, type . . .

SYSTEM (ENTER)
/1457 (ENTER)

Momentarily, the screen clears, but quickly crashes to MEMORY SIZE? Why did that happen? Notice that I said CLS was a *subroutine*. In other words, the routine must be called via some sort of GOSUB. The crash back to 'MEMORY SIZE?' occurred because there are no error messages in machine language unless you create them!

There are two ways to simulate the required GOSUB. The first is to POKE the starting

address of the clear-screen subroutine into the USR(0) command. The USR(0) command is identical to the SYSTEM command, except that it is the equivalent of GOSUB, where SYSTEM is the equivalent of GOTO.

So 01C9 (clear screen address) is the USR entry point; it must be split in two pieces (01 and C9), converted to decimal (1 and 201), and POKEd into USR addresses 16,527 and 16,526. So . . .

```
POKE 16527,1
POKE 16526,201
M=USR(0)
```

There. A screen clear without crashing to 'MEMORY SIZE?'. But there is another way. The machine language command for GOTO is called 'JUMP', and its hexadecimal code is C3. The machine language command for GOSUB is named 'CALL', and its code is CD. So here is the solution: CALL 01C9. JUMP 06CC.

Various hardware has its peculiarities. In the Z-80 microprocessor, addresses are always specified in reverse order. So CALL 01C3 is written CD (CALL) C9 (least significant byte of address) 01 (most significant byte of address). And JUMP 06CC is written C3 CC 06.

The whole process to clear the screen and jump to 'READY' is coded:

```
CD C9 01 C3 CC 06
```

Quite simple. CLS equals GOSUB 01C9, GOTO 06CC, which equals CALL 01C9, JUMP 06CC, which equals CD C9 01 C3 CC 06. The nice part of this little process is that it can be put anywhere in memory you like. Let's put it arbitrarily at 5000 to 5005 (20480 to 20485 decimal).

```
POKE 20480,205 : POKE 20481,201 :
POKE 20482,1
```

```
POKE 20483,195 : POKE 20484,204 :
POKE 20485,6
```

Convert the six hexadecimal bytes to decimal as shown earlier, and POKE them in place. Now you have the program completely under control. It is ensconced in memory, it calls the clear-screen routine, and jumps to 'READY'. Do you believe it? Just try it, entering the program at 5000. . .

```
SYSTEM (ENTER)
/20480 (ENTER)
```

There. A complete machine language program that functions from BASIC command level.

If you're new to this, take a break. The next step is to write an elementary screen-clearing routine, instead of calling one already in ROM.

I want to introduce a dilemma early on in the process of learning machine language. There is always discussion in computer programming about 'reinventing the wheel', and there is much truth to the suggestion that one should not do programming when the work has already been done. In the case of the screen-clearing program, the subroutine to do the work has already been created, and it is right there in Level II ROM waiting to be used.

Why, then, write another one? Why, in fact, even call the routine via machine language when the BASIC CLS command works so well? Indeed, with a BASIC as powerful as Level II, machine language should probably be reserved for doing things that cannot be done in BASIC at all. Among these things are upper/lower case drivers, programs to send characters to a serial printer, music-making and sound effects, telecommunications and so on. Furthermore, where machine language seems required, it often makes sense to call as many Level II ROM subroutines as possible.

Slick as this may be, it has two disadvantages: much of programming and customizing the TRS-80 requires an intense element of learning and understanding. Re-inventing the wheel is what everyone who learns must do, from the child who is forced through the memorization of 'times-tables', to the adult who has lost a job after 25 years and must learn new skills.

I can only present this from a highly personal point of view, as one who could not have learned machine language with any fluidity had I depended upon the software black boxes of others. I suggest that if you want a program to perform a certain action, try to write it in machine language. Try to make it do the same sort of error-trapping and other housekeeping that Level II's subroutines do. Take the Level II code apart and have a look. But don't deny yourself the opportunity to learn, rather than run a personal software assembly line. End of sermon. Back to clearing the screen.

The Z-80 microprocessor has internal holding latches called registers. Some are capable of holding eight bits, others sixteen. Some eight-bit registers may be paired up with others to create a single sixteen-bit register. These might be thought of as your only Z-80 machine code variables. In this screen-clearing routine, several registers will be used. The registers to be used in this experiment are: B, paired with C; H, paired with L; and A.

What's In the Memory Map

The last of these is the *accumulator*, a sophisticated register capable of doing simple arithmetic. Indirectly, the 'condition code' register (called the 'flags') will also be needed, and will be described when it is used.

First, the program in its entirety:

```
21 00 3C      START——LO      HL,3C00H
01 00 04      LO              BC,0400H
36 20          LOOP——LO      (HL),20H
23           INC              HL
00           DEC              BC
..          LD                A,B
..          OR                C
C3 XX XX      JP              NZ,LOOP
```

Listing 2-1. Simple CLS demonstration.

In the first line, the H and L register pair is prepared with the values 3C and 00. As a pair, they are capable, then, of *pointing* to memory position 3C00. This is the memory location of the first position at the top left of the video display. The next line prepares the B and C registers with 04 and 00. Although they are pointing to memory location 0400, they are going to be used as a counter in this program, 400 is 1024 decimal, the number of places on the screen.

In the next line, the parentheses around HL mean 'the memory location defined by'. That is, the command is saying 'store the value 20 in the memory contents defined by the H and L register pair' – in this case, 3C00. So LD (HL),20 stores a blank space (20 is the ASCII value for such a space) at the first place on the screen. Whatever

was on the screen before is turned into a space.

Notice that this has the name 'LOOP' next to it. This 'label' is for the programmer's eyes, not for the program's use. In the next two lines, the HL pair is *incremented*. That is, 3C00 is incremented to 3C01. Correspondingly, the BC pair is *decremented*, from 0400 to 03FF. You may see the pattern emerging – BC is being used in a kind of machine language FOR-NEXT loop.

In BASIC, a FOR-NEXT loop is sort of self-testing. The programmer doesn't have to put in anything that checks the value of the loop variable; BASIC does it automatically. But in machine language, a test has to be made. The next two lines of the program make that test. The accumulator is prepared not with an absolute numerical value, but rather with whatever the B register reads at the moment. After one pass through, it will be 03.

Then the accumulator is asked to make a logical judgment. Recall that earlier in this Chapter, logical OR gates were discussed. If either of two inputs were high (one), the output would be high (one). In this case, there are *eight* inputs, one for each bit in the byte. So the accumulator has been loaded with 03, thus . . .

00000011

. . . and it is being asked for the logical OR with whatever register C is now set to. In the first loop, that is FF –

11111111

What's in the Memory Map

The memory map of the TRS-80 has been designed for convenient, immediate use. It consists of seven major sections:

1. The BASIC language in read-only-memory (ROM).
2. An unassigned blank area for future expansion.
3. Special locations for cassette and disk.
4. A keyboard matrix, appearing in four locations.
5. A video memory.
6. A reserved block of RAM for BASIC use.
7. User-programmable RAM for programs, data, and variables.

The onboard jumper, X3, selects which of the various language possibilities is active in the TRS-80 (see section on electronic organization

below). Level I BASIC uses only 4,096 bytes from address 0000 to 0FFF, but Level II uses from 0000 to 2FFF. Because of hardware shortcuts, the cassette, disk and keyboard locations are incompletely decoded, appearing in 'phantom' areas beyond those strictly assigned to them. Using incompletely decoded locations in the cassette/disk area can result in unexpected results. But the phantom keyboards can be used interchangeably with the actual keyboard address.

The possible memory configurations of the TRS-80 Model I, together with reserved RAM areas, are shown below.

Figure 2-39. TRS-80 memory map in detail

Address	Function	RAM Pointer
0000	Beginning of all ROMs	
0000	Power-Up Position	
0FFF	End of Level I ROM	
2FFF	End of Level II ROM	
3000	Beginning of unassigned area	
3700	End of unassigned area	
370E	Disk Status Address	

37DF	Disk Data Address	40AD	*1	String Space Pointer (2 bytes)
37E0	Interrupt Flip-Flop	40A2		Current Line in Use (2 bytes)
37E1	Disk Drive Select	40A4	*2	Start of BASIC Program (2 bytes)
37E4	Cassette Select	40A6		TAB Position Value
37E8	Line Printer Date	40A7	*3	Input Buffer Pointer (2 bytes)
37EC	Disk Controller Chip	40A9		Input #1 Indicator
3800	Beginning of Keyboard	40AA		Seed Number for RND (3 bytes)
38FF	End of Keyboard	40AD		Reserved
3900	Phantom Keyboard	40AE		LET and DIM Scratchpad
3A00	Phantom Keyboard	40AF		Number Type Flag
3B00	Phantom Keyboard	40B0		"Flag Byte for Encoder"
3C00	Beginning of Video	40B1	*4	Top of BASIC Memory (2 bytes)
3FFF	End of Video	40B3		String Scratchpad Pointer
4000	Beginning Reserved RAM	40B5		String Workspace (30 bytes)
4000	Restart #0 Patch (RST 0)	40D3		Length of String in Use
4003	Restart #1 Patch (RST 10)	40D4		Address of String in Use (2 bytes)
4006	Restart #2 Patch (RST 18)	40D6		Next Available String Space (2 bytes)
4009	Restart #3 Patch (RST 20)	40D8		State of Print Using (2 bytes)
400C	Restart #4 Patch (RST 28)	40DA		Current DATA Line in Use (2 bytes)
400F	Restart #5 Patch (RST 30)	40DC		DIM Scratchpad (2 bytes)
4012	Restart #6 Patch (RST 38)	40DE		Print Using Scratchpad
4015	Keyboard Control Block	40DF		SYSTEM Loading Entry Point (2 bytes)
4015	Keyboard Device Type (D1)	40E1		AUTO On/Off Indicator
4016	Driver Entry Address (2 bytes)	40E2		Current Line in Use (2 bytes)
4018	Three Reserved Bytes	40E4		Size of AUTO Increment (2 bytes)
4018	Two Bytes Reading "KI"	40E6		Location of BASIC Command in Use (2 bytes)
401D	Video Control Block	40EB	*5	BASIC Stack Pointer (2 bytes)
401D	Video Device Type (D7)	40EA		ERRDR Line Number for RESUME (2 bytes)
401E	Driver Entry Address (2 bytes)	40EC		EDIT Line Number in Use (2 bytes)
4020	Location of Cursor (2 bytes)	40EE		Line Number before RESUME (2 bytes)
4022	Cursor Character	40F0		DN ERRDR GDTR Line Number (2 bytes)
4023	Two Bytes Reading "DD"	40F2		Reserved (3 bytes)
4025	Line Printer Control Block	40F5		Line Number Completed (see also 40E2)
4025	Printer Device Type (D6)	40F7		CONTINUE Line Number (2 bytes)
4026	Driver Entry Address (2 bytes)	40F9	*6	Simple Variables Pointer
4028	Total Lines Per Page	40FB	*7	Arrays Pointer
4029	Current Line Being Printed	40FD	*8	Free Memory Space (FRE(A))
402A	Reserved Byte	40FF		Pointer to DATA in Memory
402B	Two Bytes Reading "PR"	4101		Variable Type Workspace (27 bytes)
402D	Level II Workspace	4118		TRON/TROFF Indicator
402D	Unassigned RAM	411C		Arithmetic Workspace (20 bytes)
4035	End Unassigned RAM	413D		Line/Print Using Buffer (33 bytes)
4036	Beginning Keystroke Storage	4152		Disk Patch Points (see Chapter <?>)
403C	End Keystroke Storage	41A4		End Disk Patch Points
403D	Video Size / Cassette Latch	41A5		DOS Linking Patch Points
403E	Reserved for DOS Use (2 bytes)	41E7		End Linking Patch Points
4040	Storage Area for TIMES	<*3>		Keyboard/Edit Input Buffer
4047	End of TIMES Storage Area	4288		Z-80 Stack During Running Program
4048	Reserved for DOS Use	42EB		End Input Buffer
407F	End of DOS Reserved Area	<*2>		Beginning of BASIC Program
4080	Storage Area for Division			End of BASIC Program
408D	End of Division Storage Area	<*6>		Simple Variable Storage
40BE	USR Entry Point (2 bytes)			End of Variable Storage
4090	RND Storage Area (3 bytes)	<*7>		Array Storage Area
4093	INP Storage Area			End of Array Storage
4094	INPut Port Number (2 bytes)	<*8>		Free Memory Area
4096	OUTPut Storage Area			End of Free Memory Area
4097	OUTPut Port Number (2 bytes)	<*5>		BASIC Stack for NEXT, GOSUB, etc.
4099	INKEY\$ Storage Area			Top of BASIC Stack (works downward)
409A	ERRDR Code Storage	<*1>		String Storage Area
409B	Line Printer Position			Top of String Storage (works downward)
409C	Output Device Indicator	<*4>		Top of BASIC Memory
409D	Video Line Length (32 or 64)	4FFF		End of 4K RAM
409E	Video TAB Area	7FFF		End of 16K RAM
409F	Reserved	BFFF		End of 32K RAM
		FFFF		End of 48K RAM

— each bit in the accumulator is ORed with its corresponding bit in register C. If any pair of bits is 1, the accumulator's bit will be set to 1. When the ORing process is finished, the accumulator will contain the results, and the condition code register will reflect the meaning of those results.

Listing 2-2. Complete CLS demonstration.

```

00100 : #####
00110 : SIMPLE CLEAR-SCREEN DEMONSTRATION IN ASSEMBLER FORMAT
00120 : #####
00130 :
5000      00140      ORG      5000H      ; BEGIN ROUTINE AT 20400
5000 21003C 00150      LO      HL,3C00H      ; VIDEO SCREEN, AT 15360
5003 01FF03 00160      LO      BC,3FFH      ; 1,024 SPACES ON SCREEN
5006 3620    00170      LO      (HL),20H      ; SPACE (DEC 32) IN PLACE
5008 23      00180      INC     HL          ; READY NEXT SCREEN LOC'N
5009 08      00190      DEC     BC         ; DROP SPACES LEFT BY ONE
500A 78      00200      LO      A,B        ; GET CURRENT COUNT IN B
500B B1      00210      OR      C          ; OR WITH COUNT LEFT IN C
500C 20FB    00220      JR      NZ,LOOP    ; GO BACK IF NOT YET DONE
500E C9      00230      RET             ; BACK TO MAIN ROUTINE
06CC      00240      END      06CCH      ; READY AFTER TAPE LOAD
00000 TOTAL ERRORS

```

On the first pass, the accumulator will end up containing 11111111, and the condition code register's 'zero' flag will read 'not zero'. The last line of the program says 'jump if the result is not zero'. It jumps back to the part of the program marked LOOP, where it will store a 20 (space) in the new value pointed to by HL, increment HL again, decrement BC again, and go through the logical OR test once more.

If you carry the process through by hand, you will discover that only when B and C are *both zero* will the zero flag confirm a zero result. At that point, the program can shake loose from its loop.

Details of storage areas and their use can be found in *Microsoft BASIC Decoded, Supermap, Inside Level II*, and the *Level II BASIC Reference Manual*.

Setting MEMORY SIZE?

Because machine language programmers have devised many unique ways of storing their programs, the purpose of responding to 'MEMORY SIZE?' has been the cause of some confusion. A look at the summary of the full memory map may help clarify the reasoning.

Reserved Memory	(not available)
BASIC Program Text	(fills upward)
Simple Variables	(fill upward)
Array Variables	(fill upward)
***** FREE MEMORY SPACE *****	
BASIC Stack	(fills downward)
String Storage	(fills downward)
MEMORY SIZE Value	(above reserved)

Table 2-4. Memory map summary.

This table points out two important facts: array variables grow *upward* into the free memory space. The BASIC stack (which stores GOSUB return addresses, levels of parentheses, FOR-NEXT information, etc.) grows *downward* into the free memory space.

Simple variables can also bump the array variables upward; string storage space is set ahead of time with the CLEAR statement. So you can see that the free memory area is impinged upon from both sides while a program is running. Although BASIC might have been designed to bump everything upward in memory (leaving the top area of memory unmolested), it would have resulted in considerably longer running time. This is because many changes in memory would have to be made when new variables, strings, parentheses, GOSUBs, etc., were discovered during a program's run.

If any machine language program is to be used, it certainly must be stored out of the way of this frantic activity. MEMORY SIZE therefore is used as a sacred boundary, above it is 'terra incognita' as far as the BASIC program is concerned. For example, if MEMORY SIZE in a 16K machine is set to 20480, the computer *acts precisely as if it were a 4K machine!*

To make maximum memory available for a running BASIC program, this boundary should be set only *just* low enough so that the machine language program will fit above it. Most program authors will write these programs to fit as high as possible in memory, and so you will normally see memory sizes (for a 16K machine) above 30000.

Why in a 16K machine would there be a 'memory size' of about 30000 and not 16000 or so? It's simply that the prompt 'MEMORY

SIZE?' is a bad question. The memory size value is really not a size at all, but the *address* of a memory location above which the BASIC program and its variables must not go.

Why then are there machine language programs which do not require memory size to be set? That is because clever programmers write machine language programs that . . .

- may be written for Level II BASIC, and thus can reside in one of the DOS reserved areas (see memory map).
- may automatically reset the memory size value before returning to BASIC.
- may be packed into strings where they are safely protected in a program text line (see Chapter 3).
- may be short enough to reside in part of the input buffer and change its pointer.

In all of these cases, something is sacrificed for the convenience of not setting 'MEMORY SIZE?'. In the first case, DOS-like expansion programs, such as Level III BASIC, will conflict. In the second case, programs which also require the memory size to be set may be damaged when the loading program automatically resets it. Thirdly, string-packed lines may not be edited without calamitous results. And finally, a reduced input buffer makes editing long lines impossible, as they will probably run into the BASIC program text.

So the 'MEMORY SIZE?' boundary is a useful feature of BASIC, serving to protect machine language programs from the expansionist tendencies of a running BASIC program.

Comparing The Levels

Another source of confusion to a lot of users was the switch from Level I to Level II. How did this simple change of language alter the hardware? How did double-width characters, 500-baud tape loading, and key rollover suddenly appear? Why did the convenient abbreviations (P., N., M., F., etc.) suddenly go? Why were machine language programs happy with CLOAD in Level I, but needed SYSTEM with Level II?

The 4K BASIC in Level I is a compact, limited language with a few capabilities. Level II is three times as long, and much more powerful. Their authors, and hence their approaches, are different. The single hardware change in going from Level I to Level II is the installation of one ROM set in favor of another, and a minor change to allow 12K instead of 4K ROM to be accessed.

The 32-character mode hardware was already in place. The tape load speed and key rollover are all software controlled (see supplement to this Chapter). The abbreviations disappeared because Level II handled its keywords in a different manner from Level I, and such abbreviations would have increased execution time. Likewise, tape loading formats were a matter of design philosophy rather than any formal software requirements.

Level I has the advantage of being a simple, easily learned first language for computer beginners, and many TRS-80 owners learned by using that language. Level III is not a language distinct and apart from Level II, but rather an extension of the existing one. (In this sense it is much like Extended Color BASIC on the new TRS Color Computer, which does not supplant the original 8K BASIC, but merely adds another 8K to it).

What are the differences between the three levels? A command list for the three languages, with differences highlighted, follows:

Table 2-5. Comparison of Level I, II and III commands.

Command	Level I	Level II	Level III
@		X	X
A. (ABS)	X		
ABS	X		
ASC		X	X
AT	X		
ATN		X	X
C. (CONT)	X		
COBL		X	X
CHRS		X	X
CINT		X	X
CL. (CLOAD)	X		
CLEAR		X	X
CLOAD	X	*	X
CLOSE			X
CLS	X	X	
CHO			X
CONT	X	X	X
COS		X	X
CS. (CSAVE)	X		
CSAVE	X	*	X
CSNG		X	X
CVO			X
CVI			X
CVS			X
D. (DATA)	X		
DATA	X	X	X
DEFOBL		X	X
DEFFN			X
DEFINT		X	X
DEFSNG		X	X
DEFUSR			X
DEFSTR		X	X
DELETE		X	X
DIM		X	X
E. (END)	X		
EDIT		X	X
ELSE		X	X
END	X	X	X
ERL		X	X
ERR		X	X
ERROR		X	X
EXP		X	X
F. (FOR)	X		
FIELD			X
FIX		X	X
FOR	X	X	X
FRE		X	X
G. (GOTO)	X		
GET			X

GOS. (GOSUB)	X		
GOSUB	X	X	X
GOTO	X	X	X
I. (INT)	X		
IF	X	X	X
INKEY\$		X	X
IN. (INPUT)	X		
INP		X	X
INPUT	X	X	X
INSTR			X
INT	X	X	X
KILL			X
L. (LIST)	X		
LEFT\$		X	X
LET	X	X	X
LSET			X
LEN		X	X
LINE			X
LIST	X	X	X
LOAD			X
LOC			X
LOF			X
LOG		X	X
LPRINT		X	X
M. (MEM)	X		
MEM	X	X	X
MERGE			X
MID\$		X	X
MKD\$			X
MKI\$			X
MKS\$			X
N. (NEXT,NEW)	X		
NAME			X
NEW	X	X	X
NEXT	X	X	X
NOT		X	X
ON	X	X	X
OPEN			X
OUT		X	X
P. (PRINT,POINT)	X		
P.A. (PRINT AT)	X		
PEEK		X	X
POINT	X		
POKE		X	X
POS		X	X
PRINT	X	X	X
PUT			X
R. (RESET, RND,RUN)	X		
RANDOM		X	X
REA. (READ)	X		
READ	X	X	X
REM		X	X
RESET	X	X	X
REST. (RESTORE)	X		
RESTORE	X	X	X
RESUME		X	X
RET. (RETURN)	X		
RETURN	X	X	X
RIGHT\$		X	X
RND	X	*	X
RUN	X	X	X
S. (SET,STEP)	X		
SAVE			X
SET	X	X	X
SGN		X	X
SIN		X	X
SQR		X	X
ST. (STOP)	X		
STEP	X	X	X
STOP	X	X	X
STRING\$		X	X
STR\$		X	X
T. (THEN,TAB)	X		
TAB	X	X	X
TAN		X	X
THEN	X	X	X
TIMES			X
TROFF		X	X
TRON		X	X
USING		X	X
VAL		X	X
VARPTR		X	X
>	X	X	X
<	X	X	X
=	X	X	X
*	X	X	X
+	X	X	X
-	X	X	X
/	X	X	X
\$	X	*	X
()	X	X	X

* indicates that Level I and Level II operations differ for this command.

Hardware Reflects Software

You are probably familiar with the general operating characteristics of your TRS-80, including BASIC commands and how the machine responds to them. These responses are characteristics of how the software treats the hardware, and also of how aspects of the hardware act independently of the software.

The hardware inside the TRS computer can be broken into seven major sections:

1. CPU Hardware

- A. Central Processing Unit (Z-80), its clock, power-up, and reset circuitry.
- B. Decoding of CPU status signals into memory/peripheral access signals such as read, write, input, output.
- C. Buffering of address and data signals to and from the CPU.

2. Program RAM Control

- A. Refresh signals to maintain memory in dynamic RAMs.
- B. Address decoding able to distinguish 4K and 16K RAMs.
- C. Address multiplexing for dynamic RAM address lines.
- D. Read/Write signals to RAMs.

3. Video RAM Control

- A. Address decoding for video RAM.
- B. Read/Write signals to video RAM.
- C. Input to character generator, timing, blanking signals so characters do not run off the screen.
- D. Access management between display and CPU.
- E. Alphanumeric/graphic switching and graphics character circuitry.

4. Keyboard

- A. Address decoding for keyboard.
- B. Address/data buffering and read signals.

5. ROM Control

- A. Address decoding for 4K and 12K ROMs (Level I or Level II).
- B. Outboard decoding for three 4K Level II ROMs or two Level II ROMs.
- C. Read signals.

6. Output Controls

- A. Parallel-to-serial conversion from character generator.
- B. Horizontal and vertical video sync circuits, video output circuit.
- C. Cassette motor, cassette data, and 32 character video output control.
- D. Cassette audio output circuitry.

7. Power Supply

- A. Three regulated voltage outputs.
- B. Short-circuit protection.

Each of these sections plays a major role in the operation of the TRS-80, and few could be trimmed or eliminated without completely changing the character of the computer. The rest of this Chapter will be devoted to detailing those aspects of the TRS which are significant to customizing the hardware or software of the machine. For a more comprehensive examination, including timing diagrams and discussion of each signal line, turn to the *Technical Reference Handbook*.

CPU Hardware

The master clock is produced by the oscillations of a 10.6445 MHz (million cycles per second, or Megahertz) crystal. A countdown circuit (Z56) divides this by 6, producing the running frequency of the Z-80, 1,774,083 clock cycles per second. This is generally called the TRS-80's 1.77-MHz clock.

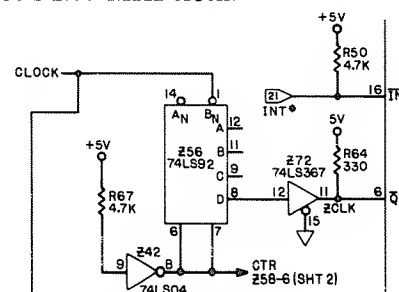


Figure 2-40. Clock circuit area of TRS-80.

It's interesting to note that there are several other clocks on the computer's board, some already wired in place, and others which can be created by interconnections. At pins 11 and 9 of Z56, 3.548 MHz is available, twice the normal TRS-80 speed. Pin 2 of Z43 clocks at 5.322 MHz, faster than the Z-80 can run, but when connected to pin 14 of Z56, a 2.661-MHz clock is available. This is 1.5 times the normal clock speed. Both the 2.661 MHz and 3.548 MHz clock rates will be used in Chapter 4.

Upon power-up, the processor's RESET input, which sets the program counter running at 0000, is triggered by Z53 in combination with Z52. The capacitor C42 takes time to charge completely, so the RESET line is held down for a few milliseconds after the rest of the system comes up and is stable. This power-on reset is a convenience so the user doesn't have to press a special reset button just to get the system going. Other computers, such as the Ohio Scientific series, demand that inconvenient action.

The TRS-80 design creates Read (RD), Write (WR), Input (IN) and Output (OUT) signals from the Z-80's Memory Use Request (MREQ), Input/Output Use Request (IORQ), Write (WR), and Read (RD) signals. These are combined by Z23, in the correct order to do that. The four Z-80 signals *are not* wired to the edge card connector, so certain functions (such as mode 0 and mode 2 interrupts – see supplement to Chapter 5) cannot be used.

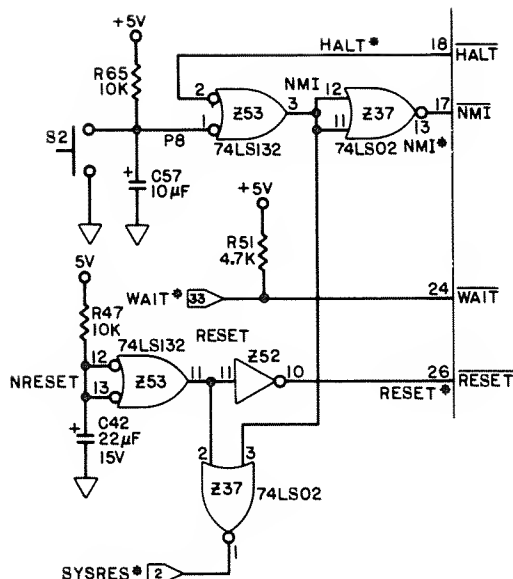


Figure 2-41. Power-up reset circuit area of TRS-80.

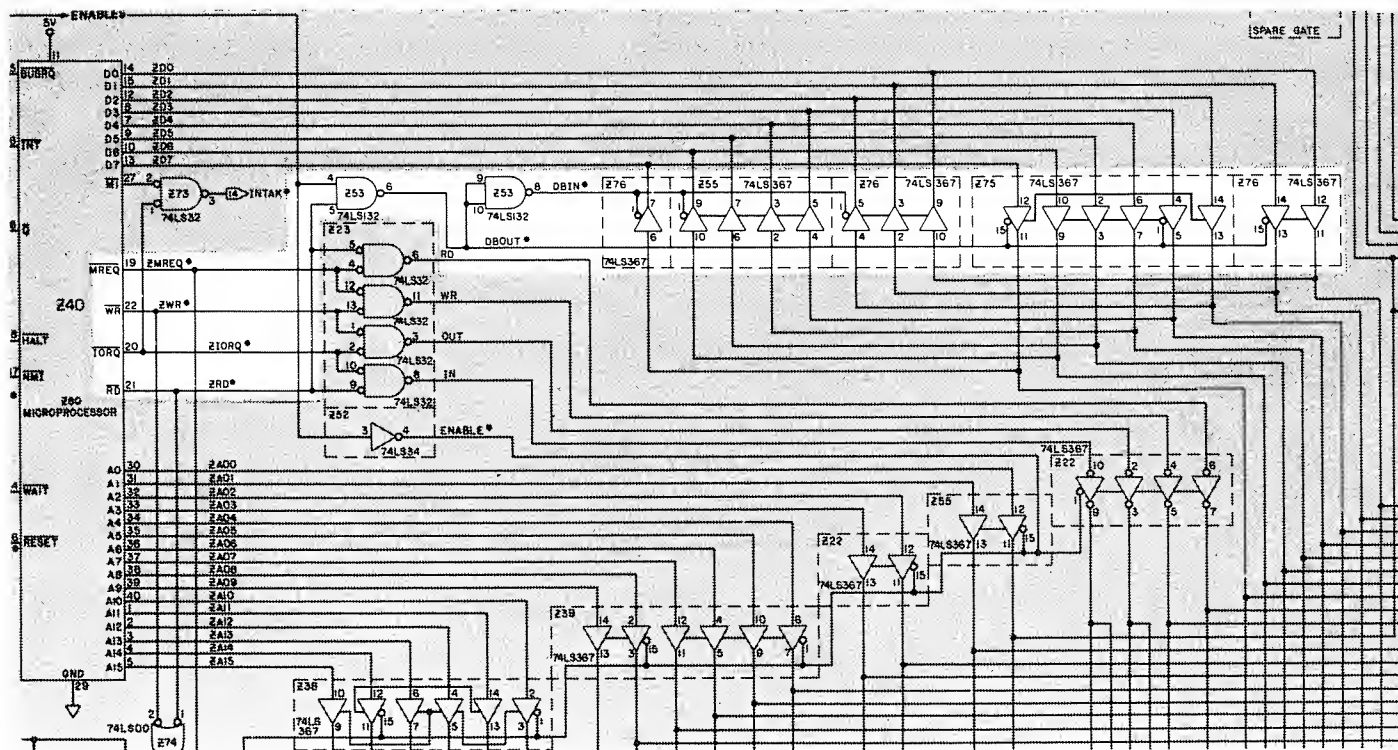


Figure 2-42. Read/write status, multiplexer of TRS-80.

Program RAM Control

A TEST line is provided to 'float the bus' – in other words, the Z-80 becomes invisible, allowing another device to take over operation of the computer. Some outboard devices which speed up the TRS-80 use this feature, essentially taking control of the memory and peripheral devices by bringing the TEST line to ground. Z52 goes high in response, electronically disconnecting the Z-80's address and data bus from the circuitry. (Note that using the TEST line without memory-refresh backup circuitry on the outside of the computer will result in loss of memory contents).

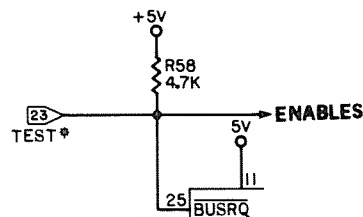


Figure 2-43. Test line circuit of TRS-80.

In normal operation, the Z-80 lines are active and buffered by Z38, Z39, Z22 and part of Z55 (address), Z75, Z76 and the remainder of Z55 (data). Except when the TEST line is used, the address buffers are always active. The data buffers are active under any circumstances, either in their READ or WRITE configuration:

Except for the memory refresh information, this completes the role of the Z-80 CPU circuitry in the TRS-80.

Program RAM Control

The CPU is also used in the creation of the signals needed to refresh the dynamic memories. Since the TRS-80 uses dynamic RAMs, (see Chapter 5 for details on this), the normal refresh (RFSH) output of the Z-80 is less than useful, at least in the minds of the computer's design engineers. That RFSH signal, which is output when the computer is not using the memory, is ignored in the TRS-80.

Instead, the processor's MREQ line, when buffered, serves as a memory address row signal. The master clock is used in conjunction with the Z-80's RD and WR lines to produce a memory address column signal (column-address strobe, CAS) and a multiplex signal (MUX) to switch from row to column. This serves a very useful double purpose: not only does it refresh memory when the processor is not specifically using the memory in the program, but it serves as the address-select lines when the Z-80 processor is using memory. Refer to the Z-80 Technical Manual for details on the timing of these signals.

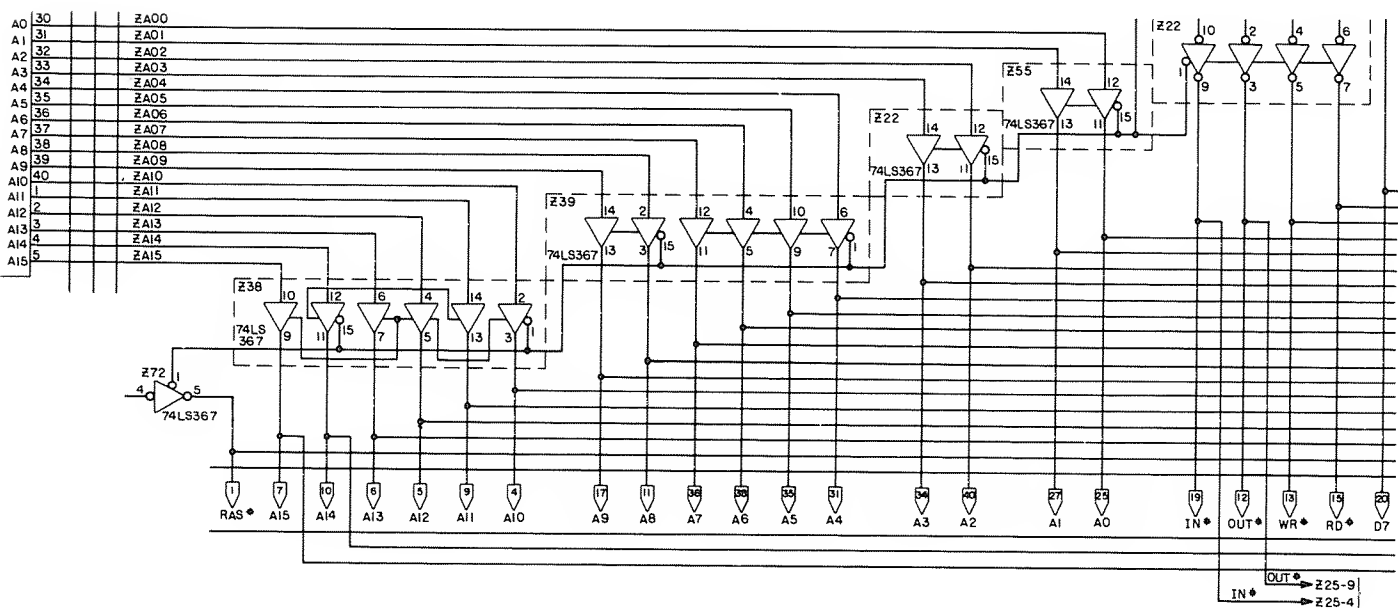


Figure 2-44. Address/data buffers of TRS-80.

Needless to say, should any of these signals fail or operate inconsistently, the memory will not retain its contents for very long *nor will a program even run which uses dynamic memory, because its address may not be selected.*

The next subject is memory management. This is basically the means by which the processor gets to the memory it wants to use. The heart of this sequence is found in Z35 and Z51, a pair of multiplexers which send the low bits of the memory address to the dynamic memories; flip from low bits to high bits according to the incoming multiplex (MUX) signal; and send the high bits to the memory. The memory, upon receiving these addresses together with the previously mentioned RAS and CAS signals, knows which address is being selected, and responds accordingly.

The DIP (Dual Inline Package) shunt Z71 plays an important role here. Specifically, it

Figure 2-45. Memory select/refresh of TRS-80.

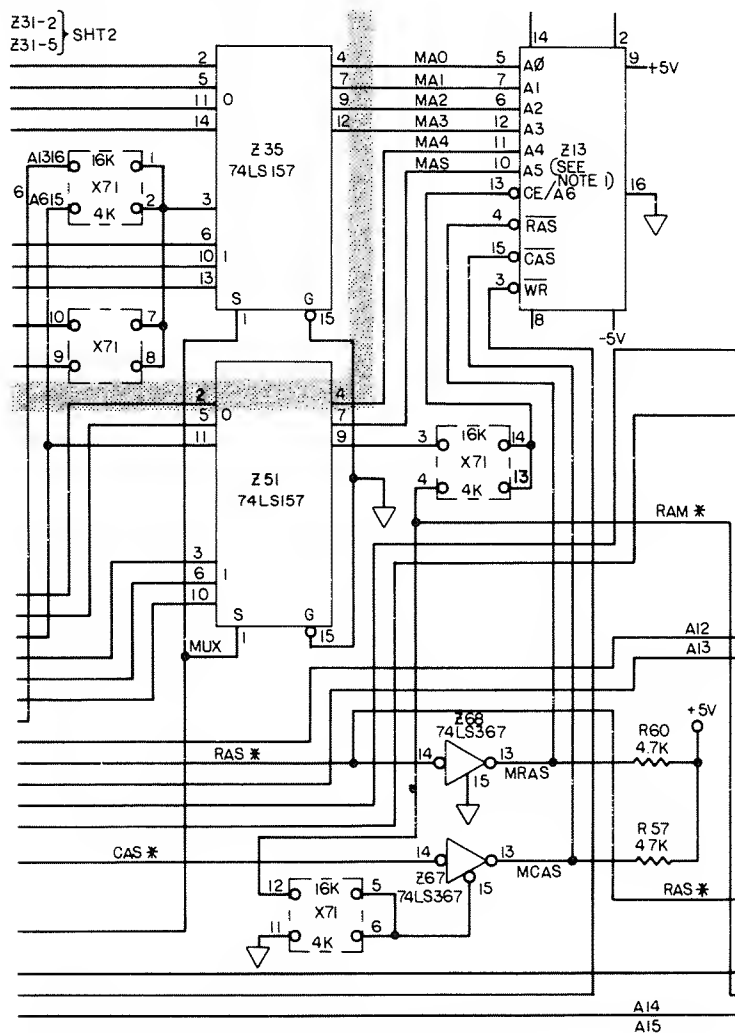


Figure 2-46. Memory multiplex, CAS, RAS & memory.

routes the signals to the multiplexers in such a manner that the computer can distinguish between 4K and 16K RAMs. Stated simply, what is a memory-chip select line for 4K RAMs is a complete address line for 16K RAMs, and a partial address line for 8K RAMs (which the TRS-80 was designed to use also). Thus, the higher address lines must be prevented from appearing at the CE (chip enable) input of the 4K RAMs. If this were to happen, phantom memory would appear, and a running BASIC program (and the power-up memory test) might try to use those phantom bytes:

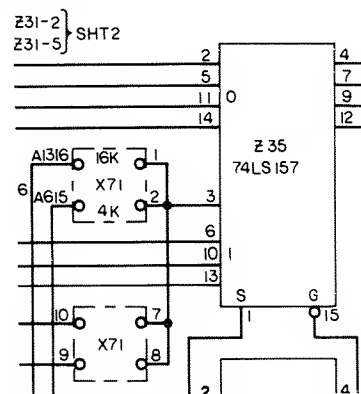


Figure 2-47. X71 and decoding scheme of TRS-80.

Only a WR signal is used to trigger these RAMs, as a high signal on their WR lines prepares them to read data. The CE (Chip Enable) determines if data should be placed on the data bus, where it is buffered by Z67 and Z68.

Video RAM

Video memory is used in three ways: it is written to by the processor, read from by the processor and read from by the *video display circuitry*. In fact, it is constantly being read by the video display circuitry *except* when the processor demands attention.

The circuitry is complicated, and if you are interested in details, turn to the *Technical Reference Handbook*. In brief, memory is accessed by the processor using the Read (RD) and Write (WR) lines in conjunction with the decoded video area address-select line (VID). The display circuitry uses the video memory in a more complex manner: characters are output to Z27 and Z28, and these in turn are fed to Z29 (a character generator) and Z8 (a multiplexer). The characters, whether alphanumeric or graphic, are fed to Z10 (a shift register), where they are fed, a bit at a time, to the video output circuitry (beginning at the input of Z30).

The character generation process is complicated by several factors. The dots that make up each letter must be fed to the video output circuitry only when the video monitor's electron beam is sweeping the visible part of the screen. The visible part of the screen does not include the upper, lower, left or right borders. The timing process must continue correctly even when the CPU is using the video memory. Because each letter is made up of twelve vertical dots, and each line is made up of 64 characters with six vertical dots each, different parts of the characters must be output to the screen at different times.

Again, the *Technical Reference Handbook* covers this in detail, but a few decoded signals are important. The output of Z30, pin 10, is the final BLANK signal; no characters are output when this signal is active. Presence of characters or graphics in the border areas points to problems with this line.

The signal to shift video bits out to the video circuit is provided by Z26. Pin 8 controls alphanumeric bits, pin 6 controls graphics bits. Mangled screen characters may be traced to here, or to any of the seven chips that select characters: Z65, Z50, Z12, Z32, Z64, Z49, and Z31. This is one of the most unpleasant areas to attempt to diagnose.

The 32/64 character mode select (MODESEL) is provided by Z59, pin 9, and changes the speed of the video clock at Z43. Failures in either Z59 or Z43 will show up as a lockup in one mode or the other.

The presence of bit 7 determines if the computer is to output graphics or alphanumerics, and that signal (DLY BIT 7) is output in normal and inverted forms from Z27 pins 2 and 3. Failure in either mode can be examined here, or at the outputs of Z26, pins 6 and 8.

Keyboard

The keyboard is very different from the video; it's just a simple key matrix. When the keyboard address area is read, the KYBD line from Z36 pin 11 triggers the keyboard integrated circuits (Z3 and Z4 on most keyboards) into action, outputting information to the data bus.

The data to be output is determined by the low eight bits of the address requested. The specific address requested is inverted, and a low signal is detected whenever a key in that matrix row is pressed. The inverting buffers to the data line provide the appropriate row information. For details on how the software interprets this switching matrix, see the supplement to this Chapter.

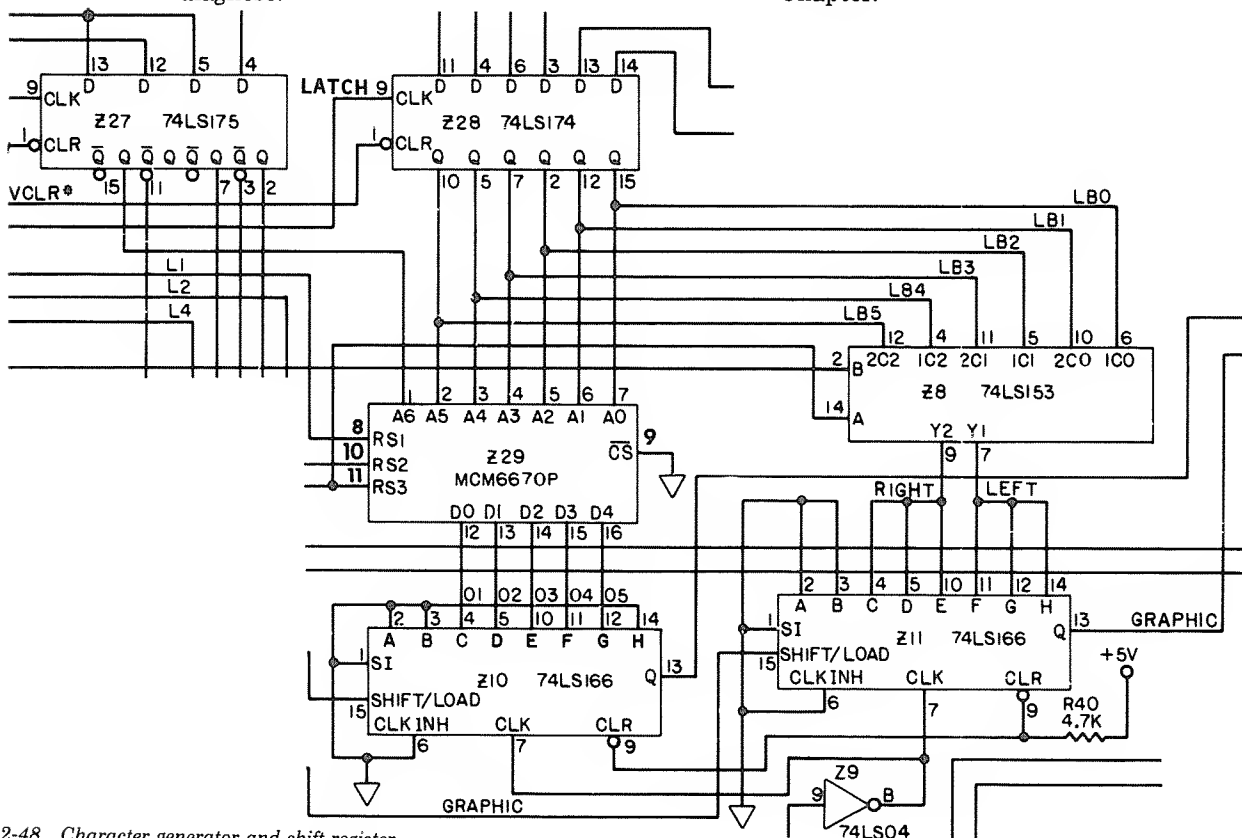


Figure 2-48. Character generator and shift register.

An interesting aspect to this is how the depression of *any* key may be detected. If address 387F is sought, all the address lines (except SHIFT) will become active, and the presence of any depressed key (except SHIFT) will appear on the address bus. Requesting the data at address 38FF will return the presence of any key *including* SHIFT. This is useful in creating a keyboard buffer, which is built from characters detected whenever the INTerrupt line is triggered.

In other words, the interrupt line triggers, and the program moves to the interrupt service routine. This routine reads address 387F, and the presence of any depressed key can be sensed. If one is pressed, it can be accepted and evaluated. Otherwise, the interrupt routine promptly returns to the program in progress.

This is also a valuable addition to INKEY\$ in some situations; see 'Peeking the Keyboard' in Chapter 3.

ROM Control

Selecting the ROMs is the biggest sticking point in the Model I. This selection is accomplished by Z3, another DIP jumper shunt, in conjunction with Z21, a 74LS156 demultiplexer. Z21 joins various address lines to produce VID, KYBD, and MEM for the video, keyboard, and dynamic RAMs and to produce variants on the ROM line.

Two-chip Level I sets were selected by a combination of methods, all of which are detailed in the *Technical Reference Handbook*. Each ROM was 2K bytes in size. Three different versions of the board were publicly released, marked 'A', 'D' and 'G'. Each had a different hard wired method of decoding ROM. An occasional 'B' or 'F' board has been reported to me, but I have never seen one. Follow *Technical Reference Manual* descriptions carefully to make sense of these Level I lines.

Some of these ROMs were EPROMs with identical pinouts, so a 'ROM A' and 'ROM B' pair of lines were needed so these memories would not conflict.

Later Level I ROMs had the selecting circuitry masked right onto the chips, which removed that conflict. Finally, a single 4K ROM was introduced to eliminate these difficulties completely.

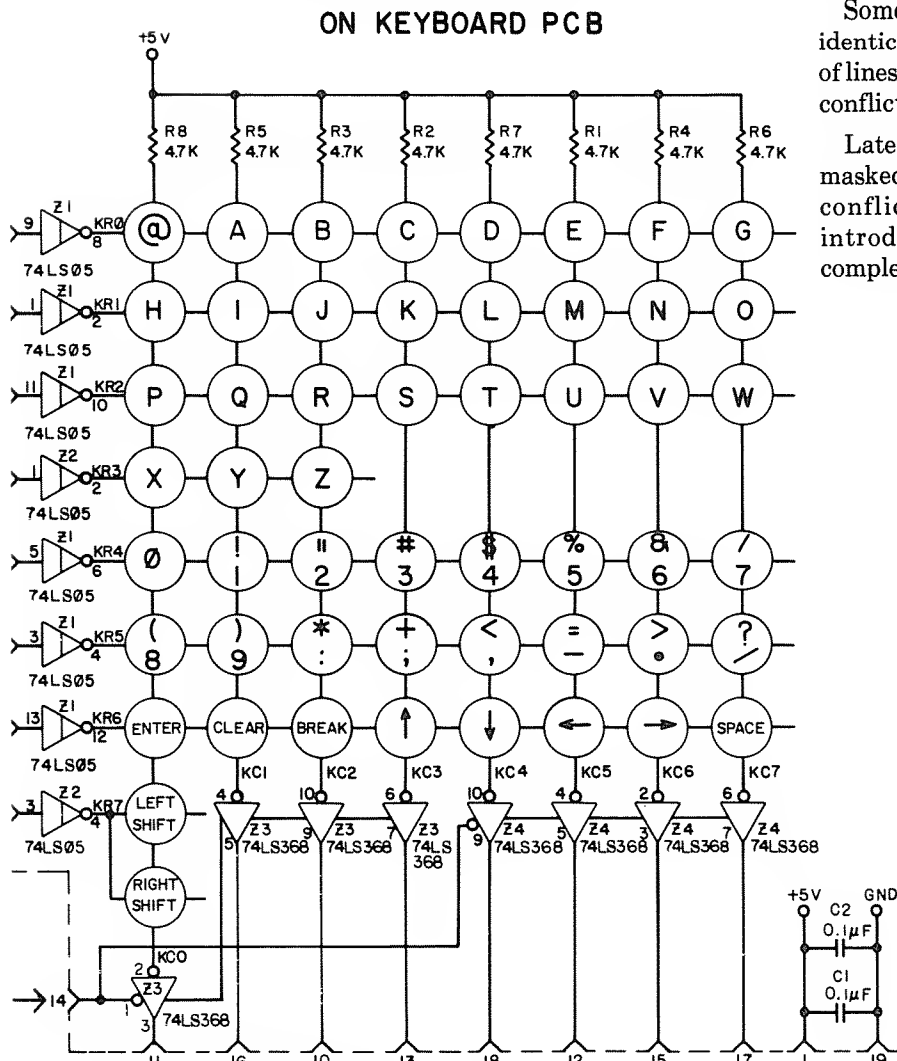


Figure 2-49. Keyboard matrix of TRS-80.

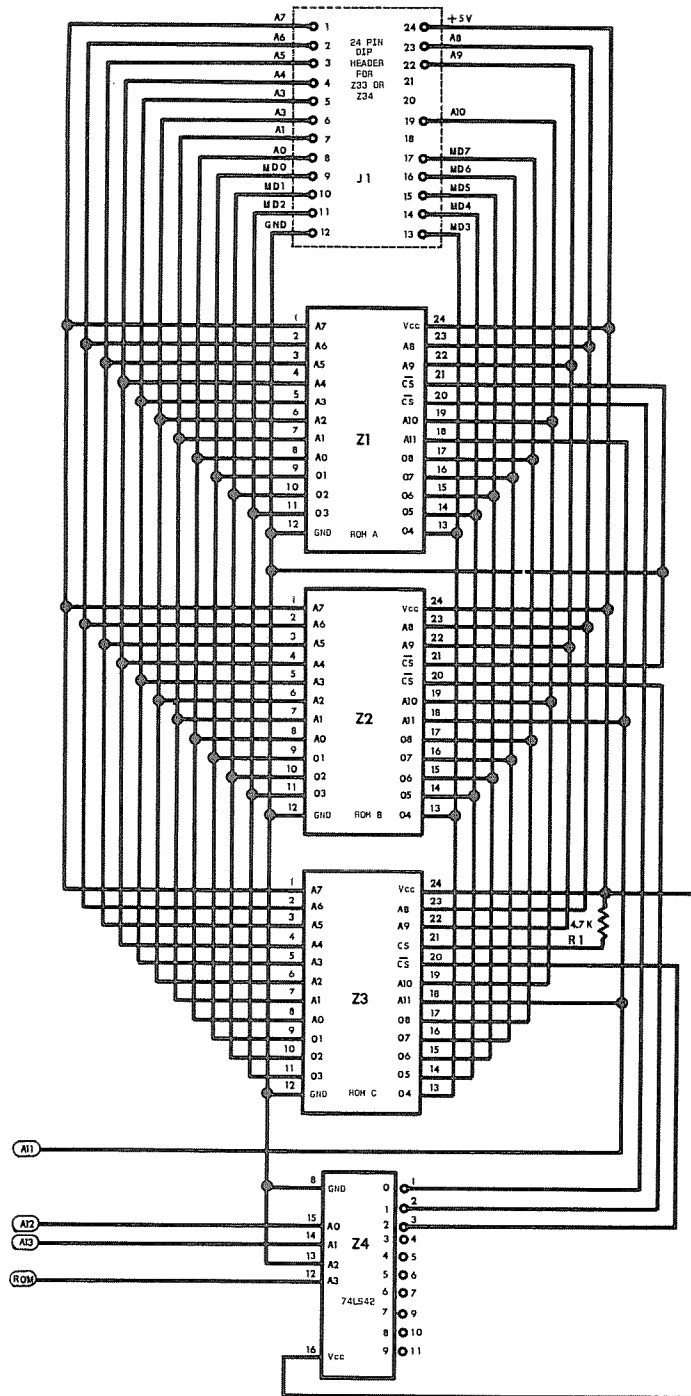


Figure 2-50. 74LS42 Level II demultiplexer of TRS-80.

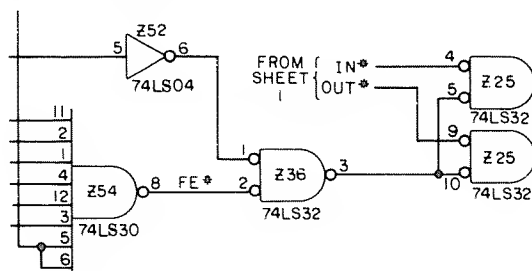


Figure 2-51. INSIG and OUTSIG area of TRS-80.

The Level II ROMs were for a long time quite consistent as a three-chip set, 4K bytes of memory in each one. A separate, piggybacked board was used. The decoded ROM line, in conjunction with address lines 11, 12 and 13, feed a 74LS42 wired as a demultiplexer. Each ROM selects a bit differently, and the combination of signals allows the memory areas from 0000 to 0FFF, 1000 to 1FFF, and 2000 to 2FFF to be selected.

In later Model I's, the three-chip set was reduced to two masked ROMs with different select characteristics. The piggyback board disappeared, and the two ROM sockets, once used for Level I, were now used for Level II. A modification to remove the two-chip ROM set, and make room for the addition of Level I, is presented elsewhere.

Output Controls

Some of the output controls have already been pointed out, including the parallel-to-serial conversion of video characters and the 32 character video control. The latter is part of a latch which is activated by information sent to port FF. Z54, Z36 and Z52 together decode port FF, and this is brought together with the computer's IN and OUT signals. The result is a pair of lines marked INSIG and OUTSIG, which are used to activate the input/output circuitry.

A latch, Z59, holds the cassette motor on or off, and also locks the video into 32- or 64-character mode. It is also latched on and off to create the 500-baud pulse rate of data to the cassette player. This output signal is purely digital, as opposed to the input signal (discussed below).

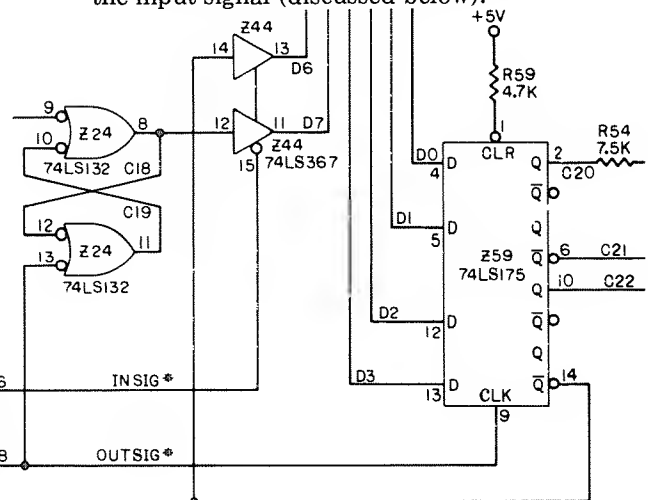


Figure 2-52. Output data latch area of TRS-80.

It would have been possible to have used Z54 to decode port FF directly, without using Z52 or Z36. Whether or not this was a design expansion consideration, you can use it as such since ports FF (255 decimal) and FE (254 decimal) are created by Z54 and the separate low-bit data line. In Chapter 4 port FE will be used for video and speed changes.

Cassette input is provided via a low-pass filter, through parts of Z4, where it is turned from low grade audio into a reasonable digital signal.

When INSIG is activated, whatever data is present at the cassette input is switched onto the data bus, and the CPU can read it. Then OUTSIG may reset the flip-flop created by Z24, when the program is ready to read the next piece of data from the cassette input. Note that the input can be any audio signal. The cassette port is not limited merely to taped data, but can be used to decode communications, shortwave, and ham transmissions, or test for the rise of voltage to a triggering level.

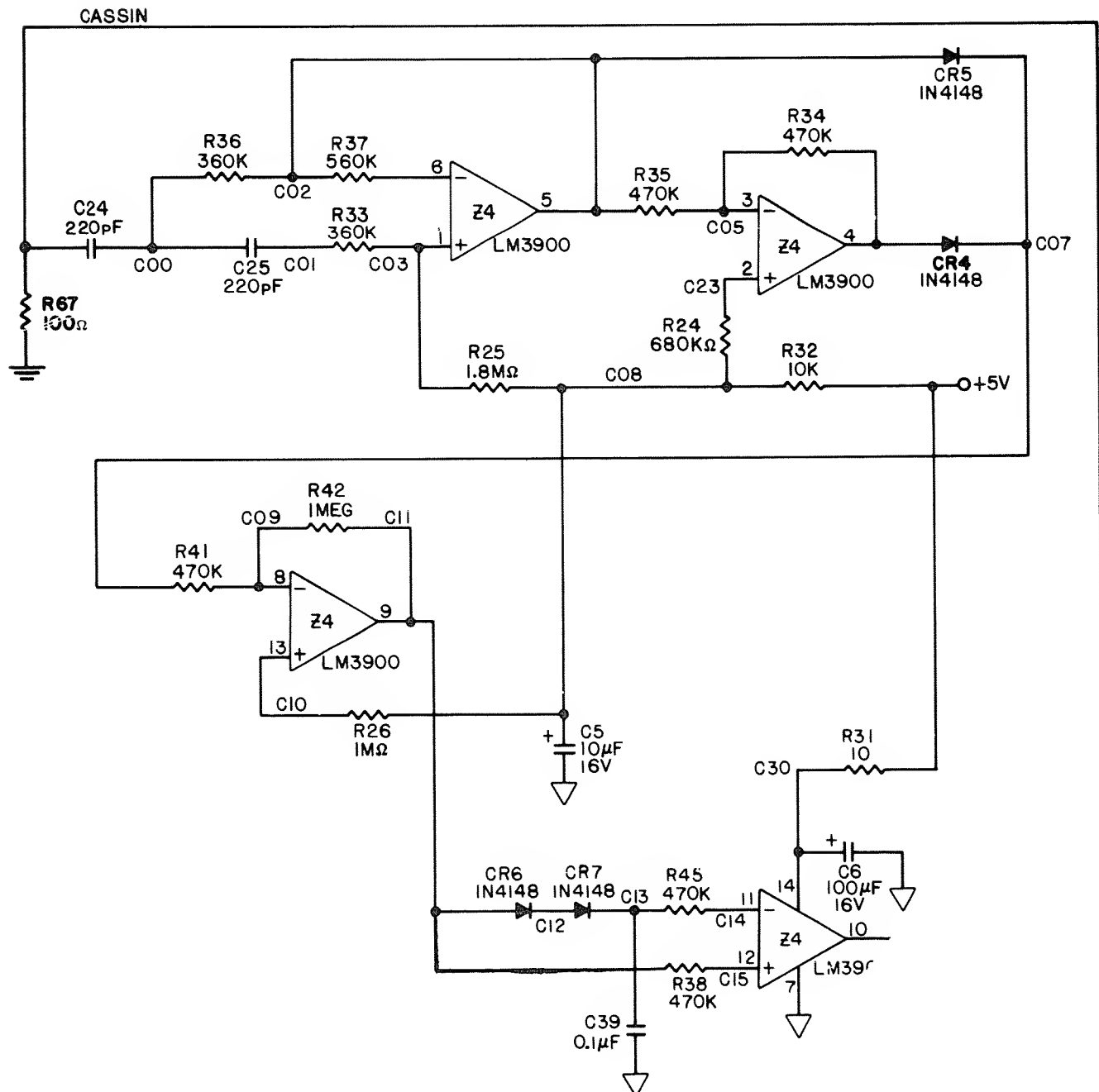


Figure 2-53. Cassette data shaping circuit of TRS-80.

The complex video divider chain provides HDRV and VDRV (horizontal and vertical drive) signals for television monitor synchronization. These signals are fed into a group of digital phase-shifting circuits which permit the signal to be adjusted on the video screen.

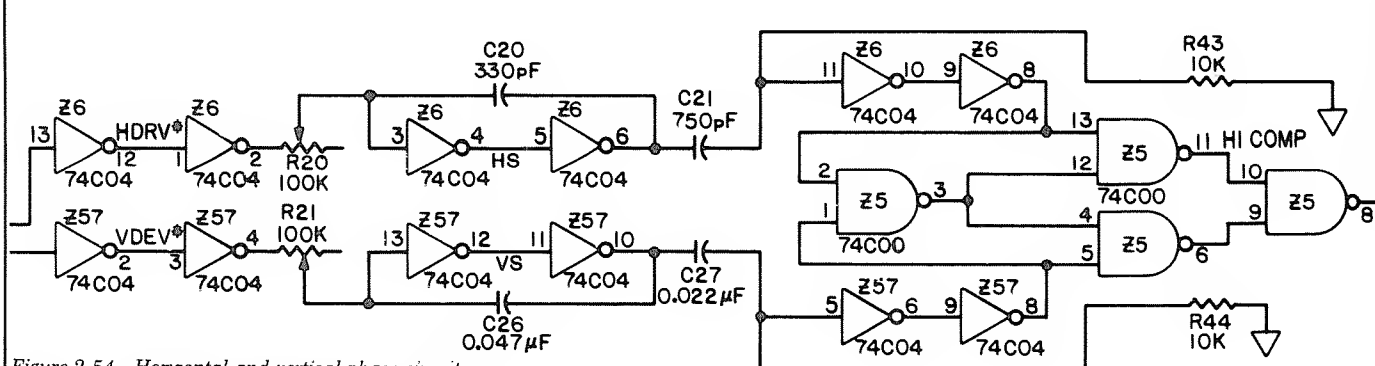


Figure 2-54. Horizontal and vertical phase circuitry.

The signals are mixed together at Z5 to provide a complete synchronization signal, and this sync signal is mixed with the video signal by Z41, Q1 and Q2. The result is a composite video signal which is capable of running a standard television monitor, or an RF modulator. The RF modulator signal can then drive an ordinary television.

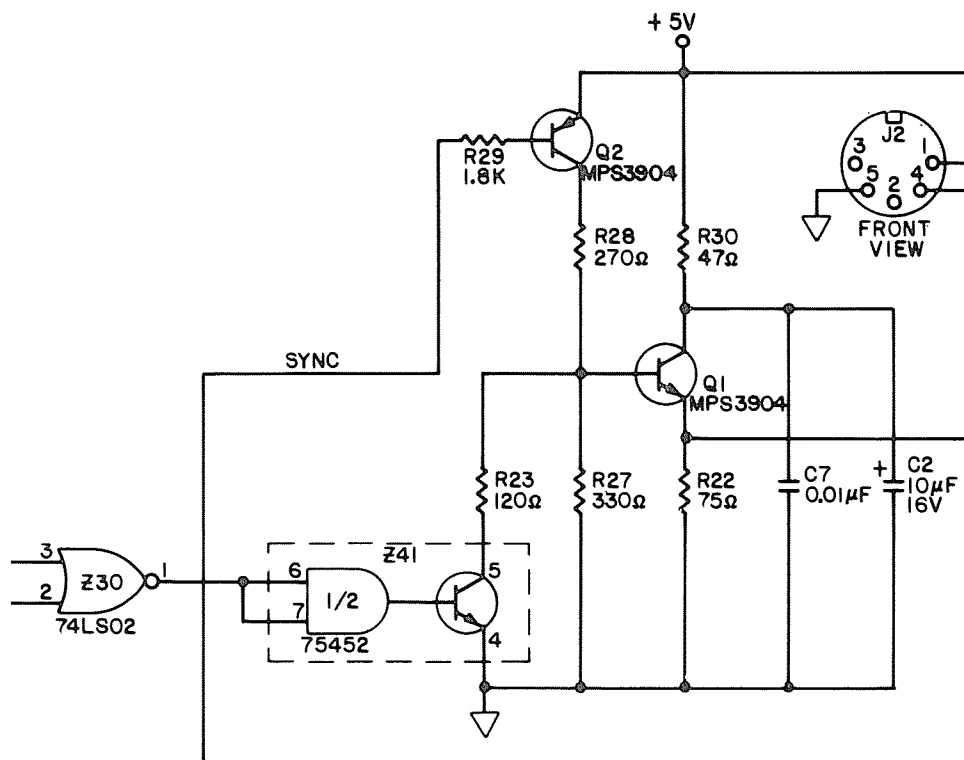


Figure 2-55. Video mixing circuit of TRS-80.

Power Supply

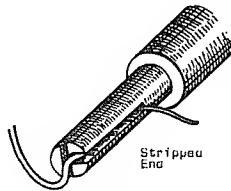
Refer to the *Technical Reference Handbook* for an excellent description of this circuitry.

Wire-Wrapping Technique

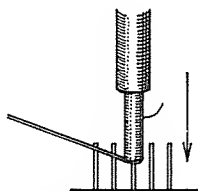
It's not without a bit of hesitation that I attacked many of the hardware projects presented in this book. Some are simple, but many, particularly those using memory circuits, need many connections. The wiring is not complicated, just tedious.

If you work carefully, all is likely to be well; but even a touch of haste will encourage confused connections. It is in these cases especially that wire-wrapping is the technique to use.

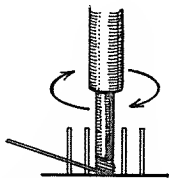
Wire-wrapping is not only easier than soldering, it is secure, simple, easier for correcting mistakes – and less costly. For wire-wrapping, you will need wire-wrap sockets, which are sold by most hobbyist supply houses including Radio Shack. Likewise, wire-wrap wire and a simple hand tool are used for the process. Here are the steps:



1. Insert stripped wire.



2. Slip over pin. Hold wire firmly, and slide fully down.



3. Spin wire-wrap tool. Wire rises along pin.



4. Finished connection has no bare wire protruding.

1. The wire, still connected to the spool, is inserted in the V-shaped stripping slot. Insert between one half and one inch of wire. Pull downward from the V, and the wire will slip out, leaving a piece of insulation in the stripper, where it can be shaken out.

2. Look carefully at the end of the wire-wrap tool. There is a small hole, meant to fit over the pins of a wire-wrap socket. Next to it is a half-circle, into which you must slide the stripped wire. The stripped portion will slide up a groove in the side of the tool, stopping where the insulation begins.

3. When the wire is in place, pull it sharply but gently upward, and slide the tool on the wire-wrap socket. Holding the wire firmly, spin the tool in your hand. The wire will wind up on the socket pin, freeing itself from the tool. Remove the tool. The wire-wrapping is complete for that end of the connection.

4. Cut the wire to a length that will comfortably reach its destination, and then some. Strip the end of it, and repeat the process above. The connection is complete. Don't forget to use different colors (white, yellow, red and blue are generally available). This will help you distinguish your connection patterns if changes become necessary.

Peripheral Addressing

The bulk of the external devices attached to the TRS-80 do their own address decoding work. Some have become standardized by conventions of their use, and others have been used somewhat haphazardly by various manufacturers.

Table 2 - 6

Addresses reserved:

3DDD-37CF	Exetron Stringy-Floppy Personal Microcomputer Festload Personal Microcomputer REX-80 Peripheral People Memory Sidecar
37DE	Communications Statue (Expansion)
37DF	Communications Data (Expansion)
37ED	Interrupt Flip-Flop (Expansion)
37E1	Disk Drive Select Latch (Expansion)
37E4	Cassette Drive Select (Expansion)
37EB	Line Printer I/O (Expansion) Percom Electric Crayon I/O Percom Speak-2-Me-2 Microcompatible Printer Buffer Floppy Disc Controller (Expansion)
37EC	Electronic Systems Serial I/O
37FB	

Table 2 - 7

Output Ports Reserved:

D (Selectable)	Alpha Product Interface Devices
1 (Selectable)	Alpha Product Interface Devices
2 (Selectable)	Alpha Product Interface Devices
7 (Selectable)	Alpha Product Interface Devices
8	JPC Poor Man's Floppy System
55	Electronic Systems Serial I/O
208	Microperipheral Microconnection
209	Microperipheral Microconnection
232	Lynx Modem Radio Shack RS-232 Board
233	Lynx Modem Radio Shack RS-232 Board
234	Lynx Modem Radio Shack RS-232 Board
235	Lynx Modem Radio Shack RS-232 Board
254	Archbold High Speed Board Mumford Micro Speed Mod Board Most Internal User Modifications
255	Cassette Data I/O (Internal) BD-Grefix (Programme) Cassette Motor Switch (Internal) Video Character Size Latch (Internal) Simutek T-Bee Addition
Addressable	Mullen Computer M-80 Interface Quent Systems PPI-80 I/O Port Orion Instruments In-Circuit Emulator Optimal Technology EPROM Programmer

Table 2 - 8

Other Peripheral Device Uses

Unaddressed (using bus control signals):

Cadnet Software/Hardware Extension, The Patch
Microgramme Programmable Graphics, Grefix-80
Microcompatible Company, The 225% Solution
Alpha Product Stick-80 Joysticks
SEL IEEE-488 to TRS-80 Interface
Xitek STD Bus I/O Card System

Unaddressed (using cassette I/O signals):

Most available Light Pens, including
Most available Cassette Data Digitizers, including
Adu-Data, Data Dubber, E-Z Loader

Unconnected (using RFI interference pickup):

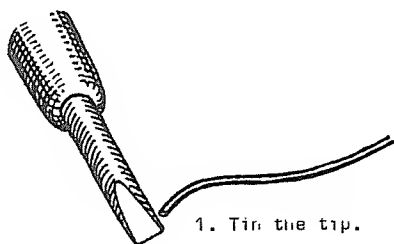
Micro-Mega CPU Monitor

Soldering Technique

For projects from scratch, soldering should be considered the final process, the actions of a self-assured, confident hobbyist. For modifications, it is a necessity. In either case, and whether you are a micro-acrobat or distinctively clumsy like me, you can solder well. The requirements are patience and good solder.

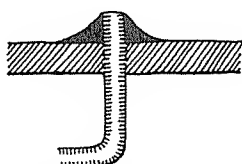
To start, make sure you are using an iron in the 25 to 40 watt range, never a soldering gun. The solder should be high quality, multicore solder. It is expensive, but will save many grief-stricken hours tracing 'cold solder joints', or removing globs of dull solder from between and under integrated circuits.

1. Clean the soldering iron tip, and heat the iron. Flow fresh solder on the tip to 'tin' the tip, which will help the solder flow from the tip of the iron to the part to be soldered. If the iron has been used, clean any encrusted material from the tip, and use coarse emery paper to shine the solder. If the tip gets deformed, bent, or very corroded, file it sharp with a fine file, and re-tin the tip.

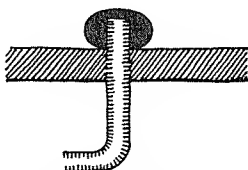


1. Tin the tip.

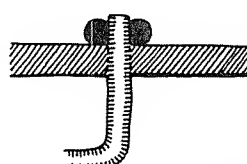
2. Keep an old sponge handy, slightly damp. Run the tip of the iron quickly over it as you solder to remove the excess flux. Always use a soldering iron holder (usually provided with an iron); if you don't, you'll wish you had the first time you burn a large hole in your imitation walnut, vinyl-topped desk.



4. Finished solder connection.

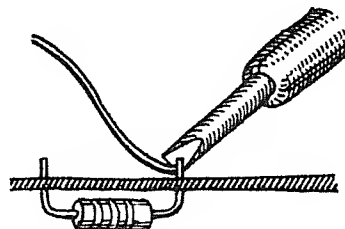


5. Bad solder connection - no contact with board (side view).



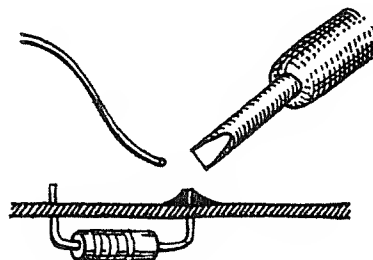
5. Bad solder connection - no contact with pin (5A side view) (5B top view).

3. In the olden days, the rule was 'heat the parts, not the solder'. Forget it. Make sure the iron is no hotter than 40 watts (and remember never to use a soldering gun) and that the parts you are about to solder are very clean. Place the iron against the part, making as much contact with it as possible along the angled tip of the iron. Place the end of the solder at the juncture of the iron and the part, and flow just enough solder to make a clean, shiny, flowing connection.



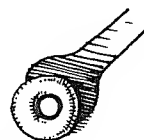
2. Bring solder, parts and iron into contact.

4. Remove the iron immediately and let the part cool. If a wire is being soldered, hold it still until the solder becomes cloudy and cool, or else an incomplete connection may result.



3. Lift iron and solder simultaneously.

5. If solder bridges develop between connections that are very close together, don't try to suck up the solder with the iron; you can only overheat the parts that way, and end up with blobs of solder and flux. Instead, use solder wick or solder-up to remove the excess solder, and start again. Let the parts cool before soldering again (a half minute should be enough).



On the Keyboard Scan

Arrows? Control codes? Autorepeat? Whatever it is you would like, that has to do with the keyboard, you can do with the TRS-80. The designers of the machine chose not to use an ASCII (American Standard Code for Information Interchange) keyboard . . . one that outputs a code for each key pressed; instead, the keyboard is a matrix of switches. Because of this decision, the TRS-80 keyboard can be extremely versatile with a minimal body of software.

First let's take a look at the keyboard matrix itself. If you have been programming in machine language, or even relatively sophisticated BASIC, this map will be familiar:

Addresses	Keys							
3801	@	A	B	C	D	E	F	G
3802	H	I	J	K	L	M	N	O
3803	P	Q	R	S	T	U	V	W
3804	X	Y	Z	RESERVED				
3810	0	1	2	3	4	5	6	7
3820	8	9	*	+	<	=	>	/?
3840	ENT	CLR	BRK	UPAR	DNAR	LFAR	RTAR	SPC
3880	SHIFT RESERVED							
Data:	01	02	04	08	10	20	40	80

At first, the arrangement of the address lines and data information may seem unappealing. What is the use of having address and data information that doubles at each change? Why not just use 3801, 3802, 3803 and so on? The reasons will eventually become apparent – in the ease and speed of the keyboard scan (and in the low cost of the hardware), but consider this binary for a moment . . .

```
00000001 00000010 00000100 00001000
00010000 00100000 01000000 10000000
```

. . . and there you have it. The bit is bumped along for each keyboard row and column, so that the presence of a single active bit instantly identifies any of the (potential) 64 keys.

The entrance to the keyboard scan is made with a jump from address 4016 in the ROM's RAM switchboard. Normally at power up, the jump address 03E3 is put in place, and the keyboard scan is entered at that point. Registers BC, DE, HL, and A are used in the scanning process:

```
03E3 21 36 40      LO      HL,4036
03E6 01 01 38      LO      BC,3801
```

The HL register pair points to the first RAM location at which the keystrokes will be stored. BC is set to look at the first row of the keyboard, whose memory map is 3801 to 3880 as noted earlier. Register D is set to zero, and it will become a 'row counter'; the process begins at

address 03EB:

```
03EB 0A          LO      A,(BC)
03EC 5F          LO      A,A
```

The accumulator reads the data at BC (recall at the outset it is pointing to the first keyboard row, 3801). The information it finds is stored in the memory location pointed to by HL (4036).

Okay so far. Now comes some of the interesting stuff that distinguishes this scan as an excellent piece of writing:

```
03E0 AE          XOR      (HL)
03EE 73          LO      (HL),E
03EF A3          AND      E
03F0 20 0B       JR       NZ,03FA
```

This short segment is responsible for the 'rollover' capabilities of Level II. The contents of the accumulator (the keystroke, if found) is XORed with the previous contents of 4036. Recalling how the XOR function works, we discover that if the key pressed was the same as the previous one at this row, the accumulator will be 'toggled' to zero. In any case, the current keystroke, whatever it is, is now saved in 4036 (so that the next time 'round, it knows if a key is still pressed).

If the key was the same, AND E will be the result of A-toggled-to-zero AND the found keystroke . . . or zero. If there was no key pressed, the result will be A XORed with HL (which is essentially irrelevant) AND E-which-is-zero . . . or zero. The test at 03F0 is for not zero. Under these conditions, it fails, so the program continues:

```
03F2 14          INC      D
03F3 2C          INC      L
03F4 C8 01       RLC      C
03F6 F2 E8 03    JP       P,03E8
03F9 C9          RET
```

The 'row counter' (D register) is incremented, and the low-order byte of HL is incremented (to storage address 4037), and the low-order byte of BC is rotated. Recalling the keyboard matrix, we can see that this command to rotate moves us from 01 to 02, from 02 to 04 from 04 to 08, from 08 to 10 and so on. That keeps track of the row that the scan is looking at, and as long as the result of the rotate is positive (bit 7 low), the loop will travel back to 03EB, where the next row will undergo the same testing as each previous one.

When RLC C shifts the row pointer to 3880, then bit 7 will be high (10000000); this is 'negative' in Z-80 architecture, and the loop falls through. Why does it fall through before checking the contents of address 3880? Because the only thing in this row is the shift key; it does not offer a decipherable code by itself, but merely modifies the information found when some other key is depressed. This explains why, among

other peculiarities, INKEY\$ does not acknowledge SHIFT alone.

When the loop falls through, the program encounters a RETURN from subroutine, which directs it immediately back to the rest of BASIC. The routine is remarkable, looping through just over 100 bytes when the keyboard is clear. Although not as time-efficient as obtaining input from a memory-mapped ASCII keyboard, it is quite speedy, and offers considerably better 'rollover' than many encoded keyboards.

When a key is pressed, the program jumps to 03FA, and is able to provide upper/lower case ASCII codes, special functions, and, incredibly enough, all of the 'missing' ASCII control codes (form feed, ring bell, etc.). Let us now follow it through:

03FA	5F	LD	E,A
03FB	7A	LD	A,0
03FC	07	RLCA	
03FD	07	RLCA	
03FE	07	RLCA	
03FF	57	LD	0,A

The position of the keystroke found has been stored in register E – recall that this is the 'column' of the keystroke. The row itself is not yet accessible, but the row counter (register D) is crucial to determining it. After E is saved, the accumulator is loaded with the value in this row counter, and rotated to the left three times. For those shaky in their binary arithmetic, this is the effect: if a decimal number is 045, a left rotation makes it 450. This is multiplication by ten. If a binary number is 010 (decimal 2), a left rotation gives 100 (decimal 4) . . . in other words, multiplication by two. Therefore, three left rotates gives us $2 \times 2 \times 2$, or multiplication by eight. That result is saved back in register D.

The purpose of this clever ploy will soon become clear:

0400	0E 01	LD	C,01
0402	79 ..	LD	A,C
0403	A3	AND	E
0404	20 05	JR	NZ,040B

Here the C register is set to 1, sucked up by the accumulator, and ANDed with E (remember E still contains that keystroke column byte). If the result is not zero (that is, if E equals 1), then the loop falls through and the program moves on. But have a look at what follows:

0406	14	INC	0
0407	CB 01	RLC	C
0409	1B F7	JR	0402

What is this about? Well, the D register, which contains 8 times the row value, is being incremented each time C is being rotated . . . making the lower three bits of D serve now as a column counter! Whoa, you say, back up there.

Okay, here it is: the original value in D could have been 0 through 6, depending on the row in use. When shifted three times, the possible values become 00, 08, 10, 18, 20, 28 and 30. Each of these possible values, when incremented through all seven possible columns, might contain 00 to 07, 08 to 0F, 10 to 17, etc., up to 37. This gives us a complete, distinct value to represent each key.

Now a fairly crude process of hunt-'n-peck begins. The status of the SHIFT key is checked, and set aside in register B. The demultiplexed keystroke value in register D is snapped back into the accumulator, and the comparisons take off:

040B	3A 80 3B	LD	A,(38B0)
040E	47	LD	B,A
040F	7A	LD	A,0

The character search can be followed through several branches; we will start with the most straightforward, and progress through some of the unique (and little publicized) aspects of the TRS-80 keyboard output.

The program adds 40 to the character value (address 0410), and checks if the result is greater than or equal to 60 (0412).

CHAR + 40 is	@	A	B	C	D	E	F	G
less than 60	H	I	J	K	L	M	N	O
(40 + 00 to 1F)	P	Q	R	S	T	U	V	W
	X	Y	Z					
CHAR + 40 is	0	1	2	3	4	5	6	7
60 or greater	B	9	;	*	;	<	=	>
(40 + 20 to 37)	ENT	CLR	BRK	UPA	ONA	LFA	RTA	SPC

If the compare finds a value less than 60, the routine rotates the SHIFT key value – which had been saved in the B register (0416). If SHIFT is released, the value in B is zero, and hence the rotate resets the carry flag (0418). The program moves directly to the terminal steps at 044B (to be discussed later). At this point, the character contained in A would be in the range 00+40 to 1F+40, the ASCII values for upper case (@, A-Z, left bracket, separator, right bracket, carat, and cursor). This is the software routine that causes the bizarre 'inverted' shift pattern on the TRS-80 . . . no shift for upper case!

If the character test at 0412 returns a value greater than or equal to 60, then 70 is subtracted (0429). No carry is generated if the test value was greater than or equal to 70, so this further separates the keyboard. See the diagram below:

60 to 6F minus 70	0	1	2	3	4	5	6	7
carries; result	B	9	;	*	;	<	=	>
is F0 to FF								
70 to 77 minus 70								
does not carry;	ENT	CLR	BRK	UPA	ONA	LFA	RTA	SPC
result is 00 to 07								

At address 043D, the value in the accumulator (00 to 07) is rotated left, producing the even values from 00 to 0E. The SHIFT byte in B is rotated right into the carry flag; if a carry is generated, the accumulator value is incremented (0442), providing the values 0+1, 2+1, 4+1 and so on – in other words, the odd values from 01 to 0F.

What follows is a classic example of machine language table look-up. HL is set to 0050, the address of the table in ROM; BC will be used as an offset, with B set to 0 and C made equal to A. When BC is added to HL, a resultant address (0050 to 005F) is produced, and the contents of that address is loaded up by the accumulator. Here is a look:

```

0443 21 50 00    LD    HL,0050
0446 4F         LD    C,A
0447 06 00      LD    B,0
0449 09         ADD   HL,BC
044A 7E         LD    A,(HL)
044B 57         LD    A,A

```

What do we find at 0050 to 005F? ASCII control codes. That result is stored in the D register (044B) before the termination sequence.

Table 2 - 9

Address	Contents	TRS-80 Action	ASCII Description	Keyboard Entry
0050	00	Carriage Ret.	Carriage Ret.	ENTER
0051	00	Carriage Ret.	Carriage Ret.	SHIFT ENTER
0052	1F	Clear Screen	Unit Separator	CLEAR
0053	1F	Clear Screen	Unit Separator	SHIFT CLEAR
0054	01	Break	Start of Heading	BREAK
0055	01	Break	Start of Heading	SHIFT BREAK
0056	5B	Up Arrow	Left Bracket	Up Arrow
0057	1B	Edit Escape	Escape	SHIFT Up Arrow
0058	0A	Line Feed	Line Feed	Down Arrow
0059	1A (00)	*See text	Substitute	SHIFT Down Arrow
005A	0B	Backspace	Backspace	Left Arrow
005B	1B	Backspace Line	Cancel	SHIFT Left Arrow
005C	09	Horizontal Tab	Horizontal Tab	Right Arrow
005D	19	32-Char. Mode	End of Medium	SHIFT Right Arrow
005E	20	Space	Space	Space
005F	20	Space	Space	SHIFT Space
00100	; *****			
00110	; SIMPLE ACTIVE KEYBOARD DISPLAY ROUTINE TO SHOW THE USE			
00120	; AND PHANTOMING OF KEYS AS THEY ARE PRESSED BY THE USER			
00130	; *****			
00140	; DENNIS BATHORY KITSZ, ROXBURY, VERMONT 05669			
00150	; *****			
00160	; *****			
5000	00170	ORG	5000H	
	00180	; *****		
	00190	; *****		
	00200	; CLEAR SCREEN, DISABLE INTERRUPTS, CLEAR ACCUMULATOR,		
	00210	; GET "0" CHARACTER, SCREEN POSITIONS, SCREEN OFFSET,		
	00220	; AND BEGINNING OF KEYBOARD FOR TEST; ALL ARE SET UP ONCE		
	00230	; *****		
	00240	; *****		
5000	00250	ENTER	CALL	01C9H ; CLEAR SCREEN TO START
5003	00260	F3	OI	; KILL ALL THEM BOTHERS
5004	00270	AF	XOR	A ; CLEAR ALL FLAGS
5005	00280	0E30	LD	C,30H ; THIS IS THE "0" CHAR.
5007	00290	211030	LD	HL,3010H ; ADDRESS NEAR SCREEN CTR
500A	00300	112400	LD	DE,0024H ; DISTANCE BETWEEN LINES
5000	00310	05	PUSH	DE ; SAVE LINE ADDRESS VALUE
500E	00320	11013B	LD	DE,3B01H ; DEFINE FIRST KEYBOARD
5011	00330	0608	LD	B,0BH ; NUMBER OF LOOPS TO 00
	00340	; *****		
	00350	; *****		
	00360	; TEST EACH ROW OF THE KEYBOARD, COLUMN BY COLUMN		
	00370	; *****		
	00380	; *****		
5013	00390	1A	LOOP	LD A,[0E] ; FIRST ROW OF KEYBOARD
5014	00400	CB47	BIT	0,A ; CHECK FIRST KEY COLUMN

Alright, we have upper case ASCII and TRS-80 control functions. How about the rest? Back up now to the test for SHIFT, at 0416. If such a shift is present, the value in A (40 to 5F) is increased by 20 (60 to 7F). These are the ASCII codes for lower case (@, a-z, left brace, separator, right brace, delete). The resultant code, as usual, is saved in the D register.

But what follows is curious:

```

0410 3A 40 3B    LD    A,(3B40)
0420 E6 10      AND    10
0422 2B 2B      JR     Z,044C

```

The keyboard is tested again, this time at row 3840, data position 10 – the down arrow. If that key is not depressed, the program skitters right to the termination routine at 044C, with the lower case ASCII code ensconced in the D register.

Why the SHIFT/down arrow combination? If the down arrow is depressed, the value in D is retrieved and placed in the accumulator (60 to 7F), then reduced by 60, becoming . . . aha! . . . 00 to 1F. The program jumps to the end sequence, with the accumulator clutching one of the complete set of 32 ASCII control codes!

(There is an anomaly in earlier Level II ROMs: the code for the down arrow at 0059 is returned before the control code. Later ROMs placed a 00 at 0059, resulting only in the return of a control code if SHIFT/down arrow was depressed.)

So where are we now? Upper and lower case, TRS-80 and ASCII control codes. We need numbers and figures, and so we shall have them. Recall the second diagram: at 042B, the command row was separated from the numbers, which were left at F0 to FF. At 042D, 40 is added, resulting in possible values of 30 to 3F. A further separation is made via a comparison with 3C:

```

.....
.      30 to 3B      0  1  2  3  4  5  6  7  .
. {CP 3C result}  B  9  :  ;  ,<  =  .>  /?  .
.....

```

If the comparison is less than 3C, a carry is generated. The usual SHIFT test is made (at 0435), and if it fails, the value in A (30 to 3B) is maintained as the program moves into the end routine. These are the ASCII codes for numbers 0 to 9, colon and semicolon.

If the test value is 3C, 3D, 3E or 3F, no carry would be generated at 042F, and these values are XORed with 10. This toggles the high nibble

On the Keyboard Scan

```

5016 71 00410 LO (HL),C ; FIRST DISPLAY A "0"
5017 2801 00420 JR Z,JUMP1 ; DON'T CHANGE IF NO KEY
5019 34 00430 INC (HL) ; MAKE IT A "1" IF A KEY
501A 23 00440 JUMP1 INC HL ; NEXT SCREEN LOCATION
501B 23 00450 INC HL ; ...PLUS TWO
501C 23 00460 INC HL ; ...PLUS THREE
501D 23 00470 INC HL ; ...PLUS FOUR
501E C84F 00480 BIT 1,A ; SECOND KEYBOARD COLUMN
5020 71 00490 LO (HL),C ; FIRST DISPLAY A "0"
5021 2801 00500 JR Z,JUMP2 ; DON'T CHANGE IF NO KEY
5023 34 00510 INC (HL) ; MAKE IT A "1" IF A KEY
5024 23 00520 JUMP2 INC HL ; NEXT SCREEN LOCATION
5025 23 00530 INC HL ; ...PLUS TWO
5026 23 00540 INC HL ; ...PLUS THREE
5027 23 00550 INC HL ; ...PLUS FOUR
5028 C857 00560 BIT 2,A ; THIRD KEYBOARD COLUMN
502A 71 00570 LO (HL),C ; FIRST DISPLAY A "0"
502B 2801 00580 JR Z,JUMP3 ; DON'T CHANGE IF NO KEY
502D 34 00590 INC (HL) ; MAKE IT A "1" IF A KEY
502E 23 00600 JUMP3 INC HL ; NEXT SCREEN LOCATION
502F 23 00610 INC HL ; ...PLUS TWO
5030 23 00620 INC HL ; ...PLUS THREE
5031 23 00630 INC HL ; ...PLUS FOUR
5032 C85F 00640 BIT 3,A ; FOURTH KEYBOARD COLUMN
5034 71 00650 LO (HL),C ; FIRST DISPLAY A "0"
5035 2801 00660 JR Z,JUMP4 ; DON'T CHANGE IF NO KEY
5037 34 00670 INC (HL) ; MAKE IT A "1" IF A KEY
5038 23 00680 JUMP4 INC HL ; NEXT SCREEN LOCATION
5039 23 00690 INC HL ; ...PLUS TWO
503A 23 00700 INC HL ; ...PLUS THREE
503B 23 00710 INC HL ; ...PLUS FOUR
503C C867 00720 BIT 4,A ; FIFTH KEYBOARD COLUMN
503E 71 00730 LO (HL),C ; FIRST DISPLAY A "0"
503F 2801 00740 JR Z,JUMP5 ; DON'T CHANGE IF NO KEY
5041 34 00750 INC (HL) ; MAKE IT A "1" IF A KEY
5042 23 00760 JUMP5 INC HL ; NEXT SCREEN LOCATION
5043 23 00770 INC HL ; ...PLUS TWO
5044 23 00780 INC HL ; ...PLUS THREE
5045 23 00790 INC HL ; ...PLUS FOUR
5046 C86F 00800 BIT 5,A ; SIXTH KEYBOARD COLUMN
5048 71 00810 LO (HL),C ; FIRST DISPLAY A "0"
5049 2801 00820 JR Z,JUMP6 ; DON'T CHANGE IF NO KEY
504B 34 00830 INC (HL) ; MAKE IT A "1" IF A KEY
504C 23 00840 JUMP6 INC HL ; NEXT SCREEN LOCATION
504D 23 00850 INC HL ; ...PLUS TWO
504E 23 00860 INC HL ; ...PLUS THREE
504F 23 00870 INC HL ; ...PLUS FOUR
5050 C877 00880 BIT 6,A ; SEVENTH KEYBOARD COLUMN
5052 71 00890 LO (HL),C ; FIRST DISPLAY A "0"
5053 2801 00900 JR Z,JUMP7 ; DON'T CHANGE IF NO KEY
5055 34 00910 INC (HL) ; MAKE IT A "1" IF A KEY
5056 23 00920 JUMP7 INC HL ; NEXT SCREEN LOCATION
5057 23 00930 INC HL ; ...PLUS TWO
5058 23 00940 INC HL ; ...PLUS THREE
5059 23 00950 INC HL ; ...PLUS FOUR
505A C87F 00960 BIT 7,A ; EIGHTH KEYBOARD COLUMN
505C 71 00970 LO (HL),C ; FIRST DISPLAY A "0"
505D 2801 00980 JR Z,JUMP8 ; DON'T CHANGE IF NO KEY
505F 34 00990 INC (HL) ; MAKE IT A "1" IF A KEY
5060 05 01000 JUMP8 PUSH OE ; SAVE THIS VALUE
5061 00E1 01010 POP IX ; PUT IT IN IX FOR A BIT
5063 01 01020 POP OE ; GET ORIGINAL OE VALUE
5064 19 01030 ADD HL,OE ; NOW START NEXT LINE
5065 05 01040 PUSH OE ; SAVE SAME VALUE AGAIN
5066 00E5 01050 PUSH IX ; STASH IT BRIEFLY
506B 01 01060 POP OE ; AND BACK INTO OE INTACT
5069 C813 01070 RL E ; DEFINE NEXT KEYBOARD ROW
506B 10A6 01080 OJNZ LOOP ; DO IT FOR EIGHT ROWS
01090 ;
01100 ; #####
01110 ; CLEAR UP POINTERS AND DELAY SO SCREEN DOES NOT JITTER
01120 ; #####
01130 ;
5060 01 01140 POP DE ; CLEAR THE STACK
506E 010008 01150 LO 8C,800H ; DELAY VALUE
5071 C06000 01160 CALL 0060H ; DELAY SUBROUTINE IN ROM
5074 186F 01170 JR START ; START THE ROUTINE AGAIN
01180 ;
01190 ; #####
5000 01200 END ENTER ; BEGIN IT ALL HERE
00000 TOTAL ERRORS
31044 TEXT AREA BYTES LEFT

ENTER 5000 00250 01200
JUMP1 501A 00440 00420
JUMP2 5024 00520 00500
JUMP3 502E 00600 00580
JUMP4 5038 00680 00660
JUMP5 5042 00760 00740
JUMP6 504C 00840 00820
JUMP7 5056 00920 00900
JUMP8 5060 01000 00980
LOOP 5013 00390 01080
START 5005 00280 01170

```

from 3 to 2, resulting in values from 2C to 2F (, - . /). If a shift key was noted at 0437, the same toggle procedure is followed, changing values 30 to 3B into 20 to 2B (these would become space ! " # \$ () * = etc.).

0420	C6 40	ADD	A,40
042F	FE 3C	CP	3C
0431	38 02	JR	C,0435
0433	EE 10	XOR	10
0435	C8 08	RRC	8
0437	30 12	JR	NC,044B
0439	EE 10	XOR	10
043B	1B 0E	JR	044B

Thus, the coding is complete: control codes (00 to 1F), punctuation (20 to 2F), numbers and figures (30 to 3F), upper case (40 to 5F) and lower case (60 to 7F). Just as an aside, the terms lower and upper case are sometimes written small and large case; old-time printers would chuckle at that. The case referred to is a printers case, which, when two were stacked one above the other, contained the capital and small letters. Thought you might like to know that.

Back to the routine, starting at the termination sequence (044C); the decoded character is saved in D, and that is the only information we need to preserve, since the bulk of the work is done.

044C	01 AC 00	LD	8C,00AC
044F	C0 60 00	CALL	0060
0452	7A	LD	A,D

A delay at 0060 is called, which was intended to wait through the bounce present with normal mechanical switch contacts – but the easily dirtied switches on the TRS-80 are abnormal! This delay uses the accumulator, and when it is free, the value in D is restored to it. This value is compared to 01 (the BREAK code), and returns directly to the main routine (0455) with any code other than BREAK.

If BREAK is discovered, the program executes a call to 0028 (RST 28), returning to Level II.

The routine is quite efficient, and is capable of returning 128 different values at a rate of better than 100 per second – ten times the speed of the world's fastest typist!

3

Software Modifications

Software makes the computer. With that in mind, it's not hard to understand the popularity of the TRS-80. Its BASIC is simple to use and immediately accessible. It reports back errors and provides clear screens and graphics with easy commands. At first, it was hard to imagine how such an elegant BASIC could be improved.

Such illusions could not last long, especially when weakly designed hardware started to exhibit keybounce, when machine language software was so fast but so difficult to access, when putting a program aside meant losing all variables, and so forth.

In this Chapter, several simple but very important software modifications will be presented :

Keyboard debounce with repeating keys, audible beep tone, and an upper/lower case driver.

Two methods of intercepting the BASIC interpreter in order to create your own commands.

Packing machine language programs in simple BASIC strings, where they can be moved about and accessed easily.

Sound and sound-effects generation routines.

Creating somewhat unlistable BASIC programs.

Auto-execution of SYSTEM programs, including an auto-load BASIC module.

A simple machine language monitor accessible directly from BASIC.

BASIC is not an incomprehensible, immutable, indivisible whole, but rather a pliant, carefully sewn, patchwork quilt of useful subroutines. These routines are accessed singly or as groups, not only whenever one of this BASIC's English-language commands is entered, but even while waiting for commands to be entered or programs to be run.

Chapter 1 contained an overview of the structures that BASIC is composed of; below is a more detailed look at the building blocks out of which Level II is created.

1. A power-up sequence: including preparation of blocks of reserved memory, clearing of internal hardware systems, and total memory examination.
2. Numeric conversion routines: single precision to integer and vice versa, numeric to string and vice versa, assignment of levels of precision.
3. Simple arithmetic operations: addition, subtraction, multiplication, division, comparison, and raising to a power (exponentiation); integer, single, and double-precision calculations.

4. Mathematical functions: activities based on the sine function, as well as logarithms, square roots, absolute value and truncation, and random number generation.

5. String operations: concatenation and truncation of strings, direct keyboard scan conversions (INKEY\$), alphabetic comparisons.

6. Variables: assignment of numeric, string, and array variable names, variable assignments (LET), updating of values, searching for variable names. Partly integrated with NEW routines.

7. Keyboard input: polling of the keyboard matrix, conversion to characters, searching for carriage return, building a keyboard buffer.

8. Cassette input/output: motor relay control, assembly of parallel data into serial form, timing of output pulses, timing of input, reassembly of data from serial to parallel form.

9. Video input/output and display management: screen clear, scrolling, tabbing, character display including line feed, carriage return, backspace, set/reset, POS, POINT, cursor control, characters per line.

10. Printer control: lines per page, top of form, output of characters, waiting for handshake.

11. A command interpreter for organizing the entry points and order of chosen subroutines.

12. Error reporting routines.

13. Program line management routines. Partly integrated with NEW routines.

14. Editing functions: Insert, delete, kill, exit, etc. Integrated in part with program line management routines.

15. Run-time management: Integrated with most of the above functions, but including subroutine handling, loop handling, etc.

A more complete rundown on the TRS-80 Level II ROM memory map can be found in *Inside Level II, Supermap, Microsoft BASIC Decoded*, and *TRS-80 Disassembled Handbook* (see Appendix II for details).

Listing 3-1. Custom BASIC interpreter patch routine.

```

00100 : #####
00110 : DEBOUNCE, AUTO-REPEAT, AND BEEP KEYBOARD PATCH ROUTINE.
00120 : THIS ROUTINE AS WRITTEN PATCHES INTO 4016. TO USE
00130 : OTHER DRIVERS SUCH AS LOWERCASE, GRAPHICS, ETC., THE
00140 : RETURN ADDRESS (NORMALLY 03E3) MUST BE REPLACED BY THE
00150 : ADDRESS OF THE OTHER DRIVER. ALSO, TO USE THIS ROUTINE
00160 : WITH THE CUSTOM INTERPRETER /ON AND /OUT COHHANOS BY
00170 : PATCHING THIS DRIVER IN PLACE WITH /ON AND THE NORHAL
00180 : (OR UPPER/LOWER, ETC.) DRIVER IN PLACE WITH /OUT.
00190 : #####
00200 :
4016 0030 00210      ORG      4016H      ; KEYBOARD SCAN PATCH
00220      DEFW      ENTREE      ; START OF DEBOUNCE ROUT.
00230      ORG      3000H      ; NOTE THAT THIS UTILITY
00240      ; IS CURRENTLY SET UP FOR
00250      ; USE WITH A MEMORY ADDRESS
00260      ; AT 3000H. IT CAN BE
00270      ; RE-ORGED AT ANY LOC'N.
06CC 00280 BASIC2 EQU 06CCH      ; BASIC "READY" DISPLAY
4036 00280 KEYHLO EQU 4036H      ; NORHAL KEYSTROKE STORE
3801 00300 KEYBRO EQU 3801H      ; FIRST KEYBOARD ADDRESS
401A 00310 HOLOER EQU 401AH      ; RESERVED BYTE FOR DELAY
4099 00320 INKEYS EQU 4099H      ; INKEY$ STORAGE BYTE
0060 00330 DELAYS EQU 0060H      ; ROW DELAY SUBROUTINE
00340 :
00350 : #####
00360 : ROUTINE BEGINS HERE, PATCHING ITSELF INTO PLACE AT 4016
00370 : #####
3000 ES 00380 :
00390 ENTREE PUSH HL      ; SAVE HL IF TRANSPARENT
3001 210930 00400 LO HL,START ; GET START OF ROUTINE
3004 221640 00410 LO [4016H],HL ; PUT INTO KEYBOARD PATCH
3007 E1 00420 POP HL      ; RESTORE HL VALUE
3008 C3CC06 00430 JP BASIC2 ; AND GO BACK TO "READY"
00440 :
00450 : #####
00460 : THIS AREA MAKES THE PRELIMINARY CHECK OF KEYBOARD ROWS
00470 : #####
00480 :
3008 213640 00490 START LO HL,KEYHLO ; SET UP STORAGE AREA
300E 010138 00500 LO BC,KEYBRO ; SET UP FIRST KBD ROW
3011 1600 00510 LO 0,0 ; SET UP COUNTER OF ROWS
3013 0A 00520 CHKKEY LO A,(BC) ; FIND IF A KEY PRESSED
3014 SF 00530 LO E,A ; SAVE VALUE IN E REG.
3015 A3 00540 ANO E ; TEST IF KEY WAS PRESSED
3016 2018 00550 JR NZ,CKPREV ; IF YES, SEE IF SAME ONE
3018 77 00560 LO [HL],A ; SAVE VALUE IN STORAGE
3019 14 00570 INCD INC 0 ; INCREMENT ROW COUNTER
301A 2C 00580 INC L ; INCREMENT STORAGE AREA
3018 C801 00590 RLC C ; SHIFT TO NEXT KBD ROW
3010 79 00600 LO A,C ; GET VALUE OF KBD ROW
301E 0680 00610 SUB 80H ; CHECK IF SHIFT KEY ROW
3020 20F1 00620 JR NZ,CHKKEY ; IF NOT THEN CONTINUE
00630 :
00640 : #####

```

Listing Continued . .

Sophisticated Debouncing

The first few hundred thousand TRS-80's were afflicted with serious keybounce problems – the appearance of double letters when only a single letter was typed. Full-scale preventive maintenance is presented in the following chapter; but there are software solutions as well. If debounce were the only criterion, though, maintenance would be the ideal solution.

But the software designers made no provision for repeating keys, nor did the hardware designers include access to lower case characters. Furthermore, the silent keyboard remains a frustration to touch-typists and others who do not refer constantly to the screen for feedback.

Thus, some sort of audible reinforcement (as with the Apple's entry-error beep) would be a thoughtful addition to the keying process.

Listing 3-1 is a complete debounce, audible beep, key repeat, and upper/lower case driver routine. The program is written as three independent subroutines, each of which may be disabled or removed before assembly.

Continued Listing

```

00660 ; CHECKING IS DONE - NOW SEE IF PREVIOUS KEYS HELD DOWN
00660 ; #####
00670 ;
3022 0607 00680      LO      B,7      ; ELSE GET NUMBER OF ROWS
3024 20    00690 DECL    DEC      L      ; MOVE BACK THRU STORAGE
3025 86    00700      ADD     A,(HL)   ; AND MAKE TOTAL OF KEYS
3026 10FC  00710      DJNZ   DECL    ; AND DO IT FOR ALL ROWS
3028 A7    00720      AND     A      ; TEST IF ANY KEYS STORED
3029 3E00  00730      LO      A,0    ; A=0, FLAGS REMAIN SAHE
302B 00    00740      RET     NZ     ; BACK IF KEY IN STORAGE
302C 321A40 00750      LU      (HOLDER).A ; SAVE NEW VALUE IN CTR.
302F C9    00760      RET     ; AND BACK TO MAIN ROUT.
00770 ;
00780 ; #####
00790 ; NEXT TEST IS FOR STATUS OF INKEYS, IF IT IS IN USE
00800 ; #####
00810 ;
3030 A6    00820 CKPREV AND     (HL)   ; SEE IF VALUE IS SAHE
3031 2B1F  00830      JR      Z,STORE ; STORE VALUE IF ZERO
3033 3A9940 00840      LO      A,(INKEYS) ; FIND VALUE AT INKEYS
3036 A7    00850      AND     A      ; SEE IF SOMETHING THERE
3037 20E0  00860      JR      NZ,INCO ; IF SO THEN GO AWAY
3039 3A1A40 00870      LO      A,(HOLDER) ; GET DELAY COUNTER VALUE
303C 3C    00880      INC     A      ; INCREMENT THE COUNTER
303D 321A40 00890      LO      (HOLDER).A ; AND SAVE VALUE BACK
3040 FEFF  00900      CP      OFFH    ; IS COUNTER AT END YET?
3042 2B09  00910      JR      Z,OECA  ; IF SO, THEN REPEAT
00920 ;
00930 ; #####
00940 ; REPEATING-KEY TIME-WASTE VALUE (FF) MAY BE VARIED
00950 ; #####
00960 ;
3044 CS    00970      PUSH    BC      ; SAVE BC FOR LATER
3045 06FF  00980      LO      B,OFFH  ; GET DELAY VALUE
3047 00    00990      NOP     ; WASTE SOME TIME
3048 10F0  01000      DJNZ   THWSTE  ; AND DO IT FF TIMES
304A C1    01010      POP     BC      ; RESTORE BC VALUE
304B 1BCC  01020      JR      INCO    ; AND GO BACK TO SCANNING
304D 321A40 01030 OECA    DEC     A      ; MAKE A BECOME FE
304E 321A40 01040      LO      (HOLDER).A ; AND SAVE IT IN DELAY
3051 78    01050      LO      A,E    ; GET KEYSTROKE FOUND
3052 73    01060 STORE   LO      (HL),E ; AND PUT IT IN STORAGE
01070 ;
01080 ; #####
01090 ; DEBOUNCE BELOW MAY BE ELIMINATED BECAUSE BEEP USES TIME
01100 ; #####
01110 ;
3053 CS    01120      PUSH    BC      ; SAVE VALUE IN BC
3054 010002 01130      LO      BC,200H ; GET DEBOUNCE DELAY
3057 060000 01140      CALL   DELAYS  ; AND CALL ROM DELAY
305A C1    01150      POP     BC      ; AND GET VALUE TO BC
305B 0A    01160      LO      A,(BC) ; GET VALUE AT KEYBOARD
305C A3    01170      AND     E      ; AND TEST IF IT'S THERE
305D CB    01180      RET     Z      ; IF NOT IT WAS BOUNCE
01190 ;
01200 ; #####
01210 ; BEEP ROUTINE PRODUCES VERY SOFT (NOT ANNOYING) SOUND
01220 ; #####
01230 ;
305E CS    01240      PUSH    BC      ; ELSE SAVE THE VALUE
305F ES    01250      PUSH    HL      ; AND SAVE THE LOCATION
3060 FS    01260      PUSH    AF      ; AND SAVE THE KEYSTROKE
3061 0640  01270      LO      B,40H  ; AND KEY BEEP DURATION
3063 3A3040 01280      LO      A,(4030H) ; AND GET SCREEN STATUS
3066 E6F0  01290      AND     0F0H    ; AND MASK OUT ALL BITS
306B 67    01300      LD      H,A    ; SAVE WAVE "0" MASK IN H
3069 F602  01310      OR      2      ; CREATE A WAVE "1" MASK
306B 6F    01320      LD      L,A    ; SAVE WAVE "1" MASK IN L
306C 70    01330      LD      A,L    ; GET THE WAVE "1" MASK
306D 03FF  01340      OUT     (OFFH),A ; AND CREATE WAVEFORM 1
306F 7C    01350      LD      A,H    ; GET THE WAVE "0" MASK
3070 03FF  01360      OUT     (OFFH),A ; AND CREATE WAVEFORM 0
3072 CS    01370      PUSH    BC      ; SAVE THE DURATION VALUE
3073 0640  01380      LD      B,40H  ; GET THE PITCH VALUE
3075 10FE  01390      DJNZ   $+0    ; AND WAIT THRU WAVEFORM
3077 C1    01400      POP     BC      ; RESTORE THE DURATION
307B 10F2  01410      DJNZ   LOOP    ; AND DO FOR FULL BEEP
307A F1    01420      POP     AF      ; RESTORE KEYSTROKE VALUE
307B E1    01430      POP     HL      ; RESTORE STORAGE VALUE
307C C1    01440      POP     BC      ; RESTORE COUNTER VALUE
307D C3FB03 01450      JP      03FBH  ; AND RETURN TO KEYS CAN
01460 ;
01470 ; #####
06CC      01480      END      BASIC2
BASIC2 06CC 00280 00430 01480
CHKKEY 3013 00520 00620
CKPREV 3030 00820 00550
OECA 3040 01030 00910
OECL 3024 00690 00710
OELAYS 0060 00530 01140
ENTREE 3000 00390 00220
HOLDER 401A 00310 00750 00870 01040
INCO 3019 00570 00060 01020
INKEYS 4099 00320 00840
KEYBRO 3801 00300 00500
KEYHLO 4036 00290 00490
LOUP 306C 01330 01410
START 300B 00490 00400
STORE 3052 01060 00830
TMWSTE 3047 00990 01000

```

This routine patches into the keyboard control block driver address at 4016, leading the program to its own entry point instead of 03E3 (see the Supplement to Chapter 2 for details on the operation of the TRS keyboard scan).

At START, the keyboard scan begins with HL pointing to the first position in a keystroke storage buffer (4036). BC points to the first keyboard row (3801). The program proceeds similarly to the normal Level II scan, except that a location (401A) has been set aside to 'count down' the time a key remains pressed. If any key or combination of keys remain pressed for the duration of the loop, the character (normally rejected by Level II's rollover capabilities) is accepted again. This is the start of the repeating process. The INKEY\$ storage area is checked (so that programs using INKEY\$ are not delayed by an unusable character acceptance), and a short debounce-delay loop is entered.

If a legitimate key is found, a different debounce-delay loop is entered, and a rapidly-fluctuating one-zero pattern is sent to the cassette port (FF). This sounds as a beep if an amplifier and speaker are connected. The routine can be exited before the beep, so only the debounce-repeat options are present; it can also be exited after the beep, returning to the main routine with the keystroke.

The final portion of the program is an upper/lower case driver program. This driver is irrelevant unless you have a lower case modification in place, and should be disabled (or not assembled) if you have not made the modification. It merely strips the conversion to upper case normally made by the Level II software, and returns to the running program with the actual key depressed instead of an upper case converted version.

Upper/Lower Case Driver

When the TRS-80 keyboard is used, all characters are automatically converted to upper case before being displayed. The keyboard itself, however, returns a full upper/lower case value (albeit inverted - shift for lower case) to the display routine. The display routine then sends this information to the screen. The screen always displays upper case because the hardware to provide lower case was not a part of the TRS-80 as sold. The addition of a single integrated circuit (see Chapter 4) provides this access.

The lower-case driver patches into the display routine just as the character to be displayed is returned in the accumulator. Control is taken

```

00100 : #####
00110 : SIMPLE LOWER-CASE DRIVER ROUTINE FOR BASIC AND 00S
00120 : #####
00130 :
7F00 00140 : ORG 07F00H : NEAR TOP OF BASIC
00150 :
00160 : #####
00170 : GET CONTROL BLOCK FOR SCREEN AND CHANGE CHARACTERS
00180 : #####
00190 :
7F00 006E03 00200 LCORIV LO L,(IX+3) : IX POINTS TO DEVICE
7F03 006604 00210 LO H,(IX+4) : CONTROL BLOCK (VIDEO)
7F06 0A9A04 00220 JP C,049AH : BACK TO SCREEN DRIVER
7F09 007E05 00230 LO A,(IX+5) : GET CURSOR CHARACTER
7F0C B7 00240 OR A : CHECK IF CURSOR IS ON
7F00 2801 00250 JR Z,GETCHR : GET ONE IF CURSOR OFF
7F0F 77 00260 LD (HL),A : PUT CURSOR INTO POSN.
7F10 79 00270 GETCHR LD A,C : GET CHARACTER TO SHOW
7F11 FE20 00280 CP 20H : SEE IF A CONTROL CODE
7F13 0A0605 00290 JP C,0506H : BACK TO DRIVER IF C.C.
7F16 FE80 00300 CP 80H : SEE IF GRAPHIC CHAR.
7F18 02A604 00310 JP NC,04A6H : BACK TO DRIVER IF SO
7F18 FE58 00320 CP 05BH : CHECK UPPER/LOWER CASE
7F10 300B 00330 JR NC,CHECK1 : IF >5B, CHECK FURTHER
7F1F FE40 00340 CP 40H : CHECK UPPER CASE
7F21 380E 00350 JR C,GOAWAY : IF <40, CHECK NO FURTHER
7F23 C620 00360 A00 A,20H : IF 40-5B, MAKE UPPER
7F25 180A 00370 JR GOAWAY : DONE - BACK TO DRIVER
7F27 FE7B 00380 CHECK1 CP 7BH : SEE IF ALPHABETIC
7F29 3006 00390 JR NC,GOAWAY : NO FURTHER IF NOT ALPHA
7F2B FE60 00400 CP 60H : SEE IF ALPHABETIC
7F20 3802 00410 JR C,GOAWAY : NO FURTHER IF NOT ALPHA
7F2F 0620 00420 SUB 20H : PLAY SWITCH TO LOWER
7F31 C37004 00430 GOAWAY JP 0470H : OUT TO DRIVER NOW
00440 :
00450 : #####
00460 : PUT VIDEO PATCH INTO PLACE UPON LOAD
00470 : #####
00480 :
401E 00490 : ORG 401EH : THIS IS VIDEO PATCH
401E 007F 00500 OEFW LCORIV : PUT LCORIV ROUTINE IN
00510 :
00520 : #####
06CC 00530 : END 06CCH : BACK TO BASIC READY
00000 TOTAL ERRORS

```

Listing 3-13. Upper/lower case driver.

from the convert-to-upper-case display function in ROM. Ideally, this ROM routine could just be entered after its convert-to-upper case code; unfortunately, this would result in the famous inverted display . . . normal upper case, shifted lower case.

To avoid this, the character is tested and converted to its proper case before being returned to the ongoing display driver routine in ROM. Notice something interesting: when programs are listed with this driver, the letters appear in lower case. That is because when the programs are entered, they are in fact being entered with the keyboard *unshifted*. Because this can be a bit disconcerting (and also quite illegible, since we all are used to upper case lists), an upper case on/off software patch is provided.

be a 'label' (such as VIDEO). Once VIDEO has been defined as 3C00 to the Editor/Assembler, it will always interpret the label as the number that was assigned to that label.

Line numbers are provided to keep things in order and to insert or edit pieces of code, and there is space on every line for comments.

Load the Editor/Assembler tape under the SYSTEM command. Its name is EDTASM. When the loading is complete, enter a slash (/), and you will be presented with the prompt:

```

TRS-80 EDITOR/ASSEMBLER 1.1
*_

```

This is EDTASM's equivalent of the BASIC prompt:

```

RADIO SHACK LEVEL II BASIC
READY
>_

```

You are being asked for input. Unlike BASIC, EDTASM has only a few commands. They are (in the order you are likely to use them):

I
This command inserts numbered program lines almost exactly like the BASIC command AUTO. When I is entered alone, numbering starts with line 100 in increments of 10 line numbers. On the other hand, I15,15 will start with line 15 in increments of 15.

P
The equivalent of a list. A single P lists the next sixteen lines of the program. P10:100 lists 10 to 100. P# is the first line, P. is the current line, and P* is the last line.

N
Here is the renumber command. All lines are

Using The Editor/Assembler

The Editor/Assembler is one of the most powerful tools available to the TRS-80 customizer. It is a fast, high-level compiler which produces a block of Z-80 machine code. Its job is to provide an easily accessible substitute for the tedious creation of bytes of Z-80 coded information.

The Z-80 microprocessor is capable of responding to many hundreds of combinations of ones and zeros. Each pattern causes the Z-80 to follow a unique pattern of electronic activity, and many thousands of those activities in concert create a sophisticated language like BASIC.

Using these patterns can be very tricky and time-consuming. Long ago, computer designers learned that it was easier to remember an action like 'load the accumulator with the contents of byte counter register' as 'Load A with B', abbreviated LD A,B. This is much handier than trying to recall 01111000. These abbreviations are called mnemonics, which are what you will find in all the program listings in this book.

You will also find that, instead of specific locations in memory (such as 3C00), there may

automatically renumbered in increments of 10 starting with line 100. Again, specific lines and increments may be specified: N300,50 will renumber all lines in increments of 50, with the first line being 300.

L

Loads a source tape, but not an object tape. Up to a six-character name may be specified.

W

Writes a source tape (the program listing). Up to a six-character name may be specified.

D

Deletes the specified line or lines. D# deletes the first line, D. deletes the current line, and D* deletes the last line. Groups of lines are specified with a colon, as D40:170 or D#:90.

E

The edit function. The pound (#), period (.), and asterisk (*) represent the first, current, and last lines. A line number (as E400) may be specified. The editing functions are identical to BASIC's editing functions – except that characters to be deleted are *not* delimited by exclamation points.

R

Replaces the indicated line. The line number is presented, and new information may be entered.

F

This command finds a text string. It is not followed by a space. To search for the phrase 'ENTRY', type FENTRY (ENTER). The entire line containing the phrase will be printed. To find the next identical phrase, merely type F (ENTER).

H

This sends the source listing, unassembled, to the printer. The complete source listing, including line numbers, is printed. As usual, (#), (.) and (*) may be used to indicate first, current, and last line, and groups of lines may be printed (as H55:3000).

T

The poor person's text editor. The source code is sent to the printer without line numbers. Thus, text may be entered a line at a time, and the numberless result printed. The same functions provided with H are available.

B

The exit to BASIC. As sold, EDTASM returns only to MEMORY SIZE?, and all programs and information, including EDTASM itself, are lost. Patches are available to re-route this exit.

A

This command directs EDTASM to compile your source code into object code, make a list of all the labels (symbols) used, and check for errors. The A command may be followed by a six-letter name, as well as the 'switches' /NL (no listing), /NS (no symbol table), /NO (no object code), or /WE (wait upon error). The switches may be used in any combination, and are useful in shaking the errors out of an assembly program.

Lines are always entered into EDTASM under the I (insert) or R (replace) commands. A line number is presented, so:

```
00010*_
```

Several columns are then available, consisting respectively of an optional label, the mnemonic instruction, the 'operand' (if any), and any comments (always following a semicolon). A complete group of lines would look like this:

```
00010 VIDEO EQU 3C00H ;SCREEN TOP
00020 ORG 5000H ;START PROGRAM
00030 ENTRY LO A,B ;GET B INTO A
00040 LO HL,VIDEO ;HL AT SCREEN
```

This excerpt gives this information: the label VIDEO is an 'equate' (is defined as) location 3C00. The program starts (has its origin – ORG) at 5000. The label ENTRY is assigned to the start of the program, and that program's first action is to load the accumulator with register B. Next, the HL register is pointed to VIDEO (3C00), the start of the screen memory.

When told to assemble this (using the A command), the results will look like the following:

```
3C00 00010 VIDEO EQU 3C00H ;SCREEN TOP
5000 00020 ORG 5000H ;START PROGRAM
5000 78 00030 ENTRY LO A,B ;GET B INTO A
5001 21003C 00040 LO HL,VIDEO ;HL AT SCREEN
```

The EDTASM program evaluated all the information in the source code and created the columns at the left. The first column specifies the current address, and the second column specifies the machine language code, if any, for that particular line. Note the correct assignment of 3C00 to VIDEO in line 00040. 78 is the machine code for LD A,B and 21 is the machine code for LD HL,NNNN. In this case, NNNN is VIDEO is 3C00.

For detailed instructions, refer to the EDTASM instruction manual. One thing to note: you can conserve source code memory space by using the right arrow (tab) instead of spacing between program lines, labels, commands, operands, and comments. Each tab is a single character, but spaces are counted separately.

Patching the BASIC Interpreter

Patching the BASIC Interpreter

Each time a BASIC command is entered or a program line is being run, a section of ROM evaluates each of these commands in order, jumping to internal subroutines that will produce the desired result. This section of ROM is called the interpreter, an area which translates the commands into program action.

At address 4003, the machine language instruction C3 78 1D can be found, which means 'jump to address 1D78'. 1D78 is the main entry point to the BASIC interpreter. But the routine can be intercepted *before* going to 1D78, by patching a different jump into addresses 4004 and 4005.

This intercept is very important, because every BASIC - transparent software modification in this book will be patched into this location, leading to the master custom interpreter program below (Listing 3-2). When a command line is entered, the program in Listing 3-2 intercepts the routine at 4003, and first examines the status of the stack pointer; if it points to 1D5B, then the intercept program knows BASIC is in the interpretation mode.

Its next step is to CALL 1D78. By calling 1D78 instead of jumping to it, a 'tokenized' version of the next command in line is returned to the master custom interpreter. Tokenizing is an important, specialized process which allows the BASIC listings to use very little memory and allows the interpreter to evaluate commands at high speed.

When the token is returned to the custom interpreter, it can then be evaluated to see if it is a specially designated indicator command.

If the slash command indicator is found by the custom interpreter, it moves on to a lookup table to search for one of the specialized commands. All these commands will be tokens as well, so only a single byte comparison need be made.

There is another method of patching into the BASIC interpreter. If you are a Level II user, merely enter the command OPEN. Very promptly the computer will respond with '?L3 ERROR'. What is an '?L3 ERROR'? It refers to a 'Level III Error', the extended BASIC that is available as a part of the TRS-80 disk system.

Now enter the statement OPNE. This time a '?SN ERROR' is produced. How does the machine know that OPEN is a disk command and that OPNE is just garbage?

```

00100 ; #####
00110 ; CUSTOH BASIC INTERPRETER PATCH FOR USE WITH ALL OF THE
00120 ; ROUTINES PRESENTED IN CUSTOH TRS-80 WHICH ARE TRANS-
00130 ; PARENT TO BASIC. THESE ROUTINES ARE CALLED BY THIS
00140 ; ROUTINE, WHICH HAS ALREADY PLACED A RETURN VALUE ON THE
00150 ; STACK. ALL TRANSPARENT ROUTINES MUST EXECUTE A RETURN
00160 ; INSTRUCTION AS THEIR FINAL INSTRUCTION TO WORK WITH
00170 ; THIS CUSTOH INTERPRETER ROUTINE. THE COHMANOS AVAIL-
00180 ; ABLE WITH THIS INTERPETER ARE:
00190 ; /LOAD /SAVE /NEW /OPEN
00200 ; /ON /OFF /GET /PUT
00210 ; /STEP /MEM
00220 ; AND OTHER USER-DEFINED "/" COHMANOS AND ROUTINES.
00230 ; THIS ROUTINE HAS PUSHED THE RETURN ADDRESS {107B} ON
00240 ; THE STACK. THE ROUTINE JUMPED TO IS A PSEUDO-CALL
00250 ; IN THAT IT EXECUTES A "RET", THUS RETURNING TO 107B.
00260 ; #####
00270 ;
107B 00280 BYTE EQU 107BH ;ROH READ KEYS & TOKENIZE
00290 ;
00300 ; #####
00310 ; CHECK THAT THE BASIC STACK IS IN INTERPRETATION MODE
00320 ; #####
00330 ;
0000 00340 BEGIN EX (SP),HL ; GET SP INTO HL FOR TEST
0001 00350 LO A,L ; GET L INTO A FOR TEST
0002 00360 CP SBH ; IS LSB OF STACK SB?
0004 00370 JR NZ,NOTROY ; NOT INTERPRETING IF NZ
0006 00380 LO A,H ; GET H INTO A FOR TEST
0007 00390 CP 10H ; IS HSB OF STACK 10?
0009 00400 NOTROY EX (SP),HL ; RESTORE STACK TO SP
000A 00410 JP NZ,BYTE ; IF NOT 10H THEN TO ROH
00420 ;
00430 ; #####
00440 ; MUST HAVE BEEN AT 10SB FOR INTERPRETATION; THEREFORE,
00450 ; CHECK TO SEE IF SPECIAL SLASH (/) COHMAND THAT IS NEXT
00460 ; #####
00470 ;
0000 00480 CALL BYTE ; READ CHAR. & TOKENIZE
0001 00490 PUSH AF ; SAVE VALUE READ
0002 00500 CP 00H ; IS IT "/" COHMAND?
0003 00510 JR Z,OKSLSH ; IF SO, THEN CONTINUE
0004 00520 POP AF ; ELSE RESTORE AF VALUE
0005 00530 DEC HL ; PUT POINTER BACK ONE
0006 00540 JP 107BH ; AND BACK TO NORHAI ROM
00550 ;
00560 ; #####
00570 ; SLASH (/) HAS BEEN FOUND, THEREFORE MUST BE COHMAND
00580 ; #####
00590 ;
000A 00600 OKSLSH POP AF ; RESTORE VALUE TO AF
000B 00610 CALL BYTE ; NEXT COHMAND IN LINE
000C 00620 JR NZ,SAVE ; GO IF ONE IS IN PLACE
000D 00630 SYNERR JP 1997H ; ?SN ERROR IF LINE END
00640 ;
00650 ; #####
00660 ; SINCE SLASH & NEXT CHARACTER HAS BEEN FOUND, NOW NEXT
00670 ; CHARACTER IN LINE MUST BE TESTED FOR VALIDITY AS USER-
00680 ; DEFINED COHMAND. SEE ABOVE FOR THOSE AVAILABLE HERE.
00690 ; #####
00700 ;
0023 00710 SAVE LO DE,107BH ; GET RETURN ADDRESS
0024 00720 PUSH DE ; PLACE IT ON STACK
0025 00730 CP 0A0H ; ---- SAVE ----
0026 00740 JP Z,SAVER ; GO TO SAVE ROUTINE
0027 00750 CP 0B0H ; ---- NEW ----
0028 00760 JP Z,RENEW ; GO TO RENEW ROUTINE
0029 00770 CP 0A2H ; ---- OPEN ----
0030 00780 JP Z,OPENER ; GO TO OPEN ROUTINE
0031 00790 CP 0C0H ; ---- STEP ----
0032 00800 JP Z,STPSET ; GO TO STEPPING ROUTINE
0033 00810 CP 0C8H ; ---- MEH ----
0034 00820 JP Z,HEHSET ; GO TO MEMORY SET ROUT.
0035 00830 CP 0EAH ; ---- LOC ----
0036 00840 JP Z,RELOC ; GO TO RELOCATION ROUT.
0037 00850 CP 0A1H ; ---- ON ----
0038 00860 JP Z,KEYON ; GO TO DEBOUNCE ON ROUT.
0039 00870 CP 0A0H ; ---- OUT ----
0040 00880 JP Z,KEYOFF ; KILL DEBOUNCE ROUTINE
0041 00890 CP 0A4H ; ---- GET ----
0042 00900 JP Z,COPYIN ; GO TO READ TAPE ROUTINE
0043 00910 CP 0ASH ; ---- PUT ----
0044 00920 JP Z,DUBBER ; GO TO WRITE TAPE ROUT.
0045 00930 JP SYNERR ; ?SN ERROR IF UNDEFINED
00940 ;
00950 ; #####
00960 ; THIS ROUTINE WILL NOT ASSEMBLE AS IT STANDS. IT MUST
00970 ; BE APPENDED TO THE OTHER ROUTINES WHICH WILL BE USED
00980 ; IN CONJUNCTION WITH BASIC. ALL THE TERMS LISTED ABOVE
00990 ; MUST BE DEFINED, OR ELSE THEY MUST BE DELETED FROM THE
01000 ; ASSEMBLY LISTING.
01010 ; #####
01020 ;
0000 01030 END
00010 TOTAL ERRORS

```

```

00100 ;
00110 ; #####
00120 ; FULL-FEATURED KEY/SCREEN DRIVER - DENNIS BATHORY KITSZ
00130 ; THIS ROUTINE IS A LEVEL II KEYBOARD REPLACEMENT ROUTINE
00140 ; CAPABLE OF PROVIDING: AUTOREPEAT AFTER SELECTED DELAY
00150 ; (FOUNO IN 8 REGISTER IN DELAY SECTION); BEEP WITH ANY
00160 ; CHOICE OF PITCH; RESULTANT DEBOUNCE; CORRECTED SHIFT-
00170 ; DOWN ARROW CONTROL CODE FOR EARLIER LEVEL II ROMS; A
00180 ; SHIFT-0 SELECTABLE UPPER/LOWER CASE DRIVER AND DISPLAY.
00190 ; NOTE THAT THIS ROUTINE IS SET UP FOR USE AT 3039 HEX
00200 ; (12345 DECIMAL) FOR ENTRY IN THE MEMORY SIOECAR WHICH
00210 ; IS ADDRESSED FROM 3000 TO 3700 HEX. IT MAY BE SET TO
00220 ; ANY ORIGIN OF THE USER'S CHOICE, SUCH AS USUAL HIGH MEM
00230 ; #####
00240 ;
4099 00250 INKEYS EQU 4099H ;INKEY$ BYTE STORAGE AREA
4030 00260 PORTFF EQU 403DH ;CASSETTE OUTPUT PORT
401A 00270 KPLACE EQU 401AH ;1-BYTE KEYSTROKE STORE
4019 00280 SHIFTR EQU 4019H ;STORAGE FOR LC DRIVER
00290 ;
00300 ; #####
00310 ; PATCH KEYBOARD ROUTINE INTO 4016 AND DISPLAY INTO 401E
00320 ; #####
00330 ;
4016 00340 ORG 4016H ; START OF KEYBOARD SCAN
4016 393D 00350 DEFW KBPFIX ; PATCH KEYBOARD ROUTINE
401E 00360 ORG 401EH ; START OF DISPLAY SWEEP
401E 2A31 00370 DEFW LDWER ; PATCH UPPR/LDWR ROUTINE
00380 ;
303B 00390 ORG 3039H ; START AT MEMORY SIOECAR
00400 ;
00410 ; #####
00420 ; SET STORAGE #1, ROW #1, COUNTER #0 PARAMETERS FOR SCAN
00430 ; #####
00440 ;
3039 213640 00450 KBPFIX LD HL,4036H ; STORAGE FOR KEYSTROKE
303C 010139 00460 LD 8C,3801H ; FIRST ROW OF KEYS
303F 160D 00470 LD 0,0 ; COUNTER FOR COLUMNS
00480 ;
00490 ; #####
00500 ; CHECK EACH ROW OF KEYS IN SEARCH OF ONE THAT IS PRESSED
00510 ; #####
00520 ;
3041 0A 00530 KEYPRS LD A,(8C) ; RETRIEVE ROW CONTENTS
3042 5F 00540 LD E,A ; SAVE IT TEMPORARILY
3043 A3 00550 AND E ; SET FLAGS FOR TEST
3044 2018 00560 JR NZ,STROKE ; NOT ZERO IF KEY PRESSED
3046 77 00570 LD (HL),A ; SAVE CURRENT VALUE
00580 ;
00590 ; #####
00600 ; INCREMENT AND ROTATE PATTERN CHECKS EACH ROW IN TURN
00610 ; #####
00620 ;
3047 14 00630 RECCHK INC 0 ; INCREMENT ROW COUNTER
3048 2C 00640 INC L ; INCREMENT STORAGE AREA
3049 C801 00650 RLC C ; GET NEXT KEYBOARD COLUMN
3048 79 00660 LD A,C ; GET VALUE INTO ACCUM.
00670 ;
00680 ; #####
00690 ; CHECK IF LAST VALID ROW (I.E., NOT INCLUDING SHIFT KEY)
00700 ; #####
00710 ;
304C D68D 00720 SUB 80M ; LAST ROW IS 3880 HEX
304E 20F1 00730 JR NZ,KEYPRS ; NEXT CHECK IF NOT DONE
00740 ;
00750 ; #####
00760 ; AUTOREPEAT STATUS TEST ... CHECK IF KEYBOARD IS CLEAR
00770 ; #####
00780 ;
3050 0607 00790 LD 8,7 ; COUNTER OF KBRO ROWS
3052 20 00800 CLRMEM DEC L ; START COUNTING BACK
3053 86 00810 ADD A,(ML) ; AND ADD IT UP IN ACCUM
3054 10FC 00820 DJNZ CLRMEM ; AND DO IT FOR 7 ROWS
3056 A7 00830 AND A ; TEST FOR ANY KEY DOWN
3057 3E0D 00840 LD A,D ; A=0, FLAGS ARE INTACT
3058 C0 00850 RET NZ ; BACK IF KEYS IN USE
00860 ;
00870 ; #####
00880 ; RESET AUTOREPEAT DELAY TO ZERO IF TME KEYBOARD IS CLEAR
00890 ; #####
00900 ;
305A 321A40 00910 LD (KPLACE),A ; ELSE DELAY GETS RESET
3050 C8 00920 RET ; AND GO BACK ANYWAY
00930 ;
00940 ; #####
00950 ; IF KEYSTROKE IS FOUND, CHECK STATUS OF AUTOREPEAT LOOP
00960 ; #####
00970 ;
305E A6 00980 STROKE AND (ML) ; CHECK KEYSTROKE STORAGE
305F 281E 00990 JR Z,FOUND ; NEW KEY IF NOT SAME
3061 3A8940 01000 LD A,(INKEYS) ; CHECK STATUS OF INKEY$
3064 A7 01010 AND A ; TEST IF SOMETHING THERE
3065 20E0 01020 JR NZ,RECCHK ; IF THERE IS, LOOP BACK
3067 3A1A40 01030 LD A,(KPLACE) ; NOW CHECK SPECIAL STORE
306A 3C 01040 INC A ; LET STORE = STORE + 1

```

Listing Continued . . .

It is in this distinction that the other patch can be made into the BASIC interpreter. All the DOS (disk operating system) commands *already exist* in Level II BASIC! A patch point (sometimes called a 'vector', other times a 'hook') is provided for each of these commands in RAM. When the disk system is added to the TRS-80, each of these patch points is filled with a jump to a DOS parameter.

Table 3-(?) presents a list of the DOS commands and their patch points in RAM. If you are not (and do not plan to be) a disk user, and if your programs will not be sold to potential disk users, then these patch points are for you. Each one can be used for your own set of commands, and every one will be accepted by a running BASIC program.

List of DOS Patch Points

DOS COMMAND	REPLACEMENT PATCH POINTS (HEX)	REPLACEMENT PATCH POINTS (DECIMAL)
CVI	4153 - 4154	16723 - 16724
FN	4156 - 4157	16726 - 16727
CVS	4159 - 415A	16729 - 16730
DEF	415C - 415D	16732 - 16733
CVO	415F - 4160	16735 - 16736
EOF	4162 - 4163	16738 - 16739
LDC	4165 - 4166	16741 - 16742
LDF	4168 - 4169	16744 - 16745
MKI\$	416B - 416C	16747 - 16748
MKS\$	416E - 416F	16750 - 16751
MKO\$	4171 - 4172	16753 - 16754
CMD	4174 - 4175	16756 - 16757
TIME\$	4177 - 4178	16759 - 16760
OPEN	417A - 417B	16762 - 16763
FILEO	417D - 417E	16765 - 16766
GET	4180 - 4181	16768 - 16769
PUT	4183 - 4184	16771 - 16772
CLOSE	4186 - 4187	16774 - 16775
LOAD	4188 - 418A	16777 - 16778
MERGE	418C - 418D	16780 - 16781
NAME	418F - 4190	16783 - 16784
KILL	4182 - 4193	16786 - 16787
&	4195 - 4196	16789 - 16790
LSET	4198 - 4199	16792 - 16793
RSET	419B - 419C	16796 - 16798
INSTR	419E - 419F	16798 - 16799
SAVE	41A1 - 41A2	16801 - 16802
LINE	41A4 - 41A5	16804 - 16805

Creating BASIC Tokens

Here's a program to start this discussion:

```

10 CLS : REM XXXXXXXXXXXXXXXX
20 Y = 15360
30 FOR X = 17129 TO 17300
40 POKE Y, PEEK (X)
50 Y = Y + 1 : NEXT X
60 PRINT "END OF PROGRAM"
70 PRINT "LIST OF PROGRAM"
80 LIST : REM ZZZZZZZZZZZZ

```

On the screen you now have two versions of the identical information - the BASIC program. The lines of X's and Z's are there to help you locate the program amidst some of what looks like garbage. You will also see that there are some familiar elements missing: the line numbers (which have been converted to hexadecimal), and all the BASIC commands

Continued Listing

```

306B 321A40 01050 LO (KPLACE),A ; AND PUT IT BACK THERE
306E FEFF 01060 CP OFFH ; CHECK IF IT IS AT END
3070 2B0B 01070 JR Z,OECA ; IF SO, THEN HOLD THERE
3072 C5 01080 PUSH BC ; SAVE ROW COUNTER REG.
3073 06FF 01090 LD B,OFFH ; GET DELAY VALUE INTO B
3075 10FE 01100 TMSWTE DJNZ TMSWTE ; AND DELAY JUST A BIT
3077 C1 01110 POP BC ; AND RESTORE ROW COUNTER
307B 18C0 01120 JR RECHK ; AND BACK TO CHECK NEXT
307A 30 01130 OECA DEC A ; LET A = A - 1 (STORAGE)
307B 321A40 01140 LD (KPLACE),A ; AND PUT IT IN STORAGE
01150 ;
01160 ; #####
01170 ; GET KEYBOARD BYTE BACK, STORE AND PREPARE TO MANIPULATE
01180 ; #####
01190 ; FIRST CONVERT 0 FROM COUNTER TO PSEUDO ASCII EQUIVALENT
01200 ; #####
01210 ;
01220 ; LO A,E ; GET KEYBOARD BYTE BACK
01230 FOUNO LO (HL),E ; STORE IT IN STROKE AREA
01240 LO A,0 ; GET ROW COUNTER FROM 0
01250 RLCA ; AND BEGIN A PROCESS...
01260 RLCA ; ...OF CONVERTING IT...
01270 RLCA ; ...TO AN OFFSET VALUE.
01280 LO 0,A ; AND PUT IT BACK IN 0
01290 ;
01300 ; #####
01310 ; NOW PREPARE ROW COUNTER C TO COMPLETE ASCII CONVERSION
01320 ; #####
01330 ;
01340 ; LO C,1 ; GET NUMBER ONE READY
01350 BACKUP LO A,C ; ACCUM. HAS C FOR MATH
01360 ANO E ; TEST IF C = KEYSTROKE
01370 JR NZ,AROUND ; IF NOT, THEN GO AROUND
01380 INC 0 ; ELSE 0 = ROW + COLUMN
01390 RLC C ; C SET TO NEXT COLUMN
01400 JR BACKUP ; GO BACK AND TEST AGAIN
01410 ;
01420 ; #####
01430 ; SHIFT ROW IS TESTED TO DETERMINE UPPER/LOWER STATUS
01440 ; #####
01450 ;
01460 AROUND LO A,(3BBOH) ; GET SHIFT ROW FOR TEST
01470 LO B,A ; AND SAVE IT IN B
01480 LD A,0 ; GET ROW COUNTER BACK
01490 ADD A,40H ; AND CONVERT TO ASCII
01500 CP 60H ; IS IT UP/LW/GRAFIX/ETC
01510 JR NC,Z0429H ; GO OUT IF GRAPHICS MODE
01520 LD 0,A ; SAVE PARTLY CONVERTED
01530 ;
01540 ; #####
01550 ; SHIFT/ODOWN ARROW CHECKED FOR CONVERSION TO CONTROL CODE
01560 ; #####
01570 ;
01580 ; LD A,(3B40H) ; GET VALUE FOUNO 7TH ROW
01590 ANO 10H ; CHECK IF ODOWN ARROW
01600 JR NZ,CNTRLOL ; IF SO, PRODUCE CONTROL
01610 LO A,0 ; ELSE GET VALUE BACK
01620 RRC B ; B BUMPS INTO CARRY FLAG
01630 JR C,GOAWAY ; IF CARRY, THEN SHIFT
01640 ;
01650 ; #####
01660 ; LOWER CASE CONVERSION, MASK STRIPPING, FINAL TOUCHES
01670 ; #####
01680 ;
01690 ; ADD A,20H ; IF NOT THEN LOWER CASE
01700 JR GOAWAY ; AND GET OUT OF ROUTINE
01710 CNTRLOL LO A,0 ; IF CONTROL CODE, GET IT
01720 SUB 40H ; GET RID OF ASCII MASK
01730 JR GOAWAY ; AND GET OUT OF ROUTINE
01740 Z0429H JR 70H ; THE BALANCE OF THE
01750 JR NC,Z0430H ; ROUTINE BELOW UP TO
01760 ADD A,40H ; THE BEEP SECTION IS
01770 CP 3CH ; VIRTUALLY IDENTICAL
01780 JR C,Z0435H ; TO THE KEYBOARD
01790 XOR 10H ; DETERMINATION SUB-
01800 Z0435H RRC B ; ROUTINE FOUNO IN
01810 JR NC,GOAWAY ; ROM. A COMPLETE
01820 XOR 10H ; DESCRIPTION OF THIS
01830 JR GOAWAY ; SECTION OF THE KEY-
01840 Z0430H RLCA ; BOARD SCAN IS FOUNO
01850 RRC B ; IN THE CHAPTER
01860 JR NC,Z0443H ; SUPPLEMENT ON THE
01870 INC A ; ROM KEYBOARD SCAN.
01880 Z0443H LD HL,TABLET ; THIS TABLE IS CHANGED
01890 LD C,A ; FROM THE ONE FOUNO
01900 LO B,0 ; IN EARLIER ROMS, BUT
01910 ADD HL,BC ; THE ROUTINE USED TO
01920 LD A,(HL) ; ACCESS IT IS THE
01930 JR GOAWAY ; SAME.
01940 ;
01950 ; #####
01960 ; TABLE BELOW DETERMINES TRS-80 (NOT ASCII) CONTROL CODES
01970 ; SEE SUPPLEMENT ON KEYBOARD SCAN FOR DETAILS ON CODES
01980 ; #####
01990 ;

```

(CLS, FOR, TO, POKE, PEEK, NEXT, PRINT, and LIST). What has happened to them?

For two purposes – economy of memory and speed of execution – legitimate BASIC commands are converted to single-byte keys called ‘tokens’. When you enter a BASIC command line, a subroutine evaluates each character group in that line, searching through all the keywords in ROM until it finds a match. When it finds a match, it replaces the original group of characters (PRINT, for example, which is five characters) with a single byte (178 in this case). Four bytes are saved, and the lengthy process of looking up the word PRINT is eliminated at run time.

Evaluating for tokens is indeed a time-consuming process. If you type 255 characters of garbage and press (ENTER), the computer will spend nearly two seconds attempting to tokenize that line before reporting a ?SN ERROR. A line which uses the command CHR\$() very often also takes time to tokenize. You can imagine the speed difference if this process were left to be done at RUN time.

The presence of tokens in a complicated program can make the difference between a running program and an ?OM ERROR. As an experiment, return to MEMORY SIZE? and respond with 17250. This gives you about 50 bytes of program space (at least 83 are needed to run any program). Enter these lines:

```

10 PRINT"THIS IS A TEST TO FIND OUT"
20 PRINT"HOW MUCH MEMORY SPACE IS IN HERE"

```

In attempting to run this program, you will get an ?OM ERROR. Now remove the word ‘IN’ and one space from the second line. The program will run fine. Finally, insert PRINT: on line 10. In spite of the fact that it *looks like* you have inserted 6 new characters (P-R-I-N-T-:), you have really only inserted two – the PRINT token (178) and a colon.

There is a lesson in this. If your program is running quite close to the end of your system’s available memory, try cutting down the lines of text within the program. Many more BASIC commands will then open up for use.

Another interesting trick opens up. You may have a BASIC program which you would like to convert for use on a printer. This can take up quite a bit of time. As a quick fix, you might just leaf through your program, replacing all PRINTs with LPRINTs. This won’t work every time (I’ll explain later), but it’s a useful

Listing Continued . .

Continued Listing

```

30DS 000D 0200D TABLE DEFW 0000H ; CARR. RET. / CARR. RET.
30D7 1F1F 0201D DEFW 1F1FH ; CLEAR SCRN / CLEAR SCRN
30D9 0101 0202D DEFW 0101H ; BREAK KEY / BREAK KEY
30DB S81B 0203D DEFW 1BSBH ; EDIT ESCAPE / UP ARROW
30DD 0A0D 0204D DEFW 000AH ; NOP [CHANGE] / LINEFEED
30DF 0B1B 0205D DEFW 1B0BH ; BACKSP. LINE / BACKSP.
30E1 0919 0206D DEFW 1B0BH ; 32-CHAR MODE / HOR. TAB
30E3 202D 0207D DEFW 2020H ; SPACE / SPACE
0208D ;
0209D ; #####
0210D ; FINAL VALUE IS SAVED IN D; STATUS OF SHIFT-D TESTED
0211D ; #####
0212D ;
30E5 57 0213D GOAWAY LD D,A ; SAVE VALUE IN D REG.
30E6 3A103B 0214D BEEEEP LD A,(3B10H) ; GET D KEYBOARD ROW
30E9 FE01 0215D CP 1 ; SEE IF IT IS ZERO (0)
30EB 2016 0216D JR NZ,BLEEEP ; GO OUT IF NOT ZERO
30ED 3AB03B 0217D LD A,(3B0BH) ; IF 0, CHECK SHIFT ROW
30F0 FE01 0218D CP 1 ; CHECK IF SHIFT KEY
30F2 200F 0219D JR NZ,BLEEEP ; IF NOT GO OUT TO BEEP
30F4 3A194D 0220D LD A,(SHIFTR) ; ELSE GET SHIFTLCK
30F7 FE01 0221D XOR 1 ; AND SWITCH 1-0 OR 0-1
30F9 32194D 0222D LD (SHIFTR),A ; AND PUT IN SHIFTLCK
30FC 010005 0223D LD BC,500H ; GET LONGER DELAY
30FF C0600D 0224D CALL 0060H ; CALL ROM DELAY SUBR.
3102 C9 0225D RET ; AND GO BACK, NO BEEP
0226D ;
0227D ; #####
0228D ; DEBOUNCE IS ADDED; ALTERNATE: BEEP MAY BE LENGTHENED
0229D ; #####
0230D ;
3103 01B001 0231D BLEEP LD BC,1B0H ; DEBOUNCE VALUE AND
3106 C0600D 0232D CALL 0060H ; DELAY CALL TO ROM
3109 7A 0233D LD A,D ; GET STORED VALUE BACK
310A C5 0234D PUSH BC ; SAVE BC REGISTER
310B F5 0235D PUSH AF ; SAVE ACCUM AND FLAGS
310C 064D 0236D LD B,40H ; GET NOTE LENGTH VALUE
310E 3A304D 0237D LD A,(PORTFF) ; GET STATUS OF SCREEN
3111 E6FD 0238D AND 0F0H ; MASK SCREEN CHANGE OUT
3113 67 0239D LD H,A ; STORE MSB IN H REG.
3114 F602 0240D OR 2 ; SET BIT 1 TO BE ON
3116 6F 0241D LD L,A ; STORE ALT. MSB IN L REG
3117 7D 0242D BEEPER LD A,L ; GET ALT. MSB TO OUTPUT
3118 D3FF 0243D OUT (OFFH),A ; AND OUTPUT RISING WAVE
311A 7C 0244D LD A,H ; NOW GET NORMAL MSB
311B D3FF 0245D OUT (OFFH),A ; AND OUTPUT FALLING WAVE
3110 C5 0246D PUSH BC ; SAVE NOTE LENGTH REG.
311E 064D 0247D LD B,40H ; GET FREQUENCY DELAY
3120 10FE 0248D FREQCY DJNZ FREQCY ; AND WAIT A LITTLE WHILE
3122 C1 0249D POP BC ; NOW RESTORE LENGTH VAL.
3123 10F2 0250D DJNZ BEEPER ; AND GO BACK THAT LENGTH
3125 F1 0251D POP AF ; RESTORE ORIGINAL CHAR.
3126 C1 0252D POP BC ; AND RESTORE ORIGINAL BC
3127 C352D4 0253D JP 0452H ; BACK TO ROM IN PROGRESS
0254D ;
0255D ; #####
0256D ; THIS IS LOWER CASE DETERMINATION FROM STORED INFO
0257D ; #####
0258D ;
312A FS 0259D LOWER PUSH AF ; SAVE NEEDED REGISTER
312B 3A194D 0260D LD A,(SHIFTR) ; GET STATUS OF SHIFTLCK
312E FE01 0261D CP 1 ; CHECK IF STATUS = 1
3130 2B04 0262D JR Z,LOWER1 ; IF SO THEN GO TO L.C.
3132 F1 0263D POP AF ; ELSE GET ORIGINAL VALUE
3133 C35B04 0264D JP 045BH ; LEAVE TO NORMAL DISPLAY
3136 F1 0265D LOWER1 POP AF ; ELSE GET ORIGINAL VALUE
3137 0D6E03 0266D LD L,(IX+3) ; GET CURSOR LSB INTO L
313A 0D6B04 0267D LD H,(IX+4) ; GET CURSOR MSB INTO H
313D DABAD4 0268D JP C,049AH ; BACK TO ROM IF CARRY
3140 0D7E05 0269D LD A,(IX+5) ; GET CURSOR CHARACTER
3143 67 0270D OR A ; TEST IF CURSOR IS ON
3144 2B01 0271D JR Z,GETCHR ; IF NOT THEN GO DO IT
3146 77 0272D LD (HL),A ; ELSE PUT IT BACK
3147 79 0273D GETCHR LD A,C ; GET VALUE TO DISPLAY
314B FE2D 0274D CP 20H ; IS IT A SPACE OR CNTRL?
314A DA0605 0275D JP C,05D6H ; OUT IF SPACE OR CONTROL
314D FE8D 0276D CP B0H ; IS IT GRAPHICS OR TAB?
314F 02A604 0277D JP NC,04A6H ; OUT IF GRAPHICS OR TAB
3152 C37D04 0278D JP 047DH ; DO UNCONVERTED DISPLAY
0279D ;
0280D ; #####
0281D ;
06CC 0281D END 06CCH ; BACK TO BASIC READY
0000D TOTAL ERRORS
2429S TEXT AREA BYTES LEFT
AROUND 309D 0146D 0137D
BACKUP 30B7 0135D 0140D
BEEEEP 30E6 0214D
BEEPER 3117 0242D 0250D
BLEEEP 3103 0231D 0216D 0219D
CLRMEN 30S2 00B0D 00B2D
CNTRL 30AC 0171D 0160D
DECA 307A 0113D 0107D
FOUND 307F 0123D 0099D
FREQCY 312D 0248D 0248D
GETCHR 3147 0273D 0271D

```

crude first pass; enter this line from command level (no line number):

FOR X=17130 TO 32767 : IF PEEK(X)=17B THEN POKE X,17S : NEXT

ELSE NEXT

This line will search through all of memory (in a 16K machine), looking for the PRINT token (via PEEK), and replacing it wherever it finds it with the LPRINT token (via POKE).

There are occasional risks with this kind of POKEing. Occasionally, a line number or part of the BASIC line organization addresses may have the same value as the PRINT token and get changed with this process. You might end up with a line like 37549 in the middle of a nicely ordered sequence of 1000, 1010, 1020, etc. Another possible flaw is that the value stored as a pointer to the next BASIC line may also correspond to the PRINT token, and get changed. You may get into a LIST-loop or 'lose' some lines (at least to the eye - they are still in there).

If a line number is wrong, merely delete the incorrect line number and retype a new line with the correct number. If some lines seem lost or the LIST command keeps looping, then temporarily reverse the process . . .

FOR X=17130 TO 32767 : IF PEEK(X)=17S THEN POKE X,17B : NEXT

. . . and add a few REM statements in a line before the error occurred. Adding statements will alter the position of program lines in memory, and you will probably be able to perform the original conversion again safely.

As you can see, the tokenizing process has a lot of distinct advantages. As a postscript, consider the program that opened this section. The tokens were displayed as graphics characters. Why is this so?

Since there are only 256 possible combinations of bits in a byte, many times they have to serve multiple masters. In a machine language program, these bytes can be instructions. In a BASIC program, they are tokenized commands. On screen - which means in video memory - they appear as graphics. The 26 letters of our alphabet can be combined and recombined to form words, sentences, paragraphs, etc., and many words can sound alike or be spelled alike and still have different meanings. Context changes how words are understood.

Douglas Hofstadter played on this most dramatically when he wrote, "This sentence no verb".

Listing Continued . . .

Continued Listing

```

GAWAY 30E5 02130 01630 01700 01730 01810 01830 01830
INKEYS 4099 00250 01000
KBPFIX 3039 00450 00350
KEYPR5 3041 00530 00730
KPLACE 401A 00270 00910 01030 01050 01140
LOWER 312A 02590 00370
LOWER1 3136 02650 02620
PORTFF 4030 00260 02370
RECHK 3047 00630 01020 01120
SHIFTR 4019 00280 02200 02220 02600
STROKE 305E 00980 00560
TABLET 3005 02000 01880
TNWSTE 3075 01100 01100
Z0428H 3081 01740 01510
Z0435H 3080 01800 01780
Z0430H 30C5 01840 01750
Z0443H 30C8 01860 01860

```

```

00100 ; #####
00110 ; ADDING DEBOUNCE/BEEP/AUTOREPEAT TO EDITOR/ASSEMBLER 1.1
00120 ; #####
00130 ; NEW EDITASM SCAN PATCH WILL USE ROM ROUTINE IN ITS PLACE
00140 ; #####
00150 ;
4016 00160 ORG 4016H ; KEYBOARD PATCH POINT
4016 68 00170 DEF8 068H ; LSB OF NEW START
4017 50 00180 DEF8 050H ; MSB OF NEW START
5068 00190 ORG 5068H ; NEW PARTIAL KBD DRIVER
5068 211640 00200 LD HL,4016H ; FORMER KBD PATCH POINT
5068 36E3 00210 LD (HL),0E3H ; LSB OF ADDRESS IN ROM
5060 23 00220 INC HL ; GET NEXT POSITION READY
506E 3603 00230 LD (HL),03H ; MSB OF ADDRESS IN ROM
5070 C38A46 00240 JP 468AH ; JUMP TO EDITASM PROGRAM
00250 ;
00260 ; #####
00270 ; NEW KEYBOARD SCAN IS PLACED AT FORMER SOURCE CODE START
00280 ; #####
00290 ;
5000 00300 ORG 5000H ; NEW BLOCK OF CODE
5000 212140 00310 LD HL,4021H ; GET REPEAT STORAGE BYTE
5003 010138 00320 LD BC,3801H ; GET FIRST ROW OF KEYBO
5006 1600 00330 LD 0,0 ; ROW COUNTER REGISTER
5008 0A 00340 LD A,(BC) ; FIND KEY PRESSED IN ROW
5009 5F 00350 LD E,A ; SAVE VALUE IN E REG.
500A A3 00360 AND E ; TEST FOR PRESSED KEY
500B 2019 00370 JR NZ,FOUND ; JUMP OUT IF KEY PRESSED
5000 77 00380 LD (HL),A ; SAVE CURRENT KEY IN HL
500E 14 00390 REDOIT INC 0 ; INCREMENT ROW COUNTER
500F 2C 00400 INC L ; INCREMENT KEY STORAGE
5010 C801 00410 RLC C ; GET NEXT ROW INTO 8C
5012 79 00420 LD A,C ; GET VALUE OF C FOR TEST
5013 0680 00430 SUB 80H ; 80H IS SHIFT KEY ROW
5015 20F1 00440 JR NZ,LOOPX ; LOOP BACK IF NOT SHIFT
5017 0607 00450 LD 8,7 ; NUMBER OF BUMPS TO 00
5019 20 00460 LOOPY DEC L ; DECREMENT BYTE STORAGE
501A 86 00470 ADD A,(HL) ; ADD TOTAL VALUE STORED
501B 10FC 00480 DJNZ LOOPY ; LOOP BACK FOR 7 TIMES
5010 FE00 00490 CP 0 ; WERE NO KEYS PRESSED?
501F 3E00 00500 LD A,0 ; CLEAR ACC., NOT FLAGS
5021 C0 00510 RET NZ ; BACK TO MAIN ROUTINE
5022 321A40 00520 LD (401AH),A ; SAVE VALUE IF ZERO
5025 C9 00530 RET ; BACK TO MAIN ROUTINE
5026 A6 00540 FOUND AND (HL) ; TEST IF SAME CHARACTER
5027 2610 00550 JR Z,SAME ; IF SAME, JUMP OUT ...
5029 3A1A40 00560 LD A,(401AH) ; GET COUNTER BYTE INTO A
502C 3C 00570 INC A ; INCREMENT IT EACH TIME
5020 321A40 00580 LD (401AH),A ; SAVE IT AGAIN FOR NEXT
5030 FEFF 00590 CP OFFH ; IS IT PAST FULL DELAY?
5032 200A 00600 JR NZ,REDOIT ; IF NOT GO BACK FOR MORE
5034 30 00610 DEC A ; DEC A TO 0 FOR REPEAT
5035 321A40 00620 LD (401AH),A ; SAVE THAT VALUE IN CNTR
5038 78 00630 LD A,E ; GET ORIGINAL CHARACTER
5039 73 00640 SAME LD (HL),E ; SAVE THAT VALUE IN HL
503A C5 00650 PUSH BC ; SAVE ROW COUNTER
5038 010002 00660 LD 8C,0200H ; GET DELAY VALUE READY
503E C06000 00670 CALL 0060H ; CALL ROM DELAY SUBROUT.
5041 C1 00680 POP BC ; RESTORE ROW COUNTER
5042 0A 00690 LD A,(8C) ; GET VALUE AT ROW CNTR.
5043 A3 00700 AND E ; TEST IF STILL PRESSED
5044 C8 00710 RET Z ; IF NOT THEN IT'S BOUNCE
5045 C6 00720 PUSH 8C ; SAVE ROW COUNTER AGAIN
5046 E5 00730 PUSH HL ; SAVE STORAGE AREA
5047 F5 00740 PUSH AF ; SAVE CURRENT KEYSTROKE
5048 0640 00750 LD 8,40H ; GET DURATION VALUE
504A 3A3040 00760 LD A,(4030H) ; GET STATUS OF SCREEN
5040 E8F0 00770 AND 0F0H ; CLEAR THE DATA OUTPUT
504F 67 00780 LD H,BECOMES OUTPUT MASK
5050 F602 00790 OR 02 ; READY ACC. FOR BIT-SET

```

Listing Continued . . .

Packing BASIC With Machine Code

One of the most attractive aspects of interpreted code like Level II BASIC is the hide-n-seek game you can play with it. One of the most fruitful games is called 'string packing', a technique that allows machine language programs to be hidden inside ordinary program lines.

It is convenient and efficient, but once it's part of a program, it looks very obscure. There are three ways of creating machine-coded programs through BASIC strings, but all depend on the code being relocatable (see Supplement to Chapter 4 regarding relocatable code). The three ways are:

1. Packing the machine code into a string on a program line, one character at a time. This is done when the program is created.
2. Packing the machine code into a string on a program line, one character at a time, reading data on other program lines. This is done each time the program is run. The read/data lines containing the packing information are then automatically deleted.
3. Building a string in the string variable area from in-line CHR\$() commands. The strings are built each time the program is run.

First, some background on strings and how to find them. In the TRS-80, all variables can be located with a Level II command known as VARPTR (VARIABLE POINTEr). The variable pointer can find out a lot about variables - their type, location, and length. In the case of string variables, VARPTR returns a memory location in which the length of the string is stored.

Assume A\$ is a string variable in a program line. The statement X=VARPTR(A\$) assigns the address of the first part of the A\$ story to X. Here's what X can reveal:

```

PEEK (X) ...the length of A$
PEEK (X+1) ...the least significant byte of where A$ is found
PEEK (X+2) ...the most significant byte of where A$ is found

```

These values are returned in decimal form, because Level II doesn't provide a hexadecimal numbering option. To find where A\$ is actually located, then, use this formula:

$$AD = \text{PEEK}(X+1) + 256 * \text{PEEK}(X+2)$$

Do you see what is happening here? And what can be done with it? If you know, for example, that A\$="XXXXXXXXXX", then you can actually change A\$ by POKEing values into the

Continued Listing

```

5052 6F 00800 LO L,A ; L BECOMES OUTPUT MASK
5053 70 00810 SOUND LO A,L ; GET BIT-SET MASK
5054 03FF 00820 OUT (0FFH),A ; OUTPUT IT (WAVEFORM HI)
5056 7C 00830 LD A,H ; GET BIT-RESET MASK
5057 03FF 00840 OUT (0FFH),A ; OUTPUT IT (WAVEFORM LO)
5059 C5 00850 PUSH BC ; SAVE DURATION VALUE
505A 0640 00860 LO B,40H ; GET PITCH VALUE
505C 1DFE 00870 QJNZ $ ; DELAY FOR AUDIBLE TONE
505E C1 00880 POP BC ; RESTORE BEEP DURATION
505F 10F2 00890 DJNZ SOUND ; LOOP FOR FULL DURATION
5061 F1 00900 POP AF ; RESTORE KEYSTROKE VALUE
5062 E1 00910 POP HL ; RESTORE STORAGE VALUE
5063 C1 00920 POP BC ; RESTORE ROW COUNTER
5064 C30744 00930 JP 4407H ; JUMP INTO EDTASM...!
        00940 ;
        00950 ; #####
        00960 ; FOLLOWING PATCH NEW SOURCE CODE START & KEYBOARD SCAN
        00970 ; #####
        00980 ;
468A 00990 ORG 468AH ; PLACE TO PATCH END
468A 21005E 01000 LO HL,5E00H ; NEW END OF EDTASM
4710 01010 ORG 4710H ; PLACE TO PATCH END
4710 11005E 01020 LD 0E,5E00H ; NEW END OF EDTASM
4A07 01030 ORG 4A07H ; PLACE TO PATCH END
4A07 11005E 01040 LD 0E,5E00H ; NEW END OF EDTASM
4A08 01050 ORG 4A08H ; PLACE TO PATCH END
4A08 21005E 01060 LO HL,5E00H ; NEW END OF EDTASM
4850 01070 ORG 4850H ; PLACE TO PATCH END
4850 21005E 01080 LD HL,5E00H ; NEW END OF EDTASM
4039 01090 ORG 4039H ; PLACE TO PATCH END
4039 21005E 01100 LD HL,5E00H ; NEW END OF EDTASM
4D80 01110 ORG 4D80H ; PLACE TO PATCH END
4D80 21005E 01120 LO HL,5E00H ; NEW END OF EDTASM
5227 01130 ORG 5227H ; PLACE TO PATCH END
5227 21005E 01140 LD HL,5E00H ; NEW END OF EDTASM
43EF 01150 ORG 43EFH ; PLACE TO PATCH KBD SCAN
43EF C30050 01160 JP 5000H ; NEW KEYBOARD SCAN
        01170 ;
        01180 ; #####
        01190 END 5000H
5000
00000 TOTAL ERRORS
3D935 TEXT AREA BYTES LEFT

FOUND 5026 00540 00370
LOOPX 5008 00340 00440
LOOPY 5019 00460 00480
RED01T 500E 00390 00600
SAME 5039 00640 00550
SOUND 5053 00810 00890
    
```

address you have calculated.

Run the following demonstration program:

```

10 CLS
20 A$ = "XXXXXXXXXX"
30 X = VARPTR(A$)
40 Q = PEEK (X)
50 AD = PEEK (X+1) + 256 * PEEK (X+2)
60 PRINT A$
70 L = 65
80 FOR N = 1 TO Q
90 POKE AD,L
100 L = L + 1
110 AD = AD + 1
120 PRINT A$
130 NEXT N
140 LIST 20
    
```

A\$ is created in line 20, and its information storage area is found by X in line 30. Its length is discovered in line 30 and assigned to variable Q; its location is determined in line 50 (assigned to variable AD), and it is printed in its original form in line 60. Line 70 assigns the value 65 (an ASCII 'A') to variable L, and a Q-character loop is created in line 80. The first character in A\$ is POKEd with L, i.e., changed to letter 'A'. L is incremented (to letter 'B'), AD is incremented (to the next character in A\$), and the new A\$ is printed. When all characters in A\$ have been changed, it is listed. You can see that the list itself is changed because A\$ is defined and stored right there in line 20!

As a further experiment, change the value of L in line 70 to 129 and watch what happens (see Box on Tokenizing for details).

Here is the point: you can create a dummy string like A\$ which will become the residence of a machine language program. By finding A\$ and POKEing your program into it, it can be CLOAded and CSAVED at will.

Now we turn to the three string-packing methods themselves. For the sample program, the following 20-byte routine will be used:

```

21 01 3C      LO    HL,3C01H
E5           PUSH   HL
01           POP    HL
28           DEC    HL
01 FF 03     LO    BC,03FFH
36 20        LD    (HL),20H
ED 80        LDIR
21 11 01     LO    HL,D111H
CD A7 28     CALL 28A7H
C9           RET
    
```

This routine clears the screen, prepares the message 'RADIO SHACK LEVEL II BASIC', and displays it.

The first and second string-packing methods are essentially identical, except that in the first method, A\$ is created and the program saved for later use. The second method creates A\$ at every program run. Here is a simple BASIC program using the routine above packed into A\$:

Packing BASIC With Machine Code

```

00100 : #####
00110 : ELEMENTARY (AND SLOW) FOUR-VOICE MUSIC SUBROUTINE.
00120 : THIS ROUTINE IS CURRENTLY SET UP FOR FOUR INDEPENDENT
00130 : VOICES. BECAUSE OF TIMING CONSTRAINTS, VOICES ARE VERY
00140 : LOW IN PITCH. BY USING THE FOUR SUBROUTINES ALTERNATE-
00150 : LY FOR SOUND EFFECTS. A HIGHER SPEED CAN BE OBTAINED.
00160 : A TEST-AND-JUMP LOOP CAN BE INSERTED FOR THIS PURPOSE.
00170 : ALSO, AS LONG AS LABELS ARE CHANGED, VOICES CAN BE
00180 : DROPPED FROM THE SEQUENCE, RESULTING A 50%, 100% AND
00190 : 150% SPEED INCREASE, RESPECTIVELY. SQUARE WAVE OUTPUT.
00200 : NOTICE ALSO THAT, WHEN VOICES ARE DROPPED, OTHER REG-
00210 : ISTERS MAY BE SUBSTITUTED FOR IX, FOR A SPEED INCREASE.
00220 : FURTHERMORE, CHANGING THE SPEED OF THE PROCESSOR WILL
00230 : PROVIDE AN IDENTICAL INCREASE IN THE OUTPUT PITCH.
00240 : #####
00250 :
00260 :
4F60 00210050 00270 LO IX,5000H ; NEAR TOP OF MEMORY
4F64 01FF4F 00280 LO BC,4FFFH ; START PITCH & RHYTHM
00290 : ; MEMORY-MAPPED SOUND
00300 : #####
00310 : OUTER (INTER-NOTE) LOOP BEGINS HERE: T-STATES 212 - 244
00320 : #####
00330 :
4F67 09 00340 LOOP1 EXX ;04:READY DURATION REGS.
4F68 004600 00350 LO 8,(IX+0) ;19:MSB OF NOTE DURATION
4F6B 004E01 00360 LO C,(IX+1) ;19:LSB OF NOTE DURATION
4F6E 09 00370 EXX ;04:STASM REGISTER AWAY
4F6F 006602 00380 LO H,(IX+2) ;19:FIRST PITCH INTO H
4F72 006E03 00390 LO L,(IX+3) ;19:SECOND PITCH INTO L
4F75 006604 00400 LO D,(IX+4) ;19:THIRD PITCH INTO D
4F78 005E05 00410 LO E,(IX+5) ;19:FOURTH PITCH INTO E
00420 :
00430 : #####
00440 : EACH VALUE ACQUIRED FROM IX IS TESTED TO SEE IF IT IS 0
00450 : AND THE VOICE IS TURNED OFF IF IT IS (DEFINING A REST).
00460 : #####
00470 :
4F7B 0A 00480 LO A,(BC) ;07:READY TO TWEAK MEM
4F7C E60F 00490 AND OFH ;04:TURN ALL VOICES ON
4F7E 24 00500 INC H ;04:BUMP VALUE: REST TEST
4F7F 2S 00510 DEC H ;04:BUMP VALUE: REST TEST
4F80 C2854F 00520 JP NZ,REST1 ;10:ONLY 00 DEFINES REST
4F83 C8E7 00530 SET 4,A ;08:SILENCE VOICE IF REST
4F85 2C 00540 REST1 INC L ;04:BUMP VALUE: REST TEST
4F86 20 00550 DEC L ;04:BUMP VALUE: REST TEST
4F87 C28C4F 00560 JP NZ,REST2 ;10:ONLY 00 DEFINES REST
4F8A C8EF 00570 SET S,A ;08:SILENCE VOICE IF REST
4F8C 14 00580 REST2 INC O ;04:BUMP VALUE: REST TEST
4F8D 15 00590 DEC O ;04:BUMP VALUE: REST TEST
4F8E C2934F 00600 JP NZ,REST3 ;10:ONLY 00 DEFINES REST
4F91 C8F7 00610 SET 6,A ;08:SILENCE VOICE IF REST
4F93 1C 00620 REST3 INC E ;04:BUMP VALUE: REST TEST
4F94 10 00630 DEC E ;04:BUMP VALUE: REST TEST
4F95 C29A4F 00640 JP NZ,REST4 ;10:ONLY 00 DEFINES REST
4F9B C8FF 00650 SET 7,A ;08:SILENCE VOICE IF REST
4F9A 02 00660 REST4 LO (BC),A ;07:SET VOICES ON OR OFF
00670 :
00680 : #####
00690 : DECREMENT H,L,O,E (WAVEFORM DURATION FOR EACH VOICE)...
00700 : NEEDED EACH TIME THE WAVEFORM IS TOGGLED DURING LOOPS...
00710 : ...INNER LOOP BEGINS HERE. T-STATES STRICTLY EQUAL 246
00720 : MEANING MAXIMUM FREQUENCY IS APPROXIMATELY 1770000/246
00730 : OR 7195.1 HZ. USEFUL FREQUENCIES ARE CONSIDERABLY LESS.
00740 : #####
00750 : BEGIN PITCH AND RHYTHM COUNTDOWN LOOPS
00760 : #####
00770 : COUNT DOWN THE PITCH LOOP FOR VOICE NUMBER ONE
00780 : #####
00790 :
4F9B 0A 00800 LOOP2 LO A,(BC) ;07:WHAT WAVE IS LURKING
4F9C 2S 00810 DEC H ;04:COUNTDOWN FREQUENCY 1
4F9D C2AB4F 00820 JP NZ,EXIT1 ;10:SAME WAVE IF NOT 0
4FA0 EEO1 00830 XOR 1 ;07:TOGGLE WAVEFORM BIT 1
4FA2 006602 00840 LO M,(IX+2) ;19:RESTORE PITCH VALUE
4FA5 C3AE4F 00850 JP EXIT1A ;10:JUMP PAST TIMEWASTERS
4FA8 F0E5 00860 EXIT1 PUSH IY ;15:WASTE 15 T-STATES
4FAA F0E1 00870 POP IY ;14:WASTE 14 T-STATES
4FAC E6FF 00880 AND OFFH ;07:WASTE 7 MORE T-STATES
00890 :
00900 : #####
00910 : COUNT DOWN THE PITCH LOOP FOR VOICE NUMBER TWO
00920 : #####
00930 :
4FAE 20 00940 EXIT1A DEC L ;04:COUNTDOWN FREQUENCY 2
4FAF C28A4F 00950 JP NZ,EXIT2 ;10:SAME WAVE IF NOT 0
4FB2 EEO2 00960 XOR 2 ;07:TOGGLE WAVEFORM BIT 2
4FB4 006E03 00970 LO L,(IX+3) ;19:RESTORE PITCH VALUE
4FB7 C3C04F 00980 JP EXIT2A ;10:JUMP PAST TIMEWASTERS
4FB8 F0E5 00990 EXIT2 PUSH IY ;15:WASTE 15 BANANAS
4FBC F0E1 01000 POP IY ;14:ORUM FINGERS ON 14
4FBE E6FF 01010 AND OFFH ;07:USELESS ARITHMETIC
01020 :
01030 : #####
01040 : COUNT DOWN THE PITCH LOOP FOR VOICE NUMBER THREE
01050 : #####
01060 :
4FC0 1S 01070 EXIT2A DEC O ;04:COUNTDOWN FREQUENCY 3
4FC1 C2CC4F 01080 JP NZ,EXIT3 ;10:SAME WAVE IF NOT 0
4FC4 EEO4 01090 XOR 4 ;07:TOGGLE WAVEFORM BIT 3
4FC6 005604 01100 LO O,(IX+4) ;19:RESTORE PITCH VALUE
4FC9 C3024F 01110 JP EXIT3A ;10:JUMP PAST TIMEWASTERS
4FCC F0E5 01120 EXIT3 PUSH IY ;15:SCRATCH LEFT HAND
4FCE F0E1 01130 POP IY ;14:SCRATCH RIGHT HAND

```

```

5 POKE 16553,255 : REM OPTIONAL LINE
10 PRINT"THE PROGRAM IS RUNNING"
20 A$ = "12345678901234567890"
30 X = VARPTR (A$) : Q = PEEK (X)
40 AO = PEEK (X+1) + 256 * PEEK (X+2)
50 FOR N = 1 TO Q : READ A
60 POKE AO,A : AO = AO + 1 : NEXT
70 HB = PEEK (X+1) : POKE 16526,HB
80 LB = PEEK (X+2) : POKE 16527,LB
90 INPUT"ENTER TO RUN M/L":Z
100 M$ =USR (0)
110 DATA 33,1,60,229,209,43,1,255,3,54
120 DATA 32,237,176,33,17,1,205,167,40,201

```

Line 5 is optional if you have an early ROM set, where data reads could RESTORE after every READ. Line 20 contains the string to be packed, and, as before, lines 30 and 40 identify the string's length and location. The data is read and POKEd into place sequentially by lines 50 and 60.

Finally, lines 70 and 80 identify the beginning of the string and place it in the USR(X) entry points at 16526 and 16527. The program pauses for user input in line 90, and then jumps to the packed routine.

After the program has been run, list it. Note that A\$ is now packed with new information replacing the string '12345678901234567890'. The first string-packing method saves space by deleting lines 40 through 60, 110 and 120, which have done their work. The program is then CSAVED, and can be loaded and run at any time. The second string-packing method leaves all lines intact so that any future users may modify them as necessary.

There are a few disadvantages to this method of string packing. First of all, two machine language instructions or pieces of data may not be used directly: 00 and 22. 00 tells a BASIC program it has found the end of a program line; two 00's in a row indicate the end of the program. 22 is the quotation mark symbol, and will inform the program that the string has ended; a ?SN ERROR will then be produced in the rest of the line.

A second difficulty is that the line containing A\$ may not be edited. This is because when a line is edited, it is placed from the LIST into a buffer that acts exactly like the keyboard buffer; the bytes within the quotes are then converted into the individual letters. For example, code (178) is a machine command which also is the BASIC token for PRINT; when listed, it comes up on the screen as PRINT. Editing the line puts P-R-I-N-T in the edit buffer; but since it is *within the quotation marks*, it is not tokenized. The result? The string now contains five ASCII characters where it once contained a machine language instruction!

Listing Continued . .

Continued Listing

```

4F00 E6FF      01140      AND      OFFH      ;07:CHECK KITCHEN CLOCK
01150 ;
01160 ; #####
01170 ; COUNT DOWN THE PITCH LOOP FOR VOICE NUHBER FOUR
01180 ; #####
01190 ;
4F02 10      01200 EXIT3A DEC      E      ;04:COUNTDOWN FREQUENCY 4
4F03 C20E4F 01210      JP      NZ,EXIT4    ;10:SAHE WAVE IF NOT 0
4F06 EE08      01220      XOR      B      ;07:TOGGLE WAVEFORH BIT 4
4F0B 00E0S      01230      LD      E,[IX+5]  ;19:RESTORE PITCH VALUE
4F0B C3E44F 01240      JP      EXIT4A      ;10:JUHP PAST TIHEWASTERS
4F0E F0ES      01250 EXIT4  PUSH     IY      ;1S:WATER NASTURTIUHS
4FE0 F0E1      01260      POP      IY      ;14:PICK 14 ZUCCHINI
4FE2 E6FF      01270      AND      OFFH      ;07:HIX APPLES AND ORANGE
01280 ;
01290 ; #####
01300 ; CHECK FOR END OF NOTE OURATION: GET MORE NOTES IF OONE
01310 ; #####
01320 ;
4FE4 02      01330 EXIT4A LD      [BC],A      ;07:OUTPUT NEW WAVEFORHS
4FE5 09      01340      EXX      ;04:GET STASHED OURATION
4FE6 08      01350      DEC      BC      ;06:COUNT DOWN OURATION
4FE7 7B      01360      LD      A,B      ;04:SET UP B FOR TEST
4FE8 B1      01370      OR      C      ;04:CHECK AGAINST C
4FE9 09      01380      EXX      ;04:STASH OURATION AGAIN
4FEA C29B4F 01390      JP      NZ,LOOP2    ;10:GO BACK TIL NOTE ENO
01400 ;
01410 ; #####
01420 ; HOVE ALL POINTERS PAST CURRENT BATCH OF NOTES/OURATIONS
01430 ; ...THIS IS THE REHAINOER OF OUTER LOOP. T-STATES = 80,
01440 ; TOTAL T-STATES OF OUTER LOOP = 80 + 244 = 324, WHICH IS
01450 ; ABOUT 2 OF THE MAXIMUM CYCLE FREQUENCIES (.0002 USEC).
01460 ; #####
01470 ;
4FE0 110600 01480      LD      0E,6      ;10:MEMORY POS'NS TO HOVE
4FF0 0019      01490      ADD     IX,0E      ;1S:MOVE 6 PLACES FORWARD
01500 ;
01510 ; #####
01520 ; CHECK FOR END OF PROGRAH CODE {00} OR DEPRESSED BREAK
01530 ; #####
01540 ;
4FF2 007E00 01550      LD      A,[IX+0]    ;19:NEXT NOTE OURATION
4FF5 B7      01560      OR      A      ;04:SET END-OF-HUSIC FLAG
4FF6 C8      01570      RET      Z      ;0S:BACK TO BASIC IF OONE
4FF7 3A403B 01580      LD      A,(3B40H)  ;13:TEST BREAK KYBO ROW
4FFA B7      01590      OR      A      ;04:SET FLAG FOR KEY TEST
4FFB CA674F 01600      JP      Z,LOOP1    ;10:CONTINUE PIECE IF OK
4FFE C9      01610      RET      ; :TO BASIC IF BREAK
01620 ;
01630 ; #####
01640 ;
06CC      01650      ENO      06CCH      ; :READY AFTER SLASH
00000 TOTAL ERRORS

```

```

30 X = VARPTR (A$)
40 POKE 16526, PEEK (X+1)
50 POKE 16527, PEEK (X+2)
60 INPUT "ENTER TO RUN M/L";Z
70 MS =USR(0)

```

This time, A\$ has been defined in the BASIC variable area using CHR\$(). In other words, because A\$ is no longer defined within the program line, *any value may be used in the machine language program*. Hybrid strings can also be used, as:

```
A$ = "1234567890" + CHR$(0) + "12345" + CHR$(34)
```

This method uses both the POKEing via READ-DATA statments, plus concatenation with 00 and 22 where necessary. Once again, A\$ is stored in the variable area, and none of the 00's or 22's affect the BASIC program. As a final amusement, RUN listing 3-(?) above, then PRINT A\$. The A\$ in listing 3-(?) and the A\$ in this listing both contain identical machine code. RUN listing 3-(?) again. LIST line 20. Now PRINT A\$. The answer is up to you.

String-packing will likely become a very important addition to your library of BASIC tools. Here, then, is a summary of the string-packing technique:

1. Write the BASIC program.
2. Create a dummy string of any unused variable name (for example, A\$="LOONIETUNES")
3. Make the string the exact length of your machine language program.
4. Write a program line that sets another variable to point to the string's variable information.
For example, X = VARPTR(A\$).
5. Find the starting address of the string by converting the decimal bytes into a single decimal value.
AD = PEEK(X+1) + 256 * PEEK (X+2).
6. Create a set of READ and DATA lines in your BASIC program which will POKE the machine language program into place in the dummy string. (For example:

```

Line # FOR N = AD TO AD+3 : READ L : POKE N,L : NEXT
Line # DATA 33, 16, 16, 204 ..... )

```

7. Set the USR(0) entry point at 16526 and 16527 to the beginning of the string variable storage address (for example, PEEK (X+1) and PEEK (X+2), respectively.
8. If you wish, delete the READ, DATA and POKE loop lines used for that routine. CSAVE the program.

The last method of string packing is capable of overcoming both these flaws. Examine the listing below:

```

10 PRINT"THE PROGRAM IS RUNNING"
20 A$ = CHR$(33) + CHR$(1) + CHR$(60) +
      CHR$(229) + CHR$(209) + CHR$(43) +
      CHR$(1) + CHR$(255) + CHR$(3) +
      CHR$(54) + CHR$(32) + CHR$(237) +
      CHR$(176) + CHR$(33) + CHR$(17) +
      CHR$(1) + CHR$(205) + CHR$(167) +
      CHR$(40) + CHR$(201)

```

Sound-Effects Generation

The essence of sound generation has already been sneaked in: the audible beep in the keyboard routine at the beginning of this Chapter. Sound has been something of a mystery, but there could hardly been a simpler machine language program. Try this experimental program with the tape recorder running in record mode (the motor plug removed), and a cassette in place:

```
10 FORX=1TO2000:OUT255,0:OUT255,255:NEXT
```

In order that this program can operate at top speed, make sure you do not use spaces, and keep the program on one line. You'll hear the tape relay clatter and see the screen jitter. Now play back the short segment of tape you recorded. There is a buzzing on that tape, similar in pitch to the buzzing of the cassette relay. What does all this mean?

It is quite simple. There is an electronic switch inside the TRS-80 which controls several activities: whether the screen is in 64 or 32 character mode, whether the cassette relay is on or off, and whether the cassette data output is 'on' or 'off'.

I've put on and off in quotation marks for a reason: which video mode is used and whether the cassette motor is on or off are examples of states or conditions. But in audio terms, the latter output represents very swift sound wave *transitions*, not on or off *conditions*. In other words, the transition from on to off or from off to on can be heard as only a slight click; but many of them in quick succession sound like a buzz; in even faster succession, they become actual pitches.

```
0000 C07F0A 00150 CALL 0A7FH : ACCEPT BASIC VARIABLE
0003 E5 00160 PUSH HL : STASH BASIC VARIABLE
0004 C1 00170 POP BC : TRANSFER VAR. TO BC
0005 C5 00180 LOOP PUSH BC : SAVE OUTER LOOP VAR.
0006 41 00180 LO B,C : GET MSB FOR TONE
0007 10FE 00200 DJNZ S-0 : AND DELAY THAT TIME
0008 3A3040 00210 LO A,(4030H) : GET SCREEN STATUS
000C F602 00220 OR 2 : AND SET BIT ONE
000E 03FF 00230 OUT (0FFH),A : OUTPUT PART OF WAVE
0010 41 00240 LO B,C : GET VARIABLE AGAIN
0011 10FE 00250 DJNZ S-0 : DELAY WHILE WAVE HIGH
0013 E6F0 00260 AND 0F0H : AND SET OTHER BIT HIGH
0015 03FF 00270 OUT (0FFH),A : TO 00 OTHER WAVE HALF
0017 C1 00280 POP BC : RESTORE OUTSIDE VALUE
0018 10EB 00280 DJNZ LOOP : AND LOOP BACK FOR FREQ.
001A C9 00300 RET : GO BACK TO BASIC
00310 :
00320 : #####
0000 00330 END
00000 TOTAL ERRORS
LOOP 0005 00180 00280
```

The single-line BASIC program, short though it may be, cannot move fast enough to produce pleasant tones or dramatic sound effects. For that you must turn to machine language for its speed. Listing 3-(?) presents a complete sound subroutine with ten sample sounds built in.

Listing 3-(?) presents a different sound subroutine which accepts values passed from BASIC.

Again, each program has an advantage. The first allows you to tailor specific sounds and their durations and repetitions with care; the second lets you develop many different sounds easily, directly from BASIC, without altering the machine subroutine already in place.

What's in a List?

Often it's reassuring to be able to give your program to others without having to worry about gratuitous examination of your code, and finding your own, carefully developed techniques in someone else's product. The easiest thing to do, then, is UN-LIST the program.

Actually, the program still LISTS, but can't be seen; and the program still LLISTs, but uses a whole sheet of paper for each line. In either case, it's a discouragement if you've got some code you like. And protecting just a few lines might give the impression that's all there is to your program – a psychological ploy.

The trick is to add two bytes to the end of the program: the command to 'Home Cursor', and the command to 'Form-Feed'. Here are a few normal program lines:

```
0 REM&&
1 CLS:PRINT:PRINT:PRINT"CHANGING...":X=17129:REM&&
2 X=PEEK(X)+256*PEEK(X+1):PRINTX"...":REM&&
3 IFX=0THEN10ELSEPOKEX-28,28:POKEX-3,12:GOTO2:REM&&
10 PRINT:PRINT"THIS IS A TEST ":REM&&
20 PRINT"OF A TECHNIQUE TO KILL":REM&&
30 PRINT"THE FEATURE THAT LISTS":REM&&
40 PRINT"OR LLISTS THE PROGRAM.":REM&&
50 PRINT"THE PROGRAM HAS NOW BEEN CHANGED":REM&&
60 PRINT"A COMPLETE LIST FOLLOWS....":REM&&
70 LIST:REM&&
80 REM&&
```

This process uses four bytes per line (colon, REM, two ampersands), and diminishes the usable characters per line by six, but if you have the memory and think you need a little protection, this is a way to go. Delete lines 1, 2 and 3, and save the program. When LIST is commanded, nothing will list on the screen but two REM lines (0 and 80). The program will require nine sheets of paper to LLIST.

```

00100 ; #####
00110 ; BASIC AUTOEXECUTION ROUTINE PATCH FOR LEVEL II BASIC
00120 ; #####
107B 00140 BYTE EQU 107BH ; INTERPRETER CALL ADDR.
00150 ;
00160 ; #####
00170 ;
7F00 00180 ORG 7F00H ; RELOCATE IF DESIRED
7F00 2A0440 00190 START LO HL,(4004H) ; GET INTERPRETER PATCH
7F03 225B7F 00200 LO (RETURN),HL ; CHAINING PROCESS HERE
7F06 210F7F 00210 LO HL,ENTRY ; ENTRY OF AUTOEXECUTOR
7F09 220440 00220 LO (4004H),HL ; PATCH INTO INTERPRETER
7F0C C3CC06 00230 JP 06CCH ; RETURN TO BASIC READY
00240 ;
00250 ; #####
00260 ; CHECK FOR STATUS OF INTERPRETER (MUST BE AT 105BH)
00270 ; #####
00280 ;
7F0F E3 00290 ENTRY EX (5F),HL ; GET RETURN ADDRESS
7F10 70 00300 LO A,L ; GET LSB INTO A REG.
7F11 FE5B 00310 CP 5BH ; CHECK LSB OF 105BH
7F13 2003 00320 JR NZ,NOTROY ; GO OUT IF NOT AT 105B
7F15 7C 00330 LO A,H ; GET MSB INTO A REG.
7F16 FE10 00340 CP 10H ; CHECK MSB OF 105BH
7F1B E3 00350 NOTROY EX (5F),HL ; RETURN STACK POSITION
7F19 C2577F 00360 JP NZ,AWAY ; BEGONE IF NOT 105BH
00370 ;
00380 ; #####
00390 ; COMPARE PRESENT LINE POSITION WITH CLOAD (TOKEN 89)
00400 ; #####
00410 ;
7F1C C07B10 00420 CALL BYTE ; GET NEXT BUFFER CHAR.
7F1F F5 00430 PUSH AF ; SAVE PRESENT ACCUM.
7F20 FEB9 00440 CP 0B9H ; SEE IF CLOAD TOKEN
7F22 2B04 00450 JR Z,CLOAD ; SPECIAL ROUTINE IF 89
7F24 F1 00460 POP AF ; RESTORE PRESENT ACCUM.
7F25 2B 00470 DEC HL ; RESTORE HL POINTER
7F26 1B2F 00480 JR AWAY ; OUT TO NORMAL MODE
00490 ;
00500 ; #####
00510 ; IF CLOAD TOKEN IS FOUND, EXECUTE CLOAD PROCESS, BUT
00520 ; FIRST INTERCEPT KEYBOARD SCAN (CRUDE WAY OF DOING IT)
00530 ; TO GRAB PROGRAM ON ITS WAY BACK TO A READY CONDITION
00540 ; #####
00550 ;
7F2B F1 00560 CLOAD POP AF ; CLEAR STACK OF AF REG.
7F29 2A1640 00570 LO HL,(4016H) ; GET CURRENT KEYBOARD SCAN
7F2C 22557F 00580 LO (STORE),HL ; SAVE IT FOR A WHILE
7F2F 21427F 00590 LO HL,BYPASS ; GET VALUE OF ROUTINE
7F32 221640 00600 LO (4016H),HL ; PATCH INTO KEYBOARD PLACE
7F35 21E941 00610 LO HL,41E9H ; POINT TO BUFFER LOC'N
7F3B 3600 00620 LO (HL),00 ; PLACE END OF LINE CMO.
7F3A 2B 00630 DEC HL ; BUMP HL BACK TO START
7F3B 3E89 00640 LD A,0B9H ; GET CLOAD VALUE
7F3D 06B0 00650 SUB 80H ; STRIP OFFSET VALUE
7F3F C36510 00660 JP 1065H ; BACK TO EXECUTION ROUT.
00670 ;
00680 ; #####
00690 ; CLOAD ABOVE WILL EXECUTE, RESET THE STACK, AND ATTEMPT
00700 ; TO RETURN TO A READY CONDITION. AT THIS POINT, THE
00710 ; KEYBOARD DRIVER INTERCEPT WILL REDIRECT THE PROCESSOR
00720 ; TO THE BYPASS ROUTINE BELOW, WHICH WILL AGAIN SET UP
00730 ; THE BUFFER, REPATCH THE KEYBOARD DRIVER, AND RUN.
00740 ; #####
00750 ;
7F42 2A557F 00760 BYPASS LO HL,(STORE) ; GET BACK KEYBOARD SCAN
7F45 221640 00770 LO (4016H),HL ; PUT BACK INTO PLACE
7F48 21E941 00780 LO HL,41E9H ; GET BUFFER LOCATION
7F4B 36BE 00790 LO (HL),8EH ; GET RUN COMMAND TOKEN
7F4D 23 00800 INC HL ; BUMP UP BUFFER POS'N
7F4E 3600 00810 LO (HL),00 ; CLOSE OFF THE BUFFER
7F50 2B 00820 DEC HL ; BUMP BACK IN BUFFER...
7F51 2B 00830 DEC HL ; ...TO THE RUN COMMAND
7F52 C35A10 00840 JP 105AH ; AND THEN EXECUTE IT
00850 ;
00860 ; #####
00870 ; THE FOLLOWING FOUR BYTES ARE TEMPORARY KEYBOARD STORAGE
00880 ; #####
00890 ;
7F55 0000 00900 STORE 0EFW 0000 ; TEMPORARY TWO-BYTE AREA
7F57 C3 00910 AWAY 0EFB 0C3H ; JUMP COMMAND IN PLACE
7F5B 7B10 00920 RETURN 0EFW 107BH ; ORIGINAL VALUE CHANGES
00930 ;
7F00 00940 END START ; PATCH ROUTINE AT START
00000 TOTAL ERRORS
31717 TEXT AREA BYTES LEFT
AWAY 7F57 00910 U0360 00480
BYPASS 7F42 00760 00590
BYTE 107B 00140 U0420
CLOAD 7F2B 00560 U0450
ENTRY 7F0F 00290 U0210
NOTROY 7F1B 00350 00320
RETURN 7F5B 00920 U0200
START 7F00 00190 U0940
STORE 7F55 00900 00580 00760

```

Autoexecute BASIC Programs

One of the pleasures of disk operating systems is the ability to load and run programs automatically. This can be done with tape systems as well, simply because all the Level II BASIC operations are organized as subroutines. Any one may be called at any time. To autoexecute a program, then:

1. The SYSTEM command must be entered and the load begun in this mode.
2. The SYSTEM tape must load over its own return point so that it can begin execution automatically.
3. The SYSTEM program loaded must CALL the CLOAD routine, first preparing the stack to return to itself instead of command level.
4. Upon return, the SYSTEM program must prepare the stack once again for return to normal Level II operation, and jump to the RUN routine.

The process is straightforward with one exception: the CLOAD routine is terminal. That is, it forces a return to command level upon completion by clearing out the return address on the stack. It means that a program which might have been written in little more than a dozen bytes must instead play some memory hopscotch first.

Before turning to this loading routine itself, here is a look at the heart of the autorun sequence itself – a mere thirteen bytes! Enter any short BASIC program first, then the routine below:

```

5000 21 E8 41      LO    HL,41E8
5003 36 8E        LO    (HL),8E
5005 23          INC    HL
5006 36 00        LO    (HL),00
500B 2B          DEC    HL
5009 2B          DEC    HL
500A C3 5A 10     JP     105A

```

From BASIC, you can put this program in place with the following lines from command level:

```

X=20480 : POKE X,33 : POKE X+1,232 : POKE X+2,65 :
POKE X+3,54 : POKE X+4,241 : POKE X+5,35 : POKE X+6,54 :
POKE X+7,0 : POKE X+8,43 : POKE X+9,43 : POKE X+10,185 :
POKE X+11,90 : POKE X+12,29 <ENTER>.

```

This simple routine sets the HL register to point to the usual beginning of the keyboard input buffer, puts an 8E (the RUN command value) into that place, bumps the register one place forward, and puts a zero there. The HL register is bumped back to just before the beginning of the keyboard buffer, and the execution routine at 1D5A is entered.

Machine Language Monitor

```

00100 : #####
00110 : MINIATURE MACHINE-LANGUAGE MONITOR OISPLAYING HEX/ASCII
00120 : #####
06CC 00130 READY EDU 06CC      ; RETURN TO READY INTACT
1078 00140 BYTE EDU 1078H     ; ROM READ KEY & TOKENIZE
1997 00150 SYNERR EDU 1997H    ; ENTRY POINT TO SN ERROR
00160 : #####
00170 : GET REST OF DATA AND CONVERT
00180 OPENER CALL BYTE      ; NEXT CHARACTER IN LINE
0003 FE22 00190 CP 22H      ; IS IT A QUOTE MARK?
0005 C29719 00200 JP NZ,SYNERR ; OUT TO ERROR IF NOT
0008 E5 00210 PUSH HL      ; SAVE LINE POINTER
0009 F0E1 00220 POP IY     ; STASH POINTER IN IY
0008 C0F800 00230 CALL XX99 ; CONVERT CHARS TO HEX
000E C0C801 00240 CALL 01C9H ; CLEAR SCREEN (ROM CALL)
0011 184C 00250 JR NEXT99   ; JUMP PAST SUBROUTINES
00260 : #####
00270 : GET 16 SCREEN POSITIONS READY (10H)
0013 7A 00280 CONTNT LD A,D ; GET ADDRESS LOW BYTE
0014 21403C 00290 LO HL,3C40H ; GET SECOND SCREEN LINE
0017 E6F0 00300 AND 0FH      ; MASK OUT LOW BITS
0019 C0E200 00310 CALL RRRRS ; ROTATE/OISPLAY ROUTINE
001C 7A 00320 LO A,D      ; GET ADDRESS LOW BYTE
0010 E60F 00330 AND 0FH      ; MASK OUT HIGH BITS
001F C0D800 00340 CALL HEXASC ; CONVERT WORKS TO ASCII
0022 77 00350 LO (HL),A    ; DISPLAY THE CHARACTER
0023 23 00360 INC HL      ; NEXT SCREEN POSITION
0024 78 00370 LD A,E      ; GET HIGH BYTE
0025 E6F0 00380 AND 0FH      ; MASK OUT LOW BITS
0027 CDE200 00390 CALL RRRRS ; ROTATE/OISPLAY ROUTINE
002A 78 00400 LD A,E      ; GET HIGH BYTE AGAIN
0028 E60F 00410 AND 0FH      ; MASK OUT LOW BITS
0020 C0D800 00420 CALL HEXASC ; CONVERT HEX TO ASCII
0030 77 00430 LD (HL),A    ; DISPLAY THE CHARACTER
0031 21803C 00440 LO HL,3C80H ; GET NEXT SCREEN ROW
0034 0610 00450 LD 8,10H   ; GET 16 VALUE INTO 8
00460 : #####
00470 : OISPLAY CONTENTS OF ADDRESS CHOSEN
0036 1A 00480 CONT02 LO A,(DE) ; GET VALUE AT ADDRESS
0037 E6F0 00490 AND 0FH      ; MASK OUT HIGH BITS
0039 CDE200 00500 CALL RRRRS ; CONVERT/OISPLAY ROUTINE
003C 1A 00510 LD A,(DE)     ; GET VALUE AT ADDRESS
003D E60F 00520 AND 0FH      ; MASK OUT LOW BITS
003F C0D800 00530 CALL HEXASC ; CONVERT CHAR TO ASCII
0042 77 00540 LO (HL),A    ; DISPLAY THE CHARACTER
0043 23 00550 INC HL      ; GET NEXT SCREEN POSN.
0044 23 00560 INC HL      ; GO ONE PLACE MORE
0045 13 00570 INC DE      ; GET NEXT ADDRESS LOCN.
0046 10EE 00580 OJNZ CONT02 ; FULL 16 BYTES OISPLAYED
00590 : #####
00600 : OISPLAY ASCII VALUES TOO
0048 0610 00610 LD 8,10H   ; GET 16 TIMES IN 8 REG.
004A 48 00620 LO C,8      ; SAVE IT IN C FOR USE
004B 18 00630 DEC DE      ; GET NEXT LOWEST ADDRESS
004C 10FD 00640 DJNZ $-1   ; DECREMENT BACK TO START
004E 41 00650 LD 8,C      ; GET 16 TIMES IN 8 AGAIN
004F 21C03C 00660 LD HL,3CC0H ; GET NEXT LINE OF SCREEN
0052 1A 00670 888A LO A,(DE) ; GET CONTENTS OF ADDRESS
0053 77 00680 LD (HL),A    ; OISPLAY EXACTLY AS IS
0054 23 00690 INC HL      ; GET NEXT SCREEN LOCN.
0055 23 00700 INC HL      ; GET NEXT AFTER THAT
0056 23 00710 INC HL      ; VISUALLY MATCHES HEX
0057 13 00720 INC DE      ; GET NEXT ADDRESS TO SEE
0058 10FB 00730 OJNZ 888A ; DO IT FOR 16 ADDRESSES
005A 41 00740 LD 8,C      ; GET 16 INTO 8 AGAIN
005B 18 00750 DEC DE      ; GO BACK TO PREVIOUS
005C 10FB 00760 OJNZ $-1   ; AND BACK TO BEGINNING
005E C9 00770 RET         ; DONE WITH DISPLAY ROUT.
00780 : #####
00790 : SCAN KEYBOARD FOR EDIT SEQUENCE
005F C01300 00800 NEXT99 CALL CONTNT ; FIND WHICH KEYS PRESSED
00810 : #####
00820 : SCAN KEYBOARD FOR BREAK, ARROWS
0062 3A4038 00830 EDITOR LD A,(3840H) ; GET BREAK, ARROWS ROW
0065 FE04 00840 CP 4      ; IS IT BREAK KEY?
0067 2006 00850 JR NZ,ARROW ; IF NOT TEST FOR ARROW
0069 F0E5 00860 PUSH IY   ; ELSE RETRIEVE LINE PTR.
006B E1 00870 POP HL      ; SWITCH BACK INTO HL
006C C3CC06 00880 JP READY ; BACK TO BASIC READY
006F FE10 00890 ARROW CP 10H ; BEGIN ARROW COMPARES
0071 2007 00900 JR NZ,AAAA ; GO IF NOT DOWN ARROW
0073 0610 00910 LO 8,10H ; GET B READY WITH 16
0075 18 00920 DEC DE      ; GO BACK IN MEMORY
0076 10FD 00930 DJNZ $-1 ; DO IT FOR 16 TIMES
0078 1848 00940 JR STNDRD ; DONE NOW; GO OUT
007A FE08 00950 AAAA CP 8 ; CHECK IF UP ARROW
007C 2007 00960 JR NZ,AAAB ; GO OUT IF NOT UP ARROW
007E 0610 00970 LD 8,10H ; GET 16 PLACES READY
0080 13 00980 INC DE      ; GET NEXT MEMORY LOCN.
0081 10FB 00990 DJNZ $-1 ; DO IT 16 TIMES IN ALL
0083 1840 01000 JR STNDRD ; DONE NOW; GO OUT
0085 FE20 01010 AAAB CP 20H ; CHECK IF LEFT ARROW
0087 2003 01020 JR NZ,AAAC ; GO OUT IF NOT LEFT
0089 18 01030 DEC DE      ; GET PREVIOUS MEM. LOCN.
008A 1839 01040 JR STNDRD ; DONE NOW; GO OUT
008C FE40 01050 AAAC CP 40H ; CHECK IF RIGHT ARROW
008E 2003 01060 JR NZ,AAAD ; GO OUT IF NOT RIGHT
0090 13 01070 INC DE      ; GET NEXT MEMORY LOCN.
0091 1832 01080 JR STNDRD ; DONE NOW; GO OUT
01090 : #####
01100 : GET FIFTH LINE OF SCREEN AND DISPLAY CHOSEN EDITING

```

Listing Continued . . .

The routine at 1D5A bumps the HL register forward, evaluates the byte (finding 8E = RUN), then looks for a possible line number to execute. Finding a zero means the command ends there, and so a simple RUN routine is entered. Here's how to try it out once you have it in place:

```

SYSTEM      <ENTER>
/20480      <ENTER>

```

The BASIC program you had entered earlier should now run just as any other BASIC program might. So the idea is to make this autorun routine the heart of the area that the CLOAD might make its way back to.

Listing 3-(?) presents a machine language program which must precede any program to be autoexecuted. It follows the rules above by taking over control of the computer, placing a patch into the keyboard scan in order to intercept the terminal CLOAD routine's return to BASIC, and directing the computer to the usual CLOAD routine. When CLOAD gets back into BASIC, it will present a 'READY' and begin to scan the keyboard. It will, however, never get there.

Instead, the intercept now patched in place will redirect the computer to a short routine also present in the keyboard input buffer area. This routine restores the original plundered keyboard return address, and executes the automatic RUN routine. The autoloader remnants in the keyboard buffer are no longer needed, and will be wiped out at the next keyboard input of any kind.

To use this program, assemble it and save it at the beginning of each in a batch of tapes. Use these tapes to CSAVE any programs you wish to autoexecute. Whenever you wish to run one of these programs, type . . .

```

>SYSTEM      <ENTER>
*? AUTO      <ENTER>

```

. . . and the program will load, acting as if a normal CLOAD were in action, but immediately beginning execution of the BASIC routine.

Machine Language Monitor

It can be very frustrating when you need to make some quick alterations to memory, or when you need to install a short machine language program. The options are few: load a decimal-to-hex conversion program and enter the code; convert the values to hex by hand and POKE them in place; write the code into a short

Continued Listing

```

0093 210030 01110 AAAO LO HL,3000H ; FIFTH LINE ON SCREEN
0096 365F 01120 LO (HL),5FH ; GET CURSOR CHARACTER
0098 23 01130 INC HL ; NEXT SCREEN LOCATION
0099 365F 01140 LO (HL),5FH ; GET CURSOR CHARACTER
0098 28 01150 DEC HL ; BACK TO FIRST LOCN.
009C 0602 01160 LO 8,2 ; GET 2 TRIES INTO B
009E 05 01170 AAAE PUSH OE ; SAVE MEMORY LOCATION
009F E5 01180 PUSH HL ; SAVE SCREEN LOCATION
00A0 C04900 01190 CALL 0049H ; BASIC'S KEYBOARD SCAN
00A3 E1 01200 POP HL ; RESTORE SCREEN LOCN.
00A4 01 01210 POP OE ; RESTORE MEMORY LOCN.
00A5 FE47 01220 CP 47H ; CHECK IF ALPHA HEX
00A7 30B9 01230 JR NC,EDITOR ; GO OUT IF NOT ALPHA HEX
00A9 FE30 01240 CP 30H ; CHECK IF NUMERIC HEX
00AB 38B5 01250 JR C,EDITOR ; OUT IF NOT NUMERIC HEX
00AD FE3A 01260 CP 3AH ; CHECK IF OV NUMERIC
00AF 3804 01270 JR C,AAAF ; CHECK NEXT IF IN RANGE
00B1 FE40 01280 CP 40H ; CHECK IF OV ALPHAHEX
00B3 38A0 01290 JR C,EDITOR ; OUT IF OV ALPHA HEX
00B5 77 01300 AAAF LO (HL),A ; PLACE CHAR ON SCREEN
00B6 23 01310 INC HL ; GET NEXT SCREEN LOCN.
00B7 10E5 01320 DJNZ AAAE ; GO GET ANOTHER CHAR.
01330 ; #####
01340 ; CONVERT CHOSEN DATA TO HEX
00B9 28 01350 DEC HL
00BA C0C000 01360 CALL ASCHEX
00BU 4F 01370 LO C,A
00BE 28 01380 DEC HL
00BF C0EC00 01390 CALL LLLL5
00C2 81 01400 AOD A,C
01410 ;
01420 ; PUT NEW BYTE IN PLACE
00C3 12 01430 LO (OE),A
00C4 13 01440 INC OE
01450 ;
01460 ; DISPLAY REVISED LINE OF DATA
00C5 C01300 01470 STNORO CALL CONTNT
00C8 C0F400 01480 CALL DELAY
00C8 18B5 01490 JR EDITOR
01500 ;
01510 ; ASCII TO HEXADECIMAL CONVERSION
00C0 7E 01520 ASCHEX LO A,(HL)
00CE FE40 01530 CP 40H
00D0 3003 01540 JR NC,NEXT98
00D2 0630 01550 SUB 30H
00D4 C9 01560 RET
00D5 0637 01570 NEXT98 SUB 37H
00D7 C9 01580 RET
01590 ;
01600 ; HEXADECIMAL TO ASCII CONVERSION
00D8 FE0A 01610 HEXA5C CP 0AH
00DA 3003 01620 JR NC,NEXT96
00DC C630 01630 AOD A,30H
00DE C9 01640 RET
00DF C637 01650 NEXT96 AOD A,37H
00E1 C9 01660 RET
01670 ;
01680 ; RIGHT ROTATES FOR CONVERSIONS
00E2 0F 01690 RRRR5 RRCA
00E3 0F 01700 RRCA
00E4 0F 01710 RRCA
00E5 0F 01720 RRCA
00E6 C00800 01730 CALL HEXA5C
00E9 77 01740 LO (HL),A
00EA 23 01750 INC HL
00EB C9 01760 RET
01770 ;
01780 ; LEFT ROTATES FOR CONVERSION
00EC C0C000 01790 LLLL5 CALL ASCHEX
00EF 07 01800 RLCA
00F0 07 01810 RLCA
00F1 07 01820 RLCA
00F2 07 01830 RLCA
00F3 C9 01840 RET
01850 ;
01860 ; DELAY FOR SCREEN DISPLAYS
00F4 010020 01870 DELAY LO 8C,2000H
00F7 C06000 01880 CALL 0060H
00FA C9 01890 RET
01900 ;
01910 ; GET/CONVERT ASCII FROM BUFFER
01920 ; TO HEXADECIMAL ADDRESS
00F8 0604 01930 XX99 LO 8,4
00F0 C07810 01940 5555 CALL BYTE
0100 F5 01950 PUSH AF
0101 10FA 01960 DJNZ 5555
0103 F1 01970 POP AF
0104 77 01980 LO (HL),A
0105 C0C000 01990 CALL ASCHEX
0108 5F 02000 LO E,A
0109 F1 02010 POP AF
010A 77 02020 LO (HL),A
010B C0EC00 02030 CALL LLLL5
010E 83 02040 AOD A,E
010F 5F 02050 LO E,A
0110 F1 02060 POP AF
0111 77 02070 LO (HL),A
0112 C0C000 02080 CALL ASCHEX
0115 57 02090 LO 0,A
0116 F1 02100 POP AF
0117 77 02110 LO (HL),A

```

BASIC program that does the work. None of these are satisfactory. Ideally, a machine language monitor is the tool to use.

But there are disadvantages to the monitors currently available. Many are too long, and are part of other, lengthier programs. Others overlap resident BASIC programs. And none make available ASCII representations as well as BASIC graphics characters. The short monitor presented in this section provides the latter, can be executed from BASIC (using the patch table presented earlier), and sits wherever in memory you would like.

It consists of a few major sections: The first clears the display, presents the requested address, displays the hex contents of that address and the sixteen following, displays the ASCII or graphics values of that address and the sixteen following, and presents a cursor for hex code entry. The second section searches the keyboard for a valid hex character, displays the character, waits for another, displays that, and advances the address and display.

The second section also searches the keyboard for the arrows, and advances the display (a) one place forward on a right arrow; (b) one place back on a left arrow; (c) sixteen places forward on an up arrow; and (d) sixteen places back on a down arrow. Last of all it searches for the (BREAK) key, which returns it to BASIC.

This monitor, as with all the BASIC-transparent programs presented in this book, must be executed by using the special command patch table (see Listing (?)-(?)). The command used in this table is /OPEN"NNNN", where NNNN is the address to be opened for examination (in hex).

Undoing NEW

This is a much easier task than the Level II manual would have you believe. When the NEW command is entered, the program remains in place, *completely unchanged!* The only alteration is that the end-of-BASIC-program pointer in memory has been changed to the beginning of the program. Hence, the computer believes that the program has a total length of zero.

But the old end-of-program information is still intact elsewhere, and can be found very easily. In fact, to restore a program you have actually

Listing Continued . . .

Resetting MEMORY SIZE?

Continued Listing

```

0118 CODEC00 02120 CALL LLLL
0118 B2 02130 ADD A,0
011C 57 02140 LO 0,A
011D C9 02150 RET
08CC 02160 END READY
0000 TOTAL ERRORS
AAAA 007A 00950 00900
AAAB 0085 01010 00960
AAAC 008C 01050 01020
AAAD 0093 01110 01060
AAAE 009E 01170 01320
AAAF 0085 01300 01270
ARROW 006F 00890 00850
ASCHEX 00C0 01520 01360 01790 01990 02080
BBBA 0052 00670 00730
BYTE 107B 00140 00180 01940
CONT02 0036 00480 00580
CONTNT 0013 00280 00800 01470
DELAY 00F4 01870 01480
EDITOR 0062 00830 01230 01250 01290 01490
HEXASC 0008 01610 00340 00420 00530 01730
LLLLS 00EC 01790 01390 02030 02120
NEXT96 000F 01690 01620
NEXT98 0005 01570 01540
NEXT99 005F 00800 00250
OPENER 0000 00180
READY 06CC 00130 00890 02160
RRRHS 00E2 01690 00310 00390 00500
SSSS 00F0 01940 01960
STNOR0 00C5 01470 00940 01000 01040 01080
SYNERR 1997 00150 00200
XX99 00F8 01930 00230

```

```

00100 : #####
00110 : ONLY A SHORT ROUTINE IS NECESSARY TO DETERMINE WHERE
00120 : A BASIC PROGRAM WAS BEFORE ITS POINTERS WERE RESET. THE
00130 : BEGINNING OF PROGRAM POINTER MUST BE RESET, THE LINE
00140 : NUMBERS RUN THROUGH, THE END OF PROGRAM POINTER RESET,
00150 : THE STACK POINTER RESET, THE SCREEN CLEARED, AND THE
00160 : CLEAR (RESET VARIABLES AND PROGRAM CONDITIONS) EXECUTED
00170 : #####
00180 : NOTE THAT THIS ROUTINE IS ALSO CALLED FROM BASIC USING
00190 : THE FORHAT /NEW. THE CUSTOM INTERPRETER IS EMPLOYED.
00200 : #####
00210 :
0000 E058A440 00220 RENEW LO 0E,(40A4H) ; GET START OF PRGRM PTR
0004 3EFF 00230 LO A,0FFH ; GET FF RESETTNG CODE
0006 12 00240 LO (0E),A ; PLACE AT PROGRAM START
0007 C0FC1A 00250 CALL 1AFCH ; GO THRU ALL LINES TILL
; END OF PRGRM 00 FOUND
000A 23 00270 INC HL ; HL MOVED JUST PAST PRGM
000B 22F940 00280 LO (40F9H),HL ; SIMPLE VARIABLE POINTER
000E E078EB40 00290 LO SP,(40EBH) ; RESET STACK TO NORMAL
0012 C0C901 00300 CALL 01C9H ; CLEAR THE SCREEN NOW
0015 C06118 00310 CALL 1B61H ; CLEAR ALL THE POINTERS
001B C3CC06 00320 JP 06CCH ; BACK TO BASIC "READY"
00330 :
00340 : #####
0000 00350 END
00000 TOTAL ERRORS
RENEW 0000 00220

```

```

00100 : #####
00110 : RAH CHECKING PROGRAM - FIRST VERSION WITHOUT DISPLAY -
00120 : EXHAUSTIVELY CHECKS EACH TRIO OF RAH LOCATIONS WITHOUT
00130 : DISPLAYING CURRENT ADDRESS; BAD ADDRESSES ONLY SHOWN.
00140 : #####
00150 : NOTE THAT THIS PROGRAM IS * SELF-RELOCATING * AND MUST
00160 : BE ASSEMBLED AT THE ORIGIN ADDRESS SPECIFIED BELOW.
00170 : #####
00180 :
S000 00190 ORG S000H ; DO NOT CHANGE ORIGIN!
00200 :
00210 : #####
00220 : SUBROUTINE TO CLEAR THE SCREEN WITH NORMAL SPACES
00230 : #####
00240 :
S000 21003C 00250 LO HL,3C00H ; BEGINNING OF VIDEO
S003 11013C 00260 LO 0E,3C01H ; DESTINATION OF SPACE
S006 01FF03 00270 LO 8C,03FFH ; SPACES ON SCREEN
S009 3620 00280 LO (HL),20H ; DISPLAY BLANK SPACE
S00B E080 00290 LOIR ; CLEAR THE SCREEN
00300 :
00310 : #####
00320 : SET UP MIDDLE OF PROGRAM POINTER, DISPLAY STAR TO SHOW
00330 : PROGRAM IS WORKING. SET UP THE STACK POINTER IN RELAT-
00340 : IONSHIP TO THE IX REGISTER, AND SPECIFY START OF TEST
00350 : #####

```

Listing Continued . . .

wiped out, you need to invoke a few ROM routines and restore the beginning-of-program pointer.

Although variables are cleared in this process, the program is totally restored. If you wish to make this a part of a transparent operating system using the interpreter patch presented earlier in this chapter, Listing 3-(?) presents a complete routine. Enter /NEW, and the lost program reappears.

One warning is in order: if before restoring the program you cause a ?SN ERROR, the computer will jumble up the first part of the program, mess around with some other memory pointers, and the program will *really* be lost.

Resetting MEMORY SIZE?

The size of BASIC memory available can also be changed from BASIC itself, because it too is simply stored in a two-byte pointer in the RAM patch area. New values may be POKEd into place, so long as they meet two conditions:

1. The new value must be within the range of actual memory available.
2. It must not dip below the top of an existing BASIC program already in memory.
3. If no program is in memory, it must not dip below address 4414.

Here's how to do it. Convert the desired new memory size to split decimal with this formula:

```

X = NNNNN <ENTER>
Y = FIX (X/256) <ENTER>
Z = X - Y * 256 <ENTER>
PRINT Z,Y <ENTER>

```

The value of Z is the least significant byte of the new memory size, Y is the most significant byte. Now:

```
POKE 16561,Z : POKE 16562,Y : CLEAR50
```

The new memory size has been set, and 50 bytes are cleared for string space, as usual. If you haven't followed the rules about legitimate memory sizes, expect a fast system crash.

Continued Listing

```

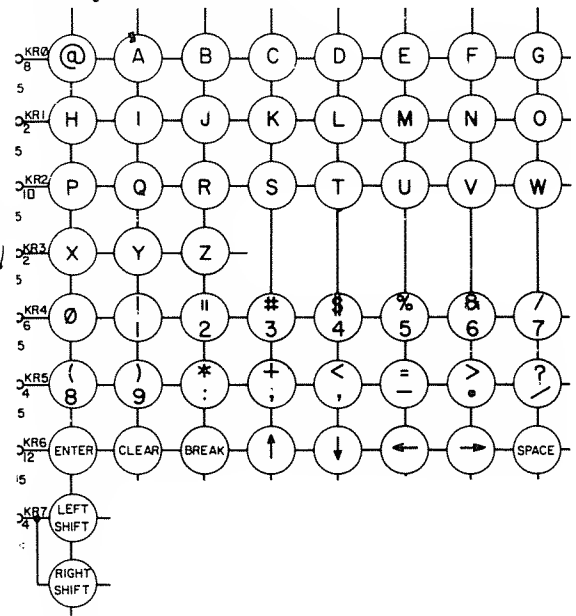
00360 ;
5000 00218050 00370      LO      IX,5080H      ; MIDDLE OF TEST PROGRAM
5011 3AFE3F    00380      LO      A,(3FFEH) ; BOTTOM RIGHT OF SCREEN
5014 E60A      00390      XOR      0AH      ; ALTERNATE SPACE & STAR
5016 32FE3F    00400      LO      (3FFEH),A ; DISPLAY IT ON SCREEN
5019 00F9      00410      LO      5F,IX    ; START STACK POINTER
501B 067F      00420      LO      B,7FH    ; NUMBER OF MOVES FOR SP
501D 33        00430      INC      SP      ; SP = SP + 1
501E 10F0      00440      DJNZ     $-1     ; DO IT 127 TIMES
5020 210060    00450      LO      HL,6000H ; BEGINNING OF TEST AREA
00460 ;
00470 ; #####
00480 ; TEST BEGINS WITH THE VALUE OF ZERO, AND PROCEEDS TO
00490 ; VALUE FF. EACH VALUE IS WRITTEN IN TURN TO EACH OF A
00500 ; TRIO OF MEMORY LOCATIONS IN ORDER TO DETERMINE THEIR
00510 ; EFFECT ON EACH OTHER AS VALID ELECTRONIC STORAGE CELLS.
00520 ; #####
00530 ;
5023 AF        00540      XOR      A      ; CLEAR ACCUMULATOR
5024 77        00550      LO      (HL),A   ; PLACE VALUE IN MEMORY
5025 4F        00560      LO      C,A      ; PLACE VALUE IN C
5026 F5        00570      PUSH     AF      ; SAVE VALUE IN ACCUM.
5027 79        00580      LO      A,C      ; GET VALUE FROM C
5028 B7        00590      LO      (HL),A   ; PLACE VALUE IN MEMORY
5029 23        00600      INC      HL      ; INCREMENT TO NEXT MEM
502A 7C        00610      LO      A,H      ; GET MSB OF ADDRESS
502B FE80      00620      CP      80H      ; IS MEM AT 8000 YET?
00630 ; NOTE THAT THE VALUE
00640 ; ABOVE IS FOR 16K; USE
00650 ; CD FOR 32K, OR FOR 48K
502D 2B30      00660      JR      Z,$+32H   ; IF 50, THEN RELOCATE
502F E5        00670      PUSH     HL      ; SAVE MEMORY VALUE
5030 00E5      00680      PUSH     IX      ; SAVE MIDDLE OF PROGRAM
5032 E1        00690      POP      HL      ; GET POSITION INTO HL
5033 BC        00700      CP      H      ; CHECK AGAINST 80 VALUE
5034 2B29      00710      JR      Z,$+2BH   ; RELOCATE IF DONE
5036 E1        00720      POP      HL      ; RESTORE ORIGINAL VALUE
5037 F1        00730      POP      AF      ; RESTORE ORIGINAL TEST
5038 77        00740      LO      (HL),A   ; PUT VALUE INTO MEMORY
5039 2B        00750      DEC      HL      ; BACK TO ORIGINAL LOC'N
503A 2B        00760      DEC      HL      ; BACK TO ONE BEFORE IT
503B 77        00770      LO      (HL),A   ; PUT VALUE INTO MEMORY
503C 47        00780      LO      B,A      ; SAVE VALUE IN B REG.
503D 7E        00790      LO      A,(HL)   ; GET VALUE AT LOC'N HL
503E 8B        00800      CP      B      ; CHECK AGAINST B VALUE
503F 2B05      00810      JR      Z,$+7     ; GO ON IF IT CHECKS OKAY
5041 161B      00820      LO      0,1BH   ; GET VALUE OF REL. SUB.
5043 05        00830      PUSH     DE      ; SAVE VALUE ON STACK
5044 1B1B      00840      JR      $+10H    ; JUMP TO SUBROUTINE
5046 23        00850      INC      HL      ; GET ORIGINAL TEST POS'N
5047 23        00860      INC      HL      ; GO ONE BEYOND IT
5048 7E        00870      LO      A,(HL)   ; GET VALUE AT THAT POS'N
5049 8B        00880      CP      B      ; CHECK AGAINST B REG.
504A 2B05      00890      JR      Z,$+7     ; GO ON IF MEMORY OKAY
504C 1610      00900      LO      0,10H   ; GET JUMP FOR REL. SUB.
504E 05        00910      PUSH     DE      ; SAVE VALUE ON STACK
504F 1B10      00920      JR      $+12H    ; JUMP TO SUBROUTINE
5051 2B        00930      DEC      HL      ; BACK TO ORIGINAL POS'N
5052 34        00940      INC      (HL)   ; INCREMENT VALUE IN MEM
5053 0C        00950      INC      C      ; INCREMENT TEST VALUE
5054 7E        00960      LO      A,(HL)   ; GET VALUE IN MEMORY
5055 47        00970      LO      B,A      ; SAVE VALUE IN B REG.
5056 FE00      00980      CP      00      ; CHECK IF 256 BYTES DONE
5058 20CC      00990      JR      NZ,$-32H ; LOOP BACK AND CONTINUE
505A 23        01000      INC      HL      ; GET NEXT MEMORY VALUE
505B 0E00      01010      LO      C,0     ; RESET TEST VALUE TO 0
505D 1B07      01020      JR      $-37H   ; LOOP BACK FOR NEXT TEST
505F 1B64      01030      JR      $+66H   ; REL.SUB. STEPPING STONE
01040 ;
01050 ; #####
01060 ; SUBROUTINE BELOW IS ENTERED WHEN A BAD MEMORY LOCATION
01070 ; HAS BEEN DETERMINED. IT CONVERTS HEX VALUES TO ASCII
01080 ; AND DISPLAYS THEM ON THE SCREEN. NOTE THE RELATIVE
01090 ; SUBROUTINE ENTRY AND EXIT METHOD USING THE 0 REGISTER.
01100 ; #####
01110 ;
5061 7C        01120      LO      A,H      ; GET VALUE FROM H REG.
5062 E6F0      01130      AND      0F0H    ; MASK OFF LOW BITS
5064 0F        01140      RRCA      ; ROTATE RIGHT FOR CONV.
5065 0F        01150      RRCA      ; ... SOME MORE ...
5066 0F        01160      RRCA      ; ... AND SOME MORE ...
5067 0F        01170      RRCA      ; ... UNTIL IT'S DONE.
5068 1622      01180      LO      0,22H    ; GET VALUE FOR REL. SUB.
506A 1B22      01190      JR      $+24H    ; AND JUMP TO SUBROUTINE
506C 7C        01200      LO      A,H      ; GET VALUE FROM H REG.
506D E60F      01210      AND      0FH     ; MASK OUT HIGH BITS
506F 161B      01220      LO      0,1BH   ; GET VALUE FOR REL. SUB.
5071 1B1B      01230      JR      $+10H    ; AND JUMP TO SUBROUTINE
5073 70        01240      LO      A,L      ; GET VALUE FROM L REG.
5074 E6F0      01250      AND      0F0H    ; MASK OUT LOW BITS
5076 0F        01260      RRCA      ; ROTATE FOR CONVERSION
5077 0F        01270      RRCA      ; ... SOME MORE ...
5078 0F        01280      RRCA      ; ... SOME MORE ...
5079 0F        01290      RRCA      ; ... AND IT'S DONE.
507A 1610      01300      LO      0,10H   ; GET VALUE FOR REL. SUB.
507C 1B10      01310      JR      $+12H    ; AND JUMP TO SUBROUTINE
507E 70        01320      LO      A,L      ; GET VALUE FROM L REG.
507F E60F      01330      AND      0FH     ; MASK OUT HIGH BITS
5081 1609      01340      LO      0,9     ; GET VALUE FOR REL. SUB.
5083 1B09      01350      JR      $+0BH    ; AND JUMP TO SUBROUTINE
5085 01        01360      POP      DE      ; RESTORE VALUE TO DE
5086 3E03      01370      LO      A,003H ; GET RETURN POS'N VALUE
    
```

Listing Continued . . .

Peek that Keyboard

One of the handiest functions for fast-running, convenient BASIC programs is the INKEY\$ command. This allows the latest keystroke to be transferred to a program variable, where it can then be evaluated.

For action games, BASIC word processors and other speed-conscious programs, INKEY\$ can be too long because it requires considerable evaluation and juggling of memory space. There is a faster way, and it involves examining the keyboard's memory contents directly. Below is the keyboard matrix:



You can see that various memory locations refer to various rows of keys. Depending on the key pressed, different values will be discovered by PEEKing. For example, A=PEEK(14400) returns a value of 1 if ENTER is pressed, and a value of 2 if CLEAR is pressed, 4 if BREAK is pressed, 8 for down arrow, 16 for up arrow, 32 for left arrow, 64 for right arrow, and 128 for a space. If only a few keys are being sought – perhaps the arrow in an action game – a very tight loop can be constructed that is many times faster than INKEY\$:

```

10 A = PEEK(14400) : IF A=0 THEN 10
20 IF A=8 THEN 100 : IF A=16 THEN 200
30 IF A=32 THEN 300 : IF A=64 THEN 400
40 GOTO 10
    
```

For top-speed operation, this should all appear on one line; try it in comparison with testing for CHR\$(8) or CHR\$(9), etc., etc., to determine the arrows.

Continued Listing

```

5088 92      D1380      SUB      D      ; SUBTRACT RETURN DIFF.
5089 DD77DD D1390      LD      (IX+DDH),A ; AND HAKE JR OPERAND
509C 18FF    D1400      JR      S+1      ; IRRELEVANT OPERAND
50BE 05      D1410      PUSH     DE      ; SAVE VALUE ON STACK
50BF F5      D1420      PUSH     AF      ; SAVE VALUE ON STACK
509D 11DD3C D1430      LD      DE,3CDDH ; GET TDP LEFT OF SCREEN
5093 1A      D1440      LD      A,(DE)   ; GET VALUE ON SCREEN
5094 FE2D    D1450      CP      2DH      ; IS IT A SPACE NOW?
5096 2B03    D1460      JR      Z,S+5    ; IF 50, GO AHEAD SOME
509B 13      D1470      INC      DE      ; INCREMENT SCREEN PDS'N
5099 1BF8    D1480      JR      S-6      ; AND GO AHEAD PAST TEST
509B 7B      D1490      LD      A,E      ; GET VALUE OF SCREEN
509C FED4    D1500      CP      D4       ; IS IT 4 PDS'N5 DVER?
509E 2DDA    D1510      JR      NZ,S+DCH ; IF NDT, GO ON AHEAD
50AD 3E2D    D1520      LD      A,2DH    ; IF 5D GET A SPACE READY
50A2 18      D1530      DEC      DE      ; GO BACK SOME....
50A3 12      D1540      LD      (DE),A   ; AND FILL WITH A SPACE
50A4 18      D1550      DEC      DE      ; AND BACK SOME MDRE....
50A5 12      D1560      LD      (DE),A   ; AND INSERT ANOTHER ONE
50A6 18      D1570      DEC      DE      ; AND BACK A BIT MDRE....
50A7 12      D1580      LD      (DE),A   ; AND STASH ANOTHER SPACE
50AB 1B      D1590      DEC      DE      ; AND BACK ONE MORE TIME
50A9 12      D1600      LD      (DE),A   ; AND STUFF A SPACE THERE
50AA F1      D1610      POP      AF      ; RESTORE VALUE TO AF
50AB FED4    D1620      CP      DAH      ; IS VALUE LESS THAN 1D?
50AD 3DD4    D1630      JR      NC,S+6    ; IF LESSER, THEN JUMP
50AF C63D    D1640      ADD      A,3DH    ; CONVERT HEX TO ASCII
50B1 1B02    D1650      JR      S+4      ; AND GO ON PAST THE REST
50B3 C637    D1660      ADD      A,37H    ; CONVERT HEX TO ASCII
50B5 12      D1670      LD      (DE),A   ; AND STASH ON THE SCREEN
50B6 C5      D1730      PUSH     BC      ; SAVE VALUE IN BC REG
50B7 D6FF    D1740      LD      B,0FFH   ; GET DELAY VALUE
50B9 1DFF    D1750      DJNZ     S-D      ; AND DELAY JUST A LITTLE
50BB C1      D1760      PDP      BC      ; RESTORE BC VALUE
50BC D1      D1770      PDP      DE      ; RESTORE DE VALUE
50BD 3EC9    D1780      LD      A,DC9H   ; GET VALUE TO RETURN
50BF 92      D1790      SUB      D      ; SUBTRACT JUMP OFFSET
50CD DD7744 D1800      LD      (IX+44H),A ; PLACE A5 JR OPERAND
50C3 18FF    D1810      JR      S+1      ; IRRELEVANT OPERAND
50C4 18FF    D1820      ;
50C5 DDE5    D1910      PUSH     IX      ; SAVE PROGRAM POSITION
50C7 DDE5    D1920      PUSH     IX      ; SAVE PROGRAM POSITION
50C9 E1      D1930      PDP      HL      ; TRANSFER TO HL REG.
50CA D1      D1940      PDP      DE      ; TRANSFER TO DE REG.
50CB 7A      D1950      LD      A,D      ; GET VALUE FROM D REG.
50CC EE2D    D1960      XDR      2DH     ; TRANSFER TO HIGH MEM.
50CE 57      D1970      LD      D,A      ; PUT BACK IN D REG.
50CF 05      D1980      PUSH     DE      ; AND STASH ON STACK
50D0 D68D    D1990      LD      B,BDH    ; DEC. TO PROGRAM START
50D2 28      D2000      DEC      HL      ; AND BEGIN DECREMENTING
50D3 1B      D2010      DEC      DE      ; FOR BOTH THE REGISTERS
50D4 1DFC    D2020      QJNZ     S-2     ; UNTIL IT'S ALL DONE
50D6 D1FF00 D2030      LD      B,C,0DFFH ; AND GET READY TRANSFER
50D9 ED8D    D2040      LDIR      ; AND THEN DO IT!
50DB D0E1    D2050      POP      IX      ; RESTORE NEW VALUE
50DD 3E2D    D2060      LD      A,2DH    ; VALUE TO MOD. ADDRESSES
50DF DDAEA2 D2070      XDR      (IX+DA2H) ; MODIFY IX+A2 ADDRESS
50E2 DD77A2 D2080      LD      (IX+DA2H),A ; AND STORE IT IN PLACE
50E5 3EFD    D2090      LD      A,DEFFH ; VALUE TO MOD. ADDRESSES
50E6 3EFD    D2100      LD      A,DEFFH ; CHANGE OPERAND ABOVE TO
50E7 DDAEAC D2130      XDR      (IX+DACH) ; TO 70 FOR 4BK MEM TEST
50EA DD77AC D2140      LD      (IX+DACH),A ; MODIFY IX+AC ADDRESS
50ED DDE5    D2150      PUSH     IX      ; AND STORE IT IN PLACE
50EF E1      D2160      POP      HL      ; STORE THE PROGRAM PTR.
50F0 2E11    D2170      LD      L,11H   ; AND TRANSFER IT TO HL
50F2 E9      D2180      JP      (HL)   ; SET LSB OF HL REGISTER
50F3 E9      D2190      ; JUMP TO PROGRAM START!
50F4 E9      D2200      ;
50F5 E9      D2210      ;
50F6 E9      D2220      ;
50F7 E9      D2230      ;
50F8 E9      D2240      ;
50F9 E9      D2250      ;
50FA E9      D2260      ;
50FB E9      D2270      ;
50FC E9      D2280      ;
50FD E9      D2290      ;
50FE E9      D2300      ;
50FF E9      D2310      ;
5000 D2210  TOTAL ERRORS

```

There is yet another use for this PEEK function. As noted in Chapter 2, location 387F reveals if any key is pressed at all. This can be used in a BASIC text editor, for example, and is much faster than INKEY\$ to check for keypressing:

```
10 A = PEEK(14463) : IF A=0 THEN 10
```

This is one of the fastest keystroke-detectors in BASIC, and from there a group of PEEKs could be done, checking the most-used rows of characters first. It's both a matter of taste and programming skill whether INKEY\$ or PEEK is used for keyboard input, but in some situations (such as testing for carriage returns and other control codes), PEEK is the winner.

	1	2	4	8	16	32	64	128
DATA:	1	2	4	8	16	32	64	128
ADDRESS:								
14337	@	A	B	C	D	E	F	G
14338	H	I	J	K	L	M	N	O
14340	P	Q	R	S	T	U	V	W
14344	X	Y	Z					
14352	D	1	2	3	4	5	6	7
14368	8	9	:	*	+	<	=	>
14400	ENT	CLR	BKR	UPAR	DNAR	LFTAR	RTAR	SPACE
14464	SHIFT							

Figure 3-(?) is a chart of all the characters and their respective PEEK positions and data returned. Remember, when more than one key is depressed at a time, the row of the data at that address is returned. For example, if PEEK (14400) returns 129, then both ENTER and SPACE are being held down. This is significant because INKEY\$ returns only the last character pressed, and there is no real way of getting to multiple characters. Try this:

```

10 A = PEEK (14400)
20 IF A = 129 THEN 100
30 PRINT "PRESS ENTER AND CLEAR"
40 GOTO 10
100 PRINT "DONE!"

```

Simple as this may appear, it offers an 'out' for programs that need special, unusual input for some functions. It can be used as the sort of double-interlock protection switch found on industrial machinery to keep both hands out of blades and moving parts.

Furthermore, when the keyboard is PEEKed, the values returned can be changed and mutated at will - for example, a BASIC letter-writer could have two easily written routines to make the keyboard the standard 'QWERTY' type or changing it to the faster Dvorak type. Of course, even with PEEK statements, BASIC is not likely to keep up with that sort of speed typing!

Make-'Em-Sweat Memory Test

Many memory tests are available for testing the dynamic memory in the TRS-80 and expansion interface, including the simple test done by the computer itself on power-up (see Supplement to Chapter 1). Most have a significant disadvantage: they are *software* tests rather than *electrical* tests. Certainly, since memory involves operating software, a software test seems to be a logical solution.

On the other hand, memories are electronic devices, run by electricity and influenced in their failings by electrical and physical forces. Barring removing each memory and running sophisticated electrical tests on it, then, there is only one serious memory test option: test each bit in combination with each and every other bit in the device. Unfortunately, that is impractical, since there are 393,216 bits in a 48K system, and testing every one would result in over 150,000,000,000 separate tests – a task that would take nearly a solid month running at the TRS-80's 1.77 MHz clock rate!

The remaining option, then, is to make reasonable electrical tests that exercise the neighboring bits in a memory chip, and read and write to them about as fast as the TRS-80 is likely to do it in real time. The process I have chosen is twentyfold:

1. Write a value to memory.
2. Read the stored value and check its accuracy.
3. Change the surrounding bits from zeros to ones.
4. Change the surrounding bits from ones to zeros.
5. Read the stored value and check its accuracy.
6. Change the surrounding bits from zeros to ones.
7. Change the surrounding bits from ones to zeros.
8. Write the value to memory again.
9. Read the stored value and check its accuracy.

10. Change the surrounding bits from zeros to ones.
11. Read the stored value and check its accuracy.
12. Change the surrounding bits from ones to zeros.
13. Read the stored value and check its accuracy.
14. Change the surrounding bits from zeros to ones.
15. Write the value to memory again.
16. Read the stored value and check its accuracy.
17. Change the surrounding bits from ones to zeros.
18. Write the value to memory again.
19. Read the stored value and check its accuracy.
20. Increment the value to be written and repeat.

This process is repeated 256 times, writing values from 00 to FF, and producing the electrical switches noted above. The program then moves itself to another area of memory, and checks the area in which it was just residing.

The entire memory test is still very time-consuming, since the address under test is displayed while the process is continuing.

In the program presented in Listing 3-(?), the test displays the memory location under test; any failed memory location is displayed on the screen, along with the bits which have failed. The test is presented in a 16K version with changes for 32K and 48K systems.

Cassette I/O

One of the most maligned aspects of the TRS-80 is its cassette loading procedure. Interestingly, it is a lengthy and skillfully designed piece of coding, a victim of a combination of poor hardware (an inexpensive cassette recorder), the inclination personal computer owners have to purchase the least expensive tapes they can find, and the lack of foresight on the part of the engineers designing the routines. But there's no question that with a good tape recorder and reasonable tape, it works well. Here's how.

The routine to read and accept serial information is fairly convoluted, collapsed to about a dozen major subroutine CALLs. We will start with the SYSTEM command; since BASIC programs have other bytes to juggle (looking for out-of-memory errors, etc.), we won't tackle its major routines.

The SYSTEM module

The SYSTEM command is evaluated by the BASIC interpreter, and its control routine is entered at 02B2. If you don't want to know how this command gets to work, then skip right to the tape loading routine two paragraphs below. An initial CALL is executed to DOS link 41E2, which in Level II merely executes a RETURN. The stack is set up at 4288, and another CALL executed to 20FE, which checks the DOS link at 41C1, picks up the 'device type' – video, tape, or printer – (video at this time), displays a carriage return, checks and saves port FF status (32 or 64 character mode and cassette state), clears the accumulator, and returns. This is preparatory housekeeping.

The accumulator is set up with a star, it is displayed (with more housekeeping), and the INPUT routine is CALLED from location IBB3. This is the same routine used for INPUT statements, and it displays a question mark, evaluates the input line, discards everything after certain punctuation, and returns the evaluated line to the CALLing program. If a BREAK is discovered, the program returns to READY. Spaces, line feeds, tabs, etc., are cleaned out, and a syntax error is declared if no alphanumeric characters are found. If a slash (/) is found, the SYSTEM program jumps past its loading routines, picks up the start address from 40DF (more about that later), cleans out blanks

again, and evaluates the string after the slash as an interger (a CALL to 1E5A). The whole business starts over if a non-numeric string is found. If, at last, the program does discover that a number was input, the SYSTEM module is executed from the starting address stored at 40DF.

Build-a-Byte

The first major loading call is to 0293, which searches for a synchronization byte. Since this will eventually call the 'build-a-byte' routine, let's move there first. It begins at 0241; BC and AF registers are saved. Then:

0243	0B FF	IN	A, (FF)
0245	17	RLA	
0246	30 FB	JR	NC, 0243
0248	06 41	LD	B, 41
024A	10 FE	DJNZ	024A

Port FF is checked repeatedly by inputting the value to the accumulator and rotating that value into the carry flag. If no carry is found – i.e., no 'one' bit has yet triggered port FF – the program loops back to 0243. Once a bit is found, the B register is loading with 41, and a 'waste time' loop is executed at 024A (a total of just under 500 microseconds). A CALL is then executed to 021E. Let's have a look at that:

021E	21 00 FF	LD	HL, FF00
0221	3A 30 40	LD	A, (4030)
0224	A4	AND	H
0225	B5	OR	L
0226	03 FF	OUT	(FF), A
0228	32 30 40	LD	(4030), A
022B	C9	RET	

This curious subroutine seems to stumble through checking port FF for its video state, then resetting the OUTSIG flip-flop (see the *Technical Reference Handbook* for details on this circuitry). Isn't a byte ANDed with FF and ORed with 00 merely itself? True enough, but since this is also called as a subroutine entering at 0221, with a different value for HL, the complex AND/OR strategy makes sense.

So at this point we have picked up a bit from tape, delayed, and reset the flip-flop, readying it for the next bit to trigger it. Another delay loop follows (over 850 microseconds), and a byte is input to A from port FF:

0253	0B FF	IN	A, (FF)
0255	47	LD	B, A
0256	F1	POP	AF
0257	CB 10	RL	B
0259	17	RLA	
025A	F5	PUSH	AF

The input byte is saved in the B register, and the previously saved value of A is restored from the stack. Here is a wonderful piece of serial-to-parallel conversion – a sort of software shift register. Bit 7 of port FF was input to A and saved in B, and is then rotated left into the carry flag. Then the accumulator is rotated left, bringing the state of the carry flag into bit 0 of A. The accumulator is then saved once more on the stack. Another CALL to 021E resets the port FF flip-flop, both registers are restored, and the subroutine returns to the calling program.

You'll notice that at this point we have only one bit saved in the accumulator. An eight-iteration loop would be necessary to create a whole byte . . . and it will be done. But for the moment let's see how this routine is used in the initial syncing program, which we were about to enter at 0293.

The routine's first action is to CALL 01FE. This is a detailed routine to determine the drive number and other parts of the syntax, the state of port FF (again), select the drive and get it moving. Examining the code will show that it also uses the routine entered at 0221, but with a value of FF04 in HL; this routine won't be covered here, but it is worth looking at.

The find-sync-byte routine thus turns on the tape, saves the HL register, clears the accumulator, and calls the 'build-a-byte' routine at 0241. Since this is the synchronization process, no loop value is specified:

0297	AF	XOR	A
0298	CD 41 02	CALL	0241
0299	FE A5	CP	A5
0290	20 FB	JR	NZ,0298

It continually seeks bits, endlessly rotating the accumulator until it assembles a serial stream which matches A5 (i.e., binary 10100101 – nice and symmetrical). This routine is so accurate, in fact, that whenever tape motor start-up is not a consideration, the leader consisting of zero bytes would be unnecessary. The leading '1' of A5 serves as a kind of serial 'start bit' – and the routine at 0241 handles it from there.

Any kind of match to sync byte A5 might be found, tough, since the serial stream coming in from the tape does not distinguish start and end of byte. For example, the byte pattern DD 28 also contains an A5 embedded in it. As a serial stream, DD 28 is

```
1101110100101000
....10100101....
```

- with the A5 appearing at the junction of DD and 28. So once the matching A5 is found, a return is executed to the main SYSTEM loading module. That module then CALLs a subroutine at 0235, which is a gussied-up bit reader. BC and HL are saved, then:

0237	06 08	LD	B,08
0239	CD 41 02	CALL	0241
023C	10 FB	DJNZ	0239

There's the byte read . . . read a bit with eight iterations. HL and BC registers are restored, and the subroutine returns to the main program.

Loading the Code

The SYSTEM module now compares the byte it created with the value 55, the code assigned to machine language programs. It loops until it finds that code, then proceeds:

0208	06 06	LD	B,6
020A	7E	LD	A,(HL)
020B	B7	OR	A
020C	28 09	JR	Z,020E
020E	CD 35 02	CALL	0235
02E1	BE	CP	(HL)
02E2	20 E0	JR	NZ,02C1

Above, the B register is loaded with the number of characters to be found in the SYSTEM program's name. The accumulator is set up with the first character of the name as entered on the *? command line. The accumulator is tested for zero, and skips out of the loop when the end of the entered name is found. Each character following the name is read into the accumulator (CALL 0235) and compared with each letter of the entered name. If at any point the entered name does not match the name on tape, the program goes back to searching for 55 (machine program indicator) and the name search begins again.

There is a minor flaw in this process. Let's look at the succeeding lines of code:

02E4	23	INC	HL
02E5	10 F3	DJNZ	020A

This coding increments the HL register to the next character and loops back, looking for a total of six letters in the name. But what if the machine program code (55) is found, and one or more characters of the name match, but the rest do not match? There is no provision in this routine to decrement the HL register pair . . . which means that, if only part of a correct name has been found, the program will begin its search anew until it finds a program that matches only the last part of the entered name! This is the

reason the SYSTEM routine is not always able to search until it finds the correct program, the way the BASIC load does.

Let's assume the best – that a machine program was found with the name as entered from the keyboard. A CALL is then made to 022C, where the star or space at 3C3F is toggled (XORed) with 0A. Star XOR 0A is a space, and space XOR 0A is a star; easily done.

The SYSTEM Module Continues —

D2EA	CD 35 D2	CALL	D235
D2ED	FE 78	CP	78
D2EF	2B 8B	JR	Z,D2A9
D2F1	FE 3C	CP	3C
D2F3	2D F5	JR	NZ,D2EA

– searching for either 78 (end of program code) or 3C (beginning of data block code). If 78 is found, the program skips back to 02A9, where a CALL is executed to 0314. This subroutine merely reads the last two bytes on tape into the HL register, preparing the start address. This is saved at 40DF, the cassette recorder is turned off (CALL 01F8), and the SYSTEM module is re-entered from the start at 02B2. This module is a continuous loop, allowing a group of machine-language programs to be entered sequentially. Only the presence of the slash-start address combination will break out of the loop.

If a 3C is found, the beginning of a block of machine code is assumed. (If neither is found, the program loops until it finds one or the other). Here's a snippet of code:

D2F5	CD 35 D2	CALL	D235
D2F8	47	LD	B,A
D2F9	CD 14 D3	CALL	D314
D2FC	85	ADD	A,L
D2FD	4F	LD	C,A

A byte is read and saved in B. At 0314, two bytes are read and saved, respectively, in the HL register pair. These three bytes are, first, the number of bytes to read, and second, the two-byte starting address of the block. The 0314 subroutine leaves the value transferred to H in the accumulator; to it is added the value in L, and this number, sans carry value, is saved in the C register. The C register will be used to calculate the checksum for the block being read.

Curious Checksum

Each succeeding byte is read from tape and placed at the address now specified by HL. That byte is also added to the C register to update the simple checksum. HL is incremented to the next contiguous address, and the loop is iterated until B (the number of bytes to read in the block) reaches zero.

When the block is fully read, another byte is read from tape. This is the checksum byte, and should match the last updated value in the C register. If it does match, the program loops back, toggles the star, and begins anew the search for end-of-program (78) or block header (3C).

A correct checksum byte, curiously enough, is not a necessary element of the SYSTEM module. If the checksum is incorrect, the program will display a 'C' at video location 3C3E, and loop back regardless to continue reading the program from tape. I first noticed this action when a gentleman from New Hampshire called; he had been using a tape duplication routine to make a corrected copy of a machine language program. He had loaded the tape, returned to BASIC, then POKEd in a few byte changes. He then continued with the duplication. When he loaded the tape later on, he got a 'C' error message on the screen... but the program continued to load and did execute properly. The checksum was wrong because of the byte changes he had made, but the program, checksum notwithstanding, was read and loaded completely.

Let's take a look at that final portion of code:

D2FE	CD 35 D2	CALL	D235
D3D1	77	LD	(HL),A
D3D2	23	INC	HL
D3D3	81	ADD	A,C
D3D4	4F	LD	C,A
D3D5	1D F7	DJNZ	D2FE
D3D7	CD 35 D2	CALL	D235
D3DA	89	CP	C
D3DB	28 DA	JR	Z,D2E7
D3DD	3E 43	LD	A,43
D3DF	32 3E 3C	LD	(3C3E),A
0312	18 06	JR	02EA

Overall, these routines give the appearance of being reasonable and reliable, and they should be. What, then, gives rise to the tape problems? Mostly the timing loop in the 0235/0241 subroutine. The values placed in the B register at 0248 and 024F are too short for low-grade audio processing. Simply stated, the audio waveform coming in from tape 'rises' too slowly for the fast bit-check loop at 0251 to catch. A 'one' might come through, but it comes through too laggardly for port FF to have flipped into place.


```

00100 ; #####
00110 ; VOICE INPUT/OUTPUT ROUTINE USING THE CASSETTE PORT AND
00120 ; AMPLIFICATION. CAN BE USED WITH CTR TAPE RECORDERS AND
00130 ; BUILT-IN MICROPHONES OR PREFERABLY EXTERNAL CRYSTAL
00140 ; MICS. SMALL SPEAKER OUTPUT INCREASES INTELLIGIBILITY.
00150 ; #####
00160 ;
4300      00170      ORG      4300H      ; LOW POINT IN MEMORY
00180 ;      ORG      6500H      ; USE WITH OISK BASIC
06CC      00190 MONITR EQU      06CCH      ; BASIC EXIT (OR OTHER)
4300 F3    00200 START OI      ; NO BOTHERSOME STUFF
4301 C0C901 00210 CALL      01C9H      ; CLEAR THE SCREEN
4304 3A3040 00220 LD      A,(4030H)      ; START BY RESETTNG PORT
4307 03FF    00230 OUT      (OFFH),A      ; TO CLEAR INCOMING BITS
00240 ;
00250 ; #####
00260 ; KEYBDARO ROUTINE FOR ENTER (INPUT), CLEAR (OUTPUT), OR
00270 ; UP-ARROW (BASIC). UP-ARROW GOES TO EXIT IF NOT BASIC.
00280 ; #####
00290 ;
4309 3A403B 00300 KEYST LD      A,(3B40H)      ; GET ENTER/CLEAR ROW
430C FED1    00310 CP      1      ; CHECK IF ENTER PRESSEO
430E 2B0B    00320 JR      Z,INPUT      ; GO TO INPUT ROUTINE
4310 FED2    00330 CP      2      ; CHECK IF CLEAR PRESSEO
4312 2B50    00340 JR      Z,OUTPUT      ; GO TO OUTPUT ROUTINE
4314 FE0B    00350 CP      B      ; CHECK FOR UP-ARROW
4316 CACC06 00360 JP      Z,MONITR      ; OUT TO BASIC OR MONITOR
4319 1BEE    00370 JR      KEYST      ; BACK FOR A VALID KEY
00380 ;
00390 ; #####
00400 ; INPUT FROM PORT FF (255 DECIMAL) AND STORAGE IN MEMORY
00410 ; #####
00420 ;
431B 21A343 00430 INPUT LD      HL,MSG01      ; GET THE "INPUT" MESSAGE
431E C0A72B 00440 CALL      2BA7H      ; AND OISPLAY ON SCREEN
4321 3A3D40 00450 LD      A,(403DH)      ; GET VALUE FOR PORT MASK
4324 4F      00460 LD      C,A      ; SAVE MASK IN C REGISTER
4325 210044 00470 LD      HL,4400H      ; BEGIN VOICE STORAGE
00480 ;      LD      HL,6700H      ; BEGIN STORAGE (DISK)
432B 160B    00490 LOOP1A LD      0,B      ; NUMBER OF BITS IN BYTE
432A 0BFF    00500 LOOP2 IN      A,(OFFH)      ; GET VALUE AT THE PORT
432C CB17    00510 RL      A      ; STASH IT IN CARRY BIT
432E CB13    00520 RL      E      ; BUMP IT INTO E REGISTER
4330 79      00530 LD      A,C      ; GET VALUE OF PORT MASK
4331 3A403B 00540 LD      A,(3B40H)      ; CHECK ENTER/CLEAR ROW
4334 FE80    00550 CP      80H      ; CHECK IF SPACE PRESSEO
4336 C25343 00560 JP      NZ,ESCAPE      ; OUT IF KEYBDARO CLEAR
00570 ;
00580 ; ### NOTE: DELAY VALUE USED IN THE B REGISTER IS ###
00590 ; ### CHOSEN FOR OPTIMUM INTELLIGIBILITY WITH THE ###
00600 ; ### CTR TAPE RECORDER AND HARDWARE MODIFICATION. ###
00610 ; ### A LONGER VALUE CAN BE USED IF HIGH-FIDELITY ###
00620 ; ### INPUT IS PROVIDED. FOR EACH INCREASE IN THE ###
00630 ; ### B-REGISTER DELAY VALUE, ALSO INCREASE THE B- ###
00640 ; ### REGISTER BY THE SAME AMOUNT FOR PLAYBACK. ###
00650 ; ### LIKEWISE, A DECREASE IN THE DELAY VALUE MAY ###
00660 ; ### INCREASE FIDELITY AT A SACRIFICE OF MEMORY. ###
00670 ;
4339 0604    00680 LD      B,4      ; GET SHORT DELAY VALUE
433B 10FE    00690 DELAY1 OJNZ      DELAY1      ; AND DELAY A WHILE
433D 03FF    00700 OUT      (OFFH),A      ; MUST RESET PORT INPUT
433F 15      00710 DEC      0      ; DECREMENT TOTAL BITS
4340 C22A43 00720 JP      NZ,LOOP2      ; CONTINUE IF MORE TO DO
4343 73      00730 LD      (HL),E      ; SAVE FULL BYTE IN MEM.
4344 23      00740 INC      HL      ; GO ON TO NEXT BYTE
4345 7C      00750 LD      A,H      ; GET VALUE OF M.S. BYTE
4346 FE00    00760 CP      00H      ; USE FOR 4BK MACHINE
00770 ;      CP      0C0H      ; USE FOR 32K MACHINE
00780 ;      CP      0B0H      ; USE FOR 16K MACHINE
4348 C22B43 00790 JP      NZ,LOOP1A      ; IF NOT DONE THEN MORE
434B 21C743 00800 LD      HL,MSG02      ; GET "INPUT COMPLETE"
434E C0A72B 00810 CALL      2BA7H      ; AND OISPLAY THE MESSAGE
4351 1B86    00820 JR      KEYST      ; DONE - BACK TO KEY TEST
00830 ;
00840 ; #####
00850 ; PAUSE CHECK DURING ENTRY; SPACEBAR = GO, OTHERWISE STOP
00860 ; #####
00870 ;
4353 E5      00880 ESCAPE PUSH      HL      ; SAVE CURRENT POINTER
4354 CDAF0F 00890 CALL      0FAFH      ; OISPLAY CURRENT MEN.
4357 213944 00900 LD      HL,MSG05      ; GET "WORD START" MESS.
435A C0A72B 00910 CALL      2BA7H      ; AND OISPLAY THE MESSAGE
435D E1      00920 POP      HL      ; RESTORE MEMORY PTR.
435E 3A403B 00930 RECHK LD      A,(3B40H)      ; ENTER/CLEAR KEYBRO ROW
4361 FE80    00940 CP      80H      ; CHECK IF SPACE AGAIN
4363 2B03    00950 JR      Z,LOOP1A      ; BACK TO MAIN LOOP
4365 FE04    00960 CP      4      ; CHECK IF BREAK KEY
4367 20F5    00970 JR      NZ,RECHK      ; KEEP LOOKING ENT OR BRK
4369 21C743 00980 LD      HL,MSG02      ; GET "INPUT COMPLETE"
436C C0A72B 00990 CALL      2BA7H      ; AND OISPLAY THE MESSAGE
436F 1B8B    01000 JR      KEYST      ; AND BACK TO KEY MENU
01010 ;

```

Listing Continued . . .

Special Loaders

This was initially one of the mysteries of TRS-80 operations. *Microchess* was produced with a loader, then others quickly followed, mysteriously taking control of the machine and locking it up completely.

Let's now take a look at some of these special loaders, which will be designated Loaders A, B, C, and D in order to help them continue to do the job they were supposed to – protect software.

Loader A sets up a stack at 5000, clears the accumulator, and calls ROM to turn on the tape recorder and find the sync byte. It places a star on the bottom of the screen, sets up the HL register to receive the program, and prepares register C to perform simple checksum. A byte is read, it is saved in memory, and the checksum is created as in the SYSTEM mode. Then:

```

4025 7C      LD      A,H
4026 1F      RRA
4027 23      INC      HL
4028 3E 2A    LD      A,2A
402A DA 2F 40 JP      C,402F
4020 3E 20    LD      A,20
402F 32 F0 3F LD      (3FF0),A
4032 3E 4C    LD      A,4C
4034 BC      CP      H
4035 C2 1F 40 JP      NZ,401F
403B 3E FF    LD      A,FF
403A B0      CP      L
403B C2 1F 40 JP      NZ,401F
403E B9      CP      C
403F C2 00 00 JP      NZ,0000
4042 C0 FB 01 CALL      01FB
4045 C3 B0 47 JP      47B0

```

The strange appearance of RRA has nothing to do with rotating incoming bits. Rather, since the accumulator contains the H register value, each page (256 bytes) of information will change the high page value by one. Consequently, the high page will alternate between odd and even values, and the least significant bit, rotated into the carry flag, will trigger the display-star or display-space routines at 4D2F.

Finally, this somewhat awkward loader does a pair of compares to see if it has yet reached 4CFF, the end of the program load. If not, it loops back and continues; if so, it examines the checksum in C. Amazingly enough, it goes back to MEMORY SIZE? if there is a checksum error! There's no tampering with this program. A successful load is followed by a jump to the program's beginning at 4780.

Loader B is virtually identical to Loader A, except that the beginning of the program is found at 41FD instead of 4780.

Loader C is of a more interesting variety. It is written entirely without calls to ROM, because it

Continued Listing

```

01020 ; #####
01030 ; OUTPUT FROM MEMORY OF RECORDED VOICE TO CASSETTE PORT
01040 ; #####
01050 ;
4371 21E743 01060 OUTPUT LO HL,MSG03 ; GET "BEGIN OUTPUT"
4374 C0A728 01070 CALL 2BA7H ; AND DISPLAY THE MESSAGE
4377 3A3040 01080 LD A,(4030H) ; PORT FF OUTPUT MASK
437A 4F 01090 LD C,A ; SAVE OUTPUT MASK IN C
437B 210044 01100 LD HL,4400H ; START VOICE STORAGE (*)
437E 1608 01110 LOOP3A LO 0,B ; NUMBER OF BITS IN BYTE
4380 7E 01120 LD A,(HL) ; GET VALUE FROM MEMORY
4381 5F 01130 LD E,A ; SAVE IT IN E REGISTER
4382 AF 01140 XOR A ; CLEAR ACCUMULATOR TO 0
4383 C813 01150 LOOP4 RL E ; SEND BIT TO CARRY FLAG
4385 C817 01160 RL A ; AND ROTATE 'ROUND TO A
4387 B1 01170 OR C ; USE THE PORT FF MASK
4388 03FF 01180 OUT (OFFH),A ; AND SEND OUT THE VALUE
01190 ;
01200 ; #### NOTE: PLAYBACK VALUE BELOW MUST BE CHANGED ####
01210 ; #### TO MATCH SAMPLING DELAY IN THE INPUT SECTION ####
01220 ; #### OF THIS I/O PROGRAM. THIS VALUE IS ROUGHLY ####
01230 ; #### TWO TIMES THAT IN THE 8-REGISTER DURING THE ####
01240 ; #### INPUT SAMPLING. VARIOUS DUMMY OPCODES MAY ####
01250 ; #### BE INSERTED WHERE NECESSARY TO KEEP VOICE ####
01260 ; #### AT THE PROPER PITCH AND QUALITY. USING THIS ####
01270 ; #### PROGRAM, THERE IS A QUARTER-TONE DIFFERENCE. ####
01280 ;
438A 0606 01290 LO 8,6 ; GET SHORT DELAY VALUE
438C 10FE 01300 DELAY DJNZ DELAY ; AND DELAY SHORT WHILE
438E AF 01310 XOR A ; AND CLEAR ACCUM. BACK TO 0
438F 15 01320 DEC 0 ; BITS = BITS MINUS ONE
4390 C2B343 01330 JP NZ,LOOP4 ; AND BACK FOR SOME MORE
4393 23 01340 INC HL ; GET NEXT BYTE FROM MEM.
4394 7C 01350 LD A,H ; GET VALUE OF M.S. BYTE
4395 FE00 01360 CP 00H ; FOR 48K MACHINE
01370 ; CP 0C0H ; FOR 32K MACHINE
01380 ; CP 0B0H ; FOR 16K MACHINE
4397 C27E43 01390 JP NZ,LOOP3A ; AND GO BACK FOR MORE
439A 210E44 01400 LD HL,MSG04 ; GET "OUTPUT COMPLETE"
439C C0A728 01410 CALL 2BA7H ; AND DISPLAY THE MESSAGE
43A0 C30843 01420 JP KEYTST ; AND BACK WHEN DONE
01430 ;
43A3 48 01440 MSG01 DEFB 'HOLD SPACE BAR AND BEGIN SPEAKING.'
43C5 00 01450 DEFB 00H
43C6 00 01460 DEFB 00H
43C7 49 01470 MSG02 DEFB 'INPUT COMPLETE OR MEMORY FULL.'
43E5 00 01480 DEFB 00H
43E6 00 01490 DEFB 00H
43E7 42 01500 MSG03 DEFB 'BEGINNING PLAYBACK; BREAK IS IGNORED.'
440C 00 01510 DEFB 00H
440D 00 01520 DEFB 00H
440E 50 01530 MSG04 DEFB 'PLAYBACK COMPLETE; PRESS CLEAR TO REPEAT.'
4437 00 01540 DEFB 00H
4438 00 01550 DEFB 00H
4439 20 01560 MSG05 DEFB ' = WORD SEPARATION POINT.'
4452 00 01570 DEFB 00H
4453 00 01580 DEFB 00H
01590 ;
01600 ; #####
01610 END START
4300
00000 TOTAL ERRORS
29121 TEXT AREA BYTES LEFT

DELAY 438C 01300 01300
DELAY1 4398 00690 00690
ESCAPE 4353 00880 00560
INPUT 4318 00430 00320
KEYTST 4309 00300 00370 00820 01000 01420
LOOP1A 432B 00490 00790 00950
LOOP2 432A 00500 00720
LOOP3A 437E 01110 01390
LOOP4 4383 01150 01330
MONITR 06CC 00190 00360
MSG01 43A3 01440 00430
MSG02 43C7 01470 00800 00880
MSG03 43E7 01500 01060
MSG04 440E 01530 01400
MSG05 4439 01560 00900
OUTPUT 4371 01060 00340
RECHK 435E 00930 00970
START 4300 00200 01610

```

is capable of loading into a Level I or Level II TRS-80. Less fortunately, the ROM timing errors are not corrected, so the chances of loading this program on a marginal machine are not at all improved. The stack is prepared, and a block of memory is cleared from 5800 to the end of potential RAM at FFFF. My only guess as to the reason for this is that the authors wish to wipe out any programs such as monitors or disassemblers, as the clearing byte (A5) does not strike me as otherwise meaningful.

The tape is then turned on, and a pattern of three asymmetrical and two symmetrical sync bytes is found (B1, 83, 79, 5A, 00). Again, the choice strikes me as arbitrary, and may be the authors' way of identifying their own code. If these bytes are found, the program continues; if not, the entire five-byte pattern is sought again.

As in the other loaders, register C is set to zero for use as a checksum byte. The program load point is set high in memory (747F), and a byte is read. Here is a part of the code:

```

4330 C0 BF 43 CALL 43BF
4340 77 LO (HL),A
4341 32 3F 3C LO (3C3F),A
4344 B1 ADD A,C
4345 4F LO C,A
4346 2B DEC HL
4347 70 LD A,L
4348 3C INC A
4349 C2 53 43 JP NZ,4353
434C C0 8F 43 CALL 43BF
434F 89 CP C
4350 C2 66 43 JP NZ,4366

```

The secret to this portion of code rests in address 4346. Unlike most other loaders, this one loads (and displays) the last byte of code first, moving backwards through memory. (438F is the location of the byte-read subroutine). When the page is crossed (4346-4348), the checksum is evaluated; if the checksum is incorrect the program jumps to 4366, where an error message is displayed and the machine locks up.

The user's display is worth noting:

```

4353 7C LO A,H
4354 32 3E 3C LO (3C3E),A

```

This loader actually displays the ASCII equivalent of the page of memory being loaded with data . . . and it looks like an alphanumeric countdown as the program is fit into place.

Finally, Loader C does a comparison for the end of the first major load block, changes the value of H, and loads the next block. It then overwrites critical portions of the load routine, effectively obscuring the loading and entry point of the program. Interrupts are disabled, and the

Conclusions

process moves out of the loader into the main program. Interestingly, the authors forgot to turn the tape recorder off.

Finally, Loader D is of an entirely different sort. First, some code:

BEFE	3E 04	LD	A, 4
BF00	03 FF	OUT	(FF), A
BF02	0B FF	IN	A, (FF)
BF04	17	RLA	
BF05	30 FB	JR	NC, BF02
BF07	06 XX	LD	B, XX
BF09	10 FE	DJNZ	BF09
BF0B	06 09	LD	B, 9
BF0D	3E 04	LD	A, 4
BF0F	03 FF	OUT	(FF), A
BF11	0B FF	IN	A, (FF)
BF13	17	RLA	
BF14	00	NOP	
BF15	3B 0C	JR	C, BF23
BF17	23	INC	HL
BF18	2B	DEC	HL
BF19	10 F6	DJNZ	BF11

This remarkable loader is written for high-speed operation, setting up the output ports

BEFE and BF0D), clocking itself with start bits (BF0D), and then reading a nine-bit serial stream. Careful timing and self-clocking are essential in high-speed data I/O, and this routine is capable of reading and writing on ordinary audio cassettes, with excellent reliability, at better than 2000 baud. The only point to the instructions at BF17 and BF18, for example, is the delay introduced by executing them; yet that timing is very important. The actual timing value at BF07 has been dropped for a measure of protection of this author's fine software.

Conclusions

In sum, the tape read/write routines of the TRS-80 are efficient and, especially now with special loaders and a corrected ROM, quite reliable. Different levels of user prompts, particularly those used by the reverse-loading module described above, are probably more satisfactory than flashing stars. A checksum process for BASIC similar to the SYSTEM module would have been valuable. Finally, by careful attention to clocking details, a reliable, higher speed loader could have been included in the TRS-80.

For those especially interested in high-speed loaders, I recommend examining the *Exatron Stringy-Floppy* operating system, which shows what can be done with equipment designed for digital operation. It is capable of reliable loading and saving at rates exceeding 11,000 baud.

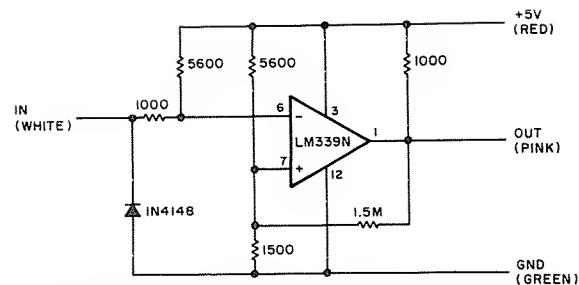


Fig. 2. Full schematic for the Model I cassette modification for speech input. It should be switched out when cassette programs are being loaded (see Fig. 4).

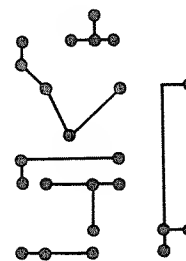


Fig. 3a.

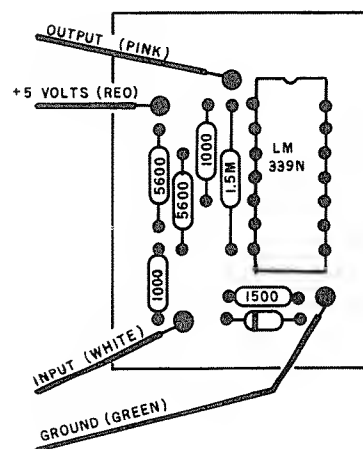


Fig. 3b.

NOTES

4

Simple Modifications

Why hardware modifications? Simply because they help a limited computer broaden its ability to serve our needs. Looking at glowing blue-white letters on a dark screen is comfortable for only a limited time . . . waiting through many minutes for the computer to make a game play is frustrating . . . attempting to imagine which characters are upper case and which are lower case is nearly impossible . . . emphasizing individual words and characters is only possible with arrows and stars . . . and so forth.

In this chapter, eleven simple modifications to the standard TRS-80 will be presented:

- Installing 16K memory to the keyboard unit, and to the expansion interface.

- A change to enable recovering Resets when there is an expansion interface attached.

- Moving the Reset button to a more accessible position.

- Adding both an extra keyboard and an extra video output for extending the computer's portability.

- Adding an RF modulator to further increase the TRS-80's portability.

- Installing a modification to increase the computer's computational speed.

- Installing a modification to reverse the video display for black letters on a white background, more in keeping with real text.

- Adding the feature of dual-language operation: Level I and Level II in the same computer.

- Accessing the lower case capabilities built into the TRS-80 by adding a single circuit.

- Reversing individual characters on the screen for emphasis.

- Adding a hexadecimal keypad for faster entry of machine code information.

Almost all of these modifications will require you to open the cabinet of your TRS-80. Should you do it? Or should you not? First of all, consider that your computer has but a three month warranty; after that, you will have to pay repair charges should something go wrong. Radio Shack has, in response to public protests, changed its former policy and will now repair machines with the most popular modifications installed, at no extra charge. They will not remove those modifications if you say so, but may refuse repair if (a) you have mangled the board, or (b) if you cannot document your modifications.

Yes, there is a chance that you can damage your TRS-80 when you make modifications. This will not be the fault of the modifications themselves, though. My own TRS-80 has been modified many times, and the only damage has been to the notoriously flakey cable that connects the keyboard to the main circuit card, and the failure of one RAM chip. The latter would have happened anyway.

Keep It Clean!

How many TRS-80's have been damaged during modification? That's a hard question to answer, but I will describe briefly the failed 80's I have seen:

1. A damaged integrated circuit in the data section was caused when heavy wire was used to make the modification.

Lesson — use the parts specified.

2. Two RAM chips were blown when the user dropped a soldering iron into the unit when he was modifying it.

Lesson — always turn the power off when making the physical changes.

3. A blown power supply regulator. The user had left the unit under an open window.

Lesson — close the windows in case of rain!

4. High levels of garbage causing keyboard lockout and blown programs. The computers were being used next to heavy electrical equipment.

Lesson — treat the computer as if it were a very sensitive piece of electronic equipment. It is.

5. Constant keyboard lockout or odd characters. The TRS-80 cable had been flexed too many times, causing breaks. These breaks had destabilized the keyboard circuits, causing two ICs to blow.

Lesson — handle the keyboard cable, that most delicate of hardware, with utmost care.

6. Constant lost memory and programs. The RAM chips were not being 'refreshed' because the refresh multiplex line was not working. The user had plugged in a peripheral device upside down, with the power on to both.

Lesson — make all interconnections with the power off.

In sum, all the above damage was caused by carelessness, haste, or attempting to use inappropriate materials to do the job. The solutions? Resolve not to do all the work in one evening. Turn back to the introduction to make sure you have the right tools. Work slowly and take breaks often. Buy parts from reliable suppliers. Read and understand the instructions *and the theory* before you start.

And finally, if you have any serious doubts about the accuracy of the printed information, contact the author. That's me. And if you have any problems that won't go away, write. If you write, you *must* enclose a self-addressed, stamped envelope, a complete description of any problem, and all tests you have made. Also, I cannot cover much beyond the scope of this book, which includes the myriad competing disk operating systems and support software.

Expanding the Memory

The simplest modification, once you have braved opening the cabinet, is expanding the unit's memory from the 4K supplied to a full 16K RAM. At this writing, 16K of reliable RAM memory can be purchased for less than \$20.

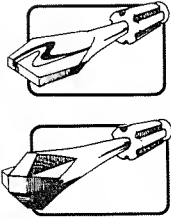
Keep It Clean!

If you examine the contacts on the keyboard unit's edge-card connector, as well as on the five connectors (expansion in, expansion out, printer, RS-232, disk) of the Expansion Interface, you will see a major cause of the TRS-80's instability: solder-plated connectors. In the interest of economy, Radio Shack did not use gold to ensure a good cable contact. This turned out to be a serious error in judgment.

There are two options in dealing with these solder-plated connectors. The first is to keep them clean. Remove the cables regularly and vigorously rub the contact surfaces with a dollar bill or talc buffing wheel. Bring the

solder to a bright shine, and spray it with contact cleaner. Reinstall the cables and memory crashes should lessen.

The other option is expensive and time-consuming, but much more reliable. The contacts can be freed of solder plating with solder wick, cleaned with flux remover, brushed to a high shine, and a soft silver compound can be flowed onto the edge contacts. Fuller Software (see Appendix I) sells a contact plating kit; price varies with the price of silver, but at this writing it is more than \$20. The process is very tedious because silver melts at a much higher temperature than lead, but the results are noise-free connections and greater reliability.



There are several types of so-called '16K memories', so when you set out to upgrade your TRS-80 memory, make sure you order the correct type. 16K is actually a shorthand term for 16,384 bytes of memory. In the TRS-80, the 16K memories are integrated circuits containing 16,384 single-bit memory cells each. To create an entire byte, then, eight integrated circuits are needed. Furthermore, the TRS-80 memory needs very little power, and must take only a small amount of space in its cabinet. The only small, low-power memories made are 'dynamic' memories. The 16K-by-one-bit, dynamic memory is industry type 4116, also called type 416.

There is one other consideration in purchasing memory expansion chips, and that is a popularly misunderstood quality called *access time*. Access time can be thought of as the time it takes the computer to inform the memory chip that it needs information (or will give it information) until the memory chip is electronically ready to respond. This figure is usually given in *nanoseconds* (see Table 4-1).

Second	1
Millisecond (mS)	0.001
Microsecond (uS)	0.000001
Nanosecond (nS)	0.000000001

Table 4-1. Relative time units.

This time is very small but quite critical to the operation of the computer. If the memory is not ready, programs and data will be recorded or reported incorrectly, and the computer will have no chance of working properly. The minimum access time for memory in an unmodified TRS-80 is 450nS; with the speed-up modifications presented in this book, that figure drops to 300 nS or less.

Most of the current crop of memory chips easily meet the 450 nS requirement, and many of those sold as 450 nS chips can be run reliably at 300 nS. A problem arises with the expansion interface, as older units had a bit of trouble with 'hotter' (faster) memory. So as a general rule, refer to Table 4-2 when looking for memory to upgrade your TRS-80 or its expansion interface.

Equipment	Access Time
TRS-80 keyboard unit, unmodified	450 nS mex, 350 nS optimum
TRS-80 keyboard unit with speed-up	300 nS mex, 250 nS optimum
Expansion Interface made before 1/80	300 nS min, 450 nS mex
Expansion Interface made after 1/80	450 nS mex, 350 nS optimum
Expansion Interface made before 1/80 with speed-up	300 nS min optimum
Expansion Interface made after 1/80 with speed-up	300 nS mex, 250 nS optimum

Table 4-2. Memory access time for various TRS-80s.

Opening and Closing the Case.

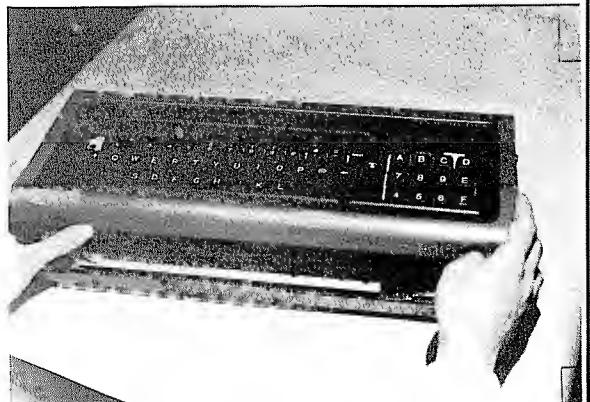
1. Locate a spacious workplace, and set a soft towel on it. You will need a Phillips screwdriver and a small box in which to set the screws and spacers.
2. Remove the power and disconnect all cables to the keyboard unit. Set it face down on the towel.



Photo 4-1. Photos showing opening computer.

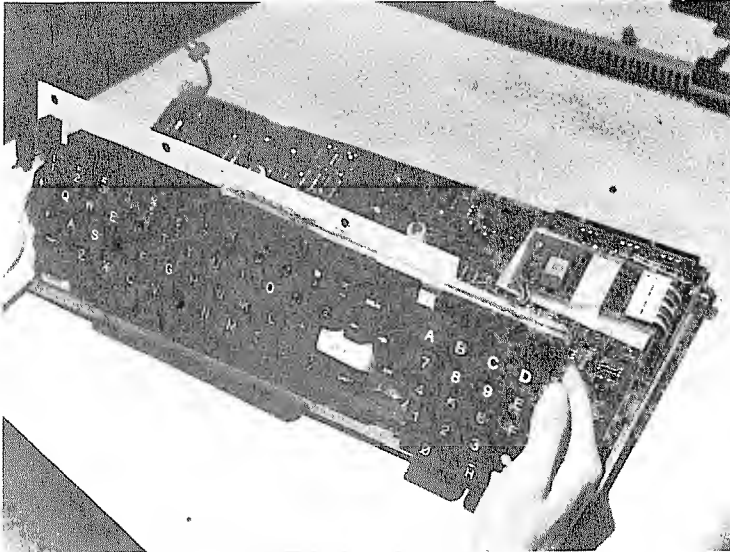
Unit is placed face down and screws are removed. On newer units one screw is beneath a warranty notice.

3. There are six screws of three different sizes which hold the case together. In later TRS-80's, one of these is covered by a label warning you not to go inside. This label just won't peel off for later use. You'll have to punch through it to remove the last screw.
4. Hold the case together and place it face up on the towel. Gently lift off the top cover. Some TRS-80's have a 'flying lead' LED (light-emitting diode) power indicator, meaning you will have to pull it gently out of the hole in the top cover. Other power indicators are fastened to the keyboard, and the cover lifts off directly.



Computer is placed face up and the top removed by lifting back and up.

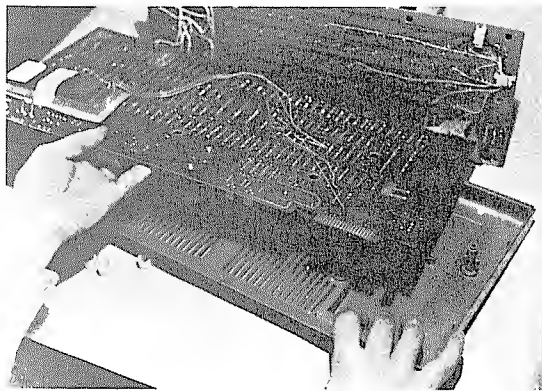
5. This step is the most delicate. Lift the keyboard slightly upward, and then swing it toward you. A cable attaches from the keyboard to the main circuit card at the bottom left side. This cable is made up of flat copper bands which have a tendency to break when flexed.



Keyboard is rocked gently forward to show white plastic spacers, which are then removed.

6. While holding the keyboard up at about a 90-degree angle, look at the main circuit card. There are five soft plastic white spacers which cushion the keyboard. On later models, a sixth, central spacer is hard plastic. Note their positions, remove them, and set them in the parts box.

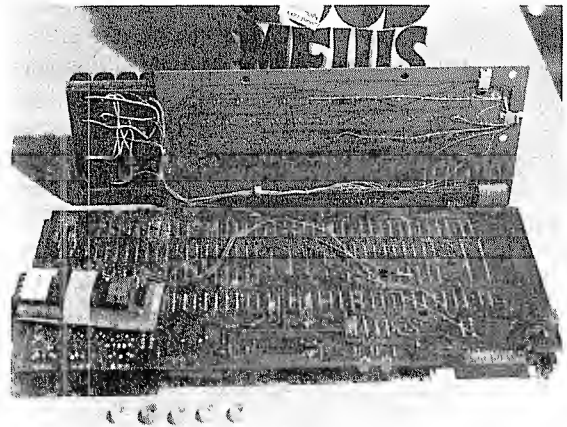
7. Swing the keyboard downward gently, and lift the entire computer out of the base of the cabinet. Set the computer down and put the base aside.



With the keyboard supported upright against a firm backing, the CPU board is lifted from the case bottom.

8. Most modifications will be done to the integrated circuit side of the circuit board, so turn the entire unit face down again. Unless the modification calls for work on the solder side of the board, do not flex the keyboard cable again.

9. If the modification calls for work on the bottom of the board, get a heavy box or other support. Set the keyboard face up, and swing the keyboard out and up a bit wider than 90 degrees, leaning it against the support. Never open the unit fully like a book, as this badly deforms the interconnect cable.



The complete unit is ready for work with the keyboard supported at a 90-degree angle.

10. To close the case, fold the unit back together, and set it into the base. The unit may have to be jostled gently to get it to fit over the plastic support posts. Make sure any added wires do not get caught and cut by the support posts.

11. Lift the keyboard slightly and restore the white spacers to their former positions. Fit the keyboard back into place.

12. If any long wires or cables have been added, make sure they all clear the cabinet edges. Restore the LED to its place on the keyboard top if it is a flying-lead type, or straighten the LED on the circuit card so it fits into the cover hole.

13. Fit the cover into place lightly, making sure there are no newly installed parts being crushed or bent in the process. Be sure no leads creep out the joints on either side.

14. Holding the unit together firmly, flip it on its face. While holding it with one hand, drop the screws into the holes, longest one towards

the top. Fasten one screw on either side; this will hold the computer together while you tighten the remainder.

15. Remember that the case is soft plastic, so just tighten enough to pull the two sides together. Otherwise, the plastic may be stripped and the screws will fall out. If that happens, drop in some clear acrylic cement (not white glue), and insert the screw. Remove it just before the glue hardens, and replace it about an hour later.

16. Flip the computer on its back, restore cables, and turn on the power. If the slightest problem seems apparent, open up and try again!

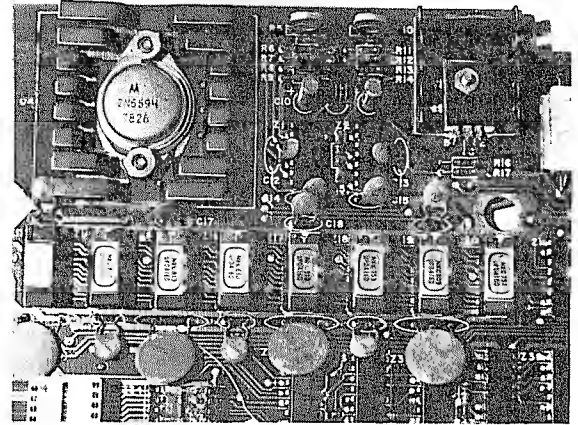


Photo 4-2. Memory chip area in TRS-80.

Power Supply and Memory: 2N5594 transistor handles 5-volt supply; adjustments are seen at top of photo. 16K RAM chips are plugged into sockets Z13 to Z20, unusually close to the power transistor's heat sink. Replacement memory chip in socket Z13 attests to the degrading capacities of excess heat.

Tips on Handling Integrated Circuits

In the early days of microcomputers, there was a lot of user hesitation about installing memory chips because of warnings about static electricity damaging the memory devices. At that time the fear was reasonable; but today (with just a little caution) there need be no problem.

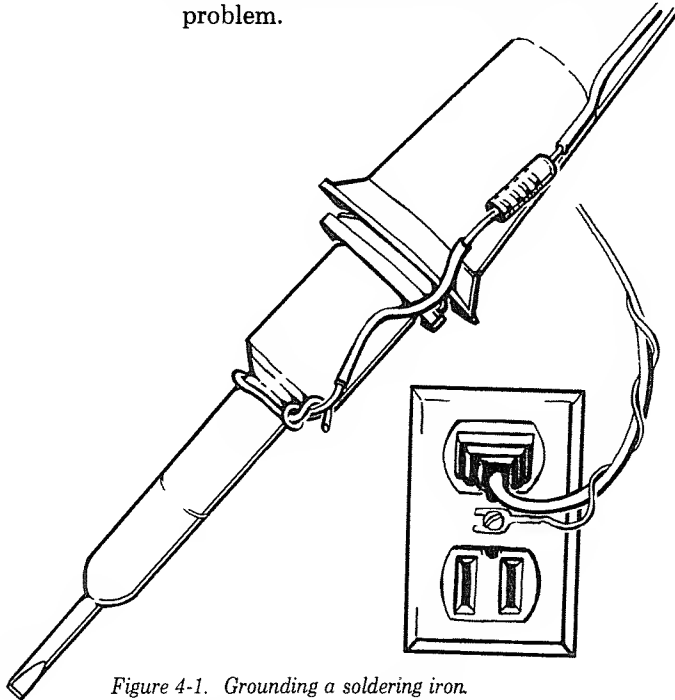


Figure 4-1. Grounding a soldering iron.

1. Never place any integrated circuit on highly charged plastic material, especially styrofoam.

2. Handle memory chips, CPUs (such as the Z-80), LSI devices (large-scale integrated circuits, usually those with 28 or 40 pins), or any marked MOS, CMOS, or NMOS (metal-oxide semiconductors), with care. Hold them by their ends, never by the connection pins.

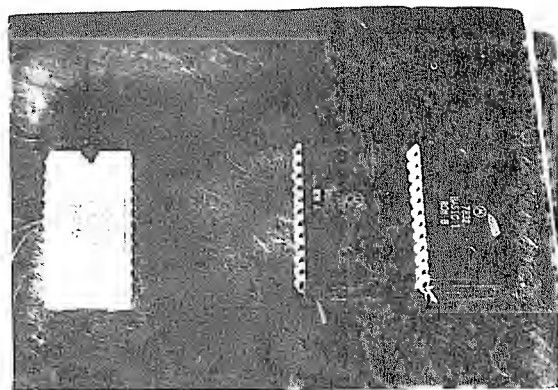
3. Purchase a static-free workbench, which is a conductive cloth sheet with a wrist strap and safe grounding cable. These can be obtained from Wescorp for about \$18.

4. Ground your soldering iron to an earth ground *but only through a series-connected one-megohm resistor — never directly!* The grounding is not absolutely essential, but helps if you live in a very dry, static-producing environment.

5. Work with any integrated circuits with the power off. Make sure the integrated circuit's ground and power pins are all connected (soldered or in sockets) before turning on the juice! A difference of a mere half a volt between certain pins can kill an IC.

6. Use high-quality sockets for integrated circuits wherever you can. This will not only keep excessive heat away from them, but will also save the day if one is damaged. Unsoldering a 40-pin integrated circuit is not pleasant.

7. Above all, work slowly and carefully. By far the greatest villain is haste. Oh yes — do keep furry animals out of the area!



Level I ROMs: Rockwell single-chip ROM and Motorola 2-chip set are pushed into aluminum-foil-covered vegetable tray.

Before installing your new memory chips, take a styrofoam meat or vegetable tray, trim off the curved ends, and cover the center with aluminum foil. This will be your static-free storage for the 4K memory chips you will be removing. To install the 16K memory in your keyboard unit, turn off the power, open the case and find the 4K memory chips.

Slide a thin-bladed screwdriver under the end of one of these chips, and rock it slightly upwards. Slide the blade under the other side, and rock. Move back and forth gently until the chip is free, but don't spring it out of the socket. Lift it by the ends and press it into the foil-covered tray. As you are doing this, notice that each of the chips has a notch or dot at one end. Keep this position in mind; the 16K chips will be installed in the same direction.

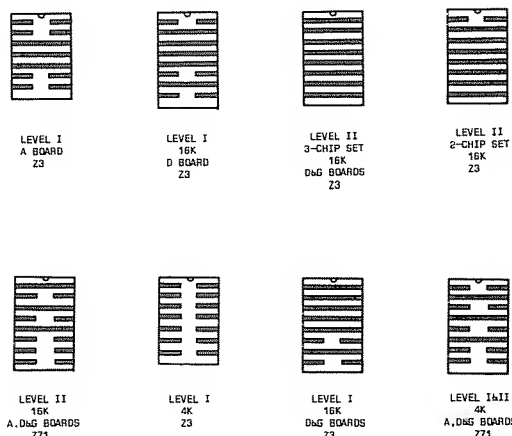


Figure 4-2. 16K RAM expansion shunt versions.

Once all the 4K chips have been removed, lift the 16K chips one at a time and press them in the empty sockets. If the pins are spread too wide, set the chip sideways on the foil-covered tray and press gently. Try again to insert the chips. Be sure none of the pins are bent underneath, or slide outside the socket.

Now turn your attention to a pair of on-board sockets marked Z3 and Z72. 16K chips will need more attention from the address lines, so these decoding shunts will have to be changed. If you didn't receive these shunts with your order of 16K chips, have no fear — just slide wires into the sockets. Or, obtain an 8-position DIP switch and push it into the socket. Turn the switches *on* where the shunt is to be connected, *off* where the bars are broken. Figure 4-1 shows the old and new positions of the shorting bars or wires of these shunts.

When you have placed the new, corrected shunts, switches or wires into sockets Z3 or Z72, and all the 16K chips are installed, the keyboard upgrade is complete. Reinstall the computer in its case, turn on the power, and press (ENTER) in response to MEMORY SIZE? Then PRINT MEM, and you should see a value of 15572 (15570 on newer machines). If you don't get that value, turn the machine off and troubleshoot:

1. MEMORY SIZE? reads much less than 15570. Turn off the computer and try again.

If you have no expansion box, type SYSTEM (ENTER) /0 (ENTER), and keep trying. If the value never changes, you may have either a bad memory chip or incorrect shunt wiring. Go back to MEMORY SIZE?, only this time enter 15560. If you get the READY message, this might point to a failed memory chip or an incorrect DIP shunt (especially if PRINT MEM? reads the same as that for 4K). To test memory, run the RAM test printed in Chapter 3. If instead you get the flash of an ?OM ERROR and a return to MEMORY SIZE?, then suspect that you've wired the shunt or shorting wires incorrectly.

2. You get a partial RADIO SHACK LEVEL II BASIC (or R/S LEVEL 2 BASIC) message, with or without READY, and with or without incorrect characters, but it only lasts for a short time before crashing back to MEMORY SIZE?

Suspect that a memory chip is very balky, is inserted only partly or with pins bent, or that you've lifted the Level II interconnect cable (if your unit has one). You may also have damaged some other circuitry, but this is very unlikely.

3. The screen never gets past a pile of garbage.

You may have lifted the Level II interconnect cable (if your unit has one), one or more memory chips may be completely dead, inserted backwards or only partly, or you may have forgotten to reinstall (or have reversed) either of the two shunts at Z3 and Z72.

4. Unexpected characters are displayed after MEMORY SIZE?, sometimes acting as if they were 'entering' themselves.

You have broken one or more wires of the keyboard interconnect cable. You can look for cracks, or just replace the whole cable.

5. The machine responds correctly, but only for a short while; it often crashes; occasionally PRINT MEM will give a smaller number than 15570, but not always.

This is probably balky memory or memory that is the wrong speed (usually older, slower memories that some discount houses may sell). For starters try reinserting the memory in case of a bad contact; run the RAM test; or just buy new memory.

6. The computer displays a screen full of 999, etc.

You have lifted the Level II interconnect cable out of its socket. Replace it very carefully.

Adding to the Expansion Interface is a much easier task. Your keyboard unit must have 16K in it already in order that the memory map be complete from 4000 (decimal 16384) to the start of expansion box memory at 8000 (decimal 32768). And, sadly, you cannot use your 4K chips in the expansion box without a hardware modification.

First, remove the cover over the power supplies and remove them; this will prevent them from tumbling all about when you open the bottom cover. Now flip the expansion box over and remove and set aside the six screws that fasten the cover. Also disconnect the power cable inside the expansion interface case.

Inside, you will find two rows of empty sockets for memory expansion. The first 16K of expansion memory goes in the sockets marked Z9 to Z16, and the second 16K into sockets Z1 to Z8. The memory must be inserted in this order, unless you want a permanently protected, 16K, high-memory block (which might be useful). Use the same procedure for installing these memory chips as for the keyboard unit, facing them in the direction of the notch on the sockets. Once again, check carefully for bent pins or pins

out of the sockets, reinstall the cover, and power up the interface and keyboard.

Press ENTER in response to MEMORY SIZE?, and your 32K machine should read 31956, and the 48K machine will read 48340 (two bytes less each in later models). Early expansion boxes, because of design flaws in memory timing, are significantly more sensitive to memory speed. If expansion memory is occasionally balky or shows frequent glitches when peripheral devices are attached, make the hardware changes to Z69 recommended in the 200% speed modification (later in this Chapter). Seeming memory failures can most often be attributed to these timing problems, although earlier interfaces (particularly those with the bulbous buffered cable) had hardware difficulties which made them extremely sensitive to noise and vibration.

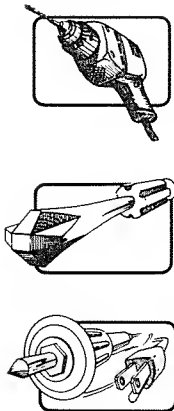
Most of these earlier units had their circuit board layout and plating done in such a way that a sharp tap on the box, board, or cable would cause a 'microphonic' reaction. That is, the vibration would be transmitted along power supply and signal lines, interfering with the actual data. The result would be frequent memory crashes. Likewise, a noisy environment (nearby washers, mixers, fluorescent lights, transformers, and even printers) can cause electronic interference which would disrupt memory.

Rescuing the RESET

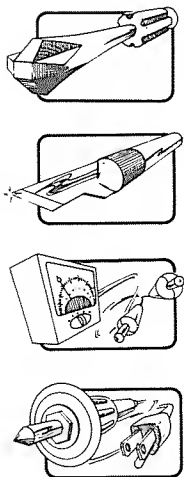
Among the conveniences of the TRS-80 keyboard computer is the Reset button. A program, especially one with machine language components, may cause the computer to 'hang'. The Reset button conveniently recovers control of the machine and returns it to you.

Once the Expansion Interface is connected, though, things begin to change. The Reset button becomes a Reboot button, causing any operating programs in memory to be lost and the complete system to restart from initialization of the disk operating system (see Supplement to Chapter 1, on the power-up sequence).

As an aside, let me note that the Z-80 HALT instruction does not have the effect of a true HALT. Instead, the CPU's Halt Acknowledge output line is tied in with the Reset button. The result is a READY in Level II and a disastrous reboot with an expansion box connected. This is another good reason for reasserting the Level II reset function with this modification.



Up-Front RESET



The solution to this problem is to disable the disk controller whenever disk access is not expected. Open the expansion box, and locate Z32, near the power switch. If you have a newer expansion interface, this circuit will be marked Z39. This is a 16-pin circuit, type 74LS155. Identify the circuit trace that runs from pin 4 underneath the IC and out the opposite side. Use an ohmmeter if necessary to make sure you have the right trace. This signal activates the disk controller chip's output to the CPU; when it is cut, the keyboard unit cannot 'see' the disk controller.

Take a sharp blade and cut this trace. Solder a 10K ohm resistor from the *far side* of this trace to pin 16 of Z32 (or Z39). This is the +5 volt lead, and will hold the pin high.

Next run a pair of fine wires from either side of the cut trace to each connection of a small toggle switch. When the switch is on, the cut trace is bridged, and the disk controller buffer can be activated normally; when the switch is off, the Reset button sends the software to a routine that checks for a disk controller. Since it does not 'see' the controller, it acts as if it were simply in Level II BASIC and returns to READY.

Be sure to mount the switch as close as possible to the trace cut, preferably right on the front of the expansion box, as shown in Photo 4-1. This will prevent noise from creeping in to an already somewhat noisy box.

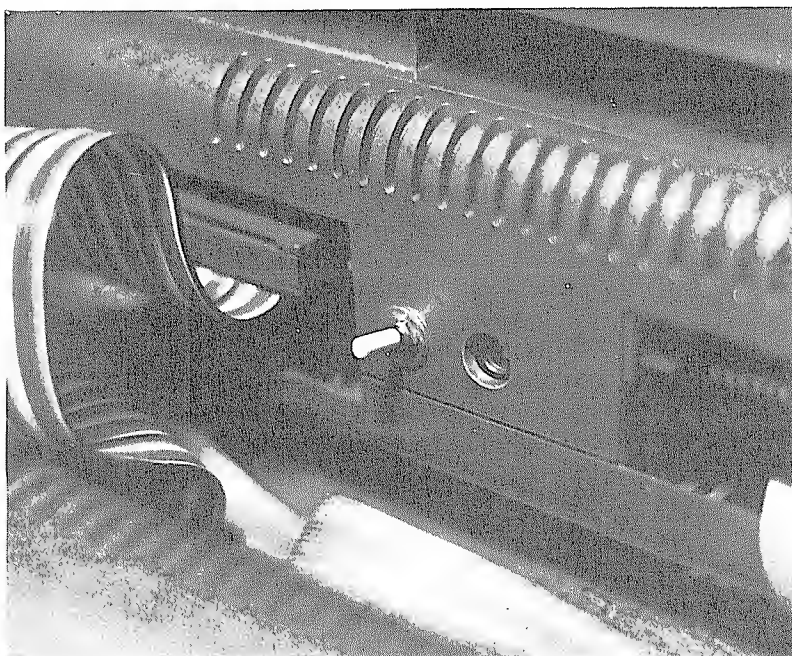


Photo 4-3. Expansion box reset modification.

Visible in this photo are a 12-inch replacement expansion cable, and the disk-defeat switch to recover the reset function.

The LNW expansion interface change would normally be identical to that for the Radio Shack box, except for the integrated circuit numbers. One additional change is necessary.

The trace from U19 pin 4 leading to U8 and U15 is cut, and a 10K resistor wired from the trace end closest to U8 and U15 to +5 volts (found at U19) pin 16).

However, the pullup/pulldown resistors in the LNW expansion box can still give an 'on' reading to the CPU. To avoid this, change the pulldown resistors from 220 ohms to 470 ohms (or, if they already are 470 ohms, from 470 ohms to 680 ohms). This will result in the 'high' reading needed to avoid picking up the disk controller signal.

Up-Front RESET

If you are a frequent user of the Reset button in Level II — and once again a user of it with your expansion box — then you will want to get the button out of the area of the sensitive interconnect cable, and well within reach. It will be a welcome relief from clawing at the silver port cover, or totally wiping out your program by jostling the cable.

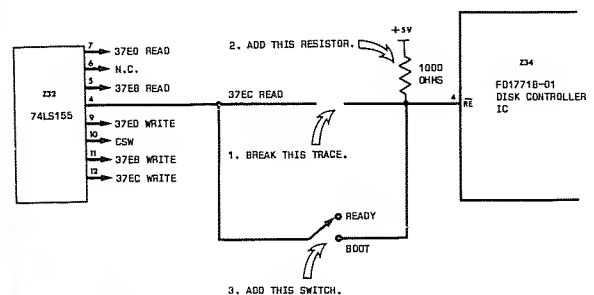


Figure 4-3. Expansion box reset modification.

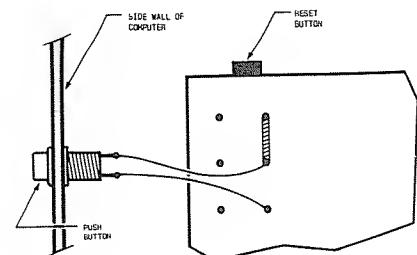


Figure 4-4. Up-front reset addition.

A momentary-on pushbutton (such as Radio Shack part number 275-1547) can be added to the cover of the keyboard unit. Photo 4-2 shows the position of the Reset button on the left side of the computer. Run two wires from the pushbutton to a small cable connector (a submini plug), and run two wires from the Reset

button (see Figure 4-2) to the other end of the cable connector (a submini jack). The project can be completed in ten minutes, carpentry and all, and reset will be less frustrating.

Note : Don't risk being a victim of the Apple syndrome! Apple's Reset button is placed much too close for comfort to the user's work area, and many a program has disappeared into the electronic stratosphere when inadvertently pressed while typing. So keep that Reset button just out of reach!

Working by the Woodstove

There comes a time when sitting up straight by the computer is no longer fun. Or when the neighborhood kids howl because they can't all reach the keyboard at once during an action game. Or when those kiddie hands are just too sticky for your sacred micro. Or, in my case, when the computer's room is just too cold to share with my typing fingers. That's when you need a keyboard and monitor in the room by the woodstove. Or an extra keyboard for the young'ns. Or a keyboard for the lap in an easy chair and a monitor on the mantle.

The additional keyboard is mostly a matter of carpentry, because there's nothing special about the TRS-80 keyboard. It's merely a matrix of switches, eight by eight. Each position in the matrix is identified by the computer's software and turned into a character.

Start by obtaining two high-quality wire-wrap integrated circuit sockets, and one good solder-tail type. These are 16-pin sockets. You will also need fine wire, a 16-wire jumper cable with plug attached, and a keyboard.

The keyboard can be any style you like, from a complete alphanumeric keyboard (\$40 to \$120), to a \$10 numeric keypad if you work mostly with numbers. Whichever you choose, it must consist of individual keys, each with a single-pole, single-throw (SPST) contact pair. Many small calculators have a prearranged matrix which is incompatible with the TRS-80. If you choose a matrix keyboard, check that it will work with the TRS keyboard pattern shown in Table 4-3.

@	A	B	C	D	E	F	G
H	I	J	K	L	M	N	O
P	Q	R	S	T	U	V	W
X	Y	Z	...	Unassigned...			
0	1	2	3	4	5	6	7
8	9	:	;	,	-	.	/
ENT	CLR	BRK	UPR	ONR	LFR	RTN	SPC
SHIFT	Unassigned.....					

Table 4-3. Keyboard matrix.

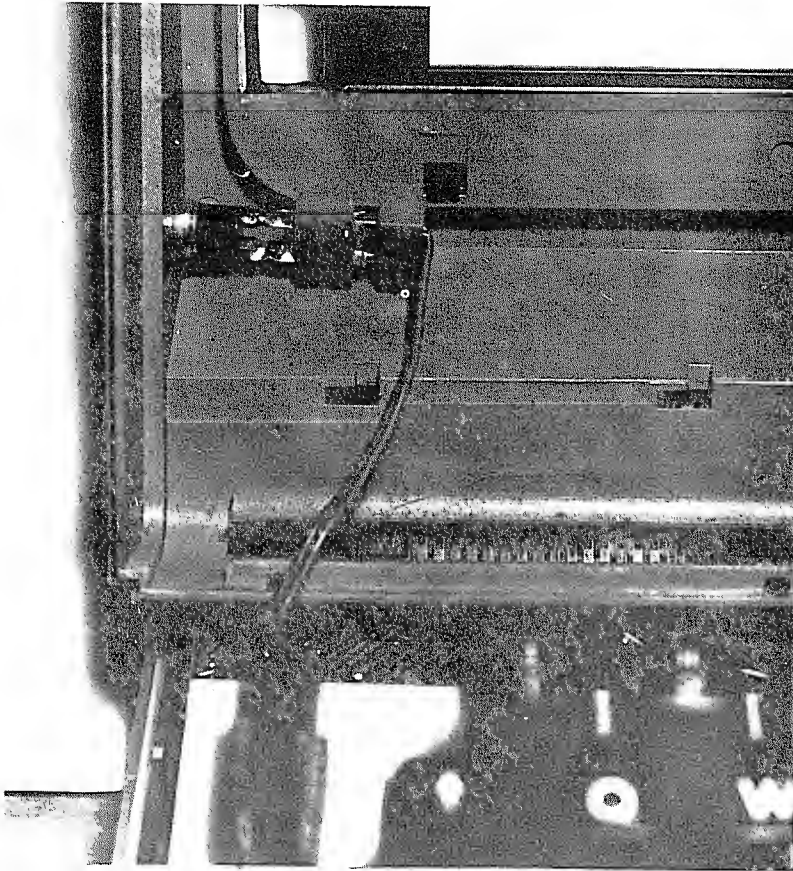
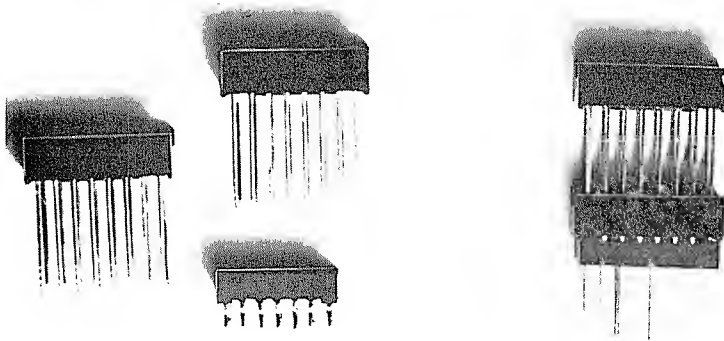
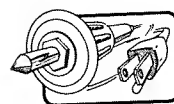
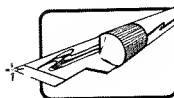
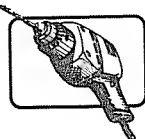
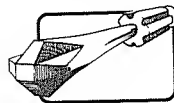
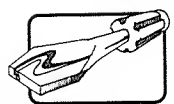


Photo 4-4. Up-front reset addition.

Reset switch added to the top cover. A connector is added so the cover can be removed easily.



Piggybacking sequence: pins are removed from short solder-tail socket, and it is used as a grommet for a wire-wrap socket. After this is soldered to the baseboard, a third socket is plugged in place.



Since depressed keys are identified in software by the row and column, you only need to know which column-row combinations produce the letter you want. That way, you can reassign your keyboard's characters for any purpose that suits you — including the Dvorak keyboard. Thus, you need make no software modifications to your favorite machine language programs to use them with different keyboard combinations.

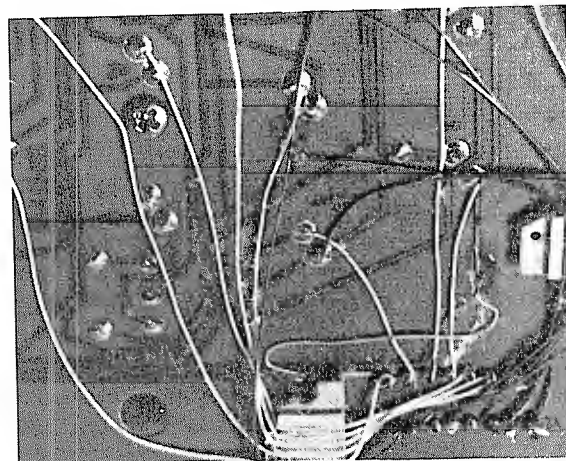
Furthermore, the attachment of a 64-key musical keyboard can open the door to direct compiling of music as you play it.

The physical layout of the TRS-80 keyboard unit is fairly compact, leaving only a space on the far left or far right for the added keyboard connection. I have chosen the left side for that addition. Inside the computer, this location is directly above a blank part of the keyboard's circuit card.

Use a strong flat screwdriver to snap out the black portion of the keyboard cover. Six tabs hold it in place at the top and bottom. Mark precisely where the free area can be found on the baseboard.

You will be using the three IC sockets to make a standoff-style keyboard connector. Pull all the pins from the solder-tail socket, and use this socket as a guide to drill 16 holes in the baseboard. Use a very fine hobby drill — #68 is good. When you have the holes completed, slide

one of the wire-wrap sockets into the disemboweled solder-tail socket, and feed the wire-wrap pins through the circuit board. Fasten with fast-drying epoxy; do not use white glue, as this will react badly with metal.



Plastic carriers from flat-pack integrated circuits make excellent 'bridges' to hold wires in place. A drop of glue holds them there.

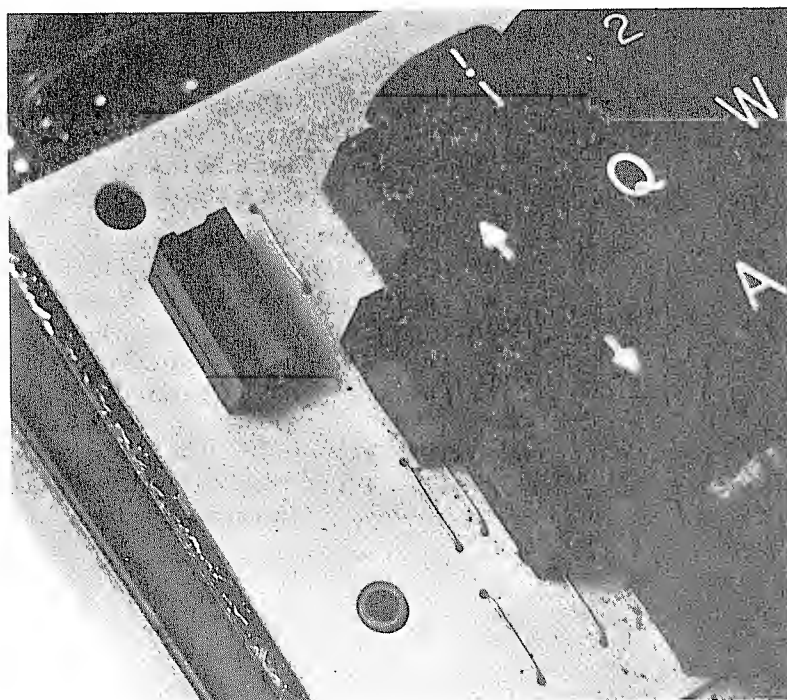
When the glue is set, remove the entire keyboard cover, turn the baseboard over, and identify pin 1 of the newly installed socket. This pin will attach to column one of the keyboard matrix (see Table 4-4). On most versions of the TRS-80, you can use the keyboard's resistors to identify the columns; I recommend this, because there were at least three separate runs of keyboards, each with a different layout.

Column 1	R8	@ H P X O 8 ENTER SHIFT
Column 2	R5	A I Q Y 1 9 CLEAR
Column 3	R3	8 J R Z 2 : BREAK
Column 4	R2	C K S 3 ; UPARROW
Column 5	R7	O L T 4 , DOWNARROW
Column 6	R1	E M U 5 - LEFTARROW
Column 7	R4	F N V 6 . RIGHTARROW
Column 8	R6	G O W 7 / SPACE

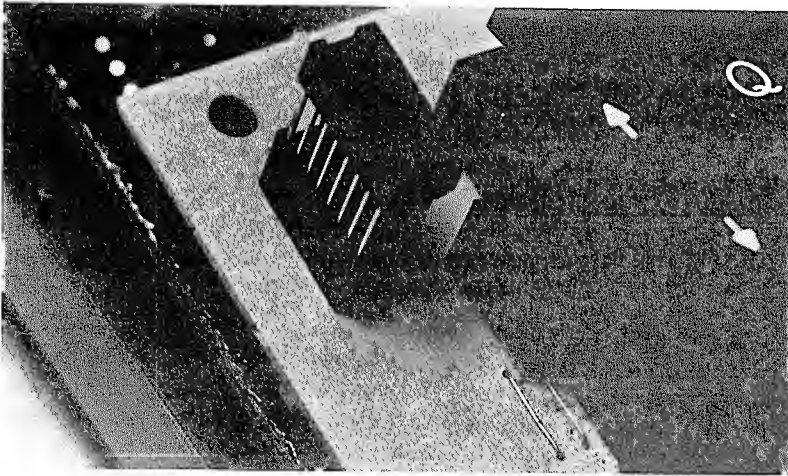
Table 4-4. Keyboard column assignments.

Match column one, then, with socket pin 1; column two with socket pin 2; column three with pin 3; etc. Solder a separate wire to each of the resistors, and wire-wrap or solder the other end to their respective socket pins. Make sure you solder to the end of the resistors which are connected to the keyswitches, not the other ends, which are all connected together.

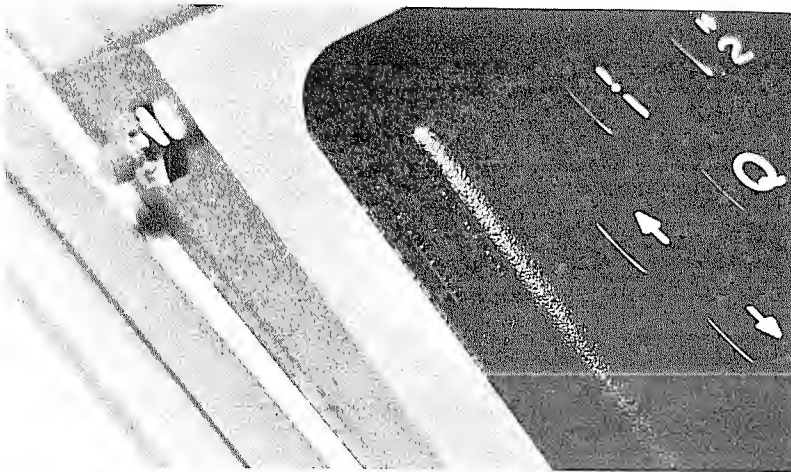
The keyboard matrix rows are found at the input pins of the on-board ICs, but because of the many versions of the TRS-80 keyboard which have been issued, this sequence is inconsistent. The technical manual identifies the rows as shown in Table 4-5, but it's better to check for yourself. Look for the traces that connect



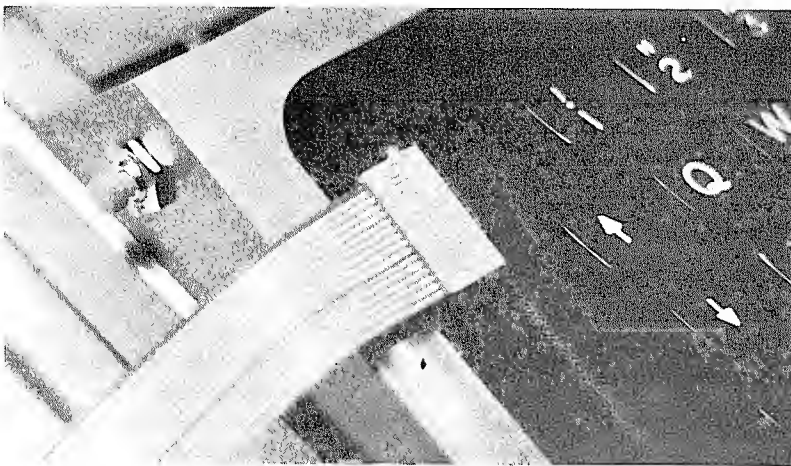
The extra keyboard socket is soldered in place, with the gutted solder-tail socket used as a grommet.



The extra keyboard socket is soldered in place, with the gutted solder-tail socket used as a grommet.



Socket rises perfectly to the height of the outer shell. Note reset button extension at left.



Completed extension cable plugs discreetly into the socket. Any type of keyboard, from a small numeric pad to a full 64-key musical keyboard, can be used.

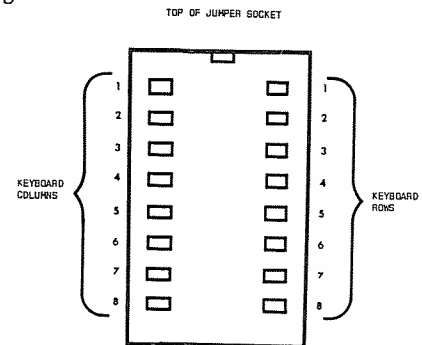
together @, A, B, C, D, E, F and G. These are all in row one. Solder a wire to some point in this row, and run it to pin 16 of the new keyboard socket.

Row 1	Z1 Pin 8	@ A B C D E F G
Row 2	Z1 Pin 2	H I J K L M N O
Row 3	Z1 Pin 10	P Q R S T U V W
Row 4	Z2 Pin 2	X Y Z
Row 5	Z1 Pin 6	0 1 2 3 4 5 6 7
Row 6	Z1 Pin 4	8 9 : ; , - . /
Row 7	Z1 Pin 12	ENTER CLEAR BREAK UPARROW DOWNARROW LEFTARROW RIGHTARROW SPACE
Row 8	Z2 Pin 4	SHIFT

Table 4-5. Keyboard row assignments.

Locate row two by using Table 4-5, and solder a wire from somewhere in this row to pin 15. Likewise, identify rows three through eight, and solder them to pins 14, 13, 12, 11, 10 and 9, in order. When viewed from the top, the pin arrangement is as shown in Figure 4-4.

Figure 4-4.



Once the wiring is complete, clip the pins on the added socket very short, turn the board over, and put everything back in the case. Power up and check the operation of the computer.

Now clip a small length of bare wire about an inch long, and bend it in the shape of a 'U'. At the newly installed socket, jumper each row across to each column, one at a time. You should produce all the non-shifted keyboard characters on the screen, including the previously inaccessible four arrows and the cursor character.

Clip an additional jumper, and cross row eight with column one. This simulates the pressing of the SHIFT key. Repeat the column-row jumpering, and note that all the shifted characters now appear. Any unusual behavior, such as repeated letters or groups of unrelated letters produced from a single jumpering, indicates a wire may be shorted, attached to the wrong column or row, or left out completely.

Finally, as with all modifications, make the cosmetics pretty. Snap the black plastic cover off again, and in it cut a rectangular hole the size

of the 16-pin socket, using a hot, sharp X-acto knife or razor blade. Work slowly, filing or smoothing, and rub the finished hole with a marble. This will result in a professional-looking addition.

The second wire-wrap socket now piggybacks into the first one, and the black cover snaps back on. The socket should fit perfectly, rising about 1/16 inch above the surface of the cover. The 16-wire cable plugs into it a comfortable distance from the typing area, well above and to the left of the up-arrow key.

For each keyboard you wish to add, work out the row-column matrix using the table. A jumper cable may be an integral part of each keyboard, or an IC socket/lug arrangement similar to the main unit can be included with each added keyboard. You can even chain keyboard to keyboard by including two sockets on each one — just be sure all the sockets and plugs are identically wired!

Working by the Woodstove — II

Once you've got a new keyboard in your lap, you'll probably want a nearby screen to glance at. There are two ways to do this: by using a video monitor or by using an ordinary television.

There are advantages and disadvantages to both methods.

A video monitor is the ideal tool because the image is crisp and clear, and your TRS-80 provides a 'composite video' (video with both image and synchronization signals) output. But a video monitor also costs somewhat more than a new black and white television, and means an added expense in any case.

A television on the other hand has limited 'bandwidth'; that is, it was made for fluctuating images, and not for the precise on-off quality of white letters on a dark background. If you've ever noticed that it's sometimes hard to read telephone numbers, addresses, or credits for television programs, you've got an idea of how hard it can be for some sets to reproduce crisp computer lettering.

Furthermore, most televisions accept only radio frequency (RF) input, meaning your TRS-80 output has to be converted to RF before your television can make sense of it. The last complication is that such a close and strong RF signal can overload your television's automatic gain control (AGC) resulting in an unstable, twisting, rolling, or badly contrasted picture.

But chances are you already have a television, and chances are even better that the television is

Making it Look Manufactured

One of the worst curses of a customized anything is how it tends to look — homemade. Now I have absolutely nothing against something looking homemade, but whenever I do that, somehow it also *acts* homemade — that is, just a bit too eccentric to be reliable as a computer!

Instead, attend to the cosmetic aspects of the TRS-80. Since the plastic case is very pliable and 'works' easily, these touches are easy. The silver coloring is a flake paint, and tends to wear off, particularly below the shift keys where the typing hands rest. The black plastic cover is very soft and can be scratched; its pebbled surface makes such scratches stand out.

If the silver flake paint wears off, it can be resprayed with the kind of paint used on model cars. Work the spray can valve for a while until it is spraying evenly, and then spray the silver cover from a slight distance. If the color match is not perfect, the bottom can be sprayed as

well. Use two or three extremely light coats for a good effect.

The black plastic cover can even have deep, obvious scratches repaired by rubbing it with a glass marble. The scratch will smooth over, making it different from the area around it. Next, pick at the smooth area very lightly with a needle, making pebble-size marks similar to the rest of the case. Rub with the marble again until the scratched area looks exactly like the surrounding area. It really works.

When cutting holes for switches, buttons, sockets, jacks, keypads, etc., always cut the holes slightly smaller and use rattail (for round holes) or triangular (for rectangular holes) file to expand it to the correct size. This way, no unsightly cut marks will extend away from the area of the modification. Bevel cut rectangular surfaces with a flat file, smooth them with a letter opener or librarian's 'bone', and touch up the corners. The result will be almost precisely like the manufacturer's molded cutouts.

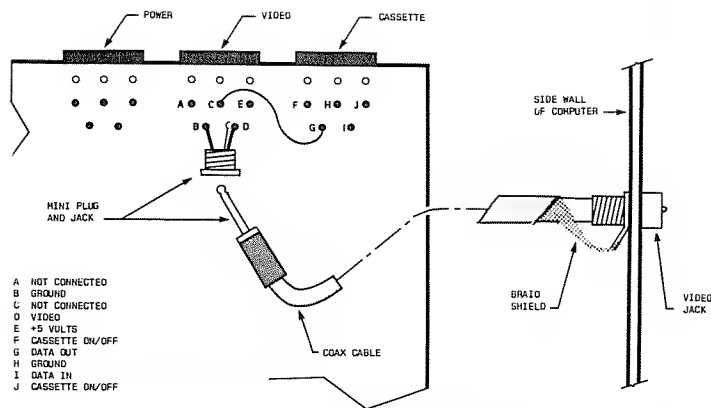


Figure 4-5. Extra video jack.

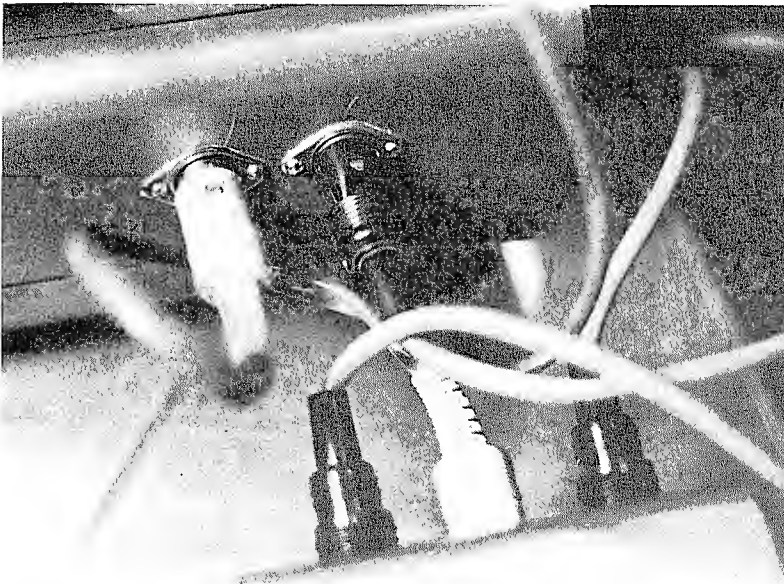


Photo 4-7. RF modulator hookup.

Additional jacks added in order to feed both the video monitor and an RF modulator installed in the expansion box.

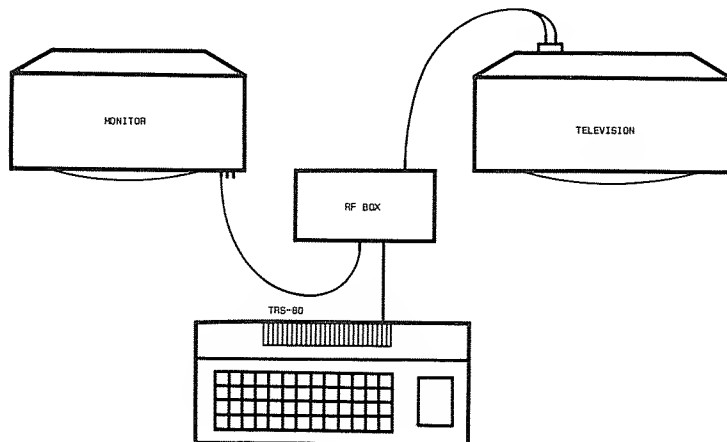


Figure 4-6. RF modulator hookup.

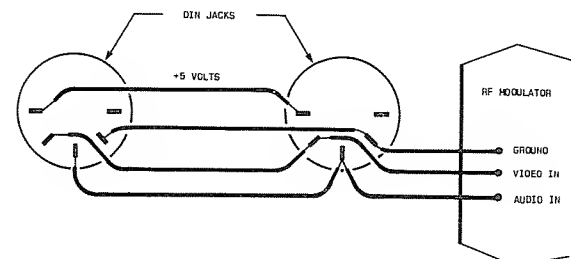


Photo 4-6. Extra video jack.

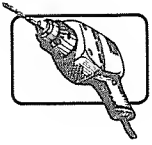
just where you were thinking of using your computer.

Before going on with these modifications, there is one important note. Working with any unknown television or monitor can crash your system if it is electrically noisy. If you plan to send a video signal directly to a monitor, or especially to a television modified for video input, make sure the set isn't 'hot' – no AC line voltage should be floating on the case. This can damage your computer . . . or you! If you are considering a monitor or direct video input, and you are not familiar with your video sets, then take the television or monitor to a service person who can check them out. With appliance devices such as ordinary televisions, this is doubly important. This note does *not* apply if you will use RF input to a television.

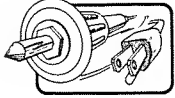
Whichever method of added video you decide upon, though, there is a solution. The first and easiest is to add an extra video output jack to your computer. There's plenty of video signal to be had, and it can be shared among several sets. Figure 4-5 shows how to wire that extra jack, and Photo 4-6 shows how mine is installed. The



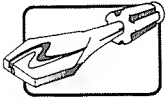
Hexadecimal Keypad



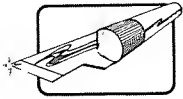
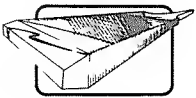
connector shown in the photo is a miniature Amphenol connector used for microphone cable, although any kind of microphone, video, CB, phono, or other shielded coaxial cable and connector can be used.



A second approach is to send the computer's video signal to an RF generator, and feed that to your television. An RF modulator is available in kit form from Radio Shack (part number 277-122, with TRS-80 installation instructions) and also from other suppliers. A few plugs and jacks are needed to complete the job.



Assemble the kit or purchase a surplus modulator, and hook it to an ordinary television. Either run the new video output to the modulator's input (use the directions with your modulator), or install a separate DIN plug-and-jack pair as in Photo 4-7.).



Actually modifying an ordinary television for video input is tricky, and I won't cover the topic here. If you are interested in doing this — and picture quality will be very much improved — refer to Don Lancaster's *TV Typewriter Cookbook*.

The results of your video modifications will be dependent on how clean your soldering is, the layout of your wiring, etc. If you add the RF modulator, you may notice something similar to herringbone on your TRS-80 monitor if the

contrast is turned all the way up. Move the modulator away, or put it in a shielded (all metal) box, and the herringbone will disappear.

If you use a long cable directly from the video output, there will be a huskiness to the characters on your TRS-80 monitor. This is actually a kind of 'smearing' introduced by the capacitance of a long cable. This is not bothersome to me; in fact, it actually seems to improve the clarity and boldness of the screen characters.

Finally, if your RF-input television addition results in an unpleasant display on the TV, try to adjust the automatic gain control (AGC) on the back of the set. Tune it in carefully as well. There will probably be less clarity in the 64-character mode than you are used to with your monitor, unless you have a very good set.

Hexadecimal Keypad

Entering machine language programs using T-Bug or another monitor is tedious enough without having to search all over the QWERTY keyboard for hexadecimal numbers. Instead, a keyboard can be added right onto the TRS-80. If you have a numeric keypad included with your Level II unit, you might want to remove it to add this one.

The addition of a hexadecimal keypad is mostly carpentry, since the connections are made in parallel to the main board, exactly like those connected to the socket addition described earlier. An unencoded hexadecimal keypad to do the job is available from Jameco Electronics (see Appendix 1.), and one of its keys can be set aside for an *Electric Pencil* or other control key, adding significant programming power to your custom TRS-80.

For this modification, you will need two 10-inch strips of 1/2-inch by 1/2-inch plastic rod (plexiglas or lucite are best, but wood strips will work as well), five-minute epoxy, wire, and the notorious hot razor blade and marble for the cosmetics.

Undo the cabinet as usual, and take the entire electronics out of the case. Later TRS-80's have an on-board, two-chip Level II ROM set, but if your Level II ROMs are the type on a separate board fastened to the end of an interconnect cable, then they will have to be moved. They are fastened to the bottom right of the circuit card with double-face tape; slit the center of the tape with a razor blade. Do not pull the ROM board off by force, as the pressure might crack either circuit board.

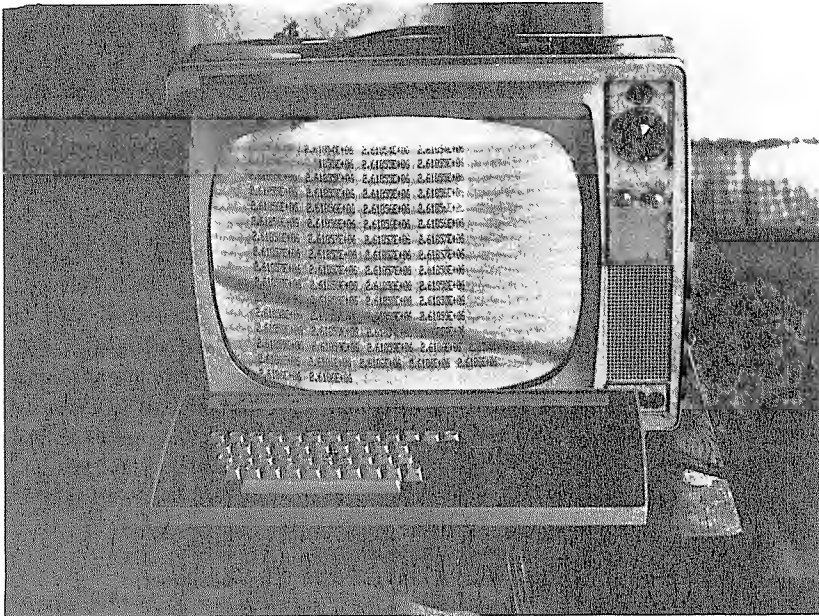


Photo 4-8. An Extra system near the woodstove.

The kitchen installation. An old RCA television and a surplus keyboard make working with the TRS-80 more convenient. Computer remains in one place, but extra monitor and keyboard can be moved closed to the woodstove on cold days. The woodstove offers no electronic interference.

The interconnect cable to the ROMs is long enough so they can be remounted inside one of the case 'feet', or above the hexadecimal keypad. Pick up a small piece of double-face tape to refasten them, or roll masking tape into cylinders (remember hanging pictures in fourth grade?).

The Jameco Electronics keypad base is identical to the TRS-80's in height and depth, so

the two plastic strips can be used to create a 'trailer hitch' arrangement with the smaller keyboard. Support both boards firmly so that they are parallel and the hex pad meets the TRS-80 printed circuit base. Cement the plastic strips in place with the quick-setting epoxy, and make sure the vertical alignment of both keyboards is identical. If you have a later style

Photo 4-10. *Black plastic cover used as template.*

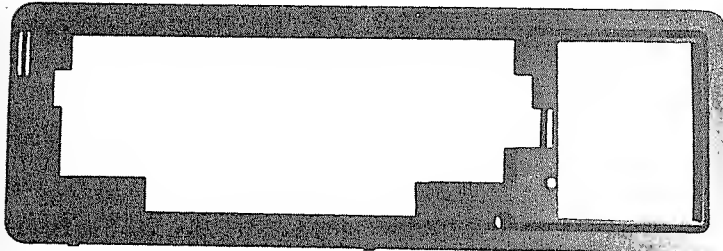
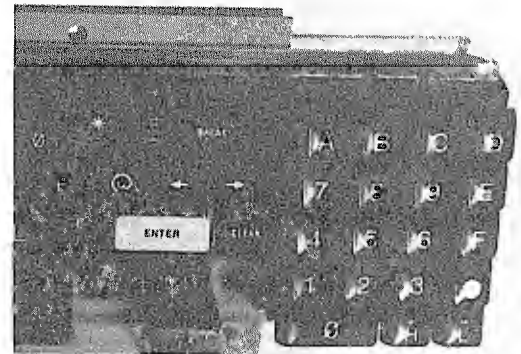


Photo 4-10. Black plastic cover used as template.

The black plastic keyboard cover can be snapped out and cut to fit. A paper template is used first to verify the position of the hex keypad.



Keyboard is attached with runners to the main keyboard, and glued in place with epoxy.

Cleaning the Keyboard

If your TRS has the old-style keyboard that was badly afflicted with keybounce, there are many ways to take care of it other than perennially loading a KBFIX routine.

The first rule is to *make sure you have the old style keyboard!* The newer keyboards have a sculptured curve to their arrangement when viewed from the side. These new boards have a contact arrangement which can be destroyed by trying to remove the keycaps. But these keyboards don't have a keybounce problem anyway.

Bend a paper clip into an ingenious keycap lifting tool, like this:



Figure 4-7. Keycap lifting tool.

Slip this clip underneath a keycap, and lift up. The plastic cap will come off, revealing a hole in which two metal plates are protruding upward.

The best way is to test the keys. Press each key quickly, gently, slowly, or sharply, until you decide whether it is a 'bouncy' key. If it still bounces after the cleaning, then take a hatpin (are there still hatpins?) or heavy sewing machine needle, and push the tines in line. This is a very delicate job; use caution and a magnifying lens. If the plates are not parallel to each other, or if they are vertically misaligned, use the pin to shift their positions.

The villains are dirt and bent tines on the plates. Brush out the dust, dirt, or hair (or much better, blow it out, using photographers' compressed air, such as 'Dust-Off'). You will be amazed at the cloud of grit that rises from the keyboard. Next, examine the tines of each key very carefully. They should be perfectly in line, so that when a key is pressed, all come into contact with their opposite (un-tined) plate.

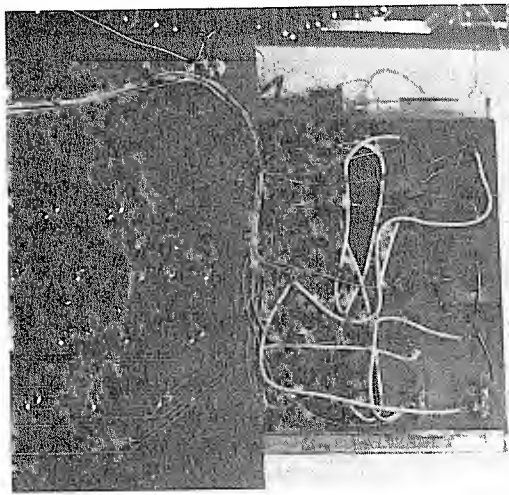
Check all the keys for bounce again, and work until it is completely cured. Some folks recommend a spray of contact cleaner at this point; I recommend against it. The cleaner tends to stay wet for a while, and dust and grit can get back into it very quickly, collecting into a dusty mudpile. Instead, give all the keys a last brushing or spraying with compressed air, and fit the keycaps back on.

Keybounce should be gone for quite a while. Monthly cleaning will keep the keyboard in shape.

Hexadecimal Keypad

keyboard (with the curved keyboard array), you will have to adjust the carpentry slightly. You may also experience a bit of keybounce on the new keypad unless you keep the keys clean.

When the glue has set, use the black plastic cover as a template for drawing your current key positions and, with the aid of a straightedge, draw extension lines horizontally across that drawing. These are the upper and lower limits of the new keypad opening. Align the template with the complete alpha/hex keyboard assembly, and mark the vertical positioning of the hex keys, allowing about 1/32 inch additional on both sides for key-travel room. This will bring you within



Back of keyboard is wired by soldering directly to the key pins. Wires are then run to resistors and integrated circuits found on the main keyboard.

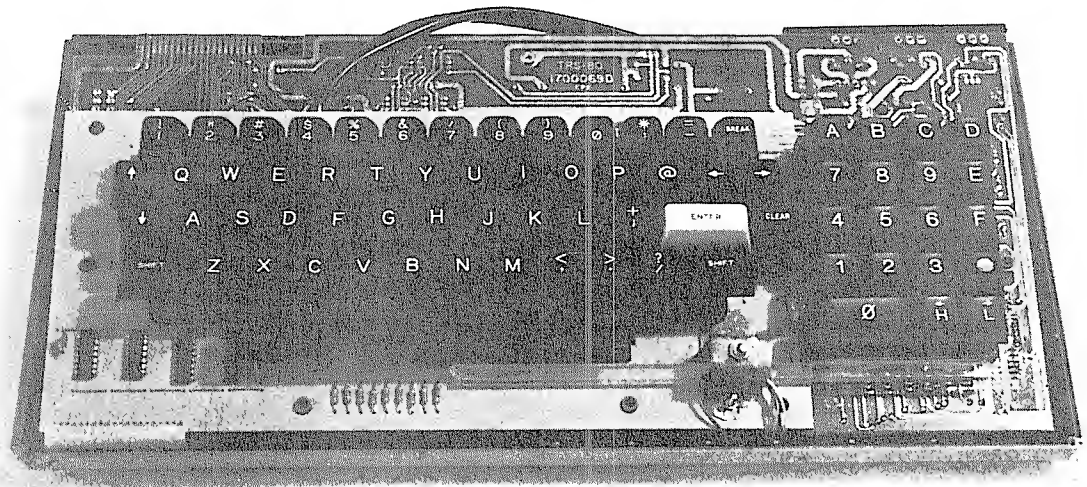
1/4 inch of the power LED.

Using the template with the black plastic cover, carefully cut an opening in the cover with the hot razor blade. This is the most time-consuming task, and should be accurate enough for the keys to travel easily (fit it atop the keyboard before re-installing the modified computer), and should look factory-finished.

It's about time to interconnect the wires from the hex pad to the main keyboard, but before that, you'll probably want to rearrange the keycaps on the hex pad. Using the lifting tool, pull off the keycaps and put them in a convenient order for hex entry; I used the accounting arrangement, bottom to top. This is the pattern used for the wiring arrangement shown in Diagrams 4-1 and 4-2.

Rest the keyboard on its face, and separate it gently from the main circuit card. Set the keyboards in an accessible position, and solder fine wires (wire-wrap type is easiest to use) to the hex pad connections shown in Diagram 4-1. Route the individual wires from the hex keys to the points on the circuit card shown in Table 4-5. As before, check to make sure this IC arrangement matches your board.

Next, solder wires to the hex pad contacts as shown in Figure 4-1, and route these wires from the hex keys to the circuit card's resistors noted in Table 4-4. Double check these too against your version of the keyboard. Once both sets of wires have been run, gather them in neat hanks and fasten them along their routes with wire ties (plastic bag ties will also work well).



Completed keyboard seats easily in the case bottom. A piece of insulating plastic should be inserted under the added keyboard to prevent pushing the keyboard pins into the main circuit board and causing a short.

Hexadecimal Keypad

Table 4-3

@	A	8	C	O	E	F	G
H	I	J	K	L	M	N	O
P	Q	R	S	T	U	V	W
X	Y	ZUnassigned....				
0	1	2	3	4	5	6	7
8	9	:	;	-	.	/	
ENT	CLR	8RK	UPR	ONR	LFR	RTR	SPC
SHIFTUnassigned.....						

Table 4-4

Column 1	R8	@ H P X 0 8 ENTER SHIFT
Column 2	R5	A I Q Y 1 8 CLEAR
Column 3	R3	8 J R Z 2 : 8BREAK
Column 4	R2	C K S 3 ; UPARROW
Column 5	R7	O L T 4 , DOWNARROW
Column 6	R1	E M U 5 - LEFTARROW
Column 7	R4	F N V 6 . RIGHTARROW
Column 8	R6	G O W 7 / SPACE

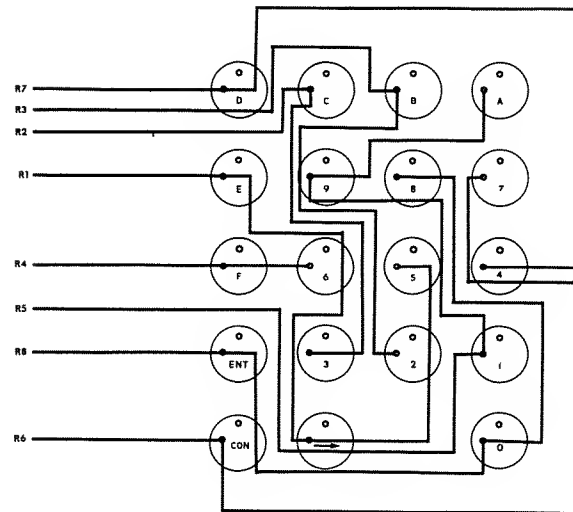


DIAGRAM 4-1

Table 4-5

Row 1	Z1 Pin 8	@ A 8 C O E F G
Row 2	Z1 Pin 2	H I J K L M N O
Row 3	Z1 Pin 10	P Q R S T U V W
Row 4	Z2 Pin 2	X Y Z
Row 5	Z1 Pin 6	0 1 2 3 4 5 6 7
Row 6	Z1 Pin 4	8 9 : ; , - . /
Row 7	Z1 Pin 12	ENTER CLEAR 8BREAK UPARROW DOWNARROW LEFTARROW RIGHTARROW SPACE
Row 8	Z2 Pin 4	SHIFT

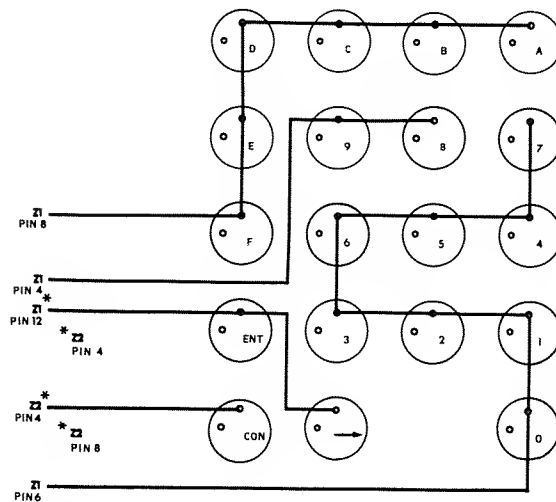
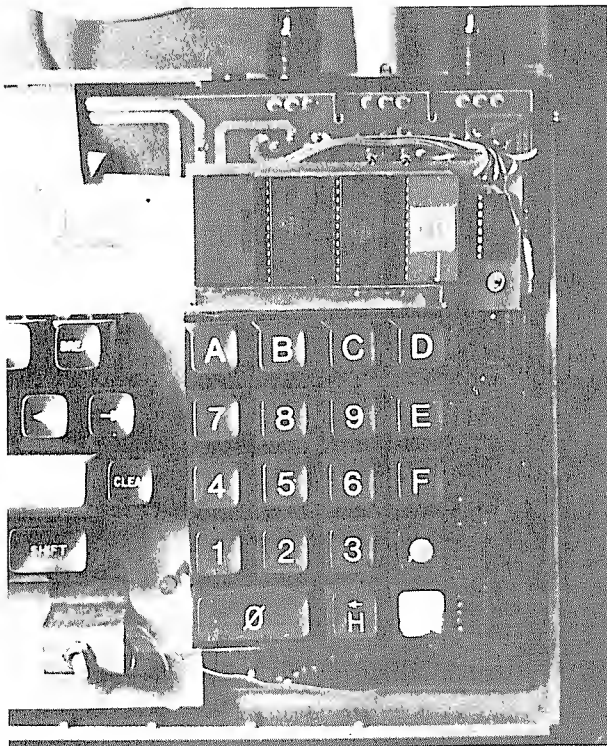


DIAGRAM 4-2

Hexadecimal Keypad



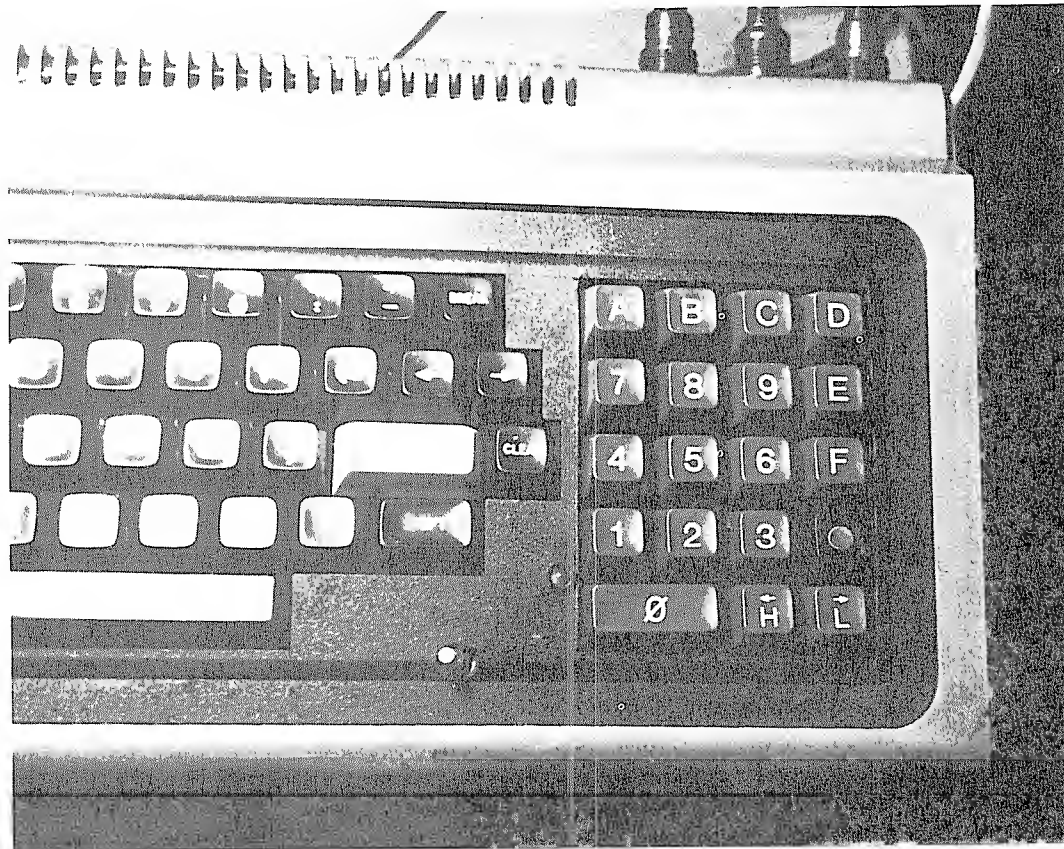
Level II ROM board can be fit immediately above the hex pad. Insulating double-face tape can be used to hold them in place.

Reassemble the keyboard in its case, remembering to orient the Level II ROM board safely in its new position. Reinsert all cables, and restore power. The keys on the main keyboard should respond normally; check them all. Now check the keys on the hex pad. All but the bottom right one should have an effect. To test its operation, enter this program:

```
10 CLS
20 PRINT PEEK (14464);
30 GOTO 20
```

The value you read should be zero unless either the shift key on the main board or the bottom right hand key on the hex pad is depressed. The shift key will return a value of 1; the new key will display a value of 16. Pressed together, they will read 17.

All keys should now be working properly. As with the socket addition described above, problems will occur in the form of incorrect letters, groups of letters on a single keypress, or dead keys. If any of these symptoms appear, recheck for shorted or unconnected wiring, or a difference in your model TRS-80 keyboard. Refer to Tables 4-3 through 4-5 if you suspect the latter.



The final modification with black plastic cover fit back in place looks manufactured (almost).

Reversing the Video

The video display of most computers including the TRS-80 suffers from a tremendous flaw – contrary to what we have known since we first learned to read, the letters are presented to us in glowing blue-white on a black background. Serious use of a computer as a day-to-day appliance, as a true adjunct to our daily lives, is limited by its formidably unappealing, tiresome, and illegible display.

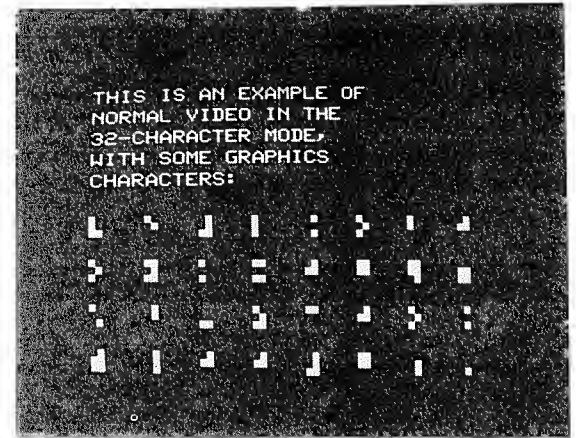
The reverse video modification is surprisingly easy, so the question arises: why is it not a standard feature of small computers? The answer is found partly in a tradition of computer monitors which have always been lighted characters on a dark background. The remainder of the answer is found in the video display itself, which is more often than not incapable of presenting a legible character in the black-on-white mode.

This problem arises with the TRS-80 as well; the video monitor has weaknesses which are emphasized by reversing the characters. But overall, and with a small change to the monitor itself, the display can be made quite legible and easy on the eyes.

For this modification, you will need three integrated circuits: 74LS02, 74LS74, 74LS368.

A 1.5K-ohm resistor will also be used, and wire-wrap wire for the interconnections. There are two ways of making this modification: on a separate board, or piggybacked atop chips already present inside the TRS-80. Since the latter approach would involve at least 15 separate wires, both this change and the high-speed modification presented below will use the piggyback system.

Open the computer's case, and remove the



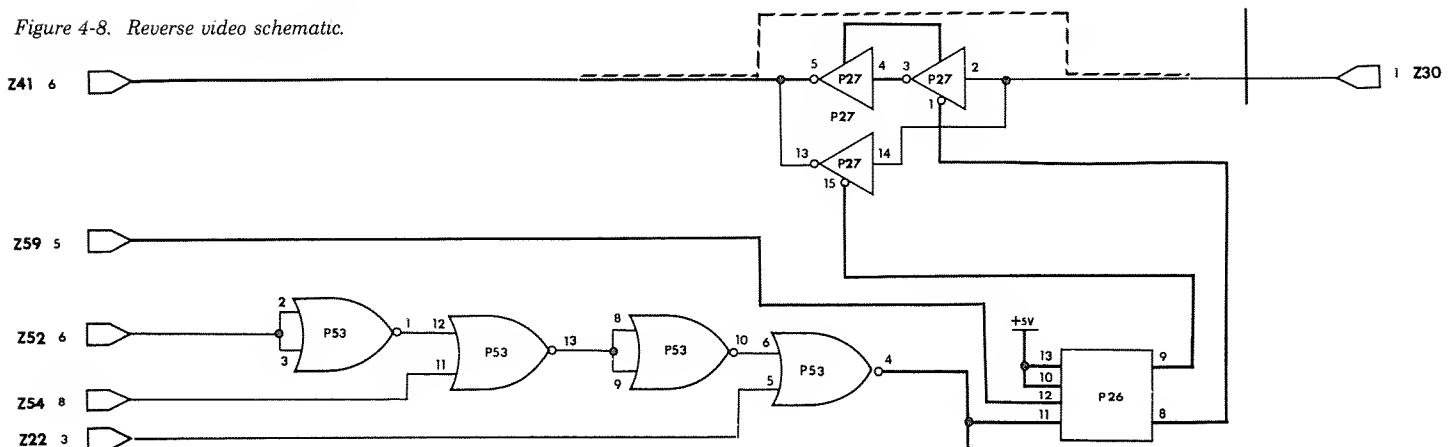
Normal screen has unnatural bright characters against a dark background, which is wearying on the eyes.

electronics so the integrated circuit side of the main card is up. Locate the following circuits (the numbers are silkscreened on the board): Z22, Z26, Z27, Z30, Z52, Z53, Z54, and Z59.

The output of Z54 (pin 8) is a decoded signal representing ports 254 and 255. Combined directly with the output of Z52 (pin 6), port 255 is selected for use by the cassette and video circuits. By inverting the output of Z52 and combining it with the output of Z54 and the computer's OUT signal (Z22, pin 3), port 254 can be selected. Figure 4-8 is the schematic for this complete decoding process.

The decoder uses the 74LS02. Prepare the integrated circuit by bending all the leads except pins 7 and 14 so that they are parallel with the IC's body. Locate Z53 on the TRS-80. Seat the 74LS02 directly atop Z53, with both notches or dots facing in the same direction. Solder power pins 7 and 14 of the piggybacked IC to corresponding pins 7 and 14 of the one below. I will refer to this piggybacked IC as ZPORT.

Figure 4-8. Reverse video schematic.



Find Z52. Run a wire from pin 6 of Z52 to both pins 2 and 3 of ZPORT, and solder it. Next locate Z54. Solder a wire from its pin 8 to pin 11 of ZPORT. Finally run a wire between pins 1 and 12 of ZPORT.

As mentioned, Z22 contains the needed OUT signal. Run a wire from pin 3 of this circuit to pin 5 of ZPORT. Solder together pins 8, 9 and 13 of ZPORT, and run a wire between pins 6 and 10 of ZPORT. Pin 4 remains unused, and it contains the complete decoded signal of port 254 as shown in Figure (?). The BASIC command OUT 254,N will activate this signal.

The next IC to be prepared is the 74LS74. Again, bend all leads parallel to the body except 7 and 14, seat this upon Z26, and solder the power pins (7 and 14) in place. This piggybacked circuit I will call ZFLOP, as it will determine which state (normal or reverse video) is flip-flopped into place when OUT 254,N is commanded.

Run a wire from the decoded signal at pin 4 of ZPORT to pin 11 of ZFLOP. Z59 has a convenient data line (bit 1) at its pin 5; run a wire from there to pin 12 of ZFLOP. Now run short wires connecting together pins 10, 13, and 14 of ZFLOP. With these connections made, OUT 254,0 will flip the circuit, and OUT 254,2 will flop it. (Is the suspense building?)

The final IC is now prepared. Bend the leads of the 74LS368 parallel to its body, except for pins 8 and 16. Seat this on Z27 and solder power pins 8 and 16 to it. For convenience, this piggybacked circuit will be called ZMODE.

Find Z30 and Z41. Pin 1 of Z30 is connected via a circuit board trace on the *underside* of the board to pins 6 and 7 of Z41. Z30 provides the characters to be output to Z41, which is part of a circuit that mixes in the synchronization information to produce 'composite video'. Cut this trace near Z30.

Why cut this trace? The TRS-80 produces characters on the screen by turning on 'dots' as the electron beam sweeps across the tube. Each dot is part of a continuous stream of pulses which might be called 'dots' and 'undots' – ones and zeros. To reverse the video, then, all you need to do is to turn the 'dots' into 'undots', and turn the 'undots' into 'dots'. We insert an electronic fork in the road at the output of Z30. When directed toward one side of the fork, the characters are made up of dots; when directed down the other fork, the characters become undots, and the background becomes dots.

Run a wire between pin 1 of Z30 and pins 2 and 14 of ZMODE. These are two inputs of an inverting buffer, the 'fork in the road'. If we invert the signal once, the video reverses . . . invert it twice, and the signal returns to normal. Connect pins 3 and 4 of ZMODE together to perform the double inversion. Pin 5 is the normal video output, and pin 13 becomes the reversed output. Connect both these outputs (pins 5 and 13) together.

Find the opposite side of the broken trace from Z30, and follow it to a hole that is plated through the board; it is at the end of a row outlined by Z29 and Z30. Be careful to select the correct hole. Run a wire from this hole to pins 5 and 13 of ZMODE. This connection feeds both normal and reverse video back into Z41 and through to the video output jack.

I have chosen ZMODE because it is a three-state circuit. That is, its outputs can be made electronically invisible. Otherwise, both normal *and* reverse video would be output at the same time. To choose between them, ZFLOP is used to *enable* one or the other of those outputs. Run a wire from pin 8 of ZFLOP to pin 15 of ZMODE. Run a wire from pin 9 of ZFLOP to pin 1 of ZMODE.

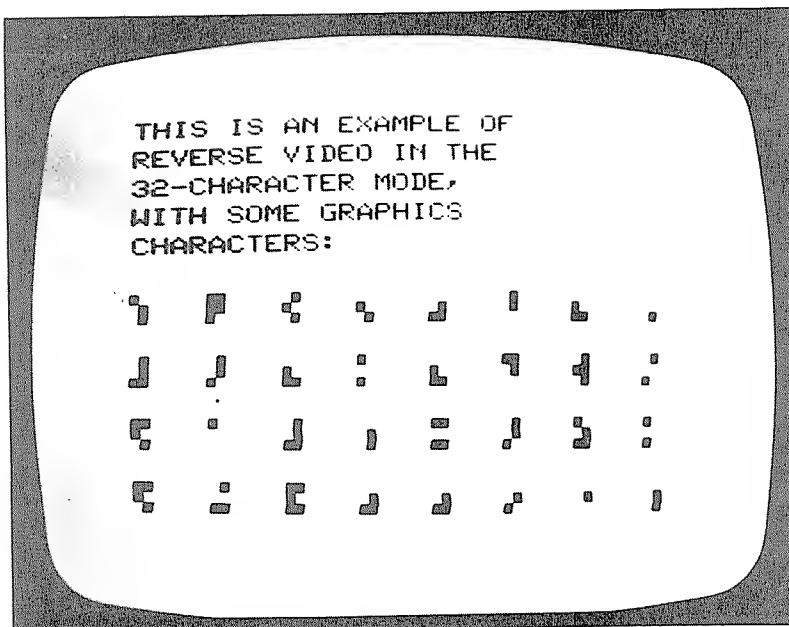


Photo 4-11. Reverse video screen example.

Illuminated background with dark characters is clear and, together with a green screen of some type, much more gentle to look at over long periods.

By commanding OUT 254,2 or OUT 254,0, data line 1 selects which output of ZFLOP will be enabled, and which video mode will be visible on the screen

Check your wiring and restore the computer to its case. Power up, and command OUT 254,2. The screen will reverse. The effect, alas, will not be as dramatic as you might expect because the video monitor is not a great piece of work.

Power down the system, unplug the monitor, and open it up. You might find that a hex-nut driver is necessary to open the monitor instead of a Phillips type; later monitors used 1/4-inch hex nuts.

When it is open, find the plug-in circuit card closest to the monitor wall (some have only one), and locate the resistor marked R14. The present value should be 3.3K ohms (orange, orange, red, silver or gold). We want to give the video signal a bit more 'oomph', so piggyback the 1.5K ohm resistor atop R14 and solder it in place. Restore the cabinet. The reverse video modification is complete.

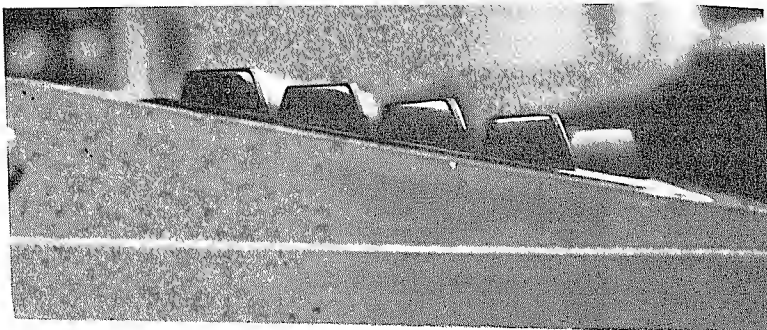


Photo 4-9. Two photos (a) & (b) of two keyboards.

Clumsy? Me Too. Do This First.

During the course of these modifications, there will be connections made to integrated circuits which are carrying piggies atop their back. In case your soldering iron is a bit obese, you might want to solder wires to these circuits in advance.

All these pins will have wires running to them:

- Z22, pin 3
- +Z25, pins 11, 12 and 13
- +Z27, pins 4 and 5
- Z30, pins 1 and 13
- Z41, pin 6
- Z52, pin 6
- Z54, pin 8
- Z59, pin 5
- Z60, pins 4 and 5
- Z63, pin 12

Only those marked with a plus sign (+) actually carry piggies, but since everything is close together, you might want to make all the connections in advance anyway; so there they are.

Lower Case with Upper

Many modifications have been proposed to obtain lower case characters already present in the TRS-80 character generator. Some are incompatible with each other, although they do provide access to a group of special control characters also burned into the character generator.

The modification provided here is simple, compatible with both the Radio Shack and Electric Pencil modifications, and should give no grief throughout its life. For this modification you will need a single integrated circuit, and 2102 AN-4L memory chip. These chips are available from several suppliers; in a pinch, the 21L02 sold by Radio Shack will do the job.

The 2102 will be piggybacked – except for pins 11 and 12 – atop Z45, also a video memory chip (it will not necessarily be marked a 2102, since Radio Shack ordered house numbered parts for a while, but it is a 2102). Bend pins 11 and 12 of the piggyback 2102 parallel with its body, and fit the integrated circuit on Z45, making sure it is positioned in the correct direction. Solder extremely carefully, pin for pin, down one side (pins 1 to 8), and up the other (pins 9 to 10, 13 to 16).

Now locate Z25. It is an integrated circuit containing four OR-gates, one of which is not used. With solder-wick, suck up the extraneous solder that is present on pins 11, 12, and 13, and with a sharp X-acto knife or razor blade, cut pins 12 and 13 free from each other and from the ground lead going to pin 7.

Now locate the trace on the circuit board running from Z60 pin 4 to Z30 pin 13. Use an ohmmeter if necessary to make sure you have the correct trace. Cut it through. Double check that you have not cut the trace that runs from Z60 pin 4 to Z27 pin 13.

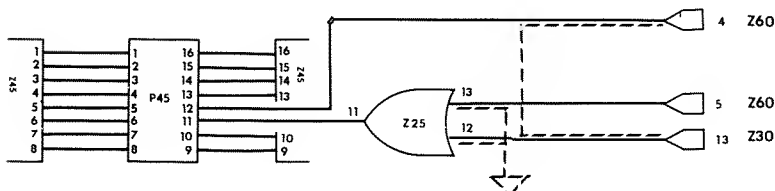


Figure 4-9. Upper/lowercase schematic.

Finally, run four wires: from Z60 pin 4 to piggy-Z45 (call it ZMEM from now on); from Z60 pin 5 to Z25 pin 13; from Z30 pin 13 to Z25 pin 12; and from Z25 pin 11 to ZMEM pin 11.

The lowercase modification is complete. A memory chip to represent the missing bit 6 has been added, and the necessary triggering information (bit 7 and bit 5) has been gated through Z25. Why is it necessary to gate bits 7 and 5 to trigger ZMEM? Because the information sent by Radio Shack's print routine does not include bit 6; in an unmodified machine the bit is generated by the presence of bit 7 and bit 5! Check the *Technical Reference Handbook* for details. In this case, the false bit 6 is generated where necessary by Z30, and ORED with a true bit 6. In other words, either case will result in bit 6 being embedded in the new bit 6 video RAM. ZMEM.

Note: Several TRS-80's in which I have installed this modification contained a Z25 with the spare gate dead. Perhaps it was a Radio Shack economy move. If so, piggyback a 74LS32 on Z25, soldering pins 7 and 14 to Z25, and continue with the remaining lowercase mod instructions.

The lowercase modification can be tested as follows; a complete lowercase driver is presented in Chapter 3:

```
10 CLS
20 FOR X = 15360 TO 15360+255
30 POKE X,Y
40 Y=Y+1
50 NEXT
60 GOTO 60
```

One by One

Reversing individual characters is similar to the process of reversing the entire video screen, except that the reversal is carried out only for a short period of time. That period is determined, of course, by the letters being displayed. (At this point, I should note that all four internal hardware modifications in this chapter are partly interrelated. Both reverse video and high-speed modifications use the same output port decoding; both reverse video and reverse characters use the same flip-flop; both reverse characters and upper/lower case use the same video memory decoding hardware).

Before installing the individual-character reverse video, the upper/lower case modification must be installed. This will provide the essential bits 6 and 7 (see that section for an

explanation of the 'phantom' bit 6). For the individual character reverse, three additional integrated circuits will be needed: one 74LS86 Exclusive-OR gate, one 74LS10 triple-input NAND gate, and one 74C04 hex inverter. The last is very important, because it is used for aligning the reversal pattern with the letter to be reversed. You will also need a small variable resistor (trimmer potentiometer, or 'trimpot'), approximately 50,000 to 100,000 ohms; two capacitors, one 330 picofarads (pF) and one 0.033 microfarads (mF).

As before, bend all pins of the 74LS86 except pins 7 and 14 parallel with the body of the integrated circuit. Affix the 74LS86 atop Z24, and solder power pins 7 and 14 to it. This gate will multiplex the combined bit 6/7 signal with the output of ZFLOP (from the previous complete reverse video modification). Call this gate ZMUXX.

Remove the wires attached to the outputs (pins 8 and 9) of ZFLOP, and run them,

respectively, to pins 2 and 5 of ZMUXX. Now run a pair of wires from pins 6 and 3 of ZMUXX to pins 15 and 1 of ZMODE (which were just disconnected from ZFLOP). In other words, ZMUXX has been inserted between ZFLOP and ZMODE.

Next position the 74LS10 correctly atop Z25 and solder the power pins in place. Call this circuit ZBITS, because it will evaluate which bits of given letters should cause the reversal to take place.

The steps below summarize all the activities to complete this modification; some have been done already, (such as the lowercase modifications) but are included for clarity:

1. Break trace from Z30, pin 13 to Z60, pin 4.
2. Break trace from Z63, pin 12 to Z42, pin 13.
3. Break trace from Z42, pin 12 to Z27, pin 4.
4. Cut loose Z25 pins 11, 12, 13, from pin 7.

Carpentry Considerations

When you first open your TRS case, you'll probably be unfamiliar with what comes out. The photo below is presented to remind you that six screws and five white spacers make up that group:



Photo 4-12. Screws and spacers in keyboard unit.

When you have been soldering in and around the circuit board, chances are that a lot of ugly, crusty brown flux residue will build up. In order to see what you are doing and make sure connections are sound, you should clean this mess. There are flux removal compounds available, and these should be applied with cotton swabs.

My own choice is a gentle but very fast acting substance which can sometimes be found in surplus – Thermo-Fax brand belt cleaner. This treats the boards and their coating without harsh chemical action, but removes the flux within seconds.

Another area of difficulty in doing these modifications is cutting traces. Place the circuit board on a very secure and stable table, cushioned just a little with a towel. Lean firmly but gently on the board, and move an X-acto knife or single-edged razor blade back and forth until the trace gives way. This may take as many as 20 or 30 scrapings.

When the trace looks cut, make sure. Cut deep into the fiberglass base so you can see a cut space, and then wipe the area clean with flux or tuner cleaner so the break is obvious.

If you must resolder a trace, scrape the green masking from both sides of the cut, wipe it clean, and flow solder on both sides, but *don't* bridge the trace with solder. Instead, take a piece of bus wire or stripped wire-wrap wire, form it into an 'L' shape, and solder the base of the L across the trace. Then cut off the excess. Bridging with solder alone is dangerous because it can look fine, but really be attached only by a glob of flux, or be attached so weakly that flexing the board will crack the solder off.

5. Piggyback on Z45 a 2102, soldering pins 1, 2, 3,
- 4, 5, 6, 7, 8, 9, 10, 13, 14, 15 and 16.
6. Piggyback on Z24 a 74LS86.
7. Piggyback on Z25 a 74LS10.
8. Piggyback on Z6 a 74C04.
9. Run a wire from Z60, pin 5 to Z25, pin 13.
10. Run a wire from Z30, pin 13 to Z25, pin 12.
11. Run a wire from Z25, pin 11 to Z45piggy, pin 11.
12. Run a wire from Z45piggy, pin 12 to Z60, pin 4.
13. Run a wire from Z45piggy, pin 12 to Z25piggy, pin 1.
14. Run a wire from Z63, pin 12 to Z25piggy, pin 13.
15. Run a wire from Z25piggy, pin 2 to Z27, pin 5.
16. Run a wire from Z25, pin 12, to Z6piggy, pin 1.
17. Connect together Z6piggy, pins 2 and 3.
18. Attach one terminal of the 100,000-ohm trimpot to Z6piggy, pin 4.
19. Attach the other terminal of the 100,000-ohm trimpot to Z6piggy, pin 5.
20. Attach one end of the 330-pf capacitor to Z6piggy, pin 5.
21. Connect together Z6piggy, pins 6 and 13.
22. Attach the other end of the 330-pf capacitor to Z6piggy, pin 12.
23. Attach one end of the 0.033-mf capacitor to Z6piggy, pin 12.
24. Attach the other end of the 0.033-mf capacitor to Z6piggy, pin 11.
25. Connect together Z6piggy, pins 9 and 10.
26. Run a wire from Z6piggy, pin 8 to Z24piggy, pin 1.
27. Connect together Z24piggy pins 1 and 4.
28. Run a wire from Z63, pin 12 to Z25piggy, pin 3.
29. Run a wire from Z42, pin 12 to Z25piggy, pin 4.
30. Connect together Z25piggy, pins 5 and 14.
31. Run a wire from Z25piggy, pin 6 to Z27, pin 4.
32. Disconnect the wires running from ZFLOP pins 8 and 9 to ZMODE.
33. Run the wire from ZFLOP, pin 8, to Z24piggy, pin 2.

34. Run the wire from ZFLOP, pin 9, to Z24piggy, pin 5.
35. Run a wire from Z24piggy, pin 3 to ZMODE, pin 1.
36. Run a wire from Z24piggy, pin 6 to ZMODE, pin 15.

The modification is complete. Using the individual reverse video is not the easiest process, but works nonetheless. Printing CHR\$(0) through CHR\$(31) results in control codes being acted upon; CHR\$(32) through CHR\$(127) now produce the full range of ASCII letters; CHR\$(128) through CHR\$(191) produce graphics characters; and CHR\$(192) through CHR\$(255) produce the 63 possible TAB positions.

To use the individual character reverse, then, you cannot PRINT the CHR\$ value. Instead, you must POKE the value onto the screen. Granted, this is a pain, but when a program is completed, the prompting can be extraordinarily effective. Besides, it's not that hard. Try this:

```
10 CLS : Y=15360
20 FOR X = 0 TO 255
30 POKE Y,X
40 Y=Y+1
50 NEXT X
60 GOTO 60
```

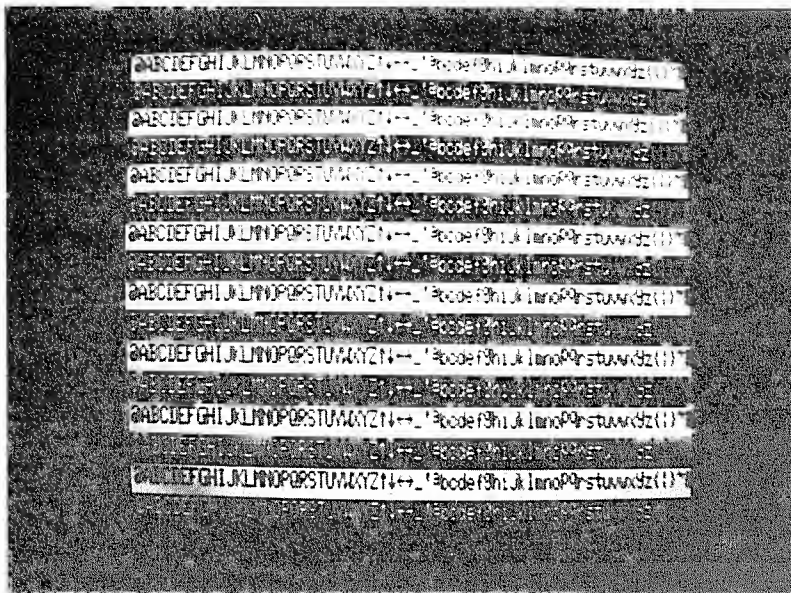
Listing 4-1. Individual character test program.

Chances are that a group of reversed characters appeared, but the reversal band didn't match up with the actual characters. That's the reason for the 74C04 and the 100,000-ohm potentiometer. Run the following test program:

```
10 CLS
20 FOR X = 16040 TO 16060
30 POKE X,255
40 NEXT
50 GOTO 50
```

Listing 4-2. Individual character alignment program.

Now adjust the potentiometer until the crosshatches are centered precisely within the white band (or rather, so the white band is in the correct background position).



Complete font of reversed characters. Because only 64 characters remained in the available set, upper and lower case were chosen as opposed to numbers and symbols.

That is a simple example of the individual character-reverse process, using BASIC. By using ordinary display routines in machine language programs, the individual character reverse is an entirely trivial matter, but really putting it to work in BASIC requires a bit of fancier footwork.

Using it through BASIC can be done without POKEing large numbers of characters in prearranged locations. Instead, the position of the cursor on the screen can be determined, and the characters POKEd in place from there. The subroutine below is useful:

```
10000 X = PEEK(16416) + 256*PEEK(16417) : RETURN
```

This subroutine determines the position of the cursor and assigns it to X. Then characters may be POKEd as follows:

```
40 GOSUB 10000
50 AS = "THESE ARE REVERSE LETTERS"
60 FOR N = 1 TO 25
70 Q = ASC (MID$(AS,N,1))
80 POKE X,Q
90 NEXT
100 GOSUB 20000
```

Listing 4-3. Individual reverse demo program.

The GOSUB 20000 executes the following one-line subroutine:

```
20000 ZH = FIX (X/256) : ZL = X-ZH*256 :
      POKE 16416,ZL : POKE 16417,ZH : RETURN
```

This short routine restores the cursor position after the POKEs have taken place. There are two very important things to note:

1. If the POKE is to take place on the last line of the screen, make absolutely certain that the value of X does not exceed 16383, because this will POKE nasty values into BASIC vectors beginning at 16384. A test for X greater than 16383 can be made so:

```
10005 IF X>16383 THEN X=96383 : GOSUB 20000
      : PRINT : RETURN
```

This will have the effect of restoring the cursor, printing a carriage return, and finishing the text to be printed.

2. The POKE feature does not include a carriage return, so this method of printing a reversed message has the effect of a PRINT; (PRINT semicolon) statement. Immediately follow the return-from-subroutine with a PRINT statement if the rest of the message is to be printed on the next line.

A few peculiarities may arise with this modification; among them:

1. The 'fill' character of Electric Pencil may change. This is because any program defining graphics characters as 192 to 255 instead of 128 to 191 is not using them according to the original Radio Shack specifications.
2. Single reversed letters scattered throughout text may not work. This is due to the extremely fast requirements of the circuit, and the fact that certain integrated circuits may not be up to it. Two characters together, however, will print properly.
3. The timing is so crucial that temperature may offset the image slightly. Use polystyrene or polycarbonate capacitors, never ceramic disc capacitors, for this modification.
4. With a screen full of reversed letters, some increase of the brightness control may be needed.
5. The full character set is not available because more hardware would be necessary to obtain the logic information necessary to select out, for example, letters and numbers or letters and symbols.

Although it sounds like I am doing a lot of warning about the limitations of this sort of modification, you should realize that the TRS-80

Stepping on the Accelerator

When you first unpacked your TRS-80 and tried a few calculations and displays, you were likely amazed at the speed with which the TRS-80 was able to respond. But certainly as you used the machine, you realized that it could spend quite a while performing complex tasks. It was then that you joined the ranks of the dozens of thousands of micro users searching for a faster running computer.

When the TRS-80 was designed, the options for higher speed were put in place. Normal and 200 percent speed options were right there on the circuit board, and 150 percent speed required the routing of but a single circuit trace; yet the final design opted for the lowest speed (1.77 MHz).

This modification releases the permanent connection of the 100 percent (1.77 MHz) clock and allows the computer to switch back and forth between this clock rate and a rate one-and-one-half times faster – 2.66 MHz – or two times faster – 3.55 MHz. There are a few pitfalls, but the actual modification is a simple one. For it, you will need the decoded port 254 explained in the reverse video section of this chapter (74LS02), the other half of the flip-flop (74LS74), and one 74LS367.

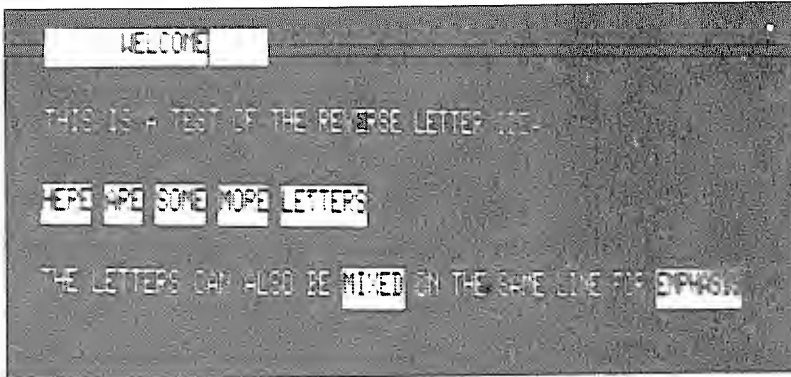


Photo 4-13. Individual reverse screen example.

Reverse lettering provides an emphasis; flashing from one to the other is extremely effective for prompts.

was never meant to do this sort of thing, and the fact that it does work is remarkable. Once it is in place, you will wonder how you could have created reasonable self-prompting programs without it. See Photo 4-13 for proof of the impressive results.

A complete circuit containing the reverse video, upper/lower case, and individual reverse video modifications appears in Figure 4-10, below:

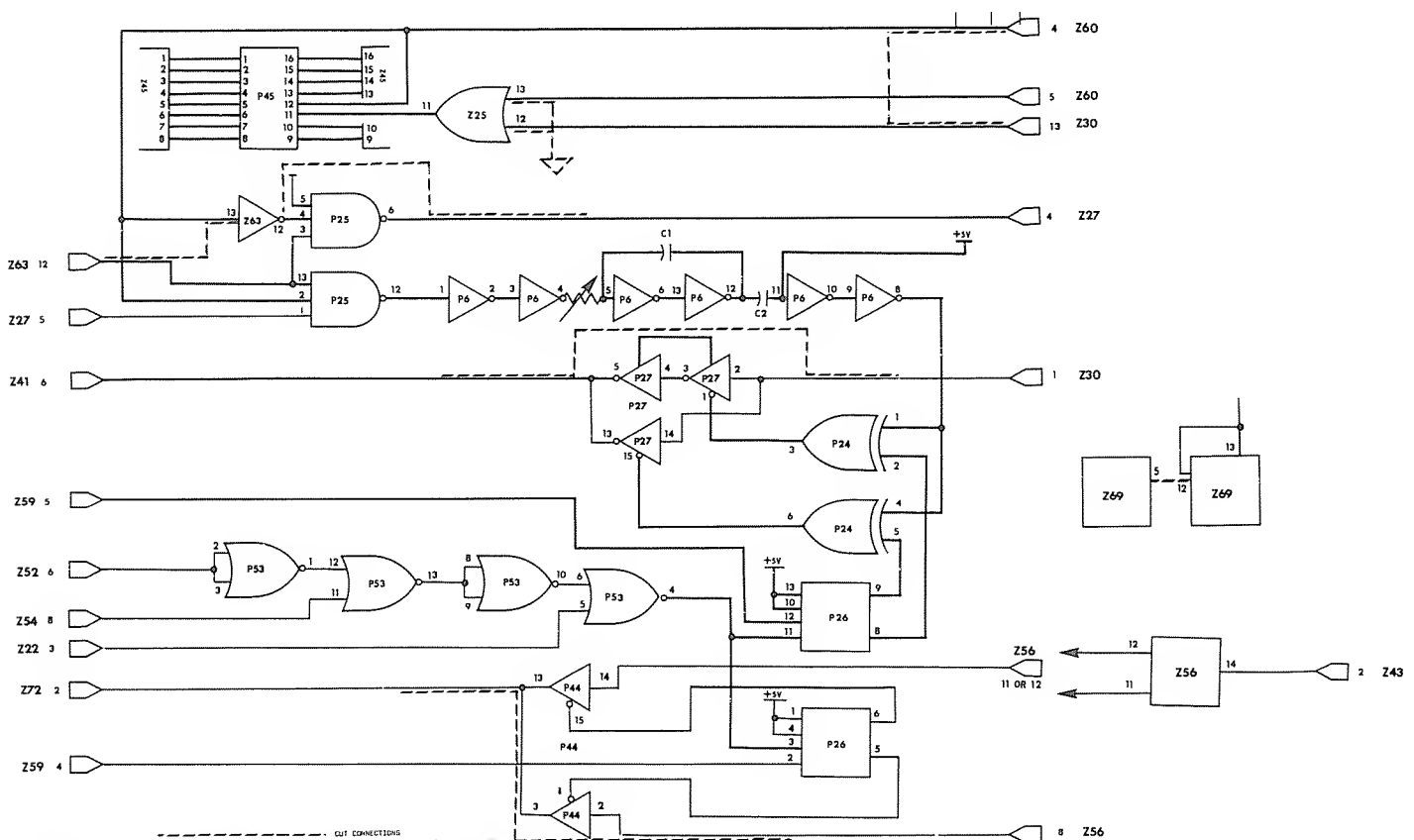


Figure 4-10. Complete video modifications schematic.

Before you perform this modification, there are a few things you should know. For a reason which I have yet to determine, some Level II machines with the *two-chip* ROM set will not accept a speed-up successfully, though most will. Second, later units have been manufactured (or earlier units may have been retrofitted) with a small board known as 'XRX III', a synchronous, 500-baud wave shaper. Cassette load at the 750- and 1000-baud rates will not work unless the modification is deactivated (for details on the cassette system, see the Supplements to Chapters 3 and 6).

The third item concerns the expansion interfaces manufactured since January 1980. Unfortunately, these interfaces must also be modified to accept a higher-speed CPU. But this modification (see Chapter 5) is quite simple and very reliable. Finally there is the question of memory speed. If your TRS-80 is *very* early and the 16K RAMs were shipped with the unit, there is a small chance that one or more of these RAMs will not be capable of running at 3.54 MHz – and an even smaller chance of not operating at 2.66 MHz.

That said, I'll turn to the modification itself. Locate Z56 on the circuit board. Cut the foil trace that leads from this pin to the hole that is plated through the circuit board. Mark that hole for future reference. By cutting this trace, you separate the 1.77 MHz output of Z56 from the clock input of the Z80 processor.

If you have not created the port 254 decoding used in the reverse video modification, you will need to piggyback a 74LS02 on Z53. Bend all the leads except 7 and 14 parallel with the body of the IC, and seat it atop Z53 with the notch or dot pointing in the same direction as the rest of the integrated circuits on the board. Solder pins 7 and 14 to Z53; this piggybacked circuit will be called ZPORT.

Locate Z52; run and solder a wire from pin 6 to pins 2 and 3 of ZPORT. Find Z54. Run and solder a wire from pin 8 of Z54 to pin 11 of ZPORT. Pins 1 and 12 of ZPORT are connected together. This completes the decoding of port 254 (hex FE). To add the necessary OUT signal, run a wire from Z22 pin 3 to ZPORT pin 5. Solder together pins 8, 9, and 13 of ZPORT; solder together pins 6 and 10 of ZPORT. The BASIC command OUT 254,X will activate the signal found at ZPORT pin 4.

Also a part of the reverse video modification was the piggybacking of a 74LS74 atop Z26. Bend all leads parallel to the body of the IC

except power pins 7 and 14. With the integrated circuit oriented in the same direction as Z26, mount it there, soldering pins 7 and 14 to the IC below. This IC is called ZFLOP.

Run a wire from the decoded signal at pin 4 of ZPORT to pin 3 of ZFLOP. Z59 has a data line (bit 0) at its pin 4; run a wire from that pin to pin 2 of ZFLOP. Now run wires connecting together pins 1, 4 and 14 of ZFLOP. With these connections made, OUT 254,0 will flip the circuit, and out 254,1 will flop it. Note that OUT 254,0 and OUT 254,2 were the flip-flop commands for the reverse video.

Now the 74LS367 is put in place. As with the other ICs, bend all leads except the power pins (pins 8 and 16 on this circuit) parallel with the body. Place the 74LS367 on Z44, and solder pins 8 and 16 to it. Call this ZFAST.

Locate pin 8 of Z56, whose trace is already cut. Run and solder a wire from this pin to pin 2 of ZFAST. Run a wire from pin 3 of ZFAST to the plated-through hole, which was previously marked for reference. Take care that this is the correct hole before soldering; it is the other end of the trace cut from Z56.

The normal clock will now be reunited with the CPU. Run a wire from pin 5 of ZFLOP to pin 1 of ZFAST and solder. When OUT 254,1 is executed, the normal speed will flip in place. The other speed output buffer is wired now. Tie pins 13 and 3 of ZFAST together. Last, run a wire from pin 6 of ZFLOP to the second gated section of ZFAST, pin 15.

The higher speed will now be selected. The options are two, and you may try either 150 percent normal speed or a hot 200 percent normal speed.

150 percent normal speed: Locate Z43, pin 2, which is the 5.32 MHz clock normally used in the video divider chain. Run a wire from this pin 2 to Z56 pin 14. That pin is the input of an unused divide-by-two segment of Z56. The output (2.66 MHz) of this divider is present at Z56 pin 12. Run a wire from there to pin 14 of ZFAST.

200 percent normal speed: Locate Z56, pin 11. This is the 3.54 MHz clock not used in the TRS-80. Run a wire from this pin 14 of ZFAST. For this super-fast mod, the memory select circuits must also be dealt with. Locate Z69, and cut the trace running from pin 5 to pin 12. Connect pin 12 to pin 13. This speeds the memory-select process (from MREQ and RD) just a tad, but enough to cope with the 200 percent modification.

Level I and Level II Together

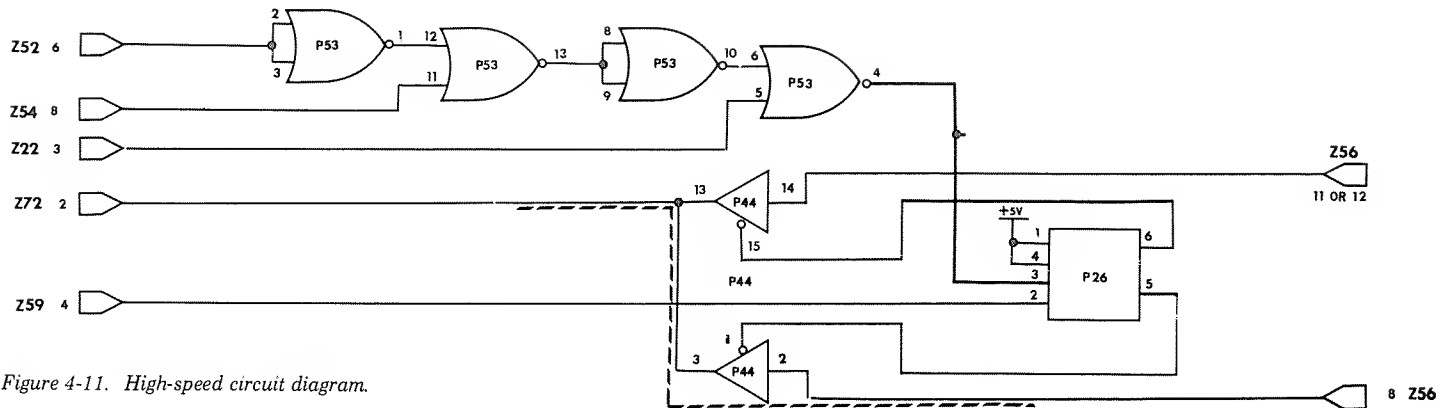


Figure 4-11. High-speed circuit diagram.

The high speed modifications are now complete. Reassemble the TRS-80 and run a test by trying OUT 254,0 and OUT 254,1, which will flip and flop the speeds. The following short BASIC program will give a good demonstration of the speed differences:

```
10 FOR B = 0 TO 1
20 CLS
30 OUT 254,B
40 FOR X = 1 TO 50
50 PRINT X;
60 NEXT : NEXT
70 GOTO 10
```

Listing 4-4. High-speed demo/test program.

In addition to the modifications themselves, it is often useful to know which speed is active. There are ways of telling after working with the higher speeds: the very slight herringbone in the video monitor (turn contrast high to see it) changes pattern, programs work faster, etc. However, a bipolar LED is ideal for this.

To pin 1 of ZFAST, attach one end of a bipolar LED. Attach the other pin to a 470-ohm resistor, and run this resistor to pin 15 of ZFAST. The light will show green for one speed and red for the other; by reversing the center pin of the LED, the color pattern selected can be reversed. I use red for the normal speed, and green for the higher speed.

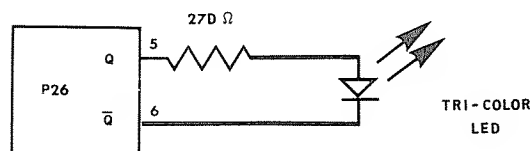


Figure 4-12. LED high-speed indicator circuit diagram.

Level I and Level II Together

Little defense of Level I BASIC has been made, but from my point of view, it's a delightful, compact little BASIC. My own familiarity with it is relatively recent, because I ordered my TRS-80 with Level II installed. Aside from the maddening lack of key rollover, Level I seems an ideal teaching language, especially for youngsters. Together with the excellent Level I manual created by David Lien, it is a fine introduction to the language, and to computers themselves.

Enough for the defense. The problem with Level I is that it is not Level II, and the bulk of we TRS users have Level II BASIC installed. Level I can nonetheless be co-resident, and there are three ways to do it: install the Level I ROM with a switch (this is the method presented here); use a disk system with the Level I-in-Level II program (offered by Apparat and others); or relocate and rewrite Level I a bit and burn it into an EPROM placed in high memory.

Level I ROMs are getting harder to obtain because of the Model I's discontinuance, but cooperative repair centers or franchise Radio Shacks ('Associate Stores', Tandy calls them) can often provide the ROM for a few dollars. Although I am not one to encourage software copying, I do feel that as a TRS-80 purchaser, you paid for your Level I ROM if you bought the unit with one. If it was not returned when Level II was installed, then Radio Shack owes you one. Try to get it.

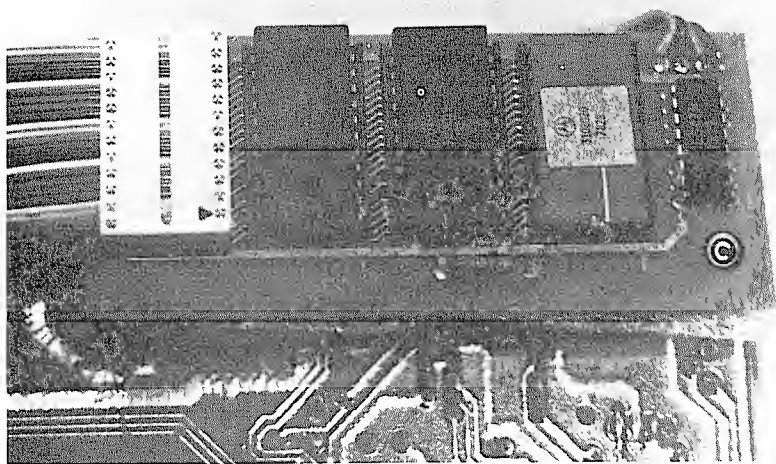
The other way is to order the ROM as a replacement part; Radio Shack's latest replacement number is MX4126, and they charge \$71.25 (!) for it. Finally, you can borrow one from someone who has it and copy it into a 2732 EPROM, but that's still about a \$40 investment.

Let's assume you're able to obtain a Level I ROM. The ones you want are marked National, and are identified with parts numbers M2316E/MMS258ET R/N and S/N, or Motorola, marked 7807 and 7804 or 7831-BASIC1-ROM A and 7832-BASIC1-ROM B. Ideal is the single-chip ROM from Motorola or Rockwell, marked 7809 and 7845, respectively (phew!). Photo 4-(?) shows the two-chip (7831/32) set and the Rockwell (7845), mounted on an aluminum-foil-covered vegetable tray.

You do *not* want the chips marked Intel; these are 2716 EPROMs, and the wiring is complex. About the only thing they are good for is scraping off the label and erasing under ultraviolet light. Then you have two spanking new \$12 EPROMs.

Check next to see if you have the two- or three-chip Level II ROM set. If you have the three-chip set, there will be a connector cable running to a separate board taped to the main circuit card. If not, both ROM sockets will be filled, and installing Level I will be more difficult. Last of all, make sure your circuit board is a 'D' board, 'G' board or later. 'A' boards won't do. (The number is part of the lettering silkscreened on the board, such as 1700069D or 1700069G).

First will be instructions for installing the ROM in the TRS-80's with outboard Level II ROMs. Mount a double-pole, double-throw switch conveniently, but discreetly enough that you won't be knocking it into Level I in the middle of a four-hour data sort.



Level II ROMs: 3-chip set on individual board at the end of a cable which is plugged into the main CPU card.

Open up the TRS-80, and note where the Level II ROM cable is plugged. Four (or six) other wires run from this Level II board to the rest of the circuit card. Find these locations:

1. The green wire on the Level II board, connected near the underside of dip shunt X3.
2. ROM socket Z33 or Z34 (whichever is empty), pins 18 and 20.

Cut the traces leading from pins 18 and 20 of the unused ROM socket. Add a short length of wire between pin 20 and the far end of the trace that used to lead from pin 18. Solder a long white wire to pin 18. Remove the far end of the Level II board's green wire from its connection point near shunt X3, and solder a red wire there. Solder a blue wire to the 5-volt supply found at Z57, pin 14. Using Figure 4-13, run the white, green, red and blue wires to the double-pole, double-throw switch.

Add two 1000-ohm resistors to the circuit to hold the ROM chips inactive when they are not in use. Without them, inadvertent selection might take place, and the running program (in either language) might crash.

If you've got the single-chip Level I ROM, you're all set to go. If not, the fun begins here. Locate the notch on each Level I ROM chip, and line up the two chips with each other, *precisely pin for pin*. Now piggyback one atop the other and solder all 24 pins so that the result is a single, hulking, integrated circuit. Solder carefully, keeping the bottom IC anchored in conductive foil. The foil will act as a static remover and heat sink, both essential in this mod.

Blobs of solder can be removed with some sort of solder-wick or a solder-removing vacuum tool. This entire chip is then inserted in the empty socket, and the unit is ready for testing. Level II gives you the expected MEMORY SIZE?, while Level I only reports 'READY'.

If you have the two-chip *Level II* ROM set, there's a bit more work to do. One ROM must be removed and piggybacked on the other, except for pin 20. Bend pin 20 out straight on ROM A (Z33), and piggyback it on ROM B (Z34). Run a wire from ROM A pin 20 to Z74 pin 9. Bend pin 20 out straight on ROM B (Z34), and run a wire from this pin to Z74 pin 12. Now insert your *Level I* ROMs in the socket for Z33.

Level I and Level II Together

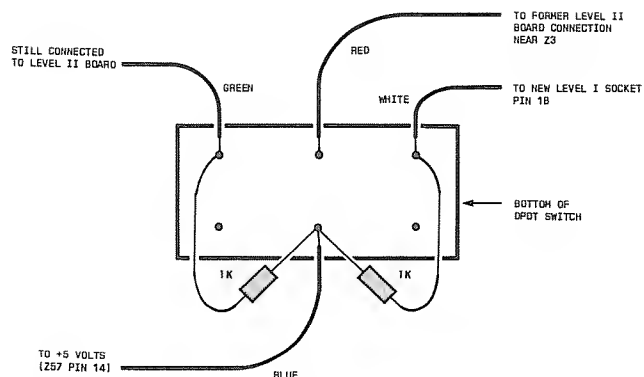


Figure 4-13. Level I and II switch wiring (3-chip set).

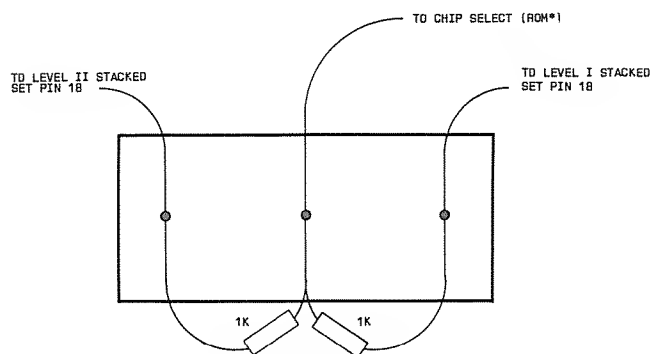
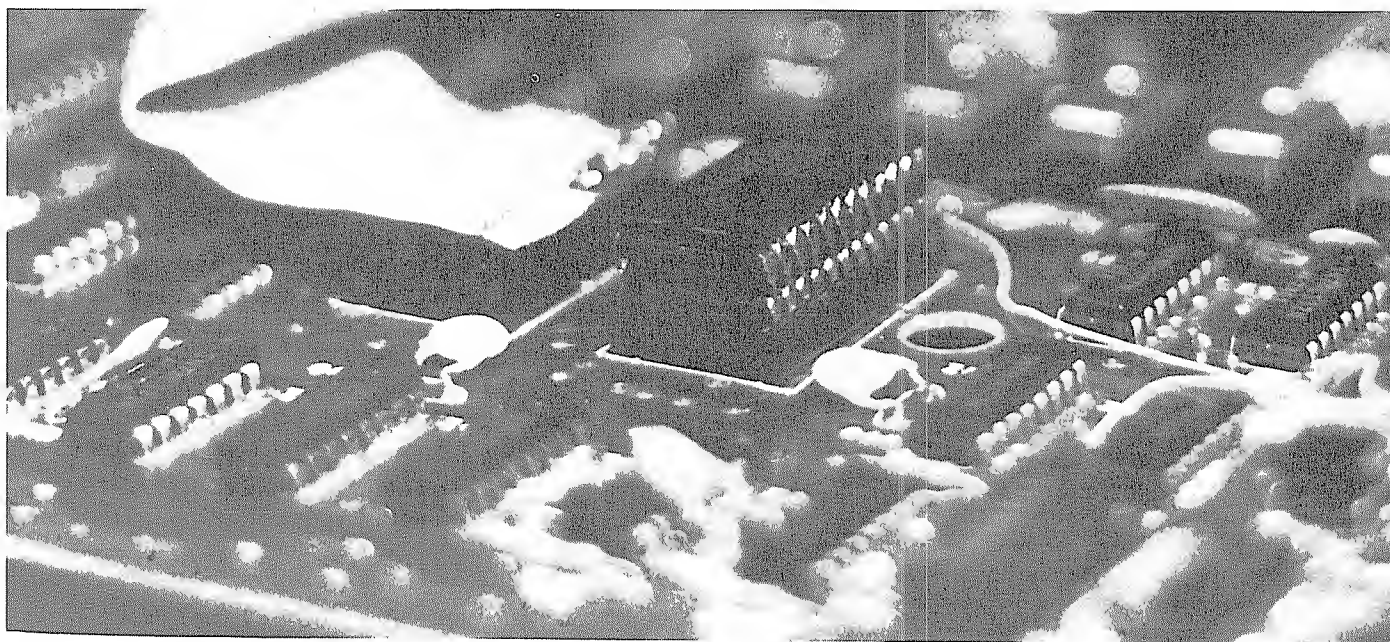


Figure 4-14. Level I and II switch wiring (2-chip set).

Since these languages use memory and pointers in such a different manner (Level I is not by Microsoft, as I understand it), they cannot be switched from one to the other directly. You must power up to the language you want.

This can present a problem. If you have installed various speed and video modifications, you may get the mode you don't want. Level I, lacking OUT statements, can't recover from this. To avoid it, solder a wire to Z53, pin 12, and attach it to one end of a pushbutton; to the other end of the button, solder a wire to ground (Z53 pin 7 is fine for that). This is the system reset (SYSRES) signal, a true restart to 0000.

To use it, power up to Level II, get in the video/speed mode you wish using OUT statements, switch to Level I (the screen will likely display a mess of @9@9@9's), and press the SYSRES button. A READY should appear. Now dig back in the Level I manual to determine if the amount of memory reads correctly.



Two-chip Level I ROM set can be added by piggybacking them and soldering them together. Together with the snaking Level II cable, they form a surrealistic electronic landscape.

On Relocatable Code

A great noise is often made by programmers and users of the TRS-80. The clamor is for an elusive programming quality which allows code to be placed anywhere in memory and to function from that place. The cry is for relocatable code, but it is often called for without understanding either the process or the sacrifices it requires. In this section I will take a look at the various levels of relocatability, and the ways it can be applied in low-level languages.

The most 'relocatable' code TRS-80 users see is BASIC itself. No matter what configuration changes the system undergoes, whether they be the use of low-memory utilities, DOS, etc., a normal BASIC program will load and run. A BASIC program's portability depends on the premise that it is *interpreted*, containing statements that are equally meaningful irrespective of the program's position in memory. Occasionally that position gains importance, as when numerical variables, strings or arrays are referenced within program lines, but normally those positions do not become significant until the program is in operation.

Unless the BASIC program is truly a hybrid (having other, lower-level languages artificially embedded in the code), it is eminently portable. All references are resolved by its interpreter, and relocatability of BASIC is assumed in the very nature of it and similar high-level, interpreted languages. In fact, that is the nature of a real 'language' as we know it as opposed to the mysteries of 'code'.

The distinction between language and code is not an accepted one, and the differences will be characterized only for descriptive clarity within the context of this article. By 'language', I intend BASIC, FORTRAN, Pascal, and others, as well as assembly language. By 'code', I suggest machine code as written in binary or hexadecimal for direct insertion in memory. The assembly language/machine code interface is curious and I will address it in terms of the conflict between understanding and using a language as opposed to a code.

At the High Level

The highest levels of languages are the so-called 'descriptive' languages, like IBM's report-generator, RPG II. A programmer need only describe a group of tasks and the order in which they should be performed, and they will be magically executed. Full-dress accounting

reports, billings, and record-keeping can be done with a language like RPG II. Relocatability is at best the task of a dour 'systems engineer'.

Likewise, BASIC normally requires no thought to its arrangement beyond the proper sequence of line numbers. Statements are arranged in the programmer's chosen order, and the interpreter selects each one for execution in that order. Line numbers are the programmer's relocatable reference. Languages containing no line numbers often use an implicit 'top-down' order; that is, once a task is completed, it is never returned to. Still, the program is created with an arbitrary arrangement selected by the programmer and evaluated by an interpreter before its execution.

Languages which use compilers rather than interpreters work in much the same manner. Whereas an interpreter selects and executes each command and operand during the program's run, a compiler produces a block of machine-coded information from high-level commands and operands. The process is completed before executing the program, so valuable space and running time can be saved. In either case, arbitrary, high-level language can be moved about, deleted, inserted, or otherwise altered with little regard for the program's ultimate position within the bowels of the computer.

In fact, the question of relocatability is implicit in the term 'high-level language', and is an essential of its high-level quality.

The reason the question arises at all is because, when writing machine-coded utilities, many programmers feel that emulating the portability of a high-level language is crucial. It becomes even more desirable when programming is done for a user rather than another programmer, a user whose contact with the machine may go little further than that of a traditional computer 'operator'.

Accomplishing relocation can be an easier task on the Model I than on other computers simply because of the extensive ROM-based operating system. Before examining some ROM uses (and its advantages and pitfalls), let's have a look at the qualities and constraints of relocatable machine code.

As a user interested in accomplishing a task, I might need to select from a group of utilities including sort routines, multiple-precision transcendental functions, mundane key-debounce and upper/lower case drivers, and so forth. Ideally, the order in which I select and load these routines should be irrelevant to their

operation, in much the same way that utilities can be 'ordered up' on a mainframe.

Microcomputer software authors have not generally respected this need; in some measure that is forgivable because the approach and intent of different utilities from a variety of vendors demonstrates that relocatability is only one of many potential conflicts to be resolved.

Hunting Through Memory

Nearly any well-written program is tailored to meet certain requirements which may not easily co-exist. Among them are:

Speed of execution. This applies in two ways. The first is the speed at which the need for the utility is evaluated - an example might be a routine scanning a keyboard for a unique combination of characters. This should be accomplished expeditiously, since no real program 'work' is being done. The second speed consideration is that, once called for, the main routine should complete its work handily.

Conservation of memory. In a system of unlimited size, speed can be optimized by replication of routines rather than use of subroutines. In appliance-level microcomputers like the TRS, however, support utilities cannot be allowed to consume memory space needed by the operation of the target program.

Resolution of Conflicts.

All operating systems make use of areas of critical reserved RAM, which are employed to tie together the major pieces of hardware and blocks of software. These patch points (I prefer this term to 'vector', which also means a disease-carrying insect) may be needed by many other utilities as well as the operating system, and must be redirected carefully.

ROMability. The traditional approach states that programs must be created so that they can be committed to ROM at any time. That is, they should neither modify themselves nor contain variable data within their bounds. This approach can certainly be argued against (I shall).

Three of these four elements - speed, size, and memory sharing - are normally considered in utility programming, and become critical in relocatable programs.

In addition, factors to be considered exclusively for program relocation are:

The placement of the program in memory. This can be a single- or multiple-step process. For example, during loading from disk or tape, a program may temporarily reside in a single block of memory, only to separate itself into smaller blocks which are shepherded to other memory fields. *Electric Pencil* is one such program.

Absolute jumps within a program. The program counter is instructed to take on a new value, a value specified as an address. These jumps are, when taken, faster than the relative branches, because the program counter need not calculate an offset to its current address.

(indent absorber) **Subroutine calls within a program.** Effective use of memory space makes subroutines very attractive, but they require that the program counter be assigned a specific new value for their duration.

A relocation block. If the program is provided to a user in object-code (SYSTEM) format, some portion of the program must be dedicated to relocating itself in another area of memory.

Assembly Programming

When programming for programmers, the most convenient way to effect relocatable code is to use assembly language, and supply that to the user. By labeling all jumps and calls, such a program becomes 'relocatable', somewhat in the sense that a program conceived in a high-level language is relocatable. In fact, in its assembly phase, machine coding is a high-level language.

	8000	ORG	8000
21003C	8000	LD	HL,3C00H
11003C	8003	LD	DE,3C01H
01FF03	8006	LD	BC,3FFH
3620	8009	LD	(HL),20H
ED80	800B	LD	HL,42E9H
21E942	8000	LD	A,(HL)
7E	8010	AND	A
A7	8011	JP	NZ,1997H
C29719	8012	JP	4300H
C30043	8015	ORG	4300H
	4300	CALL	4388H
C08843	4300	INC	HL
23	4303		

Listing 4-5. Machine coding in assembly language.

Above is a section of assembly programming. It makes several assumptions, some valid and some not. Let's consider that this program has cleared the screen (which it would in the TRS-80

configuration) after having performed some operations on a BASIC program. It checks location 42E9, the start of BASIC, for a zero. If it does not find a zero, it assumes a syntax error. Otherwise, it moves to location 4300, where it immediately calls a subroutine at 4388.

This is perfectly valid code, but here's a look at another way of creating it in this high level language:

```

04XA FIGURE 2      SETUP EQU 8000H
                     VIDEO EQU 3C00H
                     SCREEN EQU 03FFH
                     BASIC EQU 42E9H
                     PROGRAM EQU 4300H
                     SYNERR EQU 1997H
                     TESTER EQU 4388H
                     ORG $
21003C 8000          SETUP LO HL,VIDEO
11013C 8003          LO OE,VIDEO+1
01FF03 8006          LO BC,SCREEN
3620 8009            LO (HL),20H
E080 8008            LOIR
21E942 8000          LO HL,BASIC
7E 8010              LO A,(HL)
A7 8011              ANO A
C29719 8012          JP NZ,SYNERR
C30043 8015          JP PROGRAM
                     ORG $
C08843 4300          CALL TESTER
23 4303             PROGRAM INC HL

```

Listing 4-6. Machine coding in assembly language.

This assembly program, using its high-level, interpretive capabilities, can be easily relocated by a programmer changing a few equates. But it is not relocatable because it does not meet the criteria outlined above – i.e., it contains a specific origin (or rather, two), absolute jumps, calls to subroutine, and has no visible relocation block. Fortunately the start of BASIC has been defined in the table of equates, for it too can be a variable element where low-memory alterations have been made.

As a programmer's utility, this method is acceptable. As a public program, it is questionable. The solutions, however, are neither easy nor obvious.

Relocation Blocks, Time, and Space

The first option is to provide a relocation block. The amount of coding the programmer must perform is increased, but memory use is not affected because this relocation block can be deleted upon relocation. A relocation block is provided with Radio Shack's KBFIX, for example, and with other utilities. The sample above might have a relocation block which asks for a new base address in protected memory, protects that memory, adjusts its internal addresses, and automatically checks for the beginning of BASIC program storage and other patch points it may require.

For very short utilities, the relocation block may be longer than the utility itself, because it demands a user prompt and input, minimum error checking, calculation of base address, and adjustment of patch points. Even using ROM routines (such as 28A7 to display a message, 1BB3 to accept user input, 1E4A for numeric conversions), this process is lengthy.

One of the most painful aspects of this relocation is the adjustment of internal addresses; absolute jumps, calls, and table locations are among the addresses requiring such changes. Then what of the questions of speed and conservation of memory? Programs that move themselves during the course of operation must carry the excess baggage of a lengthy self-relocation block or use a plethora of relative instructions which can cost time during program execution.

Let's address these speed and space questions. Table I lists a number of the relative instructions in the Z-80 set, along with the bytes, machine cycles and time required for their execution. In contrast, Table II presents similar absolute Z-80 instructions, their execution speeds, and the percentage of time and space they require as opposed to the relative instructions.

Z-80 INSTRUCTION	BYTES	T-STATES	TIME AT 1.77 MHz
LD r, (IX+d)	3	19	10.7 uS
LD (IX+d), r	3	19	10.7 uS
LD (IX+d), n	4	19	10.7 uS
ADD A, (IX+d)	3	19	10.7 uS
INC (IX+d)	3	23	13.0 uS
RLC (IX+d)	4	23	13.0 uS
BIT b, (IX+d)	4	20	11.3 uS
SET b, (IX+d)	4	23	13.0 uS
JR e*	2	12	6.8 uS
DJNZ e	2	13	7.3 uS

When this instruction is conditional, and the condition is not met, it uses 7 T-states (3.9 uS).

Note:

'r' = registers A, B, C, D, E, H or L.

'd' = displacement byte 00 to FF

'n' = single-byte integer 00 to FF

'b' = bit position 0 to 7

'e' = offset byte 00 to FF

Z-80 INSTRUCTION	BYTES	T-STATES	TIME AT 1.77 MHz	% TIME & SPACE VS. REL. INSTR.
LD r, (HL)	1	7	3.9 uS	37% 33%
LD (HL), r	1	7	3.9 uS	37% 33%
LD (HL), n	2	10	5.6 uS	53% 50%
AUD A, (HL)	1	7	3.9 uS	37% 33%
INC (HL)	1	11	6.2 uS	48% 33%
RLC (HL)	2	15	8.5 uS	65% 50%
BIT b, (HL)	2	12	6.8 uS	60% 50%
SET b, (HL)	2	15	8.5 uS	65% 50%
JP nn*	3	10	5.6 uS	83% 150%
JP NZ, nn	3	10	5.6 uS	77% 150%

When this instruction is conditional, and the condition is not met, it still uses 10 T-states, 3 more than the relative branch instruction. In programs where such a condition is generally not met, the JR e instruction will save both time and memory.

Note:

'nn' = two-byte integer 0000 to FFFF

Relative instructions in many cases command considerably more time than absolute ones. How does this reflect on program speed? The design engineers for the 6809 chip reported in *BYTE* a survey of 6800 instruction class usage based on static analysis (i.e., with the program not running) of 25,000 lines of source code. Because of the considerable differences in chip architecture between the Z-80 and the 6800, especially in regard to the Z-80's multitude of registers, this information (see Table III) is not directly applicable. Yet it is instructive, particularly in the percentage of subroutine calls and branches, instructions similar in the two microprocessors.

TABLE III

Use of 6800 Instruction Types

Loads (movement from register to register, and from memory to a register)	23.4%
Stores (movement from a register to memory)	15.3%
Calls and returns (absolute addressing)	13.0%
Conditional branches (relative)	11.0%
Unconditional branches (relative) and jumps (absolute)	6.5%
Others	30.8%

(Adapted from *BYTE*, 4:1, January 1979, p. 26)

A remarkable 30 percent of program space is dedicated to motion of the program counter. This might suggest that a program using absolute instructions would be 15 percent faster overall than one using relative branches and calls. Since relative calls and returns are complex, however (see below), the time savings might be even greater. Absolute jumps use more space than relative ones, so space savings, though significant, would not be as impressive.

What are the practical implications of this speed differential? A 1,000-byte subroutine (long by microprocessor standards) with no internal loops might execute in four milliseconds, as opposed to five milliseconds for such a routine written exclusively with relative functions. On the average, it might be 200 bytes longer. This means 200 iterations per second as opposed to 250. For a complex mathematical function which iterates through a routine many times, it could mean the difference between a four-hour and a five-hour run. For a patch to a

keyboard scan, it could mean adding only a few seconds to an hour's program time. And for routines involved in occasional string handling, printing, low-speed input/output, error reporting, or a host of other applications, the loss of time is insignificant.

The application itself is the answer to the viability of relative coding. In time-sensitive applications (music-generation routines, for example), and where every byte of code is crucial, relative instructions can interfere with a program's effectiveness. In the bulk of TRS-80 applications, which are seldom so time-conscious, relocatable coding through use of relative instructions can be very successful.

Creating Relocatable Code

The Relative Branch. The first coding choice is the two-byte instruction for changing the program counter. No matter where the program counter is currently pointing, it can be shifted by specifying a relative branch. When the high bit of the specified offset is zero, the program counter is moved ahead (for example, 18 37 adds 37 to the program counter). When the high bit of the offset is one, the program counter is moved back, in effect subtracting 80 before adding the offset:

5236	5236	:	5236	5236
+ 00	+ 80	:	+ 7F	+ FF
5236	5186	:	5285	5235
(5236+0)	(5236-100+80)	:	(5236+7F)	(5236-100+FF)

Especially gratifying is the notion that the page of the address (the high byte) is incremented or decremented automatically by this instruction:

PC = 52C5	(page 52)
Execute 18 37	
PC = 52C7+37 = 52FE	(page 52)
Execute 18 37	
PC = 52FE+37 = 5335	(page 53)
Execute 18 A7	
PC = 5335+CA = 52FF	(page 52)

The branching can be done unconditionally, or on the basis of the zero or carry flags (JR, JR Z, JR NZ, JR C, JR NC). This does not provide the options available when using the absolute jump (which can act on the status of the parity and sign flags), and thus presents additional coding demands.

A special relative branch instruction, DJNZ e, automatically decrements the B register, and performs a branch to the indicated offset whenever B is not decremented to zero. Although this instruction can have many uses, its two-byte brevity makes it most attractive when executing loops.

The problem of branches outside the range of the +7F/-80 byte offset can be solved by inserting 'stepping stones' in the main program flow. At the end of any logical program section, two unconditional branches may be inserted: the first branch serves to skip past the second branch to prevent disturbing the program flow in that area. The second serves as a stepping stone for some branch which is too far (more than 127 bytes) from its eventual destination. The program can be ordered judiciously so as to avoid overusing such memory-gobbling stepping stones.

Indexed Addressing. One of the best features of the 6500 and 6800-series CPUs is their ability to address and manipulate memory without continually reloading a register or using an absolute address in an instruction's operand. In fact, using these processors would be incredibly cumbersome without this mode since they contain merely two storage registers (X and Y) as opposed to the Z-80's sixteen (B, C, D, E, H, L, B', C', D', E', H', L', R, I, IX, IY).

The value of the indexing features was recognized by the Z-80's designers, who added them to the 8080's primitive instruction set. As with relative branch instructions, a single-byte offset may be added to the value specified by registers IX or IY in order to produce the address of the desired memory location:

	IX = 6AA1 + D1	IX = 6AA1 + 7F	IX = 6AA1 + A1
resulting address:	6AA2	6B2D	6A42

The address of the index registers themselves does not change; the resultant value is stored in temporary internal Z-80 registers while an instruction is being executed. For example, when IX = 6AA1 and the instruction –

INC (IX + A1)

– is specified, the CPU temporarily creates the resulting address 6A42 and increments its contents. The value 6A42 itself is then discarded.

When a desired value is outside the range of an indexed register, the instruction ADD IX,pp comes to the rescue. The designation 'pp' is for any of registers BC, DE, IX or SP, which can be assigned specific values representing a full 16-bit offset, rather than the 8-bit offset of the indexing itself. For example:

```

IX = 57EB
DE = 41CC
Execute ADD IX,DE
IX = 99B7
DE = 41CC

```

Since it is an offset rather than a specific address, and since it remains unchanged after the execution of ADD IX,pp, it becomes a perfect candidate for use in relocatable code.

Instructions using the index registers can be particularly valuable when lookup tables must be accessed again and again. With these commands, any registers used as pointers may remain unaltered; they need not be redefined each time the lookup routine is used. An excellent application of the index registers for five-character string comparisons is found in William Barden's *Z-80 Microcomputer Handbook*.

Relative Subroutines. No, there aren't any secret relative call-and-return instructions that Zilog never told us about. Rather, there is a way to create the effect of a relocatable call by using an index register in combination with relative branches.

This method was detailed in my article, 'Relative Subroutines for the Z-80', *BYTE*, 4:12, December 1979. It is a bit cumbersome at first, and makes it impossible to produce 'ROMable' code. In summary, here is how it works (see also Table IV):

1. An index register is set to some point in the program. This becomes a reference point.
2. The call is prepared by determining the offset from the *end* of the subroutine *back* to the main program flow.
3. During program execution, the offset is assigned, (using the indexed instruction LD (IX+d), n) as the *operand* of a relative branch placed at the *end* of the subroutine.
4. A relative branch is executed to the beginning of the subroutine.
5. When the subroutine reaches the end of its execution, it moves back to the program in progress.

TABLE IV

Relative Subroutine Flow

```

7000 00218070 LO IX,7080 ;IN PROGRAM AREA
7004 .....< MAIN PROGRAM HERE >.....
7006 .....< PROGRAM CONTINUES >.....
.
.
7040 003608BF LO {IX+B},BF ;RETURN OFFSET
7044 1820 JR S+22 ;JUMP TO SUBR.
7046 .....< RESUME PROGRAM HERE >.....
7047 .....< PROGRAM CONTINUES >.....
.
.
7066 .....< SUBROUTINE IS HERE >.....
7067 .....< SUBROUTINE CONTINUES >.....
.
.
7087 1880 JR S-41 ; "RETURN" INSTR.
< NOTE THAT "80" IN OPERAND ABOVE WAS PUT
IN PLACE BY CALLING PROGRAM AT 7040. >

```

This method requires careful program preparation in that the beginning and end of each subroutine must be within the range of the calling program, and the offset to IX or IY must be within the range of the end of the subroutine. If the offset is outside the range of these registers, a 16-bit offset can be created as described above under indexed addressing.

There is another option to create relative subroutines using the Z-80, but it is not time-effective. This involves a ROM call to 000B.

Jerry Lindsly of West Chester, Ohio, describes it this way:

When faced with the problem of relocatable code when using subroutines . . . you can use a very useful routine in ROM. Essentially what it does is this:

```

POP HL      ; Get return address to HL
JP (HL)     ; Return to calling program

```

So after you call 000B the PC is in HL. A relative call is a little more complicated:

```

CALL 00BH
LO OE,0BH
ADD HL,DE
PUSH HL
LO OE,NNNN ; Offset from first byte
           ; past JP (HL).
           ; Use XOR A and SBC HL,OE
           ; if the call is negative.

ADD HL,OE
JP (HL)

```

(Those interested in the various uses of this routine may write to Jerry at 8106 Quailwood Court, West Chester, Ohio 45069).

One could use the equivalent to such a ROM call in the program itself. This might be desirable. Certainly, if programmers had not learned it over the past two years, then the recent ROM changes (and expansions in the Model III) might dissuade overuse of calls to the Level II ROM.

5

How the System Expands

Unlike automobiles, refrigerators, or shovels, computers are hardly ever completely self contained. Even if their manufacturers pretend they are making a complete unit, the machines themselves seem to desire bursting forth from their shells, commanding us to create a tabletop electronic octopus.

The reasons almost sound metaphysical: it needs more memory; it needs communications capabilities; it needs mass storage; it needs hardcopy . . . to achieve these goals, expansion capabilities were made available for the TRS-80 by means of an electronic hodgepodge called an Expansion Interface.

The Radio Shack Expansion Interface

The Radio Shack interface originally consisted of eight major sections :

1. An extension of the keyboard's edge-card connector, initially reserved for a screen printer, but containing all the signals on the keyboard's connector.
2. Two groups of eight sockets, together with address decoding, to read and write two 16K banks of dynamic memory.
3. A decoded input/output latch for parallel printer control. A 'Centronics compatible' electronics scheme is used, with a separate edge-card connector dedicated to the printing task.

4. A set of decoded memory addresses and a LSI (large scale integration) chip for disk drive and input/output control. The disks are accessed through an independent edge-card connector.

5. A decoded output latch for dual cassette selection and control. One input and two output jacks are integral to this circuitry.

6. A decoded, input/output port for serial communications control, with accompanying terminal for an RS-232 circuit board. The RS-232 board is accessed via a separate edge-card connector.

7. A crystal and divider circuitry to provide disk access, real-time clock, and other interrupt functions.

8. Power supply circuitry.

If it seems like a lot of unrelated material to pack into a single box, then you're on the right track - the track that, for quite a while, meant trouble. The electronic clamor inside this box was to create a system-crashing din during the first year or so of expansion interface manufacturing.

The problems were manifested by memory crashes (return to MEMORY SIZE?, keyboard lockup, or, in a disk system, complete reboot and so on). There were several culprits, but the gangleader was the misconceived design itself. Putting all that material on a single board was inviting trouble, and the actual execution of the circuitry compounded it. The main difficulties

were memory refresh/select, noisy and dirty connectors, weak buffering lines, microphonics, and susceptibility to external interference.

The memory refresh/select lines were noisy and unreliable. The purpose of the refresh lines is to read the entire memory in intervals of two milliseconds or less. The memory-select lines provide the signals to select one of the 65,536 possible addresses when needed by the CPU.

In many computers, selecting a memory address is a fairly straightforward process. The CPU signals a memory request of some kind, and provides an address on the address bus. The selected memory responds by providing or accepting data according to the direction of the CPU. On the TRS-80, however, type 4116 dynamic memories are used. These memories do not have enough external pin connections to accept a complete memory address. Instead, hardware breaks the address into two parts, which are transferred sequentially.

Briefly, this is how it works. The address is sent out on the address bus. The CPU also provides a 'memory request' (MREQ) signal which triggers special circuitry. This special circuitry stands between the address lines and the dynamic memories. This circuitry produces a multiplex signal - MUX - which chooses the low seven bits of the address. It then sends those bits together with a 'row address strobe' - RAS - to the memories and reverses the MUX signal to choose the high seven bits of the address. Then the special circuitry sends those bits together with a 'column address strobe' - CAS - to the memories.

This sequence of operation provides all the address information needed by the dynamic memories to select an address. The three items most crucial for smooth memory action are the RAS, CAS and MUX signals, and as such, these should be clean, noise-free lines. Within the confines of the keyboard unit, this is the case. But once forced to travel through a cable (via two solder-coated edge connectors) into the expansion interface, the signals pick up some measure of noise. Once inside the expansion box, the noise is increased by the surrounding electronic din (remember that a separate crystal is onboard), and the requirement that these signals feed many other devices reduces their effectiveness. They become electronically tired, and memory access becomes erratic and susceptible to external electrical influences.

Thus, memory selection was impaired and memory refresh was not necessarily executed successfully in less than two milliseconds.

The next culprit in the expansion box was the lack of 'buffering'. In the *Technical Reference Handbook*, the author points out that among the requirements for hardware are that it :

4. Contains a separate power supply.
5. Does not contain more than 1 LS TTL load on any one output from the Computer.

Points 4 and 5 are very important, if you want to guarantee proper operation of your Computer.

Radio Shack, not taking its own advice, hangs five LS TTL loads on data line 0, and four on each additional data line. If an RS-232 board is installed, that becomes another load. The address lines feed as many as three devices. With a screen printer or other device connected to the expansion interface this load increases, respectively, at least once for each additional device attached!

The third system-crashing factor is an unusual one, avoided by good circuit board layout, but not always anticipated. This is covered by the rather unusual term *microphonics*. In other words, when the expansion box circuit board is tapped or vibrates, those vibrations are amplified and communicated throughout the entire circuitry by the power supply, data and address lines.

The keyboard unit can take a hefty bounce. If you're brave, give it a try by dropping it six or more inches. Chances are a running program won't crash. But rap sharply on the expansion box with the program running, and . . . you can guess. These strong pulses are not in the high-frequency range which the 'bypass' capacitors (the small disc capacitors scattered throughout the computer) can squelch; rather these are heavy, low-frequency surges which can be electronically interpreted as signal changes. Hence, the crash.

Finally, the electronics, being spread out in a plastic case, are subject to the electrical whims of the outside world. Seated directly below the video monitor, the expansion box is susceptible to any strong noises contributed by the monitor, or through the power lines and rebroadcast by the monitor's circuitry. Printer heads and motors, disk drive motors, and what have you all let the expansion box know they are operating. So it responds by crashing its sensitive and overworked memory circuits.

Several solutions have been created. The first was that bulge between expansion box and keyboard unit: the buffered cable. The second was the 'twisted pair' modification, another bulge between the boxes. The last was a

rethinking of the expansion circuitry and redesign of the board. The last was the only one that worked.

The idea behind the buffered cable was sound - to strengthen the overworked signal lines by having them feed a single integrated circuit each, and then have that integrated circuit feed the expansion box. Furthermore, 'termination resistors' could then be added, which work something like this: power supply lines contain bypass capacitors to squelch noise, but placing capacitors on the signal lines would slow down important signals as well. Instead, low-value resistors could make the signal lines 'hug' the power supply and ground lines, increasing their resistance to low-level, transient, electronic noises.

First installation of buffered cables were less than successful, partly because the three memory-select signals were very fast signals delayed too long by a combination of buffering and transmutation inside the expansion box. Thus, these lines were pulled outside the 40-pin cable via a three-pin DIN connector, separately terminated. This kept them up to speed, but isolated them from the accompanying noise.

Nevertheless, these modifications were only making up for difficulties inherent in the expansion box design, and could not improve the box's susceptibility to noise all around, or its microphonic tendencies.

The solution was to put the buffers on board the expansion interface itself, and redesign the board to reduce interference. New memory-select lines were created by the RAS line alone to eliminate noise in these signals, but the method used was to create problems concerning speed modifications (which, of course, Radio Shack did not authorize and could not consider in its redesign - see the way out of this problem later in this Chapter).

These difficulties, though, should be placed in perspective. When connectors are well cleaned, good memories are used to stock the expansion box, and the computer's environment kept clean, static-free and vibration-proofed, even the first, unbuffered systems work successfully. My own system (on which this book is being prepared) is a 48K, one-disk system with RS-232 board, printer, two modems, *Exatron Stringy Floppy*, two cassettes and home-brew interfaces - all connected to an *unbuffered* expansion box of early vintage. But it must be cared for. The designers were simply not prepared for the TRS-80's ultimate home and business environment.

One increase in reliability is not mentioned in the Radio Shack redesign. Electronic parts are 'derated' according to temperature. That is, as temperature increases, their performance decreases. Some circuits, like 74LS types, are strong and steady right to the limit; others, particularly dynamic memories, become flakey as the heat builds up. Thus, reliability of the expansion box can be improved by removing the two power supply transformers from the expansion box.

Solving the Ground Problem

I call this a ground problem, but it really refers to the microphonic tendency of the expansion board. The first and easiest solution is to remove and cushion the expansion circuit board. One principle of vibration-proofing is that when different materials are sandwiched together, a certain amount of sound and vibration energy is reflected back from the junction of the different materials.

Thus, all sources of potential vibration should be cushioned. Set a thick cloth (like terrycloth) on your computer desk, and a wooden baseplate atop that. On it, place a mat made up of several thin layers of plastic pad, rubber foam, and cloth to a thickness of about 1/4 to 1/2 inch. Assure that the keyboard unit and interface seat comfortably in them (you might drill shallow depressions in the wood for stability), but that there is some air space for cooling. Remove all the plastic doors which attach to the interface and discard them. In their place put rectangles, made of alternating layers of cork and styrofoam, which snuggle firmly between the connectors and the upper and lower walls of the case. See Figure 5-1.

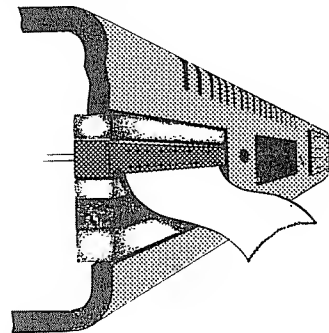


Figure 5-1. Shockproofing diagram.

Other sources of vibration are the cables that attach to the expansion box. Rather than have cables vibrate and swing with a printer, or stretch when you readjust the position of a disk drive or modem, sandwich each one under a leather (or synthetic plastic, if you prefer) bridge, and screw

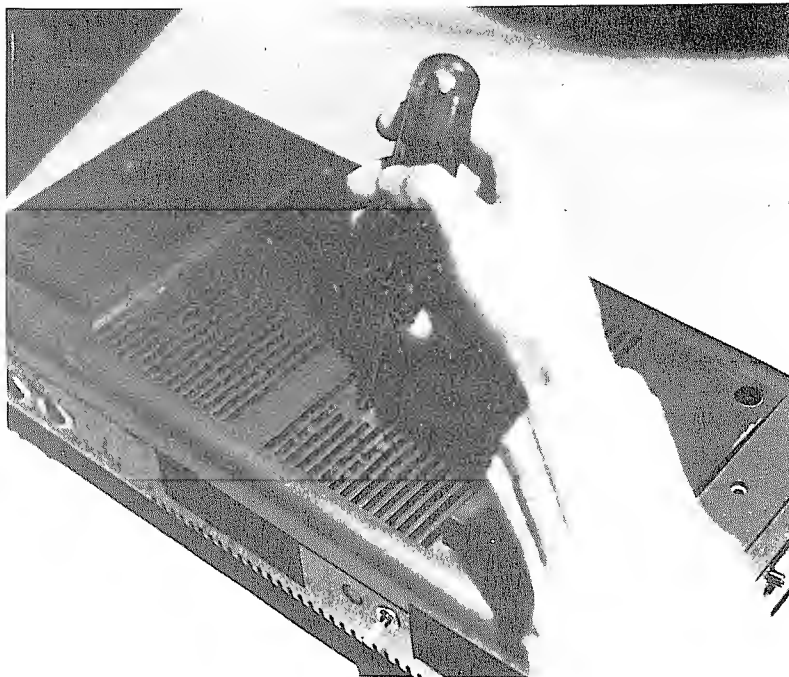
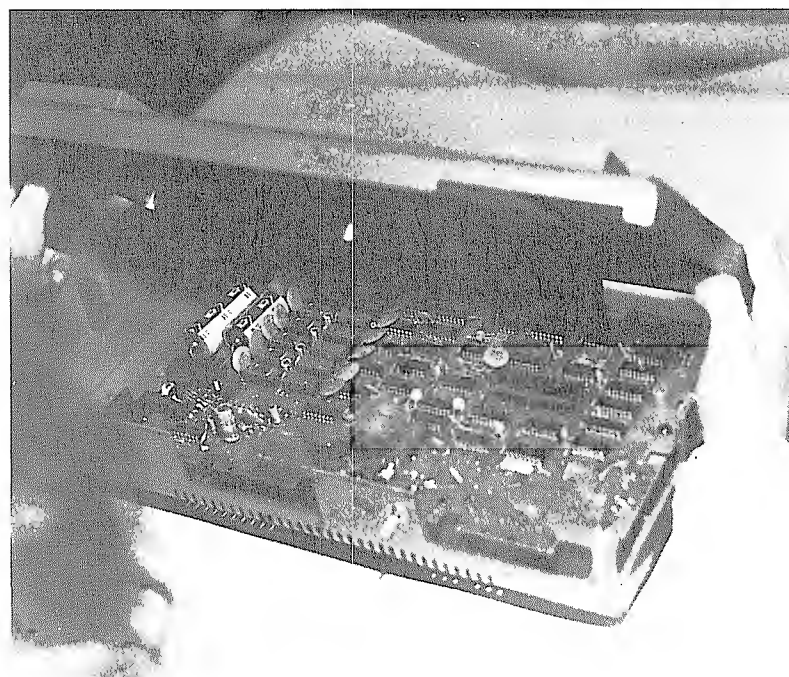


Photo 5-1. Photos on opening the expansion box.
Screws are removed from bottom; note that like the CPU, a warranty-warning label may cover one screw.



Bottom is lifted off to reveal top of expansion circuit board.
Screws holding circuit board are then removed.

the ends of this bridge through the cushion pad into the wooden base. When jostled, the cables should not move between the bridge and the expansion box.

Finally, avoid the possibility of stretching the keyboard/expansion box cable. The best way is to drill shallow depressions in the wood base, as noted earlier.

These are not specifically electronic hardware corrections, because only a thorough redesign of the expansion interface can truly cure these microphonics problems.

Expanding to 32K and 48K Memory

Memory expansion in the interface itself is quite simple. Remove power from the interface, disconnect cables, and lift out the power supply transformers. The latter action is necessary because the power supplies will only fall out anyway when the bottom comes off. Flip the expansion box over, and remove the six black Phillips screws. Lift off the bottom cover.

There are sixteen unfilled sockets on the board, reserved for memory expansion (see Photo 5-2). The *higher-numbered* sockets (Z-9 to Z-16) will be used for the first block of memory, for a 32K total RAM. The lower-numbered sockets expand the system to the complete 48K.

The memories to be used are type 4116, as used in the keyboard unit (see Chapter 4). There are some differences, though. Oddly, some sources claimed that slower memories seemed to work better in the older interfaces, but fast 250 nS RAMs are the most commonly available now. I have not come across an expansion box of any vintage that could not take fast RAMs like these, including my own, which dates to early 1978.

First, refer to Chapter 3's notes on handling static sensitive integrated circuits. Take each RAM and orient it with its notch or dot at the top facing in the same direction as the notch in the sockets (and in the same direction as the rest of the integrated circuits on the board). Fit each one of the first eight chips carefully in place, being sure that the pins slide into the sockets and neither buckle underneath the IC nor slip outside the socket edge. Press the memory chip gently but firmly in place.

Now, before installing another 16K memory block, test this group. Flip the expansion box over (the cover need not be installed), reconnect cables and power supply, and power up the expansion box. While holding (BREAK), power up the computer itself, and press (ENTER) in response to MEMORY SIZE. Now PRINT MEM, and a figure of 31956 will be displayed. If all is well, power down, remove cables, and

insert the remaining eight memory chips. Again restore cables and power, and go through the initialization procedure. Memory should now read 48340. Replace the cover and cover screws, and you're set with an expanded system.

Problems? If there are any, they will likely fall

into these categories:

1. PRINT MEM returns 15572 or some figure much below the expected 32K value.
2. The system continuously crashes back to MEMORY SIZE? or just locks up. This happens most often while running any program.
3. The system just locks up with garbage on the screen at power-up.

Well, the last first. Don't forget to hold the (BREAK) key down at power-up in a non-disk system. The other two are tougher.

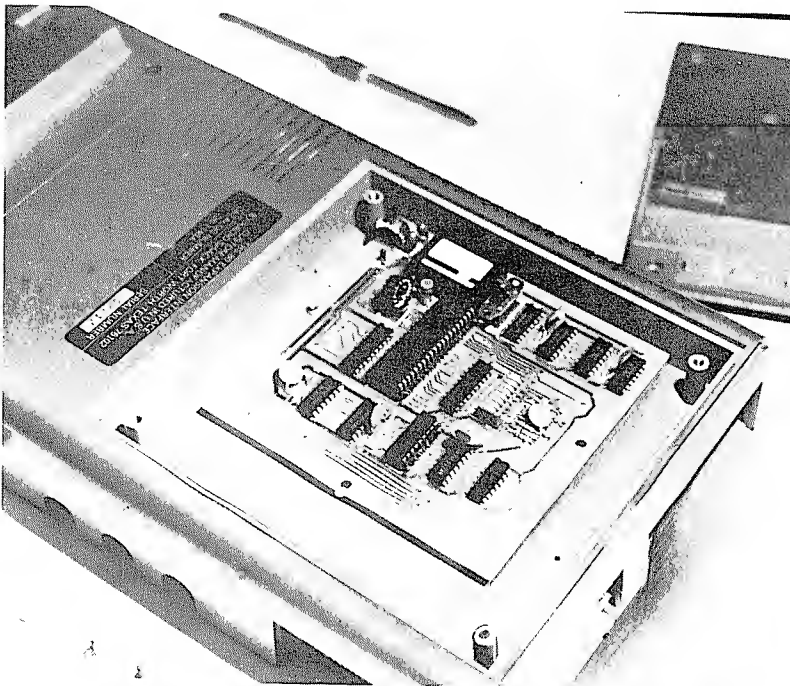
If PRINT MEM returns 15572, you might have installed the memory incorrectly. Keep booting (press the Reset button while holding down (BREAK)), and hitting (ENTER) in response to MEMORY SIZE? PRINT MEM, checking over and over. If a figure higher than 15572 is ever returned, you might have expansion box problems.

Try entering a number between 15580 and 32700 in response to MEMORY SIZE?. If it is ever accepted, chances are you've got the memory chips installed correctly. Try POKEing and PEEKing, like this :

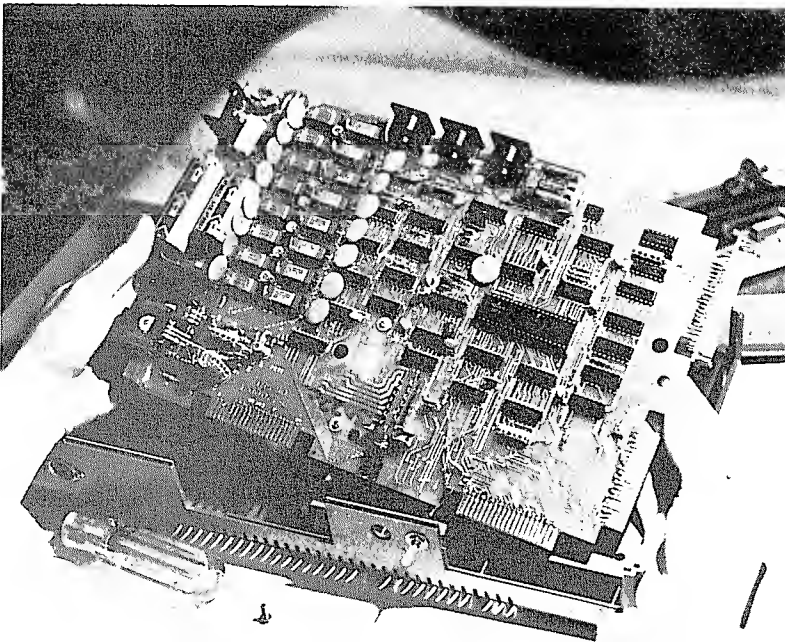
```
POKE 25000,100 : PRINT PEEK(25000)
POKE 25000,103 : PRINT PEEK(25000)
```

and so on. If the value you POKEd is the same as what you PEEK at, the memory is there, but it's not acting reliably. Here are a few solutions (see also the expansion box general solutions above) :

1. Clean cables, connectors, etc., and shock-proof the unit.
2. Remove the buffered cable if you have one, and replace it with a regular one. *Before you do this*, open the expansion box and find pin #1 on the edge connector. Cut that trace, or remove the wire you see there. This is the 5-volt power supply to run the buffering box. The matching pin on the keyboard unit is a ground!
3. If you have a newer expansion box, make the high speed alteration noted above, although you shouldn't be having these problems with the new box.
4. You actually might have some flawed memory chips, though this is unlikely. Switch those in your keyboard unit with those in the expansion box, and see if the problem goes away (or the keyboard unit doesn't work now). Run the memory test (Chapter 3) to find out which chip(s) may be bad.



If RS-232 board is in place, the two screws holding it must also be removed.



Entire expansion card lifts out easily.

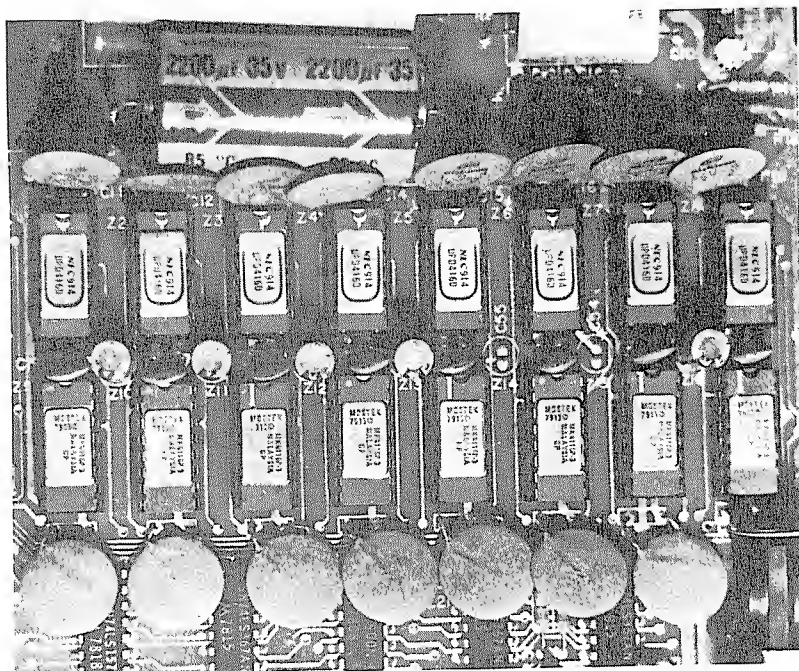


Photo 5-2. Expansion interface memory sockets. 32K memory fits into sockets Z1 through Z16. Oddly, Tandy engineers did not install power supply filter capacitors C54 and C55, nor are they included in the schematic drawing.

5. Make the small change to Z69 in the keyboard unit, as described in the 200% speed-up modification.

The options are few when dealing with an older expansion box. They will work, but you have to provide them with ideal conditions. Which points out the solution to the possible program crashes noted in problem number 3 above. Pamper the expansion interface when a program is running. Don't have the kids playing kick-the-cat near the computer, or the dog jumping about and thumping a large tail next to you, waiting to be fed. Keep the washers, driers, and mixers in the vicinity switched off.

Speeding Up Newer Expansion Boxes

Radio Shack's cures for the memory refresh/select difficulties in the expansion interface created a new problem for those who wish to increase the speed of the computer. All three select lines (RAS, CAS, and MUX) are not used in the revised expansion box; instead only RAS is used.

Technically, the RAS line is sent thrice through an inductance network which provides three differently delayed doppelgangers of the original signal. These are tweaked back into digital shape by a buffer circuit, and the three are

labeled MRAS, MMUX, and MCAS. The play works.

Trouble is, these lines don't care about the speed of the central processing unit. Fast or slow, the false MMUX and MCAS signals trip along at a fixed rate after MRAS. The only cure is to somehow speed up that rate and buy fast, high-quality memory. The speedup is accomplished by *not* delaying the RAS to produce MRAS, and creating MMUX and MCAS in the old MRAS and MMUX positions, respectively. The faster memory (200 nS) is up to you.

To provide the speedup, you'll need to cut some traces and run some wires. Remember, this adjustment is for newer expansion boxes only – the ones that the sales clerk says “don't need the buffered cable”. Here is the order :

1. Find Z37 and Z38.
2. Cut the trace leading from Z38, pin 9. Call the end furthest from Z38 ‘Trace A’.
3. Cut the trace leading from Z38, pin 8. Call the end furthest from Z38 ‘Trace B’.
4. Cut the trace leading from Z37, pin 4. Call the end furthest from Z37 ‘Trace C’.
5. Attach Trace A to Z38, pin 11.
6. Attach Trace B to Z37, pin 4.
7. Attach Trace C to Z38, pin 9.

The modification is complete. Close the box and power up again; the system should work normally (if not better as time goes by). Both higher speed modifications presented in Chapter 4 will work with these changes, so long as your memory is the faster, 200 nS type. If the memory seems at all unreliable, double-check your work. If the work is fine, then your memory may not be up to the change.

LNW, Microtek and Other Expansion Interfaces

The first of the TRS-80 products to be attacked, and rightly so, was the flakey expansion interface. By the time Radio Shack was offering its initial hopeless ‘fixes’, a trio of Californians were redesigning the box for their colleagues in a TRS-80 user group. This was to become the premier challenger to the expansion interface, the LNW System Expansion, priced at a mere \$70 for the bare board, circuit diagrams, parts placement, and a few additions to the Tandy system.

The LNW System Expansion is essentially the Radio Shack expansion interface with a new

Using Those 4K Leftovers

Did you insist on keeping the 4K RAMs which were removed from the keyboard unit when your 16K RAMs were installed by Radio Shack? Or did you do the installation yourself? If so, you've got 4K of dynamic memory worth about \$20 off the shelf. But the trouble is you can't use them in your expansion box, because it's set up for 16K memories only.

But, if you want to spend a bit of time and can't afford an expansion interface, you can extend that PRINT MEM to 19668! Almost

everything is in place already, so all you need is an address decoder and multiplexer circuit.

The circuit presented below can be wire wrapped on a small board, and will provide that extra 4096 bytes. The cost? About \$5 for circuits and sockets, another \$5 for an edge connector (use the *Texas Instruments* wire wrap type sold by *Digi-Key*, part number C6-20), and some time.

When you begin to fill your expansion interface, though, you'll have to give up the 4K RAM extension, because the first block of 16K memory in the box takes those memory addresses.

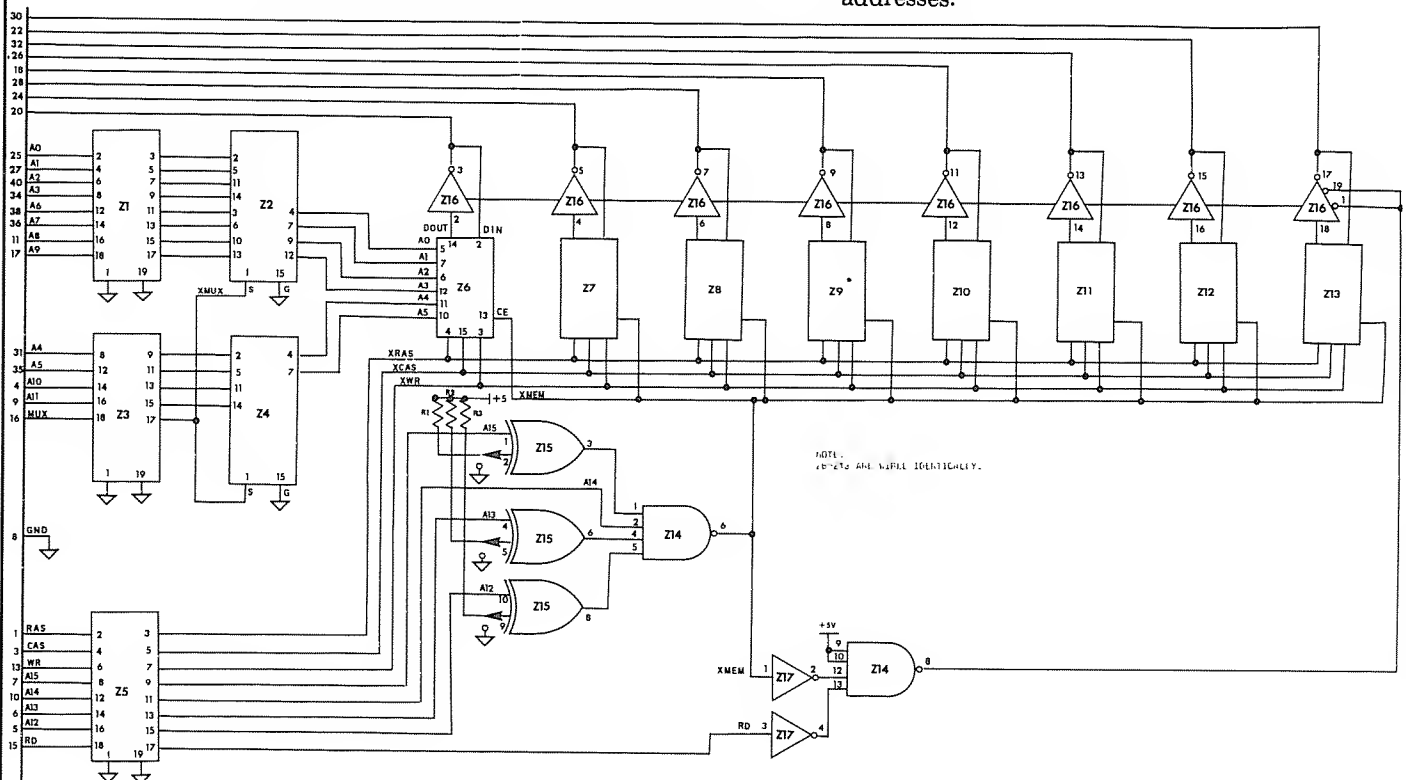
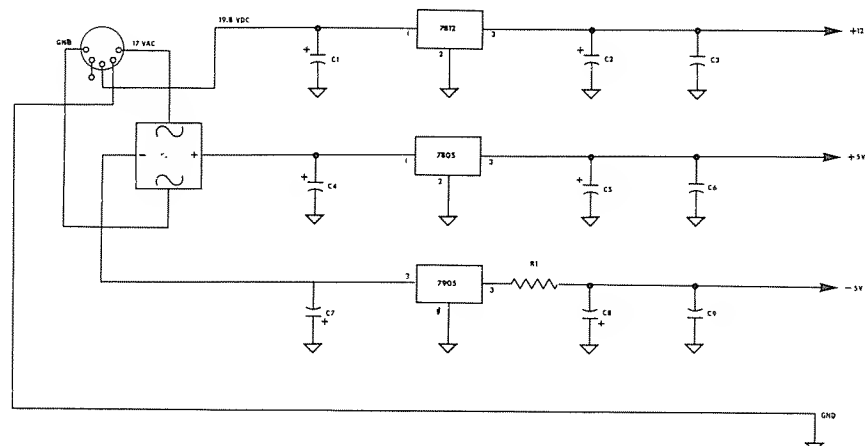
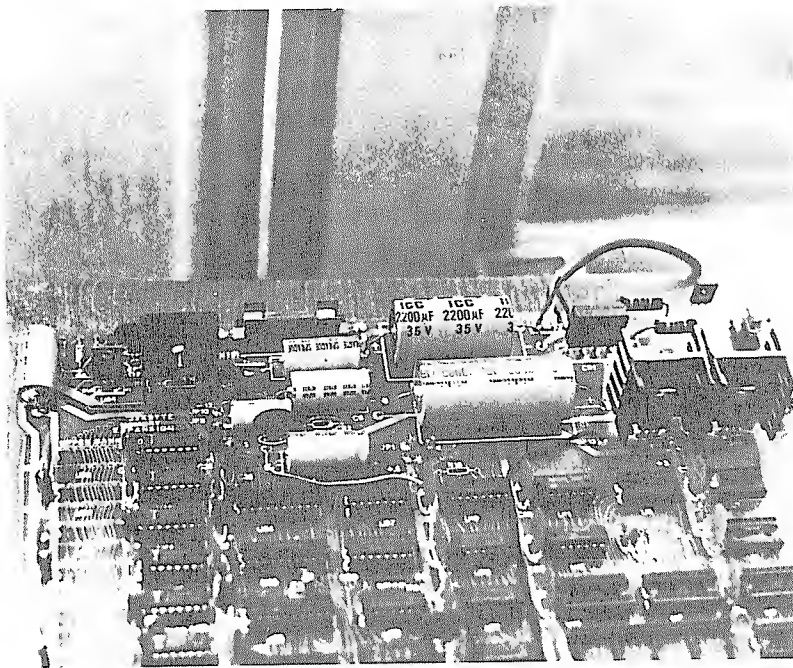


Figure 5-3. 4K dynamic RAM add-on circuit.





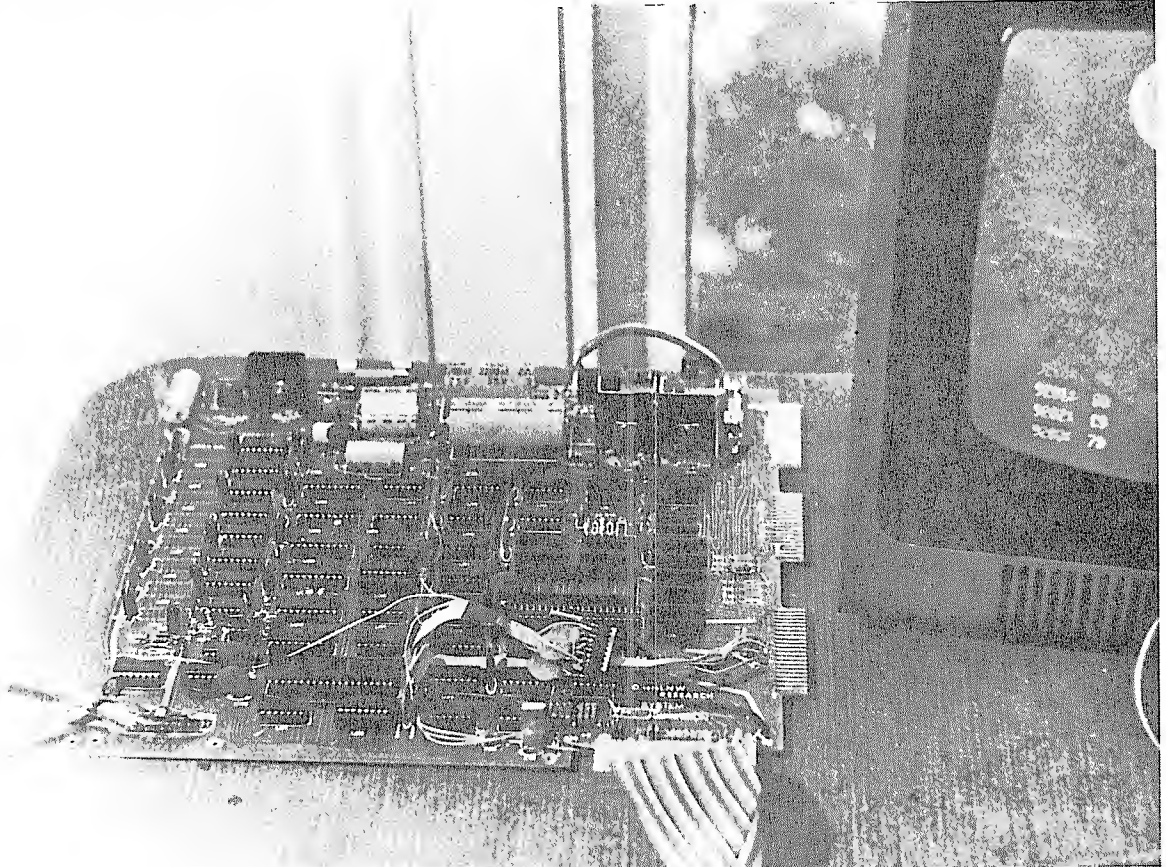
Hefty power supply includes fuse and multiple heat sinks, unlike TRS model with inaccessible fuse and light heat sinks.

layout, on-board buffers (which appeared finally on Radio Shack's own product in the spring of 1980), standard RS-232, and noise-reducing passive bus termination. This last was of particular importance in making the LNW board noise-free and consequently crash-free.

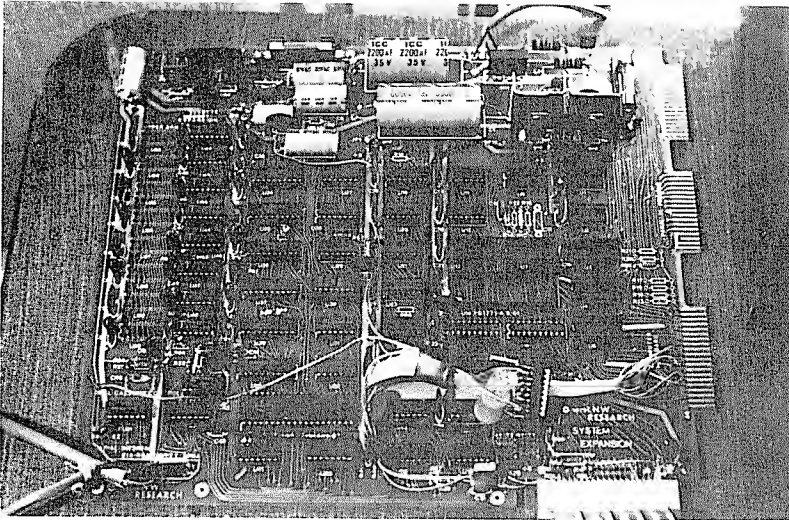
As with the Radio Shack expansion, the LNW board has sockets for 32K additional RAM, but also suggests that sockets be used for all parts – good advice, especially since chips such as the open-collector buffers used in the disk line driver need occasional replacement.

This board is also modular in a way that the complete, boxed interface from Tandy could not be. Only as much as the user needs at one time may be installed. Thus, the \$25 disk controller chip can be left out of the system until (or if) a disk drive is added; so also the disk's peripheral chips can be omitted. Where interrupts are not desired, the entire crystal controlled clock can be excluded. Less than \$150 gets the user a bare-bones 32K memory expansion.

The trick to the LNW board is dexterity. This is not a project for total novices, because it involves extensive soldering and some



LNW Expansion compared with TRS-80 monitor. Board is approximately the same size as standard TRS expansion interface.



Communications are built into LNW expansion; TR1602 UART provides serial capabilities. Note clear marking of all parts, including an industry part number at the UART socket.

knowledge about what the parts look like. There is no Heathkit-style check-off list. You get a parts layout, schematics, and encouragement. But you also get a noise free, reliable expansion interface.

For those who plan to build (or have already built) the LNW board, a few changes might be in order. If you wish your reset button to work, you must not only make a change similar to that for the Radio Shack expansion box, but you must also change the values of the pull-up and pull-down resistors. Change the pull-up resistors on the data lines to 470 ohms, and the pull down resistors to 680 ohms. This will provide the desired 11111111 signal when the reset key is hit and the disk buffer-disable modification has been made.

To make the mod, turn to the expansion interface reset recovery presented in Chapter (?), and follow those same directions, but using the integrated circuit on the LNW board marked U19.

Special Section: Two Other 80s

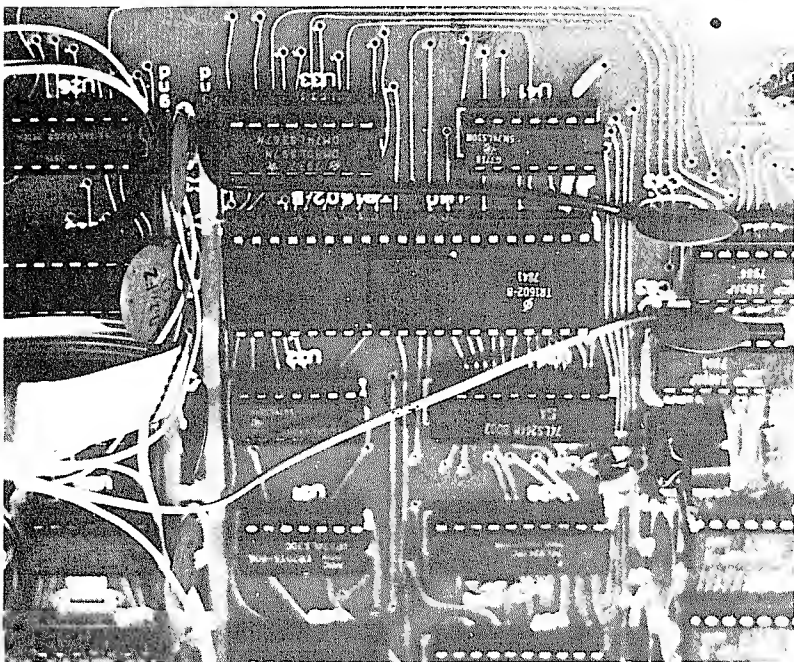
The PMC-80.

As production of the TRS-80 Model I was winding down at Tandy, the rumored 'Hong Kong Copy' made its appearance in the United States, under the name PMC-80, and sold by Personal Microcomputers, Inc., of Mountain View, California.

I was personally ready to greet the PMC-80 with much skepticism, expecting a weaker TRS-80 with little to commend it and a lot to avoid. But, tacky advertising aside, the PMC-80 is a functional, satisfactorily designed product. In fact, because of the experience of the original '80 itself, this copy is probably better designed than Tandy's product. That's probably an indirect compliment to both Fort Worth's foresight and Microsoft's Level II BASIC, which the PMC-80 contains virtually byte for byte.

First, the obvious similarities :

1. Z-80 microprocessor running at 1.77 MHz.
2. 64 x 16 screen resolution, with 128 x 48 graphics.
3. Video monitor output compatible with the Radio Shack monitor.
4. Microsoft Level II BASIC in a 3-chip ROM set and standard type 4116 16K memory.



Virtually complete LNW board in use. User has added programmable baud rate generator (bottom right) in place of hard-wired jumpers recommended by LNW Research.

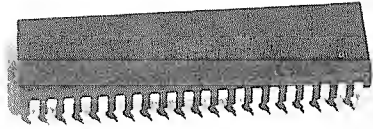
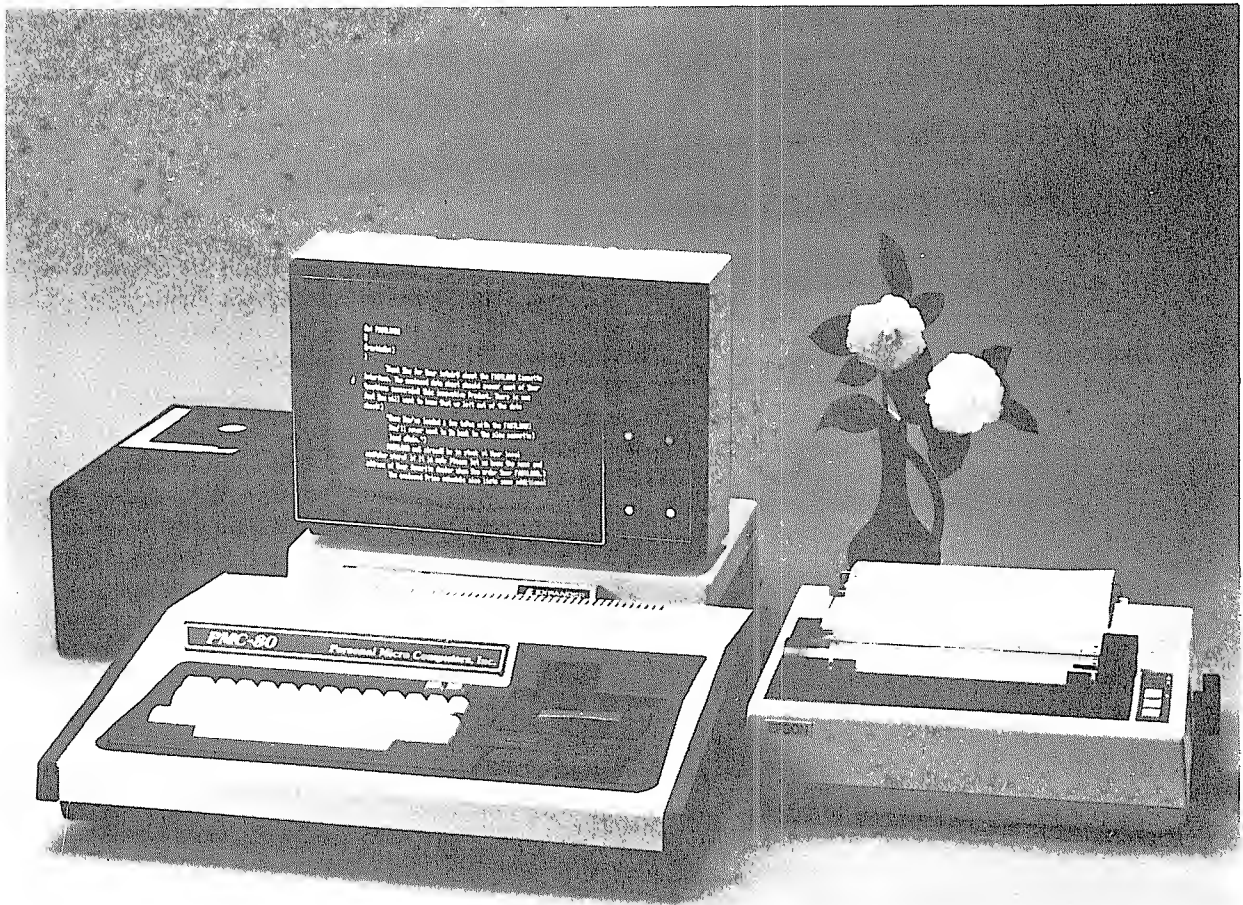


Photo 5-3. The Z-80 CPU — Used in the PMC-80.

And then, the differences :

1. Built-in cassette player with no level control necessary, with provision for two via a second cassette jack already in place.
2. Video RAMs are in sockets, all seven are installed for lower case use, and an industry-standard character generator is used. Lower case is not provided, but is available.
3. An oversized, well-ventilated power supply feeds the system.
4. Channel 3 RF video is provided, along with a connector cable and adaptor.
5. Keyboard has a cassette motor on/off control, but is missing clear and right-arrow keys. These can be installed by the user on some machines.



Complete PMC-80 system can include CPU unit with cassette, printer, expansion box, monitor, and disk drives. Photo courtesy of Personal Micro Computers, Inc.

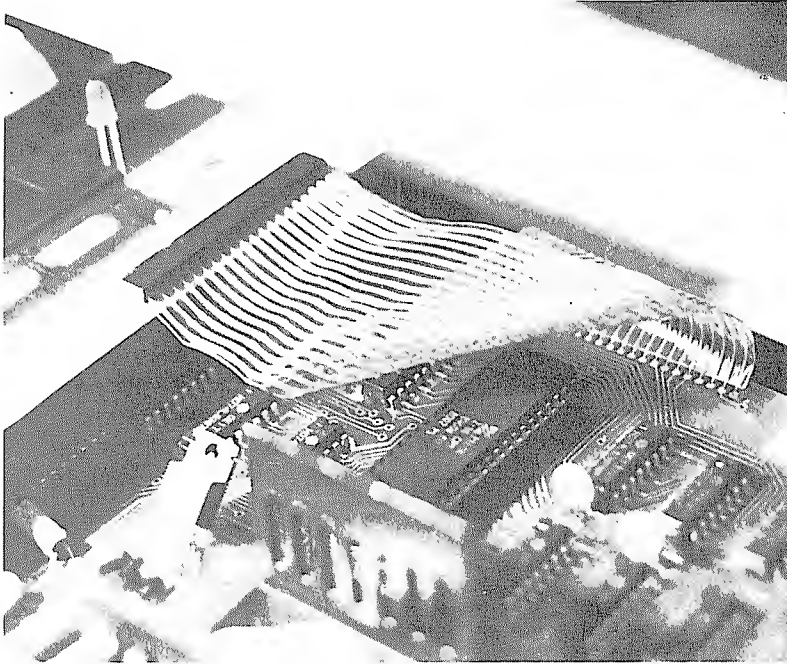
6. An odd 32-character mode, selected differently (with a switch). Instead of alternate letters in video memory, the 'left half' of video memory is used. Normal CHR\$(32) is shown instead as normal-size letters appearing alternately on the screen.
7. The keyboard types easily, with no evidence of keybounce.
8. The expansion connector is gold-plated, but the inter-board and keyboard connectors are just solid wires pushed into receptacles. Heavy duty relays are used for the dual-cassette I/O.
9. The gold-plated expansion connector, alas, is not a TRS-80 compatible 40-pin type. The PMC-to-TRS adaptor is an option.

Are the TRS-80 and PMC-80 software-compatible? From what I can tell after running both BASIC programs and machine language utilities employing calls to ROM: yes, with minor reservations. First of all, the missing clear and right arrow keys are a downright pain; users of programs such as Scripsit or the Penmod version of Electric Pencil are back in the good old days of having to add keys to the system to get these programs to work.

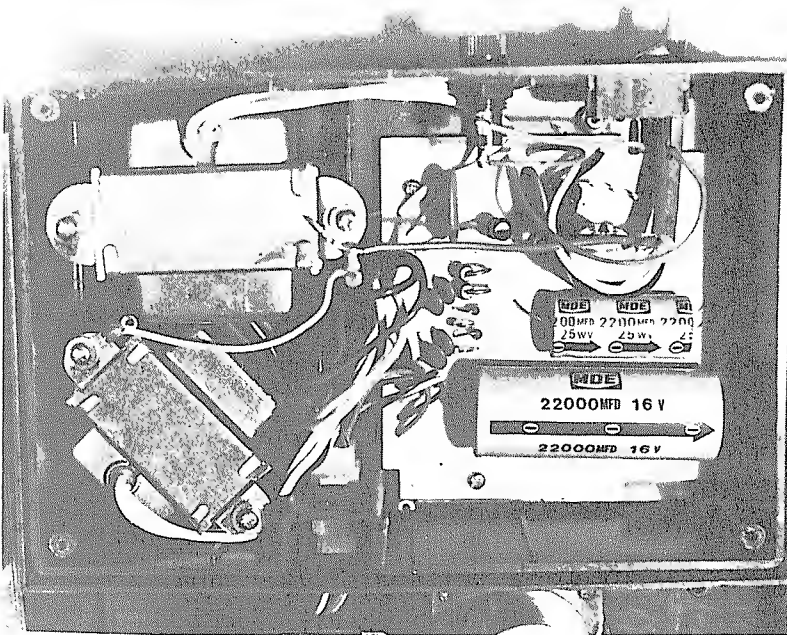
Major ROM calls, though, seem to be intact, so there's little worry about compatibility there. The only programs that seem to get lost are those that call inside routines, or that call partial routines in ROM. BASIC runs flawlessly, as do mixed BASIC/machine programs.



Inside the PMC-80. Keyboard, two electronics boards, cassette player, and power supply are attached to the baseplate. Multiconductor wire jumpers plug in place.



Keyboard connector attaches to main board via a removable cable, a distinct improvement from the permanent TRS-80 cable. Note Z-80 CPU underneath cable.



Hefty power supply provides drive for computer electronics as well as a separate supply for the cassette recorder. Fuse is user-replaceable, in a standard socket.

Other comments :

The video display is cleaner, and the letters are better formed, prettier, and with no touch of blurring, even on the inexpensive TRS-80 monitor. They are quite similar to the better quality characters of the TRS-80 Model III. The lowercase option was not installed in the unit I used, but since it uses the standard character generators and has a clean video output circuit, similar well formed letters can be expected.

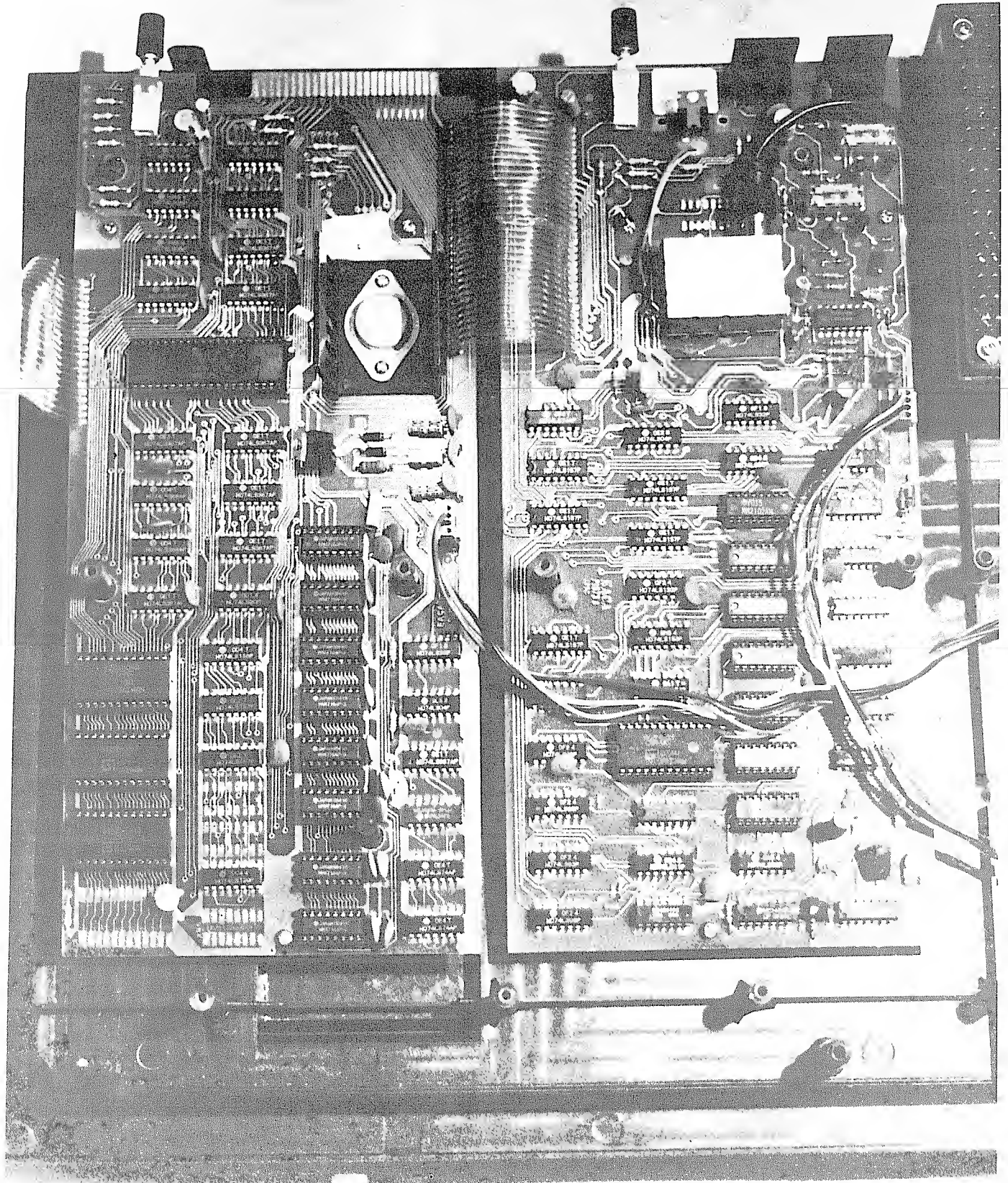
Cassette input/output is sweet. Problem programs loaded easily with this system. The playback of the cassette deck is digitized with LM324 comparators instead of 3900 Norton amps, and so the digitization is more successful.

The oversized power supply shows every sign of providing plenty of surplus current. There is no need for power supply adjustment in this system, because in place of the adjustable 723 regulator used in the TRS-80, fixed 7812 and 7805 regulators are employed.

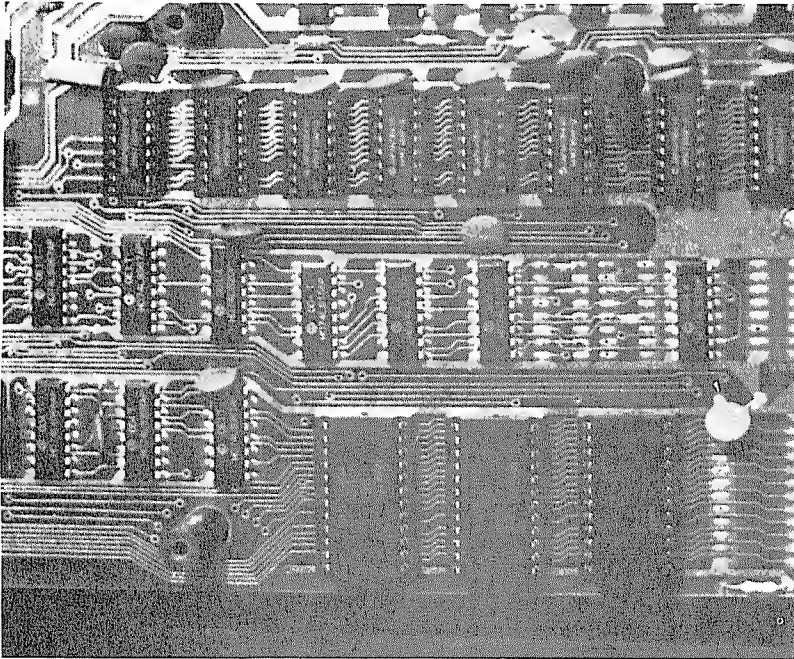
The unit, though heavier than the TRS-80, is more complete in the sense that it has a built-in cassette player and RF video output. For traveling, the PMC-80 makes much more sense, because a video monitor can be used in long sessions at home or office, and an ordinary television can be used at other times. The case is tacky, especially the fake walnut grain plasticoid sides, but then the down-to-business futura TRS-80 was no great aesthetic shakes, either. The keyboard is placed at a comfortable angle, and the added weight and size gives the unit a more responsive typing feel.

The LNW-80 Kit.

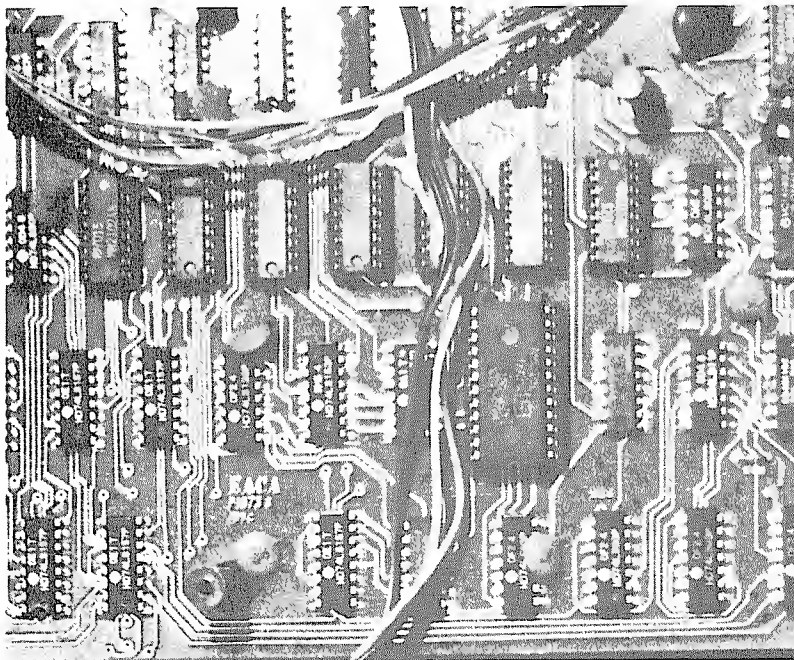
This is a welcome newcomer, providing the advantages of the Model I with LNW's well respected circuit board. As with their system expansion, this LNW board is sold naked. The user contributes the parts as required.



Under the keyboard is the complete electronics control.



CPU control card contains 3-chip Level II ROMs, 16K dynamic memory, and power supply. Circuit board design isn't pretty, but is very reliable.



Video display card includes 7-bit video memory (unlike original TRS-80 6-bit memory), industry-standard CGR-001 character generator, RF modulator, and sockets for video memory chips.

CPU Interruptus

The interrupt is a programmer's mystery tool. On large machines the interrupt controls are so sacred that they are available only to the system engineers. Without a password, assembler commands directed toward interrupts are generally ignored.

TRS-80 users might have in mind some dreams of the Z-80's interrupt power. They've heard of it. Allusions have been made to it. Yet because it operates outside familiar programming bounds, it is even now a mystery.

Furthermore, this may be one of those cases where the Tandy engineers seriously erred in the concept, (rather than in the design, which problems we know all too well). The inclusion of additional interrupt power could have been done for pennies.

Whatsis?

The activities of a computer program can be viewed in relative rather than real time. In other words, a computer performs its tasks in the order that the programmer has chosen, with a timing that is inherent to the oscillations of the computer's master clock. Change the spacing of those oscillations, and the absolute time – the time with respect to 'real' time – changes, yet the relative time within the computer program is still the same.

Unfortunately for timekeeping purposes, changing either the clock of the computer, or changing the way in which the computer handles its tasks, may actually cause a shift in the real time reported to the user. An interrupt, then, is electronic shoulder tapping, the ability of a device external to the microprocessor and its respective memory to demand the CPU's attention for imperative tasks. For example, an external pulsing clock says to the CPU, 'Okay, I don't care what you're doing, stop it now and take care of updating of the real-time clock.'

Such is the simplest use of an interrupt. And, for a change, the computer jargon term is true to its actual purpose. It interrupts the operation of the central processing unit, demanding attention.

An interrupt is akin to a ringing telephone. As long as you are within earshot, the ringing of the telephone interrupts your activity. Whether you choose to respond consciously to it is irrelevant to the actual interruption. If you can hear, and if you are within hearing range of the bell, it does

indeed influence the activities of your nervous system.

How you respond to that ringing telephone – whether you ignore it, whether you go to it and answer it, or whether you merely swear at it – is what might be called, in computer terms, ‘servicing the interrupt’.

The Z-80 microprocessor contains a wide range of ways in which an interrupt may be acknowledged, received, or serviced. You respond to a knock, doorbell, telephone, child’s cry, or gunshot with various levels of urgency. If more than one were to occur about the same time, you would be forced to mask out the one of lower priority and give your attention to the more pressing interruption.

Before describing the ways the Z-80 microprocessor might respond to an interrupt, it should be noted that an interrupt is uniquely different from the computer’s usual input/output schemes, just as the ringing telephone or gunshot are considerably different from checking the color of a traffic signal as you approach it; inspecting the mailbox as you arrive home; or scratching the cat.

In these analogies, your ears most often serve as your human interrupt line. As long as they are ‘in the circuit’, they will interrupt you. You can close your eyes to the traffic signal, ignore the presence of the mailbox by averting your glance, or cold-shoulder the cat (at your own risk), but (lacking earflaps), sound will give you an appropriate kick in the eardrums.

Back to the Z-80’s interrupt schemes. During a movie, were someone to yell “Fire!”, chances are that everyone – barring suicides and existentialists – would respond to that interruption with their personal interrupt service routine. Most likely that routine would be to leave the theater quickly.

An equivalent reaction by the Z-80 is produced by a signal applied to its ‘non-maskable interrupt’ (NMI). This interrupt is always acknowledged. The CPU cannot ignore it. On the TRS-80, it is this interrupt which is activated when the Reset button is pressed.

The NMI is hard-wired into the system in such a way that we have no control over it, so that if these sentences were a program . . .

```
READY
>_
```

. . . and not only that: the execution of a HALT instruction produces a ‘Halt Acknowledge’ output signal which the Tandy Engineers have

OR-gated with the NMI to ...

```
READY
>_
```

. . . yup, you guessed it.

Contrarily, the Z-80 has a maskable interrupt (INT). It is an interrupt line which, when activated, may either be ignored or acknowledged by the processor if the programmer has so specified.

Because their activities are not fixed, these maskable interrupts are the ones of interest to us. Although, as I’ve said, the Z-80 was designed with a wide range of interrupts, that wide range is not accessible to TRS-80 users because the Tandy engineers did not design the TRS-80 in a way that could make those interrupts available.

Nevertheless, here is a look at the Z-80 interrupt modes.

Mode 0. This mode allows the interrupting device to place a machine language instruction directly on the data bus. The CPU will then execute that command rather than the next instruction in the normal program flow. Designers of the Z-80 suggest that a restart (RST) instruction be used.

Restarts have been covered in *The Alternate Source* and other publications, but briefly, they are single byte subroutine calls especially designed for this kind of interrupt operation. Normally, subroutines require specification of a CALL instruction (CD) and two absolute address bytes. The RST instruction is a single byte command causing the CPU to stash the program counter on the stack and move to one of eight locations in low memory.

Mode 1. Upon interrupt, this mode forces the CPU to preserve the program counter and then move directly to location 38.

(line to absorb indent)

Mode 2. This is considered the most powerful interrupt response mode of the Z-80 because “the programmer maintains a table of 16 bit starting addresses for every interrupt service routine . . . When an interrupt is accepted, a 16 bit pointer (is) formed to obtain the desired interrupt service routine address from the table. The upper 8 bits of this pointer (are) formed from the contents of the I register. The I register must have been previously loaded with the desired value by the programmer. The lower eight bits of the pointer must be supplied by the interrupting device....” (from Z-80 Technical Manual).

Before you permit excitement (or even a raised eyebrow) about the last interrupt mode, it should be pointed out that neither Mode 0 nor Mode 2 may be used on the TRS-80, for this hardware reason: in order to place any instruction or address on the data bus, it must be possible to configure the data bus as an input. Within the internal circuitry of the TRS-80 itself, the data bus is fully bidirectional. This is normal, because the data bus is in an *input-to-CPU* state whenever memory is being read or input instructions are being executed, and in an *output-from-CPU* state whenever information is being stored in memory or output instructions are taking place.

Unfortunately, only the data output is brought to the edge card of the TRS-80. There is no buffer wired to the edge card which is pointed *inward* to the CPU, activated when the IORQ signal is sent from the CPU during interrupt! Therefore it is impossible to place either an instruction or an address byte on the data bus from outside the computer.

(Of all the possible under-utilization of the Z-80's power in the TRS-80, this is to me the most discouraging. You can cure something like the stingy lack of an extra video bit for lower case, but adding a second set of buffers requires a true hardware hacker. You might as well build a new computer. *Aveatquevale*, Model I.)

Thus, the balance of this section will be concerned exclusively with interrupt Mode 1 (set by executing the command IM1). As noted, the execution of this instruction forces the CPU to address 38.

Address 38 has also been usurped by the TRS-80 designers by placing it in ROM (read only memory). Because the first act of the instruction at 38 (which incidentally is the last of the RST instructions) is to jump into RAM at 4012, it can nevertheless be accessed easily.

With a Level II CPU unit alone, the edge-card connector is brought out so the interrupt line is present. However, if you are an expansion interface owner or disc user, then you should know that the interrupt itself is dominated by the expansion interface unit's clock, which places a constant 25-millisecond pulse on the interrupt line. Therefore, not only are interrupt modes 0 and 2 made unavailable by the TRS-80 design, but (unless you are willing to cut a trace) the presence of the expansion interface limits even the use of Mode 1.

Despite all these limitations (whew!), the 25 mS interrupt does allow options beyond just the

real-time clock.

Which brings us to the reason the Z-80 has so many interrupt modes. The most powerful kind of interrupt places an instruction or address on the bus, directly influencing CPU activities every time it occurs, but only when it is needed. If the hardware is well designed, this means every interruption is essential, and is handled expeditiously.

On the other hand, 'polling' interrupts for service requests is wasteful. That is, when there is a single interrupt mode which all devices must use, then the CPU is forced to ask each one, 'Did you interrupt me? No? Well, did you interrupt me? No? Well then, did you interrupt me? No? Okay, then . . .', until it finds the culprit. It not only services it, but also checks the remainder of the devices in case other interrupts occurred simultaneously with the first, or during the interrupt service routine.

We are therefore limited not only to a single interrupt mode, but also to a required polling technique, because this 40-times-per-second pulse is coming through without cease. No real 'device' is asking for service, but the effete request is present nonetheless. Yet we can still use it. Here's how: when the interrupt comes through, the CPU sends a signal to the relatively few devices in our I/O scheme, and asks them, 'do you have any information?' If there's no information, the CPU will quickly beat a path back to normal program flow. If there is information, the CPU will service the interrupt.

Whysis?

You may wonder what type of situation — within the context of the humble TRS-80 itself — might require this interrupt, other than the familiar real-time clock. Let's take an example of flawed human interaction in TRS-80 programming: the *Electric Pencil*.

As any reasonably good typist knows, this 'text editing' program is infuriatingly slow at the end of lines, particularly at the bottom of the page. Not only must it move the word in progress to the next line, but also scroll all 1,024 characters up the screen. If you are nimble-fingered, you can lose characters during this process.

The interrupt system might be used in order to build a buffer of keys pressed. A good typist's fastest keystrokes ('the', 'is', etc.) can be less than 50 mS apart, or 20 characters per second. As the interrupt comes through from the expansion box 40 times each second, location 387F can be

examined by the CPU. If any key other than 'shift' is pressed, its presence can be sensed by looking at this location. If its presence is sensed, the key-scan routine can be initiated, and the key information and any successive keystrokes are stored until the line break and/or scrolling is completed.

In this way the keyboard may be 'scanned' quickly, not only during entry of text, but also during I/O, such as to tape. A keystroke buffer may be built during this process (since the keyscan need only be initiated if a key is pressed) between output bytes. The same is true of printing; the text can continue to be output to the printer while this activity is going on.

In fact, if a key were pressed during an I/O routine, it could be accepted into a buffer, and subsequently displayed when the next interrupt occurs. In effect, you would be time-sharing your TRS-80.

Light Bulbs

Aha! Time-sharing your TRS-80! (Light bulbs flicker on in millions of brain rooms). Although a 25-mS pulse does not allow totally transparent time-sharing, it allows enough of it for your TRS-80 to be able to service several user requests at a time, in different ways.

But first, here's a short routine to prove that the interrupt is really there. You can assemble this, or just POKE it into memory.

7F00 F3	DI		;Interrupts off
7F01 3E C3	LD	A,DC3H	;Get JP command
7FD3 32 12 40	LD	(4012H),A	;Replece RET cmd
7FD6 21 14 7F	LD	HL,7F14H	;Stert service
7FD9 22 13 40	LD	(4013H),HL	;Plece after JP
7F0C 21 19 1A	LD	HL,1A19H	;Command level
7F0F E5	PUSH	HL	;Reedy to return
7F10 E0 56	IM	1	;Set to Mode 1
7F12 FB	EI		;Interrupts on
7F13 C9	RET		;Return to 1A19

The segment above is merely a setup routine, executed once. It disables interrupts during setup, places C3 (a jump instruction) in place of C9 (return) currently found at 4012, the interrupt patch point in RAM. It then places the destination address (7F14) as the jump's operand. It pushes the return to BASIC or DOS onto the stack, sets interrupt Mode 1, enables interrupts, and returns (effectively jumping to 1A19).

Note that after an EI (Enable Interrupts) instruction, acknowledgment of any interrupts is actually delayed until *after* the completion of the instruction following the interrupt enable. Here's the interrupt service routine itself:

7F14 F3	DI		;Interrupts off
7F15 F5	PUSH	AF	;All of these
7F16 E5	PUSH	HL	; registers
7F17 D5	PUSH	DE	; will be used
7F18 C5	PUSH	BC	; to do this
7F19 3A EC 37	LD	A,(37ECH)	;Read disc chip
7F1C 3A E0 37	LD	A,(37E0H)	;Reset clock F/F
7F1F 21 11 01	LD	HL,0111H	;Messege start
7F22 11 25 3C	LD	DE,3C25H	;Display space
7F25 D1 1A DD	LD	BC,DD1AH	;Messege length
7F28 E0 80	LDIR		;Display it
7F2A C1	POP	BC	;Put ell these
7F2B 01	POP	DE	; registers
7F2C E1	POP	HL	; back in
7F2D F1	POP	AF	; place
7F2E FB	EI		;Interrupts on
7F2F C9	RET		;Return whence

The instructions loading the accumulator from 37EC and 37E0 accommodate the quirks of the TRS-80 hardware — the Z-80's interrupt acknowledge signal is not used in this hardware scheme. Instead, reading 37EC (disk controller status) followed by reading 37E0 (real-time clock) resets a flip-flop used to signal input from disk controller and real-time clock, respectively. This must be done each time or the interrupt will continue to be 'on', forcing the routine ever back upon itself.

You may protect memory at 32512 if you wish, as this will prevent any accidental crashes while you're trying this demo program. Also note that if you have installed a disk-disable switch on your interface, this program may initially lock up. The lockup may be cured by momentarily flipping the disable switch on then off.

So what's the point of this? Only to show that no matter what BASIC (or machine language) program you are operating, including DOS, the sign-on message will remain in the upper-right corner of the screen. CLS, scrolling, even attempting to overwrite the area with text will have no apparent effect, because the interrupt takes priority. To disable it, put a RETURN instruction in place -

PDKE 16402,2D1

and to enable it, replace the return with a JUMP:

POKE 16402,195

Under DOS, the usual CMD"R" and CMD"T" instructions will have their intended effect. Re-initializing Level II BASIC or DOS from 0000 will wipe out the interrupt patch to 7F14.

On to Something Useful

There are a few rules of thumb for using interrupts. First, the interrogation and service routines must be as short as possible, or else the main program will get little work done between interrupts. This means, alas, very few calls to

ROM; the ROM calls, although effective, often contain a level of detail and error-checking that may be beyond the absolute minimum needs of the service routine.

Second, every register to be used by the service routine must be preserved, usually on the stack, since there is no way of knowing which registers will be in action when an interrupt occurs. The only safe bet for ordinary BASIC programs using Level II alone is to alternate register pairs (by executing EX AF,AF' and EXX), since the alternates are not used in Level II. IX, however, must still be saved.

With these thoughts in mind, have a look at the listing below, which prints a more useful message onto the screen: the BASIC memory available at all times.

Listing 5-3. Free memory constant display.

```

7F00 F3      SETUP  OI      ;
7F01A3E C3    LO      A,0C3H ;
7F03 32 12 40  LO      [4012H],A ; Setup
7F06 21 14 7F  LO      HL,SERVE ; routine
7F09 22 13 40  LO      [4013H],HL ; similar
7F0C 21 19 1A  LO      HL,1A19H ; to that
7F0F E5      PUSH  HL      ; presented
7F10 E0 56    IM      1      ; above
7F12 FB      EI          ;
7F13 C9      RET          ;
3C00      VIOE0 EQU      3C00H
40EB      STKB0T EQU     40EBH
40FB      ARRTOP EQU     40FBH
7F14 F3      SERVE  DI          ; No interrupt
7F15 F5      PUSH  AF          ;
7F16 E5      PUSH  HL          ;
7F17 05      PUSH  DE          ; Save all
7F18 C5      PUSH  BC          ; registers to
7F19 00 E5    PUSH  IX          ; be used
7F1B FD E5    PUSH  IY          ;
7F10 3A EC 37 LO      A,(37ECH) ; Clear FDC
7F20 3A E0 37 LO      A,(37E0H) ; Clear F-F
7F23 21 69 7F LO      HL,MESSG ; Get "MEM="
7F26 11 36 3C LO      DE,VIOE0+36H ; Place to go
7F29 01 05 00 LO      BC,5      ; Its length
7F2C E0 80    LOIR         ; Display it
7F2E AF      XOR      A          ; Clear carry
7F2F 2A EB 40 LO      HL,(STKB0T) ; BASIC stack
7F32 E0 5B FB 40 LO      DE,(ARRTOP) ; Array aree
7F36 E0 52    SBC      HL,DE      ; Room left
7F3B 00 21 3B 3C LO      IX,VIOE0+3BH ; Place to go
7F3C F0 21 6E 7F LO      IY,TENTAB ; Tens-table
7F40 06 05    LO      B,5        ; No. digits
7F42 AF      XOR      A          ; Clear carry
7F43 F0 5E 00 LO      E,(IY+0)   ; 10000 lobyte
7F46 F0 56 01 LO      0,(IY+1)   ; 10000 hobyte
7F49 B7      LOOP1  OR      A      ; Set flag
7F4A E0 52    SBC      HL,DE      ; 1st diffnce.
7F4C 3B 03    JR      C,JUMPO     ; If greater
7F4E 3C      INC      A          ; Additive div
7F4F 1B FB    JR      L00P1      ; Divide again
7F51 19      JUMPO  A00      HL,DE ; Was too much
7F52 C6 30    A00      A,30H     ; ASCII cnvrt.
7F54 00 77 00 LO      [IX+0],A   ; Display it
7F57 00 23    INC      IX        ; Next screen
7F59 F0 23    INC      IY        ; Move once,
7F5B FD 23    INC      IY        ; and again
7F5D 10 E3    OJNZ     L00P0     ; Oo 5 digits
7F5F F0 E1    POP      IY        ;
7F61 00 E1    POP      IX        ; Restore all
7F63 C1      POP      BC        ; the reg's
7F64 01      POP      DE        ; used to do
7F65 E1      POP      HL        ; the work
7F66 F1      POP      AF        ;
7F67 FB      EI          ;
7F68 C9      RET          ; Inte. okay
7F69 BF      MESSG  OEFB  BFH    ; BASIC awaits
7F6A      OEFM  'MEM=' ; Block
7F6B 40      7F6B 45 7F6C 40 7F6D 30
7F6E 10 27    TENTAB OEFW 100000 ; Tens-table
7F70 EB 03    OEFW 010000 ; is used for
7F72 64 00    OEFW 001000 ; division by
7F74 0A 00    OEFW 000100 ; subtraction
7F76 01 00    OEFW 000010 ;
7F00      END      SETUP

```

The 19-byte interrupt setup routine is found as usual, followed by the 100-byte service subroutine. All registers are saved — IY is optional but may be used by other utilities — a message is displayed, and a calculation is made to determine free memory. This calculation is the bottom of the BASIC subroutine stack (referenced at 40E8) minus the top of array space (referenced at 40FB).

Following that is a short, five-digit binary-to-ASCII table lookup which, in 35 bytes, converts and displays the free memory.

This routine is especially useful because it can help identify the actual area causing out-of-memory errors at run time. Heavily nested calculations can push the BASIC stack well down into memory, which is then visible on the screen. Put the program in place, and then try this :

```

10 GOSUB 20
20 GOTO 10

```

Something Even More Useful

If you by chance have a printer which prints less than 40 characters per second, the following routine will speed up your BASIC programs. Otherwise, this routine will have little effect on your program speed (perhaps a two or three percent slowdown), but will allow a program to continue without 'locking up' while the printer is in action. It is called a 'spooler', a word which is based on an acronym mercifully forgotten.

When an LPRINT or LLIST is commanded, the spooler sends the characters to an in-memory buffer, and returns to the waiting program. Since 255 characters is the maximum string length, this buffer size is adequate; increasing it will be valuable only if the main program is highly interactive, involving user, screen, calculations and printer.

Normally, the LPRINT and LLIST routines send a character to the printer, and then check to see if the printer is 'busy'. In the meantime, the line printer routine is idle, performing no other work, and ignoring the BASIC program. Those of you with Teletypes know the discouraging feeling of waiting for hard copy.

Instead of waiting for the printer to output the entire number, word, or line before returning to the program, an interrupt service routine is used as a 'despooler'. When the 25-mS interrupt is generated, the despooler checks the 255-byte buffer to see if a printable character is present. If there is, the routine sends that character to the printer, advances the buffer, and returns to the program in progress. Also, if the printer is busy,

Listing 5-4. Print spooler.

```

00100 ; #####
00110 ; THIS INTERRUPT SERVICE ROUTINE STORES UP TO 255 CHAR-
00120 ; ACTERS IN A BUFFER, WHERE THEY MAY THEN BE OUTPUT TO
00130 ; A PRINTER. THIS SPOOLER CHECKS FOR CHARACTERS AT THE
00140 ; PRINTER DRIVER ROUTINE, INTERCEPTS THEM, STASHES THEM,
00150 ; AND RETURNS TO THE MAIN PROGRAM WITHOUT DELAY.
00160 ; #####
00170 ;
7E00 00180 ORG 7E00H
00190 ;
00200 ; #####
00210 ; SETUP ROUTINE CREATES A BUFFER AND PLACES ITS ADDRESS,
00220 ; ALONG WITH A JUMP TO REPLACE THE NORMAL RET, AT 4012H,
00230 ; #####
00240 ;
7E00 F3 00250 SETUP OI ; INT. OFF DURING SETUP
7E01 3EC3 00260 LD A,0C3H ; GET JUMP VALUE INTO A
7E03 321240 00270 LD (4012H),A ; PLACE INTO INT. VECTOR
7E06 AF 00280 XOR A ; CLEAR ACCUMULATOR
7E07 47 00290 LD B,A ; PUT 100 HEX INTO B
7E08 218A7E 00300 LD HL,BUFFER-1 ; HL JUST AHEAD OF BUFFER
7E0B 23 00310 CLEAR INC HL ; NEXT POSITION IN BUFFER
7E0C 77 00320 LD (HL),A ; PLACE ZERO INTO BUFFER
7E0D 10FC 00330 DJNZ CLEAR ; DO IT FULL 256 TIMES
7E0F 213F7E 00340 LD HL,SERVE ; GET SERVICE ROUT. ADDR.
7E12 221340 00350 LD (4013H),HL ; PLACE INTO INT. VECTOR
7E15 21237E 00360 LD HL,SPOOL ; GET SPOOL ROUTINE
7E1B 222640 00370 LD (4026H),HL ; PLACE IN PRINTER DRIVER
7E1B 21191A 00380 LD HL,1A19H ; RETURN TO BASIC INTO HL
7E1E E5 00390 PUSH HL ; PLACE ON STACK
7E1F E056 00400 IM 1 ; SET INTERRUPT MODE
7E21 FB 00410 EI ; INTERRUPTS READY TO GO
7E22 C9 00420 RET ; BACK TO BASIC READY
00430 ;
00440 ; #####
00450 ; SPOOL ROUTINE STARTS HERE AND INTERCEPTS PRINTER DRIVER
00460 ; #####
00470 ;
7E23 F3 00480 SPOOL OI ; INT. OFF DURING SPOOL
7E24 E5 00490 PUSH HL ; SAVE HL REGISTER
7E25 79 00500 LD A,C ; GET CHARACTER INTO A
7E26 F5 00510 PUSH AF ; SAVE CHAR. TO PRINT
7E27 21BC7E 00520 LD HL,BUFFER+1 ; POINT TO BUFFER START
7E2A FB 00530 EI ; INTERRUPTS BACK ON NOW
7E2B 7E 00540 LOOP2 LD A,(HL) ; GET BUFFER VALUE
7E2C A7 00550 AND A ; IS IT A CHAR. OR ZERO?
7E2D 20FC 00560 JR NZ,LOOP2 ; WAIT IN LOOP IF FULL
7E2F F3 00570 OI ; INTERRUPTS BACK OFF
7E30 06F0 00580 LD B,0F0H ; GET PRESENT BUFFER SIZE
7E32 23 00590 LOOP3 INC HL ; GET NEXT BUFFER POS'N
7E33 7E 00600 LD A,(HL) ; BRING VALUE INTO A
7E34 A7 00610 AND A ; TEST FOR CHAR. OR ZERO
7E35 2002 00620 JR NZ,SAVEIT ; FOUND FREE SPACE - GO
7E37 10F9 00630 DJNZ LOOP3 ; SEARCH THROUGH BUFFER
7E39 2B 00640 SAVEIT DEC HL ; BACK OFF ONE POSITION
7E3A F1 00650 POP AF ; RESTORE CHAR. TO PRINT
7E3B 77 00660 LD (HL),A ; PUT IT IN BUFFER
7E3C E1 00670 POP HL ; RESTORE FORMER VALUE
7E3D FB 00680 EI ; INTERRUPTS BACK ON
7E3E C9 00690 RET ; BACK TO MAIN ROUTINE
00700 ;
00710 ; #####
00720 ; INTERRUPT SERVICE ROUTINE FIRST SAVES REGISTERS, THEN
00730 ; RESETS INTERRUPT FLIP-FLOPS IN E/I. PRINTER STATUS
00740 ; IS EXAMINED, AND THE ROUTINE EXITED IF PRINTER BUSY.
00750 ; #####
00760 ;
7E3F F3 00770 SERVE OI ; INT. OFF DURING DESPOOL
7E40 F5 00780 PUSH AF ; SAVE VALUE IN ACCUM.
7E41 E5 00790 PUSH HL ; SAVE VALUE IN HL
7E42 05 00800 PUSH DE ; SAVE VALUE IN DE
7E43 C5 00810 PUSH BC ; SAVE VALUE IN BC
7E44 00E5 00820 PUSH IX ; SAVE VALUE IN IX
7E46 00212540 00830 LD IX,4025H ; PRINTER CONTROL BLOCK
7E4A 3AEC37 00840 LD A,(37ECH) ; RESET INT. FLIP-FLOP
7E4D 3AE037 00850 LD A,(37E0H) ; RESET INT. FLIP-FLOP
7E50 3AEB37 00860 LD A,(37EBH) ; GET PRINTER STATUS TO A
7E53 E6F0 00870 AND 0F0H ; MASK OUT LOW BITS
7E55 FE30 00880 CP 30H ; SEE IF PRINTER IS BUSY
7E57 205A 00890 JR NZ,OUT ; GO OUT IF PRINTER BUSY
00900 ;
00910 ; #####
00920 ; WHEN PRINTER IS READY, BUFFER IS MOVED UP, THE CHAR-
00930 ; ACTER IS TESTED, AND PRINTED IF A VALID CHARACTER.
00940 ; IF IT IS CARRIAGE RETURN, FORM FEED, LINE FEED, ETC..
00950 ; APPROPRIATE TESTS ARE MADE IN THE PRINTER CONTROL BLOCK
00960 ; #####
00970 ;
7E59 21897F 00980 LD HL,BUFFER+0FEH ; GET NEXT TO LAST CHAR.
7E5C 118A7F 00990 LD DE,BUFFER+OFFH ; GET NEXT CHAR. IN QUEUE
7E5F 01FF00 01000 LD BC,OFFH ; GET TOTAL BUFFER SIZE
7E62 E0B8 01010 LDOR ; MOVE IT UP ONE POS'N

```

Listing Continued.

it returns to the program immediately.

A ten-character-per-second (cps) printer like a Teletype is 'busy' for 100,000 microseconds for each character. In that time, a BASIC program might perform many calculations. Selectrics type a maximum of 15 cps, which is also time-consuming.

The program will not necessarily benefit in time with faster printers, but the waiting period will be eliminated in favor of a more interactive program overall.

The following listing is uncommented, as the bulk of the interrupt explanations have already been presented. The spooler is entered from BASIC via LLIST and LPRINT, either of which moves directly to the printer driver address referenced at 4026 and 4027. This is replaced with the address of the spooler. The character to be printed is delivered to the printer driver in the C register. The spooler searches through the buffer for the first non-zero value and places the character to be printed immediately before that. (If your system requires nulls after a carriage return, you can back up the appropriate number of places in the buffer).

Upon interrupt, the despooler saves registers and loads the IX register with 4025, the start of the printer device control block. This block contains line and page information, which is used to determine if paging has been completed, and how line feeds, carriage returns, and form feeds will be handled. It is functionally identical to the driver in ROM, except that it takes its characters from the spooler's buffer instead of directly from the Level II routines.

If the printer is not busy, then, the character is sent to the printer address at 37E8. Following is the complete routine, including setup, spool, and despool.

All of the interrupt programs presented in this article depend on the expansion interface for their 25 mS clock pulses. However, a simple rectified and shaped 16.66 mS pulse, which will work fine with all of these routines, can be created using the 60 Hz output of a 6.3-volt, low-current filament transformer, resistor, and 7414 Schmitt trigger.

The complete circuit for a printer decoder is available from Radio Shack as 'Printer Cable Service Manual'.

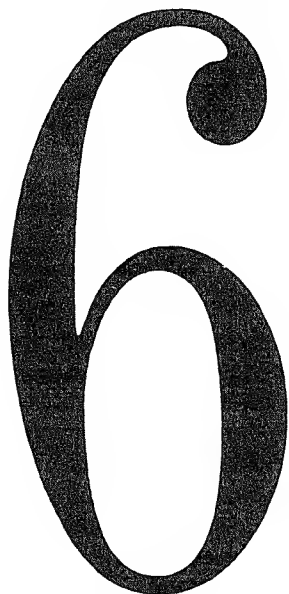
Continued Listing

```

7E64 3ABA7F 01020      LO      A,(BUFFER+OFFH) ; GET QUEUE CHAR TO PRINT
7E67 A7 01030      ANO      A      ; TEST TO SEE IF ZERO
7E6B 2B49 01040      JR      Z,OUT      ; GO OUT IF NONE LEFT
7E6A FE0B 01050      CP      0BH      ; CHECK IF TOP OF FORM
7E6C 2B0A 01060      JR      Z,TEST1    ; TEST IF CHAR. IS ONE
7E6E FE0C 01070      CP      0CH      ; TEST IF CHAR = F.F.
7E70 201F 01080      JR      NZ,TEST2   ; NEXT TEST IF NOT
7E72 AF 01090      XOR      A      ; CLEAR ACCUMULATOR
7E73 00B603 01100     OR      (IX+3)     ; GET LINES PRINTED
7E76 2B19 01110      JR      Z,TEST2     ; OUT IF VALUE = 0
7E7B 007E03 01120     TEST1  LO      A,(IX+3) ; GET LINES PRINTED
7E78 009604 01130     SUB      (IX+4)     ; GET LINES PER PAGE
7E7E 47 01140      LD      B,A      ; PUT LINES LEFT IN B
7E7F 3AEB37 01150     LOOPA  LD      A,(37EBH) ; GET PRINTER STATUS
7E82 E6F0 01160     ANO      0F0H      ; MASK OUT LOW BITS
7E84 FE30 01170      CP      30H      ; CHECK STATUS BITS
7E86 20F7 01180      JR      NZ,LOOPA    ; LOOP IF PRINTER BUSY
7E8B 3E0A 01190      LO      A,0AH      ; GET LINE FEED CHAR.
7E8A 32EB37 01200     LO      (37EBH),A  ; SEND IT TO PRINTER
7E80 10F0 01210     DJNZ    LOOPA      ; DO IT TILL FORM IS FED
7E8F 1B1E 01220      JR      EXIT      ; GO OUT WHEN DONE
7E91 47 01230     TEST2  LD      B,A      ; STASH CHAR. IN B REG.
7E92 3AEB37 01240     LDOPB  LO      A,(37EBH) ; GET PRINTER STATUS
7E95 E6F0 01250     ANO      0F0H      ; MASK OUT LOW BITS
7E97 FE30 01260      CP      30H      ; CHECK STATUS BITS
7E99 20F7 01270      JR      NZ,LDOPB    ; LOOP TILL PRINTER READY
7E9B 7B 01280      LO      A,B      ; GET CHAR. BACK INTO A
7E9C 32EB37 01290     LO      (37EBH),A  ; SEND IT TO PRINTER
7E9F FE00 01300      CP      00H      ; CHECK IF CARRIAGE RET.
7EA1 2010 01310      JR      NZ,OUT      ; IF NOT THEN GO OUT
7EA3 003404 01320     INC      (IX+4)     ; ELSE INC. LINES COUNTER
7EA6 007E04 01330     LO      A,(IX+4)   ; GET NEW LINES COUNTER
7EA9 00BE03 01340     CP      (IX+3)     ; CHECK WITH LINES/PAGE
7EAC 79 01350      LO      A,C      ; GET CHARACTER BACK TO A
7EAD 2004 01360      JR      NZ,OUT      ; GO IF NOT PAGE END
7EAF 00360400 01370     EXIT  LO      (IX+4),0 ; RESET PAGE LINES TO 0
      01380 ;
      01390 ; #####
      01400 ; REGISTERS ARE RESTORED AND THE ROUTINE IS EXISTED. THE
      01410 ; LAST 255 BYTES OF SPACE ARE RESERVED FOR PRINTER BUFFER
      01420 ; #####
      01430 ;
7EB3 00E1 01440     OUT      POP      IX      ; RESTORE IX REGISTER
7EB5 C1 01450      POP      BC      ; RESTORE BC REGISTER
7EB6 01 01460      POP      DE      ; RESTORE DE REGISTER
7EB7 E1 01470      POP      HL      ; RESTORE HL REGISTER
7EBB F1 01480      POP      AF      ; RESTORE AF REGISTER
7EB9 F8 01490      EI      ; INTERRUPTS BACK ON
7EBA C9 01500      RET      ; BACK TO MAIN ROUTINE
00FF 01510     BUFFER  OEFS      255      ; OEFINE 255-CHAR. BUFFER
7FBA 00 01520      OEFB      00      ; ENO BUFFER WITH 0 BYTE
      01530 ;
      01540 ; #####
7E00 01550      ENO      SETUP
00000 TOTAL ERRORS
29420 TEXT AREA BYTES LEFT

BUFFER 7EBB 01510 00300 00520 009B0 00990 01020
CLEAR 7E0B 00310 00330
EXIT 7EAF 01370 01220
LOOP2 7E2B 00540 00560
LOOP3 7E32 00590 00630
LOOPA 7E7F 01150 01180 01210
LOOPB 7E92 01240 01270
OUT 7EB3 01440 00B90 01040 01310 01360
SAVEIT 7E39 00640 00620
SERVE 7E3F 00770 00340
SETUP 7E00 00250 01550
SP00L 7E23 00480 00360
TEST1 7E7B 01120 01060
TEST2 7E91 01230 010B0 01110

```

More Hardware Modifications

By now your TRS-80 is an electronic intimate. In this chapter, some significant (and some sophisticated) hardware modifications will be made. A few of these are simple, and their value is not immediately obvious. Others involve major work (such as the high-resolution graphics), but the results are exciting and very rewarding.

Making Halt Work

Let's start with an easy one. The HALT instruction is a useful command. It isn't exceptional, but you will find that time spent idling in loops waiting for an interrupt is easier when the HALT instruction is used. In fact, its main advantage comes when the Z-80 microprocessor is used in interrupt-based systems. Where program work is being done, the HALT instruction is not particularly meaningful; but where program operations are suspended except for interrupts, it is very useful.

When HALT is executed on the Z-80, the processor simply ceases program counter operation and executes NOPs in order to continue memory refresh. That is, the Z-80 outputs a fetch (the M1 or machine cycle one signal) which – together with a refresh address on the address bus – constitutes enough information to keep the memory active. The information received during the fetch is ignored by the CPU, which continues to execute NOPs until an

interrupt is received. It then exits the NOP state and services the interrupt.

Whenever the HALT instruction is executed, a halt-acknowledge signal (HALT, active low) is output on pin 18. In the TRS-80, pin 18 is tied to an input of NAND gate Z53. The other input of Z53 is tied high through a 10K resistor, but can be pulled low when the Reset button is pressed. In other words, either a Reset or a HALT will redirect the CPU to the address of the non-maskable interrupt (NMI), where it will continue execution from either 'READY' or DOS reboot.

By cutting loose pin 18 of the CPU (Z40), and tying Z53 pin 2 high, with a 4.7K ohm resistor, the HALT instruction will operate as intended by the Z-80's designers. Memory will continue to be refreshed, as the 4116s' refresh cycle depends on RAS (row address strobe), which will be output as usual by Z72.

Cut the trace running from Z40, pin 18. Attach a 4.7K ohm resistor from Z53 pin 2, to Z53 pin 14. The modification is complete.

A Real Break

Later in this Chapter is a power-on monitor, which will give more importance to the minor change described here. The TRS-80 does not have a true 'break' function that resets the CPU to its power-up condition. It can be simulated by entering SYSTEM followed by /0, or by pressing the Reset button on a disk system.

This change is quite simple. Locate Z53, a 74LS132 NAND gate, which feeds (through Z52) the system reset pin of the Z-80 microprocessor. When the power is turned on, capacitor C42 takes time to charge, holding the processor in its reset mode for a few milliseconds. When the capacitor charges completely, the input to Z53 is held high by resistor R47.

Run a wire from pins 12 and 13 of Z53, through a 1K resistor, to one end of a SPST (single pole single throw) normally open pushbutton switch. The other end of the switch goes to ground (which can be found at Z53, pin 7). When this switch is pressed, the Z53 sees a low signal, and resets the processor to address 0000. A resistor is used in series so that C42 is not overexerted by shorting directly to ground.

Be sure to mount this button well out of the way; for disk systems, it's as fatal as the Reset button. For Level II it means MEMORY SIZE? and the loss of any program in memory.

Stuck Relays

There are two ways to fix a stuck relay: bang it until it unsticks, or replace it. The many published fixes to tape recorders, or additional external circuitry, simply don't take into consideration any changes you might make to your tape system; nor to the possibility you may take your CPU with you when you work elsewhere.

A replacement relay is available from *Lab Service, Inc., Box 383, Hustisford, Wisconsin 53034*. Unlike the relay normally installed in the TRS-80, it is reliable, low power and has gold contacts. I have installed this unit in my computer, and run the high current motor of a CTR-41 through 100,000 operations continuously over the course of a week. Neither the CTR-41 nor the relay failed.

The cassette relay is mounted just behind the jacks for power, video and cassette. It is a cylinder approximately the size of the LSI replacement relay, but in most TRS-80s it contains six leads instead of four.

Using solder-wick and a hot soldering iron, heat the solder connections and draw off excess solder. While doing this, slide a thin flat blade under the relay. Take care not to move the nearby small video trimmer potentiometers. Use the blade as a lever, lifting the old relay off the board. Apply very gentle pressure to this lever, and alternately melt the three solder connections on each side of the relay. Do not use force that

might crack the board or lift solder traces from the board.

When one side of the relay is completely free, grasp it with your hand and pull *gently*, while heating the remaining three connections on the opposite side of the relay. It will eventually pull free.

Remove excess solder from the holes with the solder-wick. Use a fine splinter to open the eyes of the holes, if all the solder does not flow into the solder-wick.

Now examine the circuit board, noting that on the end of the relay position nearest the computer's connection jacks there are two otherwise unused connections; no circuit traces run out from these points. Look at the replacement relay, noting that it has four wires, three on one side and one on the other. Orient it above the board so that the unused holes match the end of the relay with the single wire. This wire feeds into the center of the trio of holes.

Carefully insert the relay in place; if you have properly cleaned out the connection holes, the new relay wires will slide in, barely protruding from the opposite side of the board. You may have to bend them just a bit in order to get them to feed through.

Once the wires are in place, apply a very small amount of solder, and secure the relay on the board. Check carefully for solder splashes, shorts or cracks, and clean or repair them. Such a check is doubly important here because the relay driver circuit (unlike most other circuits in the TRS-80) cannot handle a short circuit for very long.

Refit the boards and covers of the computer together, replace the power and cassette cables; then put the cassette player into its play position, and enter this program:

```
10 PRINT#-1," "
20 FOR N = 1 TO 500 : NEXT
30 GOTO 10
```

This program will turn the cassette player on and off at regular intervals. If it does not work properly, double check all connections, especially the orientation of the relay, and try again. This fix should eliminate any further concern over a sticking relay causing skipped data or missed program loads.

High Resolution Graphics

Now for the biggie. In the past few years, there has been quite a bit of excitement generated by the idea of high-resolution graphics. Reasonably representative images can be drawn with them, and animation is considerably more exciting; especially when compared with the extremely high resolution of the type used in the latest generation of coin-operated video games. The Apple computer was advertised with a heavy emphasis on their high resolution, hard as it is to work with. The new Radio Shack Color Computer also offers several modes of resolution.

The following project can either stand alone as a plug-in peripheral, or be integrated as part of the TRS-80 keyboard unit. In either case, the specifications are:

1. Resolution of 384 dots wide by 192 dots deep.
2. Full compatability with all current software.
3. Simultaneous use *and overprint* of normal TRS-80 alphanumerics *and graphics*.
4. Addressing using six bits in contiguous memory blocks of 768 bytes each; sixteen total memory blocks are used.

The hardware involved in this project, including power supply and miscellaneous hardware, will be under \$100 (probably closer to \$70 by the time you read this), yet will compete easily with any high-resolution add-on for the TRS-80.

On the negative side, this project will involve a great deal of wire wrapping or soldering, and will eat up one chunk of 16K memory address space when it is used. It will not actually compete with or replace the top memory block in the expansion box (there is no electronic conflict) but will be addressed from C000 to FFFF. Alternatively, it may be addressed from 8000 to BFFF. In either case, you do not need the expansion box to run this memory.

The TRS-80 screen has 1,024 locations in a grid of 64 characters across by 16 lines down. Within each of these grid elements are blocks 2

pixels by 3 pixels (a pixel is a 'picture element') for the familiar coarse graphics mode accessed with SET and RESET. If you turn the contrast fully down and reduce the brightness of the screen, the individual dots which make up the graphic and alphanumeric characters can be seen with a sharp eye; a magnifying glass will make the dots very clear.

In order to produce a complete screen line of letters, the locations in video memory are handed to a circuit which actually accesses them twelve times - once for each pass the electron beam makes horizontally. At each pass, a row of dots corresponding to part of the whole line of letters is shifted out to the video beam. Each dot (or 'undot') then turns the beam on or off for the tiny fraction of a second it takes to sweep across 1/384th of the screen.

The point is clear: an electronic event takes place for every dot on every line of the screen. This means that it is possible to create an individual, addressable electronic event for each screen dot.

The process might work like this:

1. Devise a circuit whose timing characteristics are identical to the video timing of the usual TRS-80 circuits.
2. Instead of addressing the same video memory on twelve consecutive screen lines, have the addressing select *different* memory for each line.
3. Have the contents of that memory filled by the TRS-80, and displayed on the same or a different monitor. The add-on has its own video circuitry, but can be displayed on the same monitor because step 1 specifies that the timing characteristics must be identical to those in the TRS-80. It's like an auto with 4-wheel drive, where all wheels are capable of working together; or a dual-capstan tape recorder, where both capstans pull the tape to ensure steady contact with the playback head.

The circuit shown opposite presents the complete high-resolution circuitry. There are two ways of building this circuit since the areas shaded in grey *are already present within the TRS-80*. If you wish, you can solder directly to those circuits inside the computer, saving yourself some parts and perhaps a bit of time. Otherwise, the entire circuit can be built separately.

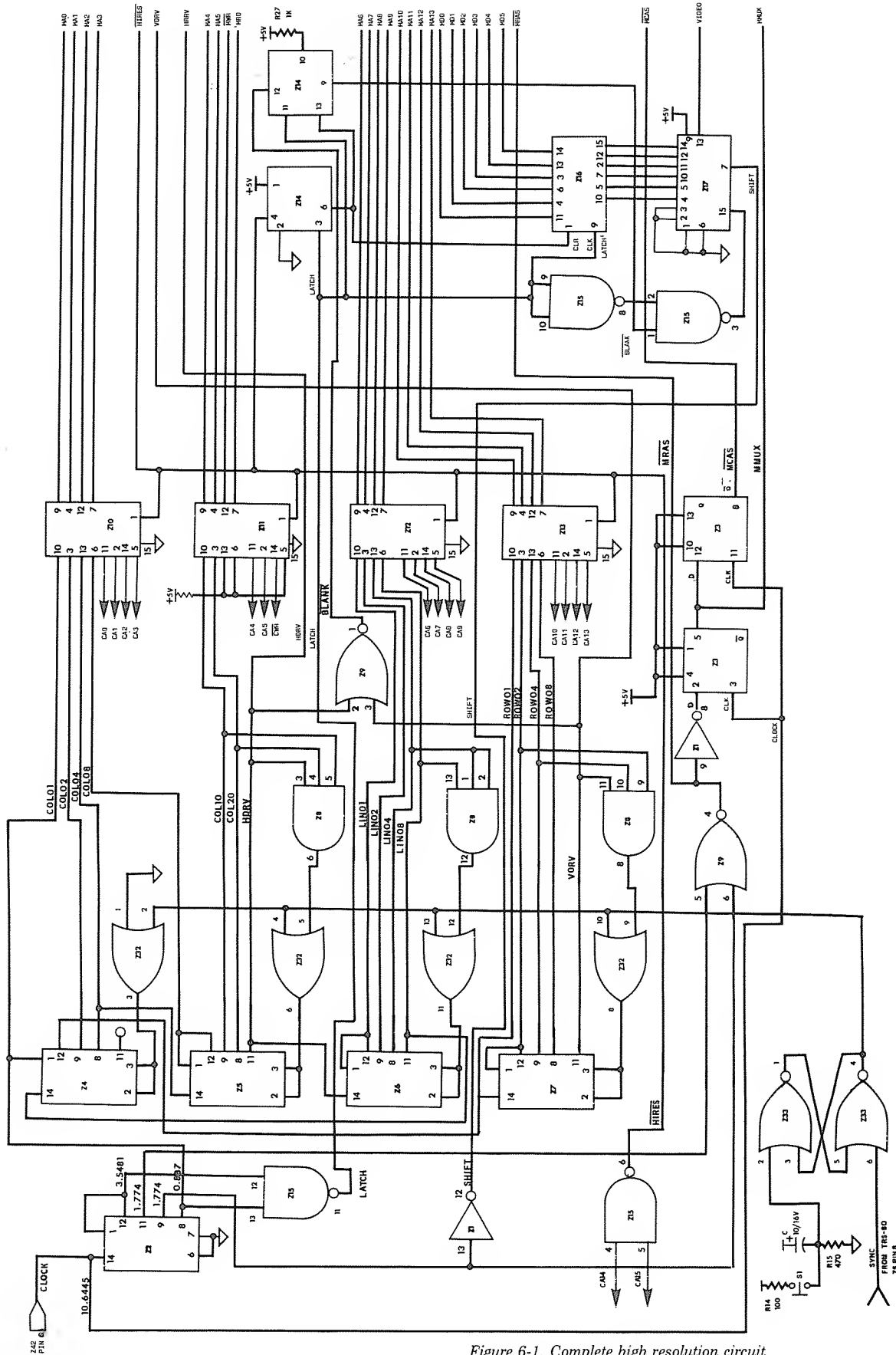


Figure 6-1. Complete high resolution circuit.

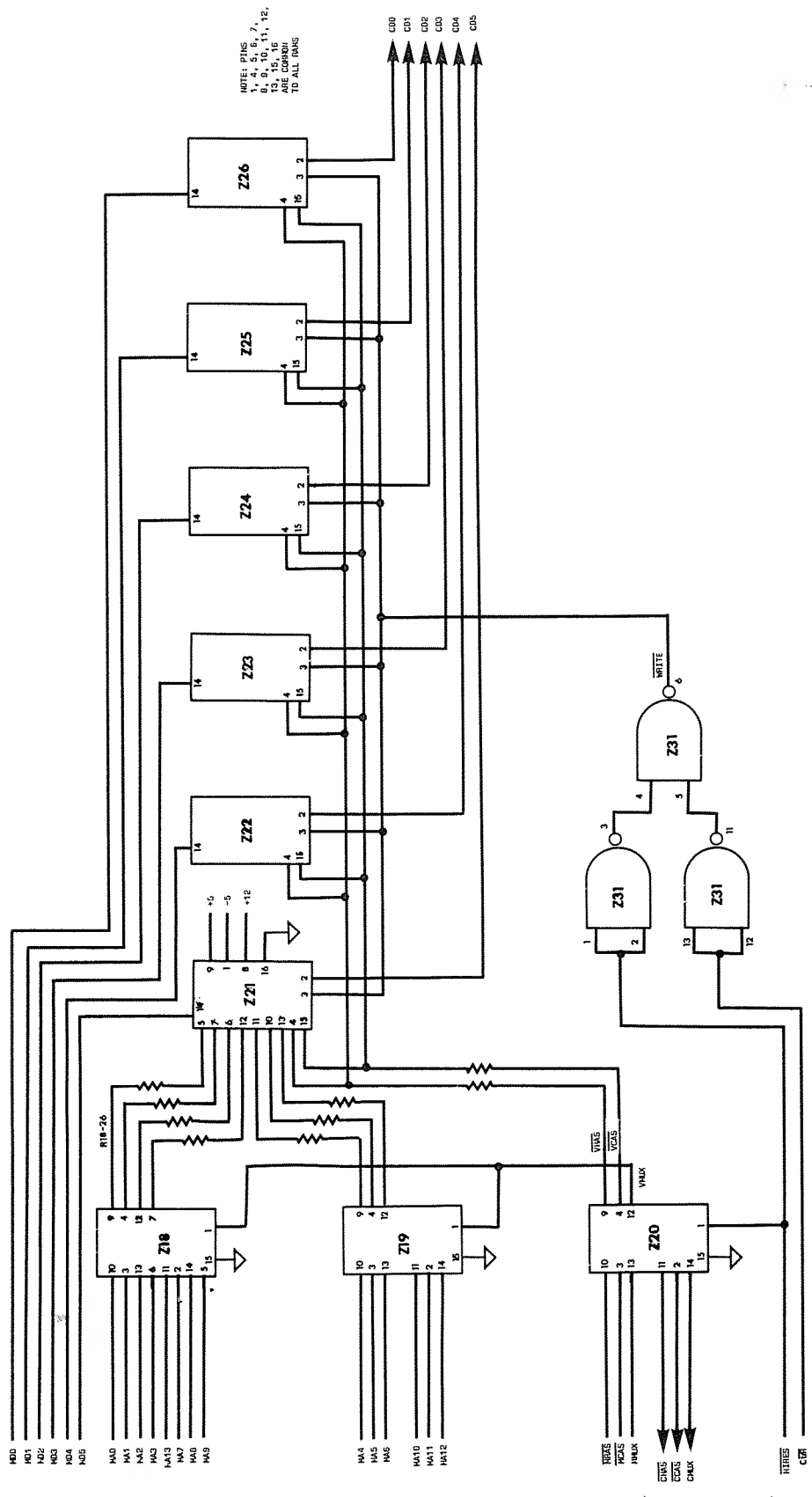


Figure 6-1b. High resolution memory.

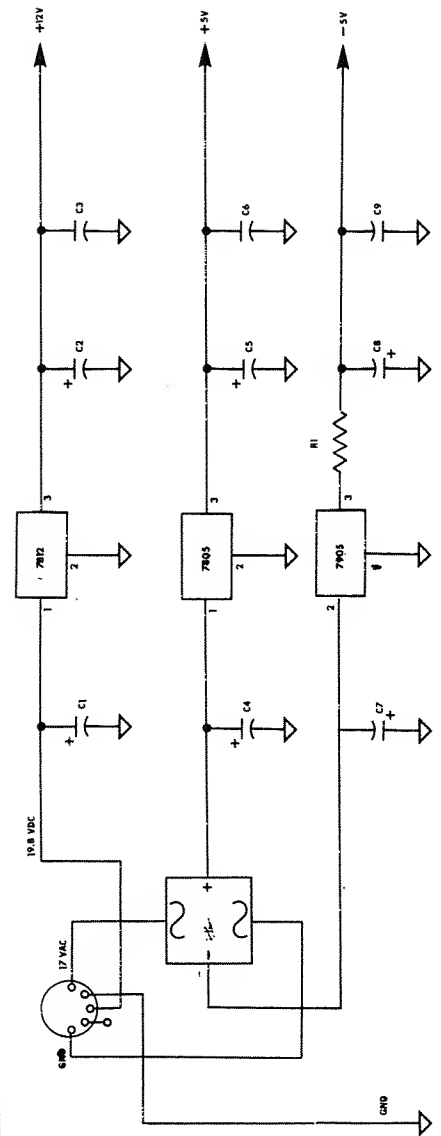


Figure 6-1c. High resolution power supply.

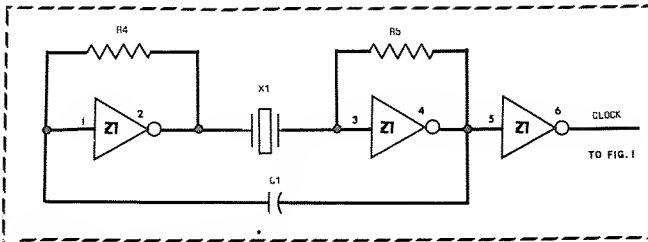


Figure 6-2. Hires clock circuit.

Figure 6-2 is the clock. The 10.6445 MHz crystal is the same one used in the TRS-80; Radio Shack sells it for \$3.95, on special order. A small trimmer capacitor is included so that the frequency of the high resolution board can be aligned identically to that in the computer.

Figure 6-3 is the master video countdown chain and timing circuitry, also nearly identical to that inside the TRS-80. There are a few exceptions. The logic necessary for 32 characters per line is not present (it is not needed in high resolution mode, although the normal alphanumeric may be displayed in 32-character mode simultaneously with the hi-res graphics screen). Also, the four outputs of Z6 do not feed any latches or character generators; instead, they become the line of dots addresses for high resolution memory. Identical (top and side) screen blanking is used.

Figure 6-1. is the high-resolution memory itself. The familiar 4116 type, 16K dynamic memories are used in this circuit (250 nS or less is essential), but with a difference. The hi-res board must generate its own memory refresh, yet hand over control to the TRS-80 when it needs to select memory into which it will write information. Thus, Z20 multiplexes the on-board refresh/select (MRAS, MCAS, MMUX) with the TRS-80 select (CRAS, CCAS, CMUX).

The switch is made by the simultaneous presence of addresses 14 and 15 on the address bus.

On-board refresh/select is generated by the clock in combination with two flip-flops (Z3a/b), producing select in this order:

1. Z2 pin 8, selects the lowest portion of the address; as such, it is the fastest changing memory select signal.
2. Two clock cycles later, Z2 pin 11, produces a signal which will be gated by Z9 and inverted to produce MRAS (row address strobe).
3. One clock cycle later, MMUX goes high, produced by the clocking of Z3a.
4. One clock cycle later, MRAS is continued as Z2 pin 11 goes low by Z2 pin 9 going high. The transition is simultaneous and virtually invisible.
5. At that time, MCAS is produced when Z3b is clocked low at the NOT Q output.
6. Memory data is stable at this time, and two clock cycles later, LATCH is issued by Z1e, latching the data into Z16 for its trip through shift register Z17.

Figure 6-3 shows the creation of horizontal and vertical synchronization signals, and the horizontal and vertical screen positioning controls. This circuitry again is identical to that in the TRS-80, as is Figure 6-4, the video mixing circuitry.

There is only one critical construction area in the device, and that is the circuitry surrounding the 10.6445 MHz crystal (Z16, R4-5, C1). The wires in this area must be very short, and all the parts clustered together. Capacitor C1 should be the only part of the circuit responsible for tuning the crystal's frequency, not random capacitance

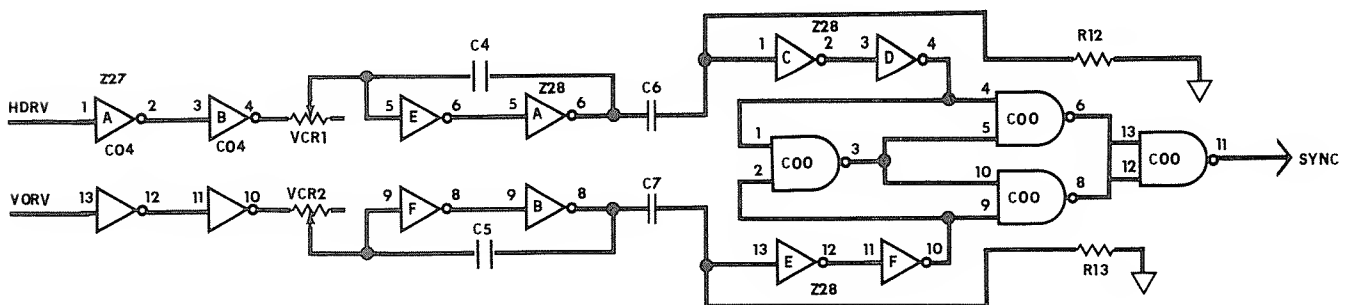


Figure 6-3. Master video circuit.

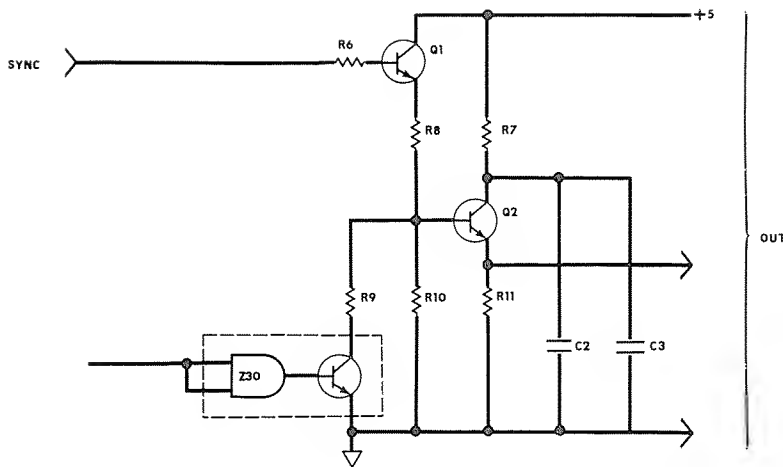


Figure 6-4. The video mixing circuitry.

introduced by a haphazard bunch of wires.

Other than the wire layout near the crystal, construction is time-consuming but straightforward. I recommend wire wrapping the entire circuit; use different colors for data, address, video, ground, etc., so that troubleshooting will be simplified.

The completed circuit will have these external controls:

1. **Power.** Three voltages, +5, -5, and +2 are necessary. -5 volts must be present first and last.
2. **Mixing.** This controls the intensity of the high-resolution board with respect to the TRS-80 alphanumerics and graphics.
3. **Fine tuning.** This adjusts the frequency of the 10.6445 MHz crystal to that of the TRS-80. Occasional adjustments will be necessary with temperature changes.
4. **Vertical and horizontal positioning.** These control the placement of the image on the screen; it should coincide with the alphanumeric screen normally produced by the TRS-80.
5. **Input.** This accepts a cable running from the TRS-80 video jack, which would normally attach to the video monitor.
6. **Output.** This accepts a cable from the video monitor, and provides an output which mixes the TRS-80 alphanumerics

and graphics with the high-resolution dots.

To use the device, attach a 5-pin DIN cable between the TRS-80 video jack and the input jack hi-res board. Connect the video monitor to the output jack of the hi-res board, then attach a 40-pin edge connector from the TRS-80 to the hi-res board. Turn the mixing control fully counterclockwise (Hi-Res Out). Power up the hi-res board, and then the rest of the system in normal order.

As usual, MEMORY SIZE? should appear; if so enter 49152 (for a 48K computer). The system should operate as usual. Enter the following program:

```
10 FOR X = 15360 TO 16383
20 POKE X,129
30 NEXT
40 FOR X = -16384 TO 0
50 POKE X,175
60 NEXT
70 GOTO 70
```

The screen will fill with small graphics blocks. There will be a pause of almost a minute while the rest of the program is running. Put an AM radio next to the computer to determine when the program is complete. Now bring the mixing control of the hi-res board clockwise until dots, herringbone, jitter and/or other interference appears on the screen. This is a good sign.

If you have a stable enough screen to see the alternating dot patterns produced by the hi-res board, then adjust the horizontal and vertical positioning controls, if necessary, to center the image with that of the TRS-80. To remove the jitter and herringbone adjust the fine tuning control.

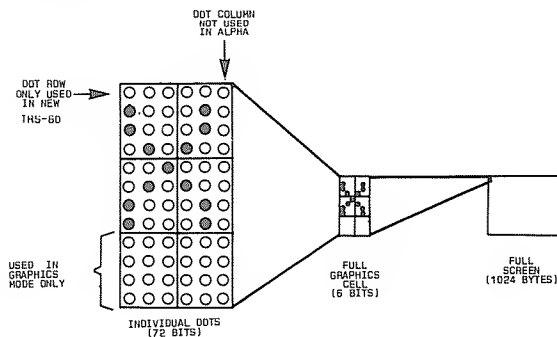
Now your screen should display an alternating pattern of TRS-80 graphics, along with an overlay of thin vertical lines of hi-res graphics dots. If you have any difficulty getting this pattern, or if there are any other problems, refer to the troubleshooting section.

Use of the Hi-Res board is simple. Addresses from C000 through C03F contain the information to create the first line of dots, addresses C040 through C07F contain the second line, etc. A contiguous block of memory from C000 through C2FF is used for the first twelve lines of dots. But since the display is twelve lines, and not sixteen, the addressing takes a jump in order to be compatible with the familiar 64 x 16 normal screen display. Thus, addresses C300 through C3FF are not used, and the second group of 12 dot lines begins at address C400 and continues

through C6FF. Here is a full memory map of the Hi-Res board:

Hi-Res Board Memory Map		
C000 - C03F	Group 1, Line 1, Screen Line 1	
C040 - C07F	Group 1, Line 2, Screen Line 2	
C080 - C0BF	Group 1, Line 3, Screen Line 3	
C0C0 - C0FF	Group 1, Line 4, Screen Line 4	
C100 - C13F	Group 1, Line 5, Screen Line 5	
C140 - C17F	Group 1, Line 6, Screen Line 6	
C180 - C1BF	Group 1, Line 7, Screen Line 7	
C1C0 - C1FF	Group 1, Line 8, Screen Line 8	
C200 - C23F	Group 1, Line 9, Screen Line 9	
C240 - C27F	Group 1, Line 10, Screen Line 10	
C280 - C2BF	Group 1, Line 11, Screen Line 11	
C2C0 - C2FF	Group 1, Line 12, Screen Line 12	
.		
C400 - C43F	Group 2, Line 1, Screen Line 13	
C440 - C47F	Group 2, Line 2, Screen Line 14	
.		
C680 - C6BF	Group 2, Line 11, Screen Line 23	
C6C0 - C6FF	Group 2, Line 12, Screen Line 24	
.		
C800 - CAFF	Group 3, Lines 1 - 12 Screen Lines 25 - 36	
.		
CC00 - CEFF	Group 4, Lines 1 - 12 Screen Lines 37 - 48	
.		
0000 - 02FF	Group 5, Lines 1 - 12 Screen Lines 49 - 60	
.		
0400 - 06FF	Group 6, Lines 1 - 12 Screen Lines 61 - 72	
.		
0800 - 0AFF	Group 7, Lines 1 - 12 Screen Lines 73 - 84	
.		
0C00 - 0EFF	Group 8, Lines 1 - 12 Screen Lines 85 - 96	
.		
E000 - E2FF	Group 9, Lines 1 - 12 Screen Lines 97 - 108	
.		
E400 - E6FF	Group 10, Lines 1 - 12 Screen Lines 109 - 120	
.		
E800 - EAFF	Group 11, Lines 1 - 12 Screen Lines 121 - 132	
.		
EC00 - EEFF	Group 12, Lines 1 - 12 Screen Lines 133 - 144	
.		
F000 - F2FF	Group 13, Lines 1 - 12 Screen Lines 145 - 156	
.		
F400 - F6FF	Group 14, Lines 1 - 12 Screen Lines 157 - 168	
.		
F800 - FAFF	Group 15, Lines 1 - 12 Screen Lines 169 - 180	
.		
FC00 - FEFF	Group 16, Lines 1 - 12 Screen Lines 181 - 192	

Only six bits of each byte are used (the least significant six); thus, six one-bit-wide, memory chips are used in the circuit. The bits fit into their respective lines and memory addresses as follows:



Before continuing, clear out the hi-res memory with:

```
FOR X = -16384 TO 0 : POKE X,0 : NEXT
```

Drawing simple lines is an easy process; for a horizontal one, just enter:

```
FOR X = -12288 TO -12224 : POKE X, 63 : NEXT
```

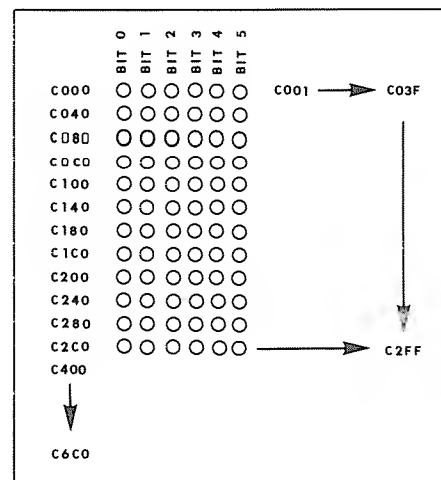
A vertical one is produced by stepping through groups:

```
5 CLS : REM KILL ALPHANUMERICS
10 FOR Y = -16352 TO 0 STEP 1024 : REM STEP THROUGH GROUPS
20 FOR X = Y TO Y+(12*64) STEP 64 : REM STEP THROUGH LINES
30 POKE X,1 : NEXT X : REM SET ONE PIXEL
40 NEXT Y : REM TO NEXT LINE GROUP
50 GOTO 50 : REM KEEP DISPLAY INTACT
```

Diagonal lines are more complicated, because more than two sets of increments must be specified; but simple diagonals can be created. For diagonals and variable-width lines, change listing to read as follows:

```
10 INPUT Q : INPUT R : CLS
20 FOR Y = -16352 TO 0 STEP 1024
30 FOR X = Y TO Y+(12*64) STEP 64+Q
40 POKE X,(R AND 63) : NEXT X
50 NEXT Y
60 GOTO 60
```

For serious graphics, assembly language programming is the only way real speed can be achieved. This is a very 'custom' type of programming, and only the simplest of subroutines will be presented here. For drawing circles, ellipses, and curves the functions will have to be stored in a look-up table and calculated. Listing 6-1 is an assembly listing to draw graphic lines on the screen, given a set of coordinates.



```

00100 ; #####
00110 ; ROUTINE TO DRAW HORIZONTAL AND VERTICAL LINES FR. BASIC
00120 ; #####
00130 ;
C000 00140 HIRES EQU 0C000H ; START OF HIRES GRAFIX
0A7F 00150 XFER EQU 0A7FH ; VARIABLE XFER ROUTINE
00160 ;
7F00 00170 ORG 7F00H ; SOMEWHERE IN MEMORY
00180 ;
00190 ; #####
00200 ; SUBROUTINE TO DETERMINE THE CORRECT BASIC USR(X) CALL
00210 ; #####
00220 ;
7F00 C07F0A 00230 ENTRY CALL XFER ; GET VALUE FROM BASIC
7F03 7C 00240 LD A,H ; GET MSB INTO ACCUM.
7F04 B5 00250 ADD A,L ; ADD LSB FROM HL PAIR
7F05 A7 00260 AND A ; TEST IF IT IS ZERO
7F06 2B08 00270 JR Z,PCLS ; CLEAR SCREEN ROUTINE
7F0B FE01 00280 CP 1 ; TEST IF IT IS ONE
7F0A 2B15 00290 JR Z,PHORIZ ; HORIZONTAL LINE ON 1
7F0C FE02 00300 CP 2 ; TEST IF IT IS A TWO
7F0E 2B30 00310 JR Z,PVERT ; VERTICAL LINE ON 2
7F10 C39719 00320 JP 1997H ; SN? ERROR IF NOT 1,2,3
00330 ;
00340 ; #####
00350 ; SUBROUTINE TO CLEAR THE SCREEN IN HIGH-RESOLUTION MODE
00360 ; BASIC FORMAT: M=USR(A). A MUST ALWAYS BE ZERO.
00370 ; #####
00380 ;
7F13 AF 00390 PCLS XOR A ; GET CHARACTER TO WRITE
7F14 F5 00400 PUSH AF ; SAVE THAT CHARACTER
7F15 2100C0 00410 LD HL,HIRES ; GET BEGINNING OF HI-RES
7F1B F1 00420 PCLEAR POP AF ; RESTORE ZERO CHARACTER
7F19 77 00430 LD (HL),A ; PUT IT IN PLACE IN MEM
7F1A F5 00440 PUSH AF ; SAVE CHARACTER AGAIN
7F1B 7C 00450 LD A,H ; GET MSB OF CURRENT LOC.
7F1C B5 00460 OR L ; GET LSB AND TEST PAIR
7F10 20F9 00470 JR NZ,PCLEAR ; LOOP BACK FOR NEXT
7F1F F1 00480 POP AF ; GET STACK BACK IN SHAPE
7F20 C9 00490 RET ; BACK TO CALLING ROUTINE
00500 ;
00510 ; #####
00520 ; SUBROUTINE TO DRAW A HORIZONTAL LINE. BASIC FORMAT:
00530 ; POKEN,B:C1=INT(C/6):C2=C-C1*6:POKEN+1,C1:POKEN+2,C2:
00540 ; 01=INT(0/6):02=0-01*6:POKEN+3,01:POKEN+4,02:M=USR(1)
00550 ;
00560 ; B=LINE NUMBER (0-11, 16-27, 32-43, 48-59, 64-75, 80-91,
00570 ; 96-107, 112-123, 128-139, 144-155, 160-171, 176-
00580 ; 187, 192-203, 208-219, 224-235, 240-251
00590 ; C=ORIGIN POSITION (0-383).
00600 ; D=DESTINATION POSITION (0-383). C MUST BE LARGER THAN 0
00610 ; #####
00620 ;
7F21 0021897F 00630 PHORIZ LD IX,N ; POINT TO THE STORAGE
7F25 C06E7F 00640 CALL FINDER ; GET STARTING POSITION
7F2B 3ABC7F 00650 LD A,(N+3) ; GET VALUE AT "N+3"
7F2B 4F 00660 LD C,A ; PLACE IN MSB OF BC
7F2C AF 00670 XOR A ; CLEAR ACCUM TO ZERO
7F2D 47 00680 LD B,A ; PLACE IN LSB OF BC
7F2E E5 00690 PUSH HL ; SAVE START LOC'N
7F2F 05 00700 PUSH OE ; READY FOR XFER BACK
7F30 E1 00710 POP HL ; TRANSFERRED BACK
7F31 09 00720 ADD HL,BC ; PERFORM THE ADDITION
7F32 E5 00730 PUSH HL ; READY FOR TRANSFER
7F33 01 00740 POP OE ; AND GET INTO DESTIN'N
7F34 E1 00750 POP HL ; RESTORE ORIGINAL VALUE
7F35 3AB87F 00760 LD A,(N+2) ; GET VALUE AT "N+2"
7F38 86 00770 OR (HL) ; ADD TO CURRENT LOC'N
7F39 77 00780 LD (HL),A ; AND PUT INTO PLACE
7F3A 23 00790 INC HL ; NEXT SCREEN POSITION
7F3B AF 00800 LOOP1 XOR A ; CLEAR ACCUM TO ZERO
7F3C 3A3F00 00810 LD A,(03FH) ; VALUE TO FILL BYTE
7F3F E052 00820 SBC HL,OE ; PERFORM SUBTRACTION
7F41 2B04 00830 JR Z,HOROUT ; OUT OF ROUTINE
7F43 77 00840 LD (HL),A ; PUT BYTE IN PLACE
7F44 23 00850 INC HL ; GET NEXT SCREEN POS'N
7F45 1BF4 00860 JR LOOP1 ; GO BACK FOR NEXT FILL
7F47 3A807F 00870 HOROUT LD A,(N+4) ; GET BACK FINAL BYTE
7F4A B6 00880 OR (HL) ; ADD TO VALUE ON SCREEN
7F4B 77 00890 LD (HL),A ; PUT IT ON THE SCREEN
7F4C C9 00900 RET ; AND BACK TO BASIC
00910 ;
00920 ; #####
00930 ; SUBROUTINE TO DRAW A VERTICAL LINE. 48K RAM IN PLACE.
00940 ; BASIC FORMAT:
00950 ; POKEN,B:C1=INT(C/6):C2=C-C1*6:POKEN+1,C1:POKEN+2,C2:
00960 ; 01=INT(0/6):02=0-01*6:POKEN+3,01:POKEN+4,02:M=USR(2)
00970 ;
00980 ; E=HORIZONTAL POSITION (0-255)
00990 ; F=LINE NUMBER (SEE HORIZONTAL LINE, ABOVE)
01000 ; G=LINE NUMBER (SEE HORIZONTAL LINE, ABOVE)
01010 ; #####
01020 ;
7F40 0021897F 01030 PVERT LD IX,N ; POINT TO THE STORAGE

```

Listing 6-1. Hires demonstration program.

Listing Continued

Without TRS-80 memory parallel to that in the hi-res board, it is not possible to read the contents of the high-resolution memory directly.

The contents must be stored in some form elsewhere. When the high 16K block is in place in the expansion box, however, six bits of each byte are identical to those on the screen. Ideally, the entire block of high resolution memory (16K by 6 bits) and TRS-80 memory (16K by 8 bits) should be cleared out by POKEing 0 in place first. Then an in-computer image of the high resolution screen can be maintained at all times.

Another interesting mode of using the high resolution board is with a separate screen. Normal alphanumerics can appear on the TRS-80 monitor, while the high-resolution graphics can be presented on a parallel screen. This way, the table calculations and information reported can be displayed on the computer's monitor for reference. The high-resolution screen will be unaffected by anything done by the TRS-80 unless its memory is being specifically written to. Not only can action games of the Startrek type be more interesting and challenging — with visuals and info displayed on different screens — but for experimentation and analysis, the high-resolution display is unbothered by program changes.

To use this mode, merely leave your TRS-80 monitor plugged into the computer. Then send the video information in the hi-res board to another video monitor, or to an ordinary television set via an RF modulator.

Troubleshooting

With a complicated project like this, there is a good chance that the system will not work perfectly the first time. The main problems together with their causes and solutions are outlined below.

1. The screen keeps tearing or jittering no matter what setting the fine tuning is on. If the fine tuning has no effect at all, it may be defective. Replace it. If the tuning gets better, but can't quite pull it in, you can put another capacitor in parallel to increase the capacitance, or replace the crystal (in either the hi-res board or the TRS-80) with one better matched to the other.
2. The high-resolution graphics cannot be changed, remaining the same as when the power was turned on. The memory write circuits are not working properly. Check the memory-select wiring at Z20; the write

Continued Listing

```

7F51 C06E7F 01040 CALL FINDER ; GET SCREEN START BYTE
7F54 E5 01050 PUSH HL ; SAVE THE START VALUE
7F55 0021BC7F 01060 LO IX,N+3 ; POINT TO THE TABLE
7F59 C06E7F 01070 CALL FINDER ; AND DO THE WORK
7F5C E5 01080 PUSH HL ; READY VALUE TO TRANSFER
7F5D 01 01090 POP OE ; AND TRANSFER TO DEST'N
7F5E E1 01100 POP HL ; RESTORE START POSITION
7F5F 014000 01110 LO BC,40H ; SCREEN BYTE LINE OFFSET
7F62 AF 01120 XOR A ; CLEAR CARRY FLAG
7F63 3AB87F 01130 LO A,(N+2) ; GET BYTE FROM STORAGE
7F66 77 01140 LD (HL),A ; STASH IT ON SCREEN
7F67 09 01150 ADD HL,BC ; MOVE UP ON THE SCREEN
7F68 E052 01160 SBC HL,OE ; CHECK IF DONE YET
7F6A 20FA 01170 JR NZ,LOOP2 ; BACK IF NOT DONE
7F6C 77 01180 LD (HL),A ; PUT LAST BYTE IN
7F6D C9 01190 RET ; BACK TO BASIC
01200 ;
01210 ; #####
01220 ; FINDER SUBROUTINE LOCATES PROPER BYTE WITHIN HIRES AREA
01230 ; #####
01240 ;
7F6E 2100C0 01250 FINDER LO HL,HIRES ; GET HI-RESOLUTION SCRIN
7F71 007E0D 01260 LO A,(IX) ; GET START BYTE OFFSET
7F74 47 01270 LO B,A ; PLACE IN B REGISTER
7F75 AF 01280 XOR A ; CLEAR THE CARRY FLAG
7F76 CB1B 01290 RR B ; DIVIDE BY TWO AND ...
7F7B CB19 01300 RR C ; ... ROTATE IN ORDER ...
7F7A CB1B 01310 RR B ; ... ACTUALLY TO ...
7F7C CB19 01320 RR C ; ... MULTIPLY BY 64
7F7E 09 01330 ADD HL,BC ; MAKE NEW SCREEN POS'N
7F7F E5 01340 PUSH HL ; READY IT FOR TRANSFER
7F80 01 01350 POP OE ; TRANSFER TO DEST'N
7F81 007E01 01360 LO A,(IX+1) ; GET START BIT OFFSET
7F84 4F 01370 LO C,A ; PLACE IN MSB OF BC
7F85 AF 01380 XOR A ; CLEAR ACCUM TO ZERO
7F86 47 01390 LO B,A ; PLACE IN LSB OF BC
7F87 09 01400 ADD HL,BC ; GET NEW START BIT
7F8B C9 01410 RET ; BACK TO CALLING ROUTINE
01420 ;
7F89 01430 N EQU $ ; MEMORY STORAGE POS'NS
01440 ;
06CC 01450 END 06CCH ; BACK TO BASIC
00000 TOTAL ERRORS
30026 TEXT AREA BYTES LEFT

ENTRY 7F00 00230
FINDER 7F6E 01250 00640 01040 01070
HIRES C000 00140 00410 01250
HOROUT 7F47 00870 00830
LOOP1 7F3B 00800 00860
LOOP2 7F66 01140 01170
N 7F69 01430 00630 00650 00760 00870 01030 010B0 01130
PCLEAR 7F1B 00420 00470
PCLS 7F13 00390 00270
PHORIZ 7F21 00830 00290
PVERT 7F40 01030 00310
XFER 0A7F 00150 00230

```

lines to the memory (Z11); and the write line from the computer (from edge card pin 13).

3. The high-resolution graphics keep changing without writing to them. The memory-select circuits may be selecting write for both read and write; check Z11. More likely, the memory refresh/select circuitry is miswired; check Z9a and Z3.

4. When creating lines of graphics, the dots do not appear in the correct place. This indicates the memory data and /or address lines are miswired; check the lines from the computer (pins 4-7, 9-11, 17-18, 20, 22, 24-28, 30-32, 34-36, 38 and 40), making sure they are in the correct order. Also check each memory circuit to be sure the address lines (pins 4-7, 10-13, and 15) are parallel in each memory IC. Finally, be sure Z4 correctly feeds Z10; Z5 correctly feeds Z11; Z6 correctly feeds Z12; and Z7 correctly feeds Z13. These four circuits are the memory count/multiplex circuits. Also check that Z10, 11, 12 and 13 correctly feed Z18 and 19.

5. No graphics are produced. This is a tough one. The fault could lie with

- (a) the clock formed by Z1a-c
- (b) the memory refresh circuits Z9b and Z3
- (c) latch and shift registers Z16 and Z17
- (d) memory circuits Z21 to Z26
- (e) video output formed by Z30, Q1 and Q2.

Check the screen display carefully, because if any of these sections are working (except the video output) the screen will be affected in some way, even if it is minor. If herringbone or some tearing is present when the fine tuning is adjusted, then the video output and sync circuits are probably okay. Also, be sure that the mixing control is not turned fully counterclockwise (TRS-80 on, hi-res off).

6. The computer crashes to MEMORY SIZE? or otherwise acts problematically. The hi-res board has no effect on the computer. No data is written to the computer from the hi-res board at any point; it only receives information. If the computer crashes, then faulty wiring is likely.

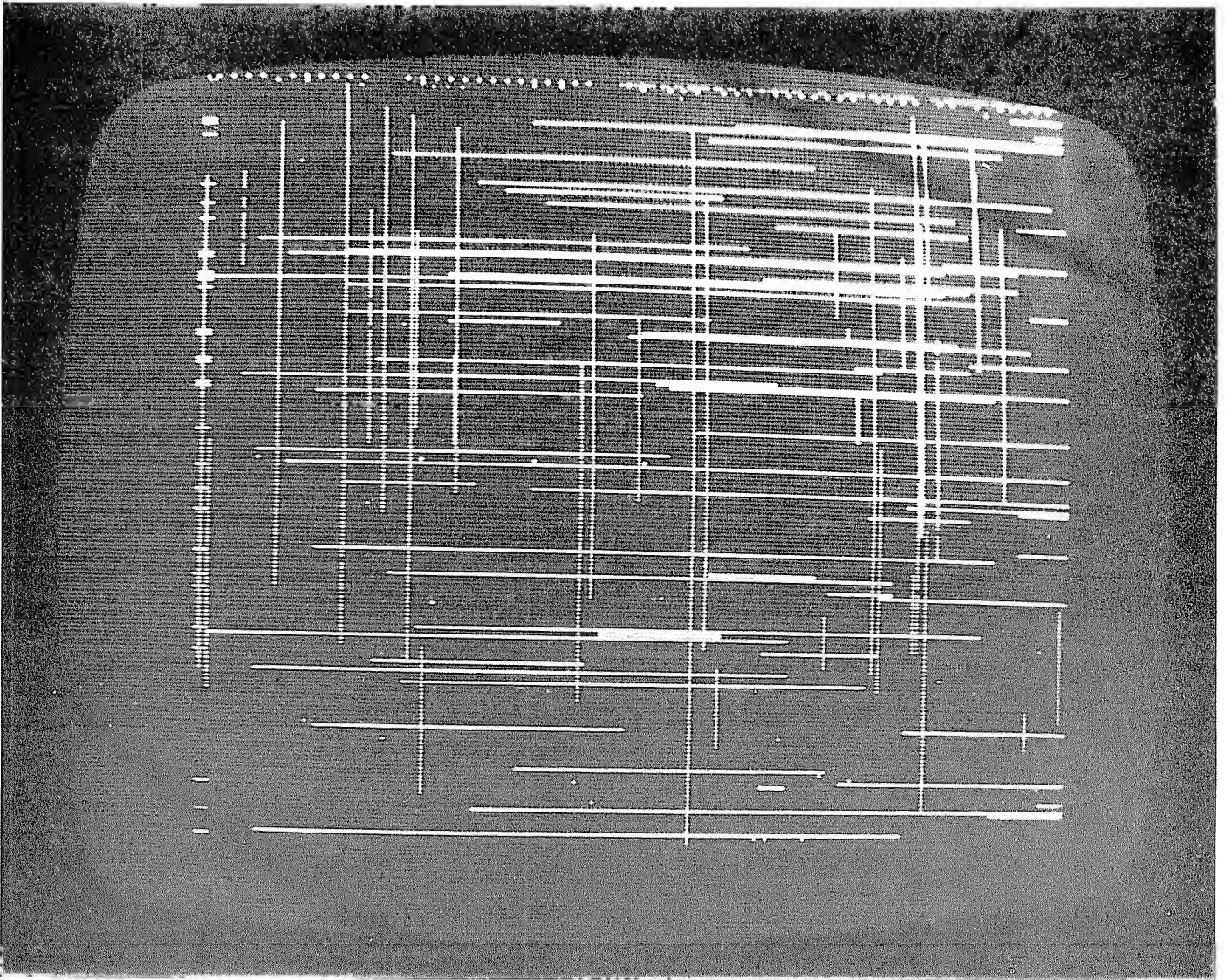


Photo 6-1. Hi-Res graphics example.

Replace BASIC ROMs

In the past few years, more and more TRS-80 users have asked me the question, "I don't need BASIC because I do all my work in machine language. How can I use that space for my own routines and language in ROM?". The answer is simple: replace your ROMs with 2716 (2K x 1 byte), 2732 (4K x 1 byte) or 2764 (8K by 1 byte) EPROMs.

The process is electronically simple. The Level II (or Level I) ROMs are removed, and a board containing the new *and old* ROMs is inserted at the edge connector. This board, then, can select either ROMs at the flick of a switch – exactly like the Apple's 'softcards' do it.

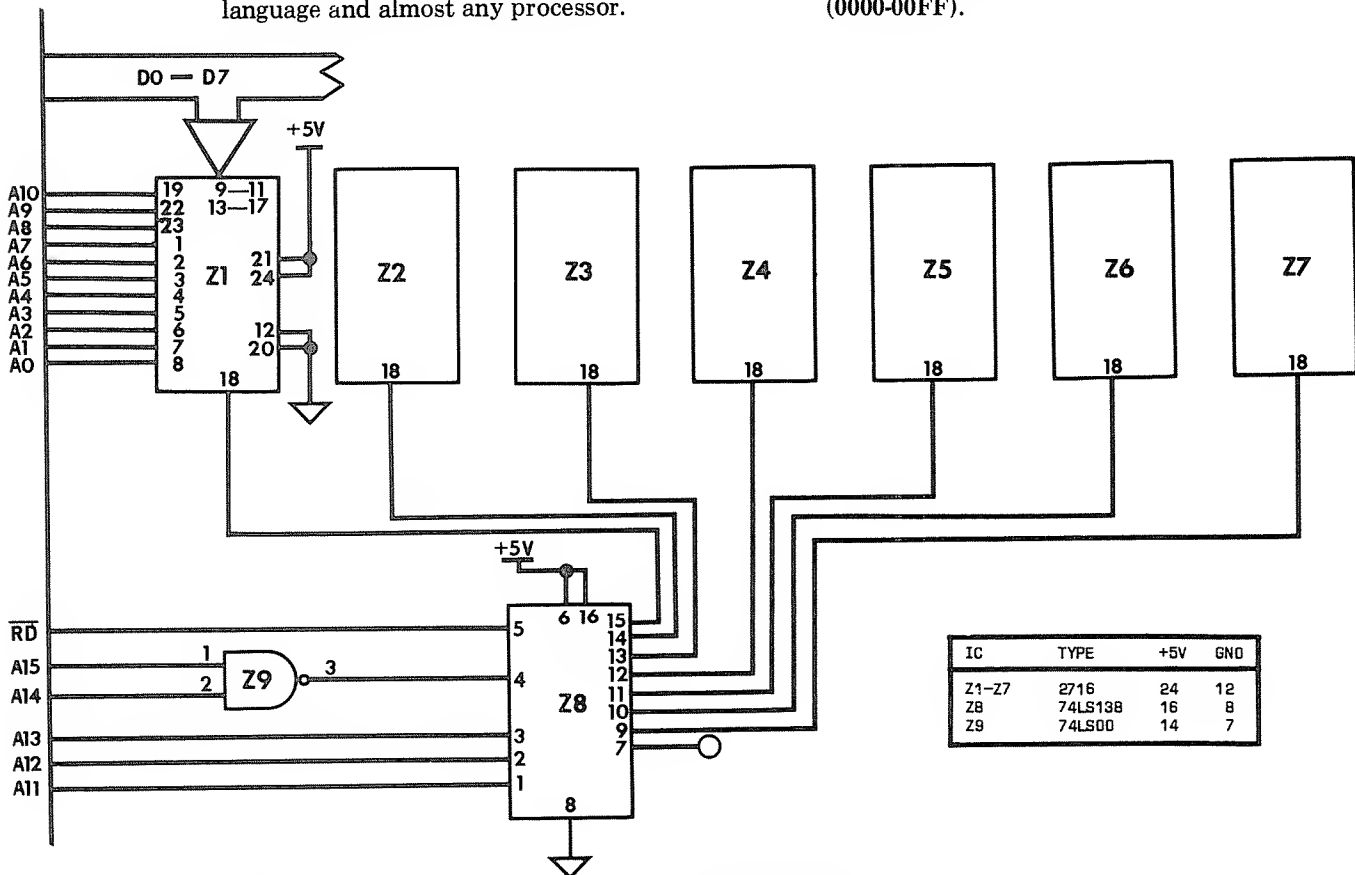
This section will present a circuit to use 2716 EPROMs and any other ROMs together and, as a bonus, a way to hand control of your TRS-80 over to another *microprocessor*! The power of such processors as the 6502, 6800 series, the 8060 (SC/MP) and others then becomes available to the TRS-80 user. Together with the appropriate bootstrap and executive programs in ROM, the TRS-80 can act with the strengths of almost any language and almost any processor.

There is no secret to adding ROM. The area from 0000 to 37C0 is free to use, and some of that method has already been presented. Figure 6-5 presents the circuitry that will handle seven 2716 EPROMs, switchable with the 3-chip Level II ROM set.

The real trick is making another processor available to the TRS-80's hardware. This other processor must be able to:

- access memory and peripherals in the TRS-80 memory map
- address at least 32K of memory
- be able to move its software stack anywhere in memory
- refresh the 4116 dynamic memories.

Of these restrictions, the last two are the trickiest. Without also placing RAM on this outboard device, the 6502 cannot be used, because it requires its stack in page 1 (0100-01FF) and many of its fastest (and hence most advantageous) instructions are limited to page 0 (0000-00FF).



1. REMOVE DIP SHUNT Z3.
2. ALL SEVEN 2716'S ARE CONNECTED IN PARALLEL EXCEPT PIN 18.

Figure 6-5. Circuit for replacing BASIC ROMs.

However, the refresh requirement is the most severe. Because the RAS-only refresh is controlled within the TRS-80, and its signals are killed when the processor is removed, the entire refresh process must be handled outside the computer by chips which were not designed for dynamic memory support.

Before turning to the details, you may be asking "How can control be wrested from the Z-

80?". The answer is found in the line marked TEST. This input line is present on pin 23 of the TRS-80 bus, and when brought low, causes all address lines (A0-A15), all data lines (D0-D7), read (RD), write (WR), input (IN), output (OUT), row address strobe (RAS), column address strobe (CAS) and memory select multiplex (MUX) to be put into three-state condition. This provides a wide-open computer bus for the outside world.

Power-Up Monitor

From my programmer (but not my user!) point of view, entering BASIC immediately upon power-up, or reset, is inconvenient and sometimes maddening. Rather, I would prefer a jump to a monitor of the type which can be found on the Ohio Scientific microcomputers: instead of MEMORY SIZE?, the user is presented with 'D/C/W/M?', which means 'Disk reboot, Cold start in BASIC (MEMORY SIZE?), Warm start in BASIC (READY), or Machine Language Monitor?'.

The user can exit from BASIC to the monitor and back at any point without jumping to a program-creating MEMORY SIZE? The cold start serves that purpose, and is almost always by choice.

There is such a possibility on the TRS-80, so long as two conditions are met:

1. A machine language monitor is resident, preferably in ROM.
2. The restart from 0000 is redirected from BASIC to that machine language monitor, and/or the NMI located at 0066 is redirected to that machine language monitor. This latter configuration preserves the last location of the program counter on the stack for examination.

The first requirement is easy to meet, and is in my opinion the most logical use of a ROM memory addition.

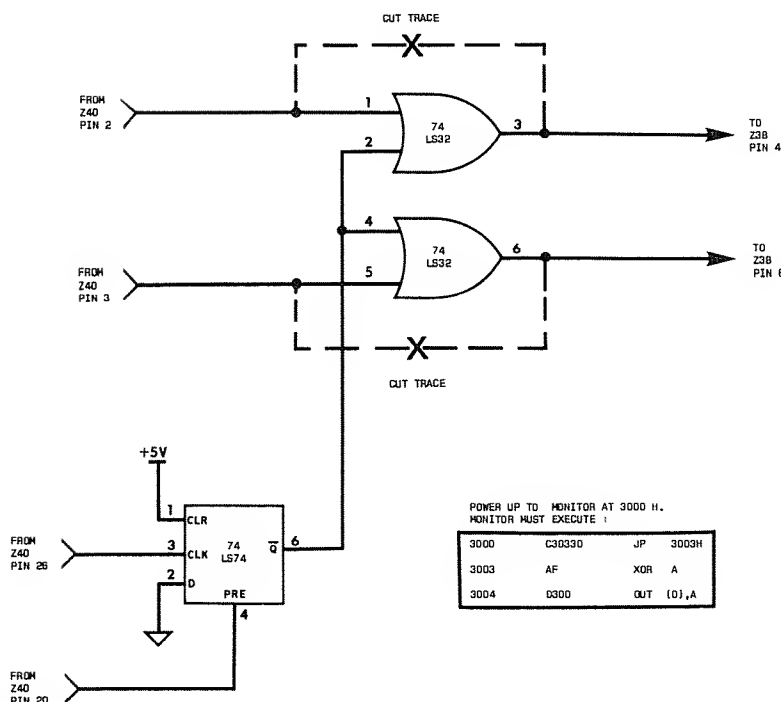


Figure 6-6. Modification needed for power-up monitor.


```

00100 ; #####
00110 ; POWER-UP PROGRAM FOR USE INDEPENDENTLY (WITHOUT BASIC
00120 ; ROMS IN PLACE) OR IN CONJUNCTION WITH EXIT AND RE-ENTRY
00130 ; TO BASIC UP RESET TO 0000 (BEFORE OISC SYSTEM REBOOT).
00140 ; #####
00150 ;
401A 00160 KPLACE EQU 401AH ; RAM STORAGE FOR STROKE
3000 00170 ORG 3000H ; MONITOR ENTRY POINT
00180 ;
00190 ; #####
00200 ; OPTIONAL SCREEN-CLEAR. IF THE EXIT CONOITION OF THE
00210 ; SCREEN IS IMPORTANT, THEN ELIMINATE CLEARING PROCESS
00220 ; #####
00230 ;
3000 F3 00240 OI ; KILL THE INTERRUPTIONS
3001 E073E242 00250 LO (42E2H),SP ; SAVE STACK POINTER
3005 31E042 00260 LO SP,42E0H ; GET NEW STACK POINTER
3008 F5 00270 PUSH AF ; SAVE REGISTERS FOR
3009 C5 00280 PUSH BC ; LATER OISPLAY
300A 05 00290 PUSH OE ; DURING THE
3008 E5 00300 PUSH HL ; MONITOR ENTRY
300C 21003C 00310 CLS LO HL,3C00H ; BEGINNING OF SCREEN
300F 11013C 00320 LO OE,3C01H ; NEXT POSITION ON SCREEN
3012 01FF03 00330 LO BC,3FFH ; NUMBER OF SCREEN POS'NS
3015 3620 00340 LO (HL),20H ; PRINT SPACE 1ST POS'N
3017 E0B0 00350 LOIR ; DO IT FOR WHOLE SCREEN
00360 ;
00370 ; #####
00380 ; OISPLAY OISK/MONITOR/WARM START/USER ROUTINE MESSAGE
00390 ; #####
00400 ;
3019 219732 00410 LO HL,MSG01 ; GET FIRST MESSAGE
301C 11C03C 00420 LD OE,3C00H ; GET OISPLAY POSITION
301F C0BF32 00430 CALL OISPLY ; AND OISPLAY ON SCREEN
3022 C0C331 00440 CALL INPUT ; GET INPUT FROM KEYBOARD
3025 FE63 00450 CP 63H ; IS IT A LETTER "C"?
3027 2B13 00460 JR Z,C0LO ; GO TO COLOSTART ROUTINE
3029 FE64 00470 CP 64H ; IS IT A LETTER "O"?
302B CA0000 00480 JP Z,0000 ; GO TO OISK REBOOT
302E FE60 00490 CP 60H ; IS IT A LETTER "H"?
3030 2B30 00500 JR Z,MONITOR ; GO TO MACHINE MONITOR
3032 FE75 00510 CP 75H ; IS IT A LETTER "U"?
3034 2B22 00520 JR Z,USER ; GO TO USER ROUTINE
3036 FE77 00530 CP 77H ; IS IT A LETTER "W"?
303B 2B2A 00540 JR Z,WARM ; GO TO BASIC WARMSTART
303A 1BE6 00550 JR MENU ; LOOP BACK IF NONE
00560 ;
00570 ; #####
00580 ; COLOSTART ROUTINE DUPLICATES LEVEL II ACTION TO 0075H.
00590 ; STACK POINTER NOT RESET HERE BECAUSE COLOSTART DOES IT.
00600 ; #####
00610 ;
303C F3 00620 COLO OI ; INTERRUPTS TURNED OFF
3030 AF 00630 XOR A ; CLEAR ACCUMULATOR
303E 03FF 00640 CLEAR OUT (OFFH),A ; BEGIN WRITE-PORT LOOP
3040 210206 00650 LO HL,0602H ; GET COMMUNICATION BLOCK
3043 110040 00660 LO OE,4000H ; GET COMMUNICATION AREA
3046 013600 00670 LO BC,0036H ; GET SIZE OF COMM. BLOCK
3049 E0B0 00680 LOIR ; BLOCK MOVE COMM. AREA
304B 30 00690 DEC A ; A = A-1 (A=0 AT START)
304C 30 00700 DEC A ; A = A-1 (FE, FC, FO...)
3040 20EF 00710 JR NZ,CLEAR ; CLEAR ALL EVEN PORTS
304F 0627 00720 LO B,27H ; GET 39 VAR. LOCATIONS
3051 12 00730 LOOP1 LD (OE),A ; BEGIN SETTING UP CHARS
3052 13 00740 INC OE ; GET NEXT MEM POS'N
3053 10FC 00750 QJNZ LOOP1 ; LOOP THROUGH THEM ALL
3055 C37500 00760 JP 0075H ; REST OF BASIC STARTUP
00770 ;
00780 ; #####
00790 ; USER ENTRY ROUTINE REQUIRES THAT A TWO-BYTE ADDRESS B
00800 ; STORED IN MEMORY AT 40BE/40BF (16526/16527), <USR(0)>
00810 ; #####
00820 ;
305B E1 00830 USER POP HL ; BEGIN TO CLEAR THE
3059 01 00840 POP OE ; STACK OF ALL
305A C1 00850 POP BC ; REGISTERS PUSHEO
305B F1 00860 POP AF ; AT ENTRY POINT
305C E07BE242 00870 LD SP,(42E2H) ; GET ORIGINAL STACK PTR
3060 2ABE40 00880 LO HL,(40BEH) ; GET USR(0) ENTRY POINT
3063 E9 00890 JP (HL) ; AND JUMP TO IT
00900 ;
00910 ; #####
00920 ; WARM START ROUTINE RESTORES STACK BEFORE GOING TO BASIC
00930 ; #####
00940 ;
3064 E1 00950 WARM POP HL ; RESTORE ALL REGISTERS
3065 01 00960 POP OE ; PUSHEO ON STACK
3066 C1 00970 POP BC ; WHEN MONITOR
3067 F1 00980 POP AF ; WAS ENTEREO
306B E07BE242 00990 LD SP,(42E2H) ; GET ORIGINAL STACK PTR
306C C3CC06 01000 JP 06CCH ; GO TO BASIC WARMSTART
01010 ;

```

Listing 6-2. Power-up monitor program.

How Interpreters Work

The concept of computer languages is far too complex for a single section, a single chapter or even a single book. They run from easy, messy, but capable languages like BASIC through hard and messy languages like FORTRAN, to neat, structured, but obtuse languages like Pascal and LISP. In between are pseudo BASICs like BASEX, plus a whole range of compiled languages and hybrid self-definers like FORTH.

All these languages have one thing in common: they must eventually be broken into subroutines which operate in the machine language of the host processor. Compiled languages are broken into these subroutines when the program is written, but the bulk of material installed in small computers is accessed most easily by means of interpreters - executing one statement at a time, during a program's run.

How does a computer get to the information to be interpreted or compiled? Here's a short rundown of how an interpreter does its work; compilers are similar, but they won't be covered here.

Once a program has been constructed (and the method varies from machine to machine) it is in place for the computer to evaluate and execute. The process that follows is consistent with most interpreters:

Once a program has been constructed, and the method varies from machine to machine, it is in place for the computer to evaluate and execute. The process that follows is consistent with most interpreters:

1. Upon an execution command, the interpreter identifies the start of the program.
2. The first command from the program is obtained. The interpreter compares this command, byte for byte, against a table of legitimate commands. In some interpreters, the commands are stored as full words; in others, they are tokenized. When a valid entry is found in the command table, this stage of interpretation is complete. If a valid entry is not found, the routine is exited, usually via a loop which can be intercepted by extensions to the interpreter.


```

01020 ; #####
01030 ; MACHINE LANGUAGE MONITOR IS SIMPLE, DISPLAYING REGISTER
01040 ; CONDITIONS AT ENTRY, AND ACCEPTING MEMORY CHANGES. ALL
01050 ; REGISTERS ARE RESTORED UPON EXIT FROM THIS ROUTINE.
01060 ; #####
01070 ;
306F 210632 01080 MONITOR LD HL,MSG02 ; GET REGISTER DISPLAY
3072 11CD3C 01090 LD DE,3CC0H ; GET SCREEN POSITION
3075 00BF32 01100 CALL DISPLY ; DISPLAY THE MESSAGE
307B 11DD30 01110 LD DE,300DH ; GET SCREEN POSITION
3078 ED4BE242 01120 LD BC,(42E2H) ; GET STACK POINTER POS'N
307F C5 01130 PUSH BC ; GET VALUE READY TO XFER
308D E1 01140 PDP HL ; TRANSFER BYTES TO HL
30B1 CD1731 01150 CALL WORDER ; CONVERT & DISPLAY BYTE
30B4 00E5 01160 PUSH IX ; GET IT READY TO USE
30B6 E1 01170 PDP HL ; AND DO IT WITH IT
30B7 CD1731 01180 CALL WORDER ; CONVERT & DISPLAY BYTES
30BA FDE5 01190 PUSH IY ; GET IT READY TO USE
30BC E1 01200 PDP HL ; AND DO IT WITH IT
30BD CD1731 01210 CALL WORDER ; CONVERT & DISPLAY IT
309D 2A0E42 01220 LD HL,(420EH) ; POINT TO AF STACK VALUE
3093 CD1731 01230 CALL WORDER ; CONVERT & DISPLAY BYTE
3096 2A0C42 01240 LD HL,(420CH) ; POINT TO BC STACK VALUE
3099 CD1731 01250 CALL WORDER ; CONVERT & DISPLAY BYTE
309C 2A0A42 01260 LD HL,(420AH) ; POINT TO DE STACK VALUE
309F CD1731 01270 CALL WORDER ; CONVERT & DISPLAY BYTE
30A2 2A0842 01280 LD HL,(4208H) ; POINT TO HL STACK VALUE
30A5 CD1731 01290 CALL WORDER ; CONVERT & DISPLAY BYTE
30AB ED5F 01300 LD A,R ; GET R REGISTER INTO A
30AA 67 01310 LD H,A ; GIVE IT FOR CONVERSION
30AB EDS7 01320 LD A,I ; GET I REGISTER INTO A
30A0 6F 01330 LD L,A ; GIVE IT FOR CONVERSION
30AE CD1731 01340 CALL WORDER ; CONVERT & DISPLAY BYTE
30B1 118030 01350 LD DE,3080H ; GET NEW VIDEO POSITION
30B4 08 01360 EX AF,AF' ; TRANSFER OTHER VALUE
30B5 FS 01370 PUSH AF ; READY TO XFER TO HL
30B6 E1 01380 PDP HL ; EFFECT TRANSFER TO HL
30B7 DB 01390 EX AF,AF' ; RETURN ORIGINAL VALUE
30BB CD1731 01400 CALL WORDER ; CONVERT & DISPLAY BYTE
30BB D9 01410 EXX ; TRANSFER BC, DE, HL
30B8 ES 01420 PUSH HL ; FIRST SLIP BC ON STACK.
30B0 05 01430 PUSH OE ; THEN PUSH OE THERE,
30BE CS 01440 PUSH BC ; AND THEN THE LAST ONE.
30B8 D9 01450 EXX ; TRANSFER REGISTERS BACK
30C0 E1 01460 PDP HL ; GET VALUE ALL READY
30C1 CD1731 01470 CALL WORDER ; CONVERT & DISPLAY BYTE
30C4 E1 01480 POP HL ; GET DE VALUE READY
30C5 CD1731 01490 CALL WORDER ; CONVERT & DISPLAY BYTE
30CB E1 01500 POP HL ; GET HL VALUE READY
30C9 CD1731 01510 CALL WORDER ; CONVERT & DISPLAY BYTE
30CC 11DD3F 01520 CHECK LD DE,3F00H ; GET NEW SCREEN POS'N
30CF CD1731 01530 CALL WORDER ; CONVERT & DISPLAY BYTE
30D2 1C 01540 INC E ; BUMP SCREEN POSITION
30D3 7E 01550 LD A,(HL) ; GET VALUE FROM MEMORY
3004 C02431 01560 CALL HEXASC ; AND DISPLAY IT IN ASCII
01570 ;
01580 ; #####
01590 ; ADDRESS MOD ROUTINE GETS KEYBOARD VALUE & LOOPS BACK
01600 ; #####
01610 ;
3007 CDC331 01620 AODMOD CALL INPUT ; GET VALUE FROM KEYBOARD
300A FE21 01630 CP '!' ; TEST FIRST IF EXECUTE
30DC 2832 01640 JR Z,EXEC ; OUT IF EXECUTE COMMAND
30DE FE2F 01650 CP '/' ; TEST SECOND IF DATA
30ED CA6E31 01660 JP Z,DATA00 ; OUT IF DATA MOD COMMAND
30E3 FE2A 01670 CP '*' ; TEST THIRD IF RET MENU
30E5 CADC3D 01680 JP Z,CLS ; OUT TO MENU IF A STAR
30E8 FE30 01690 CP '0' ; SEE IF <0 CHARACTER
30EA 3BE8 01700 JR C,AODMOD ; LOOP BACK IF <0 CHAR.
30EC FE67 01710 CP 67H ; SEE IF >F CHARACTER
30EE 30E7 01720 JR NC,AODMOD ; LOOP BACK IF >F CHAR.
30F0 FE3A 01730 CP ':' ; SEE IF <9 CHARACTER
30F2 3804 01740 JR C,NUMBER ; GO TO NUMBER ROUTINE
30F4 FE61 01750 CP 61H ; SEE IF >A CHARACTER
30F6 38DF 01760 JR C,AODMOD ; LOOP BACK IF <A & >9
30F8 21D43F 01770 NUMBER LD HL,3F04H ; GET ADDRESS SCREEN POS'N
30FB 11033F 01780 LD DE,3F03H ; GET NEXT SCREEN POS'N
30FE D1D300 01790 LD BC,3 ; GET THREE TOTAL POS'NS
3101 ED80 01800 LDIR ; AND MOVE THEM OVER
3103 2B 01810 DEC HL ; REPOSITION TO LAST CHAR
3104 FE60 01820 CP 60H ; COMPARE TO L.C. ALPHA
3106 3802 01830 JR C,ZIPBY ; IF NUMERIC, THEN SKIP
3108 0620 01840 SUB 20H ; ELSE CONVERT TO U.CASE
31DA 77 01850 ZIPBY LD (HL),A ; DISPLAY NEW CHARACTER
3108 C04331 01860 NUM2 CALL ASCHEX ; CONVERT DISPLAY TO HL
310E 188C 01870 JR CHECK ; AND LOOP BACK FOR MORE
3110 21063F 01880 EXEC LD HL,3F06H ; POINT HL TO ADDRESS
3113 C04331 01890 CALL ASCHEX ; AND CONVERT TO HEX
3116 E9 01900 JP (HL) ; HL CONTAINS ADDRESS
01910 ;
01920 ; #####
01930 ; SUBROUTINE RESPONSIBLE FOR CONVERTING & DISPLAYING TWO-
01940 ; BYTE WORDS DELIVERED TO IT IN THE HL REGISTER.

```

Listing Continued

3. If the command is identified as a legitimate one, a subroutine is called which executes the command. That subroutine in turn may further examine the command line for operands and conditions, incrementing and decrementing pointers in its search for required, and valid, information. Further subroutines are entered as necessary to evaluate and put to use this additional information. In-memory variables and pointers are set up, modified, and accessed by all subroutines, usually from a master table which defines variable types and syntax conditions. Transcendental functions are also accessed via tables within the interpreter itself. Error checking is done at all points.

4. When the execution of each subroutine is completed, it returns to the calling program. Eventually, all subroutines have returned to their upper level of subroutine 'nesting'. Then the execution routine finds itself re-entered, positioned at the next executable point in the program, where the execution process is repeated.

5. The execution routine may, depending on the language, find itself repositioned in the program out of normal execution sequence. On the other hand, some languages have an inherent structure which disallows any repositioning, demanding an inviolable first-to-last execution sequence. In these latter interpreters, repositioning will be interpreted as an error condition.

6. If any information, commands, program order or program syntax is incorrect, an error handling routine is entered, usually by direct jump rather than subroutine call. Normally, this routine prepares and presents an error message. It returns the program from execution level, arbitrarily cancelling any nested execution subroutines by readjusting stack pointers and other variable information.

This readjustment is necessary to avoid unsettling the user-interactive command-level routine, and causing the processor to lose its way in a complex of incomplete subroutines.

7. Upon completion of all program statements, the program is returned from execution level to command level. Some interpreters allow commands to be entered, interpreted, and executed from a command buffer, without actually entering a program execution condition.

Continued Listing

```

01950 ; #####
01960 ;
3117 1C 01970 WORDER INC E ; BUMP OE REGISTER ALONG
3118 1C 01980 INC E ; SOME MORE BUMPING ...
3119 1C 01990 INC E ; ... AND SOME MORE OF IT
311A 7C 02000 LD A,H ; GET FIRST BYTE OF IT
311B C02431 02010 CALL HEXASC ; CONVERT & DISPLAY BYTE
311E 70 02020 LD A,L ; GET SECOND BYTE OF IT
311F C02431 02030 CALL HEXASC ; CONVERT & DISPLAY BYTE
3122 1C 02040 INC E ; AND POSITION SCREEN
3123 C9 02050 RET ; AND BACK TO CALLER

02060 ;
02070 ; #####
02080 ; HEXADECIMAL TO ASCII CONVERSION FOR TWO-BYTE WORDS
02090 ; #####
02100 ;
3124 F5 02110 HEXASC PUSH AF ; SAVE BYTE IN AF REG.
3125 E6F0 02120 AND OF0H ; AND MASK OUT LOW NYBBLE
3127 0F 02130 RRCA ; BEGIN ROTATING BIT ...
3128 0F 02140 RRCA ; ... IN ORDER TO TEST...
3129 0F 02150 RRCA ; ... THE LOW NYBBLE ...
312A 0F 02160 RRCA ; ... ALL BY ITSELF.
312B C03931 02170 CALL HXAC ; BYTE RANGE EVALUATION
312E 12 02180 LD (OE),A ; DISPLAY ON THE SCREEN
312F 1C 02190 INC E ; AND MOVE TO NEXT POS'N
3130 F1 02200 POP AF ; RESTORE ORIGINAL BYTE
3131 E60F 02210 AND OFH ; AND MASK OUT HIGH BITS
3133 C03931 02220 CALL HXAC ; RANGE EVALUATION ROUT.
3136 12 02230 LD (OE),A ; AND DISPLAY ON SCREEN
3137 1C 02240 INC E ; MOVE TO NEXT VIDEO POSN
3138 C9 02250 RET ; AND BACK TO CALLER

02260 ;
02270 ; #####
02280 ; HEX-ASCII RANGE EVALUATION FOR 0 TO 9 AND A TO F AREAS
02290 ; #####
02300 ;
02310 ;
3139 FE0A 02320 HXAC CP 0AH ; CHECK AGAINST 10 DEC.
313B 3003 02330 JR NC,NEXTX ; IF >10 THEN ALPHA CHAR
3130 C630 02340 A00 A,30H ; ELSE NUMERIC - CONVERT
313F C9 02350 RET ; AND BACK TO CALLER
3140 C637 02360 NEXTX A00 A,37H ; ALPHABETIC - CONVERT
3142 C9 02370 RET ; AND BACK TO CALLER

02380 ;
02390 ; #####
02400 ; ASCII TO HEXADECIMAL CONVERSION FOR TWO-BYTE WORDS
02410 ; #####
02420 ;
02430 ;
3143 C06331 02440 ASCHEX CALL ACHX ; RANGE EVALUATION ROUT.
3146 4F 02450 LD C,A ; SAVE PART OF CONVERSION
3147 2B 02460 DEC HL ; AND GET NEXT BYTE
3148 C05B31 02470 CALL LLLLS ; EXECUTE LEFT ROTATES
314B B1 02480 A00 A,C ; AND A00 TO MAKE BYTE
314C 4F 02490 LD C,A ; SAVE IT BACK IN C REG.
3140 2B 02500 DEC HL ; AND MOVE TO NEXT BYTE
314E C06331 02510 AX2 CALL ACHX ; 00 RANGE EVALUATION
3151 47 02520 LD B,A ; AND SAVE IT IN B REG.
3152 2B 02530 DEC HL ; GET FINAL BYTE READY
3153 C05B31 02540 CALL LLLLS ; AND EXECUTE LEFT ROTATE
3156 B0 02550 A00 A,B ; CREATE COMPLETE BYTE
3157 47 02560 LD B,A ; AND SAVE IT IN B REG.
315B C5 02570 PUSH BC ; GET 2 BYTES TO XFER
3159 E1 02580 POP HL ; TRANSFER TO HL REGISTER
315A C9 02590 RET ; AND BACK TO CALLER
315B C06331 02600 LLLLS CALL ACHX ; EVALUATE RANGE OF CHAR
315E 07 02610 RLCA ; BEGIN LEFT ROTATES ...
315F 07 02620 RLCA ; ... WHICH WILL POS'N...
3160 07 02630 RLCA ; ... BYTE READY FOR ...
3161 07 02640 RLCA ; ... CONVERSION.
3162 C9 02650 RET ; AND GO BACK TO CALLER
3163 7E 02660 ACHX LD A,(HL) ; GET BYTE FROM SCREEN
3164 FE40 02670 CP 40H ; CHECK AGAINST ALPHA
3166 3003 02680 JR NC,NEXTZ ; IF NUMERIC THEN JUMP
316B 0630 02690 SUB 30H ; ELSE NUMBER TO HEX
316A C9 02700 RET ; AND BACK TO CALLER
316B 0637 02710 NEXTZ SUB 37H ; THEN IT'S ALPHA TO HEX
3160 C9 02720 RET ; AND BACK TO CALLER

02730 ;
02740 ; #####
02750 ; ROUTINE TO ACCEPT DATA FOR INPUT TO DISPLAYED MEMORY
02760 ; LOCATIONS, WITH EXIT TO MENU UPDATE OR MEMORY INCREMENT
02770 ; #####
02780 ;
316E C0C331 02790 OATMOO CALL INPUT ; GET CHAR. FROM KEYBRO.
3171 FE2E 02800 CP ' ' ; IS IT A PERIOD?
3173 C8031 02810 JP Z,A0EXIT ; IF SO, EXIT TO A0OR MOO
3176 FE00 02820 CP 00H ; IS IT A CARRIAGE RET?
317B 2B2F 02830 JR Z,NEXT0 ; IF SO, GO TO NEXT DATA
317A FE30 02840 CP 'G' ; BEGIN TESTING FOR RANGE
317C 3BFO 02850 JR C,OATMOO ; BACK IF <0 CHARACTER
317E FE67 02860 CP 67H ; CHECK AGAINST LC ALPHA
3180 30EC 02870 JR NC,OATMOO ; IF >F THEN GO BACK

```

Listing Continued . . .

It was natural that Radio Shack would choose an inexpensive storage medium, cassettes, to accompany their low cost microcomputer. Because ordinary cassette players are audio devices, the tape saving and loading routines were designed to be slow but sure. With care, the CTR series of recorders can be as reliable as any other storage system designed for the TRS-80.

The weakness of the tape process comes from the obvious mismatching of an audio device, of very limited precision, with a digital device of unyielding high-precision. Portable cassette recorders are intended to reproduce audio signals with a reasonable level of fidelity. What constitutes a reasonable level of fidelity is disputable, and only a person with a true tin ear would not be able to pick out a portable player, from amongst a group of high fidelity tape decks.

But even with this 'reasonable' fidelity, much of what we recognize as harmonies and instrumentation is perceptible only because we already have an acculturated comprehension of sound; and this directly influences what we believe we are hearing. Our minds, in conjunction with our ears - average, fill in, smooth over and forgive minuscule failings. We have internal mechanisms which remember our experiences.

The cassette load/save system consists of seven major elements:

1. Serialization.

The individual bytes of computer data are converted into a stream of individual bits. This is a completely digital process, and the timing is provided by the computer.

2. Audio Processing.

The signal is converted into a 'digital audio' wave for recording on tape decks of unknown polarity. In other words, a digital, one to zero, signal is converted to an audio, one to minus one to zero, signal. In this way, an 'upside down' signal looks the same to a computer as the original.

3. Recording.

The signal goes through the tape recorder's electronics, and is recorded on a thin strip of magnetic tape. The audio electronics round the wave's edge, and the limitations of the tape contribute noise to the signal.

Continued Listing

```

31B2 FE3A      028B0      CP      '!'      ; TEST FOR NUMERIC >9
31B4 3B04      02890      JR      C,INQATA ; IF <9 THEN GET DATA
31B6 FEB1      02900      CP      61H      ; TEST FOR ALPHA <A
31B8 3BE4      02910      JR      C,0ATM00 ; IF <A THEN GO BACK
31BA 210A3F    02920      LD      HL,3FOAH ; READY THE SCREEN PTR.
31B0 F5        02930      PUSH     AF      ; SAVE VALUE FROM KEYBRO
31BE 7E        02940      LD      A,(HL) ; GET LOWER ASCII NYBBLE
31BF 32093F    02950      LD      [3F09H],A ; AND PUT IN HIGH POS'N
3192 F1        02960      POP      AF      ; RESTORE VALUE FROM KB0
3193 FE60      02970      CP      60H      ; CHECK AGAINST LC ALPHA
3195 3B02      02980      JR      C,ZIPZY ; IF LC, THEN SKIP PAST
3197 D620      02990      SUB      20H      ; ELSE CONVERT TO U.C.
3199 77        03000      LD      (HL),A ; AND DISPLAY ON SCREEN
319A C04E31    03010      CALL    AX2 ; SINGLE BYTE CONVERSION
3190 7C        03020      LD      A,H      ; SAVE ONE BYTE IN A REG.
319E F5        03030      PUSH     AF      ; SAVE THAT NOW ON STACK
319F 21063F    03040      LD      HL,3F06H ; POINT HL TO SCREEN POSN
31A2 C04331    03050      CALL    ASCHEX ; EVALUATE ADDRESS THERE
31A5 F1        03060      POP      AF      ; RESTORE VALUE TO SHOW
31A6 77        03070      LD      (HL),A ; AND PUT ON THE SCREEN
31A7 1B07      03080      JR      EXIT1 ; NOW GO OUT AND STORE
31A9 21063F    03090      LD      HL,3F06H ; POINT HL TO SCREEN POSN
31AC C04331    03100      CALL    ASCHEX ; CONVERT ADDRESS TO HEX
31AF 23        03110      INC      HL ; HOVE OVER TO NEXT POSN
31B0 11003F    03120      LD      DE,3F00H ; POINT DE ADDR. DISPLAY
31B3 C01731    03130      CALL    WORDER ; AND CONVERT/DISPLAY IT
31B6 1C        03140      INC      E ; GO TO NEXT SCREEN POSN
31B7 7E        03150      LD      A,(HL) ; GET VALUE FROM MEMORY
31B8 C02431    03160      CALL    HEXASC ; AND CONVERT FOR DISPLAY
31B8 1B81      03170      JR      DATM00 ; AND BACK FOR SOME MORE
31B0 21063F    03180      LD      HL,3F06H ; POINT TO PRESENT ADDR.
31C0 C30B31    03190      JP      NUM2 ; AND BACK TO UPDATE MENU
31C0          03200      ;
31C0          03210      ; #####
31C0          03220      ; COMPLETE KEYBOARD ROUTINE BELOW MAY BE CALLED BY NOT
31C0          03230      ; ONLY THE MONITOR ROUTINES, BUT ANY ROUTINES NEEDING AN
31C0          03240      ; UPPER/LOWER CASE, AUTOREPEAT, BEEP DRIVER. FOR FURTHER
31C0          03250      ; DESCRIPTION, SEE SUPPLEMENT TO CHAPTER ON KEYBOARD I/O
31C0          03260      ; #####
31C0          03270      ;
31C3 213640    03280      LD      HL,4036H ; SAME AS LII BUFFER
31C6 01013B    03290      LD      BC,3B01H ; FIRST KEYBOARD ROW
31C9 1600      03300      LD      D,0 ; COUNTER FOR COLUMNS
31CB 0A        03310      LD      A,(BC) ; RETRIEVE ROW CONTENTS
31CC 5F        03320      LD      E,A ; SAVE IT TEMPORARILY
31CD A3        03330      AND      E ; SET FLAGS FOR TESTING
31CE 201B      03340      JR      NZ,STRDKE ; NOT ZERO IF KEY PRESSED
31D0 77        03350      LD      (HL),A ; SAVE CURRENT VALUE
31D1 14        03360      INC      D ; INCREMENT ROW COUNTER
31D2 2C        03370      INC      L ; INCREMENT STORAGE ADDR.
31D3 CB01      03380      RLC      C ; GET NEXT KEYBOARD COL.
31D5 79        03390      LD      A,C ; GET VALUE INTO ACCUM.
31D6 06B0      03400      SUB      B0H ; LAST ROW IS 3BBD HEX
31DB 20F1      03410      JR      NZ,KEYPRS ; NEXT CHECK IF NOT DONE
31DA 0607      03420      LD      B,7 ; COUNTER OF KEYBOARD ROWS
31DC 20        03430      LD      L,CLRMEM ; START COUNTING BACKWARD
31DD 86        03440      ADD      A,(HL) ; AND ADD IT UP IN ACCUM.
31DE 10FC      03450      LD      CLRMEM ; AND DO IT FOR 7 ROWS
31E0 A7        03460      AND      A ; TEST FOR ANY KEY DOWN
31E1 3E00      03470      LD      A,0 ; A=0, FLAGS ARE INTACT
31E3 C0        03480      RET      NZ ; BACK IF KEYS IN USE
31E4 321A40    03490      LD      (KPLACE),A ; ELSE DELAY GETS RESET
31E7 C9        03500      RET ; AND GO BACK ANYWAY
31E8 A6        03510      AND      (HL) ; CHECK KEYSTROKE STORAGE
31E9 2B1B      03520      JR      Z,FOUND ; NEW KEY IF NOT SAME ONE
31EB 3A1A40    03530      LD      A,(KPLACE) ; NEW CHECK SPECIAL STORE
31EE 3C        03540      INC      A ; LET STORE = STORE + 1
31EF 321A40    03550      LD      (KPLACE),A ; AND PUT IN BACK THERE
31F2 FEFF      03560      CP      0FFH ; CHECK IF IT IS AT END
31F4 2B0B      03570      JR      Z,DECA ; IF SO, THEN HDLO THERE
31F6 C5        03580      PUSH     BC ; SAVE ROW COUNTER REG.
31F7 06FF      03590      LD      B,0FFH ; GET DELAY VALUE INTO B
31F9 10FE      03600      DJNZ     TMWSTE ; AND DELAY JUST A BIT
31FB C1        03610      POP      BC ; RESTORE ROW COUNTER
31FC 1B03      03620      JR      RECCHK ; AND BACK TO CHECK NEXT
31FE 3D        03630      DECA     DEC A ; LET A = A - 1 (STORE)
31FF 321A40    03640      LD      (KPLACE),A ; AND PUT IT IN STORAGE
3202 7B        03650      LD      A,E ; GET KEYBOARD BYTE BACK
3203 73        03660      LD      (HL),E ; STORE IT IS STROKE AREA
3204 7A        03670      LD      A,0 ; GET ROW COUNTER FROM D
3205 07        03680      RLC      A ; AND BEGIN PROCCSS DF...
3206 07        03690      RLC      A ; ...CONVERTING IT...
3207 07        03700      RLC      A ; ...TO AN OFFSET VALUE
3208 57        03710      LD      D,A ; AND PUT IT BACK IN D
3209 0E01      03720      LD      C,1 ; GET NUMBER ONE READY
320B 79        03730      LD      A,C ; ACCUM. HAS C FOR MASK
320C A3        03740      AND      E ; TEST IF C = KEYSTROKE
320D 2005      03750      JR      NZ,AROUND ; IF NOT, THEN GO AROUND
320F 14        03760      INC      D ; ELSE D = ROW + COLUMN
3210 CB01      03770      RLC      C ; C SET TO NEXT COLUMN
3212 1BF7      03780      JR      BACKUP ; GO BACK AND TEST AGAIN
3214 3AB03B    03790      LD      A,(3BBDH) ; GET SHIFT ROW FOR TEST
3217 47        03800      LD      B,A ; AND SAVE IT IN B REG.

```

Listing Continued

4. Storage.

The tape sits on the shelf, affected by temperature and humidity, where its oxide coating may 'creep'. The tape may stretch or buckle or warp, and its upper and lower edges may become slightly feathered.

5. Playback.

The recorded signal, including warps, dropouts, feathering, creep and noise, is fed to the playback electronics. This audio circuit contributes further noise, providing a purely low grade audio signal to the computer.

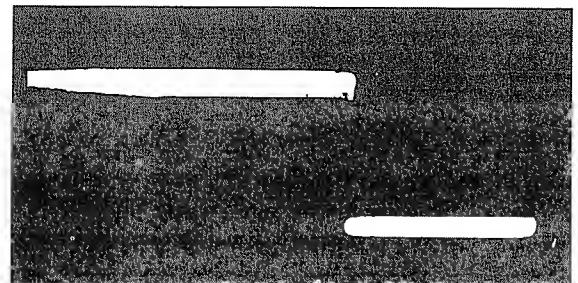
6. Digital Processing.

The signal is received, its top and bottom edges are squared, and it is returned to the one to zero digital state. The timing is provided by the recorded tape in cooperation with timing loops provided by the computer.

7. De-Serialization.

A completed group of bits is assembled into an 8 bit byte for use by the CPU in determining synchronization, type of program, loading location, etc.

An oscilloscope representation of a digital signal being generated during the CSAVE routine is shown below. The signal changes from one to zero very crisply, spending only a few nanoseconds (not visible at all on the photo) in the transition area between one and zero. This true digital signal is measured just before the audio processing section sent to the cassette.



The audio processing is actually no more than a digital signal whose polarity is reversed. Two outputs of a latch are tied together, and as one goes high, the other is forced low. Between pulses, the signal floats to the middle. This is done by data bits 0 and 1 of Z59, which are alternately switched by the CPU during CSAVE. The digital signal as it leaves the cassette port is shown in the second photograph.

The sharp edges of the signal have been ever so slightly blurred, partly due to the capacitance

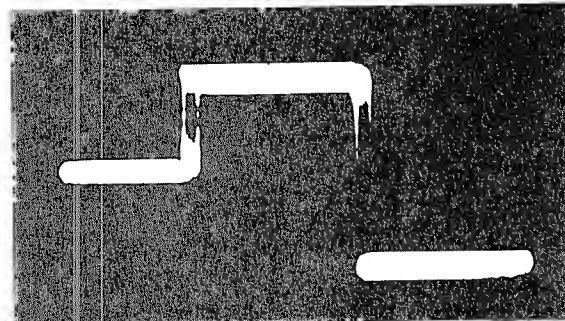
The CLOAD Problem

Continued Listing

```

3218 7A 03810 LD A,0 ; GET ROW COUNTER BACK
3219 C640 03820 ADD A,40H ; AND CONVERT TO ASCII
3218 FE60 03830 CP 60H ; IS IT UP/LW/GRAFIX/CTRL
3210 3016 03840 JR NC,JUMP1 ; GO OUT IF GRAPHICS
321F S7 03850 LD 0,A ; SAVE PARTLY CONVERTED
3220 3A40B8 03860 LD A,(3840H) ; GET VALUE FOUND 7TH ROW
3223 E610 03870 ANO 10H ; CHECK IF DOWN ARROW
3225 2009 03880 JR NZ,CTRL ; IF SO, PRODUCE CONTROL
3227 7A 03890 LD A,0 ; ELSE GET VALUE BACK
3228 C808 03900 RRC B ; B BUMPS INTO CARRY FLAG
322A 3830 03910 JR C,GOAWAY ; IF CARRY, THEN SHIFT
322C C620 03920 ADD A,20H ; IF NOT THEN LOWER CASE
322E 1839 03930 JR GOAWAY ; AND GET OUT OF ROUTINE
3230 7A- 03940 CTRL LO A,0 ; IF CONTROL CODE, GET IT
3231 0640 03950 SUB 40H ; GET RID OF ASCII MASK
3233 1834 03960 JR GOAWAY ; AND GET OUT OF ROUTINE
3235 0670 03970 JUMP1 SUB 70H ; THE BALANCE OF THE
3237 3010 03980 JR NC,JUMP2 ; ROUTINE BELOW UP TO
3239 C640 03990 ADD A,40H ; THE BEEP SECTION IS
323B FE3C 04000 CP 3CH ; VIRTUALLY IDENTICAL
3230 3802 04010 JR C,JUMP3 ; TO THE KEYBOARD
323F EE10 04020 XOR 10H ; DETERMINATION SUB-
3241 C808 04030 JUMP3 RRC B ; ROUTINE FOUND IN
3243 3024 04040 JR NC,GOAWAY ; RM. THIS ENTIRE
3245 EE10 04050 XOR 10H ; ROUTINE IS REPEATED
3247 1820 04060 JR GOAWAY ; BECAUSE IT CAN
3249 07 04070 JUMP2 RLCA ; SERVE AS AN EXACT
324A C808 04080 RRC B ; REPLACEMENT FOR
324C 3000 04090 JR NC,JUMP4 ; THE LEVEL II BASIC
324E 215832 04100 JUMP4 LD HL,TABLET ; ROUTINES WHEN (IF)
3251 4F 04110 LD C,A ; THE ROMS HAVE BEEN
3252 0600 04120 LD 8,0 ; REMOVED, REPLACED,
3254 09 04130 ADD HL,BC ; DISABLED, ETC.
3255 7E 04140 LD A,(HL) ; SEE KEYBOARD ROUTINE
3256 1811 04150 JR GOAWAY ; RUNDOWN ELSEWHERE
3258 3C- 04160 INC A ; IN THIS BOOK.....
3259 0000 04170 TABLET OEFW 0000H ; CARR. RET. / CARR. RET.
3258 1F1F 04180 OEFW 1F1FH ; CLEAR SCRIN / CLEAR SCRIN
3250 0101 04190 OEFW 0101H ; BREAK KEY / BREAK KEY
325F 581B 04200 OEFW 1858H ; EBIT ESCAPE / UP ARROW
3261 0A00 04210 OEFW 000AH ; NO CHANGE / LINEFEED
3263 0818 04220 OEFW 180BH ; BACKSP. LINE / BACKSP.
3265 0919 04230 OEFW 1909H ; 32-CHAR MODE / HOR TAB
3267 2020 04240 OEFW 2020H ; SPACE / SPACE
3269 57 04250 GOAWAY LO 0,A ; SAVE VALUE IN 0 REG.
326A 018001 04260 LO BC,180H ; OEBOUNCE VALUE AND
3260 08 04270 DELAYS DEC BC ; IMPLEMENTED AT THIS
326E 7B 04280 LO A,B ; POINT FOR ACTUAL
326F 81 04290 OR C ; KEYBOARD SCAN WHEN
3270 20F8 04300 JR NZ,DELAYS ; CHARACTERS ARE
3272 7A 04310 LD A,0 ; GET STORED VALUE BACK
3273 FS 04320 PUSH AF ; SAVE ACCUM. & FLAGS
3274 0640 04330 LD B,40H ; GET BEEP LENGTH VALUE
3276 3A3040 04340 LD A,(4030H) ; GET STATUS OF SCREEN
3279 E6F0 04350 ANO 0FH ; MASK SCREEN CHANGE
327B 67 04360 LO H,A ; STORE MSB IN H REG.
327C F602 04370 OR 2 ; SET BIT 1 TO BE ON
327E 6F 04380 LD L,A ; STORE ALT. MSB IN L REG
327F 70 04390 BEEPER LO A,L ; GET ALT. MSB TO OUTPUT
3280 03FF 04400 OUT (OFFH),A ; AND OUTPUT RISING WAVE
3282 7C 04410 LO A,H ; GET NORMAL MSB CHAR.
3283 03FF 04420 OUT (OFFH),A ; AND OUTPUT FALLING WAVE
3285 CS 04430 PUSH BC ; SAVE NOTE LENGTH VALUE
3286 0640 04440 LO B,40H ; GET FREQUENCY DELAY
3288 10FE 04450 FREQUENCY OJNZ FREQUENCY ; AND WAIT A LITTLE WHILE
328A C1 04460 POP BC ; NOW RESTORE LENGTH VAL.
3288 10F2 04470 OJNZ BEEPER ; AND GO BACK THAT LENGTH
3280 F1 04480 POP AF ; RESTORE ORIGINAL CHAR.
328E C9 04490 RET ; AND BACK TO CALLER
04500 ;
04510 ; #####
04520 ; ROUTINE BELOW TAKES SIMPLE MESSAGES AND DISPLAYS THEM
04530 ; #####
04540 ;
32BF 7E 04550 DISPLAY LO A,(HL) ; GET FIRST MSG CHARACTER
3290 A7 04560 ANO A ; TEST IF 0 OR CHARACTER
3291 CB 04570 RET Z ; BACK TO CALLING ROUTINE
3292 12 04580 LO (OE),A ; DISPLAY THE CHARACTER
3293 23 04590 INC HL ; GET NEXT MSG CHARACTER
3294 13 04600 INC OE ; GET NEXT DISPLAY POS'N
3295 1BF8 04610 JR DISPLAY ; AND TEST NEXT CHARACTER
04620 ;
04630 ; #####
04640 ; MESSAGES FOR DISPLAY FOLLOW. IN A STRIPPED SOFTWARE
04650 ; MONITOR CONFIGURATION, THESE MESSAGES MAY BE ELIMINATED
04660 ; AND DISPLAY PARAMETERS CHANGED ELSEWHERE IN THE LISTING
04670 ; #####
04680 ;
3297 3C 04690 MSG01 OEFM 'C>OLOADSTART, <O>ISKBOOT, <M>ONITOR, '
3298 3C 04700 OEFM 'U>SERJUMP, <W>ARMSTART? '
3204 BF 04710 OEFB BFH
3205 00 04720 OEFB 00
3206 5A 04730 MSG02 OEFM 'ZBO REGISTERS AT MONITOR ENTRY: '

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introduced by the cables and tape input. This is the first (and least significant) step in the extensive route of signal degradation.

```

32F6 20 04740 OEFM '
3316 53 04750 OEFM 'SP= IX= IY= AF=
3336 42 04760 OEFM 'BC= OE= HL= RI=
3386 41 04770 OEFM 'ALTERNATE REGISTERS AT MONITOR E'
3376 4E 04780 OEFM 'NTRY (NOT USED IN LEV II BASIC):'
3396 41 04790 OEFM 'AF= BC= OE= HL=
3388 20 04800 OEFM '
3306 53 04810 OEFM 'USE PERIOD {.} TO ENTER AN ADDR'
33F6 53 04820 OEFM 'SS, SLASH (/) TO ENTER DATA. TO'
3416 45 04830 OEFM 'EXIT TO MONITOR MENU, TYPE STAR'
3436 2B 04840 OEFM '(*), TO EXECUTE AT ADDRESS SHOWN'
3456 4F 04850 OEFM 'ON SCREEN, TYPE EXCLAMATION POIN'
3476 54 04860 OEFM 'T {}'
3486 00 04870 OEFB 00
3497 04880 XXXXXX EQU $
04890 ;
0000 04900 ENO
ACHX 3163 02660 02440 02510 02600
AODMOO 3007 01620 01700 01720 01760
AOEXIT 3180 03180 02810
AROUND 3214 03790 03750
ASCHEX 3143 02440 01860 01890 03050 03100
AX2 314E 02510 03010
BACKUP 3208 03730 03780
BEEPER 327F 04390 04470
CHECK 30CC 01520 01870
CLEAR 303E 00640 00710
CLRMEM 310C 03430 03450
CLS 300C 00310 01680
CNTRL 3230 03940 03880
COL 303C 00620 00460
DATMOO 316E 02790 01660 02850 02870 02910 03170
DECA 31FE 03630 03570
DELAYS 3260 04270 04300
DISPLY 328F 04550 00430 01100 04610
EXEC 3110 01880 01640
EXIT1 3180 03120 03080
FOUND 3203 03660 03520
FREQUENCY 3288 04450 04450
GOAWAY 3269 04250 03910 03930 03960 04040 04060 04150
HEXASC 3124 02110 01560 02010 02030 03160
HXAC 3139 02320 02170 02220
INGATA 318A 02920 02890
INPUT 31C3 03280 00440 01620 02790
JUMP1 3235 03970 03840
JUMP2 3249 04070 03980
JUMP3 3241 04030 04010
JUMP4 324E 04100 04090
KEYPRS 31CB 03310 03410
KPLACE 401A 00160 03490 03530 03550 03640
LLLLS 3158 02600 02470 02540
LOOP1 3051 00730 0075U
MENU 3022 00440 00550
MONTOR 306F 01080 00500
MSG01 3297 04690 00410
MSG02 3206 04730 01080
NEXT0 31A9 03090 02830
NEXTX 3140 02360 02330
NEXTZ 31EB 02710 02590
NUM2 3108 01860 03190
NUMBER 30FB 01770 01740
RECHKEK 3101 03360 03620
STROKE 31EB 03510 03340
TABLET 3259 04170 04100
THWSTE 31F9 03600 03600
USER 3058 00830 00520
WARM 3064 00950 00540
WORDER 3117 01970 01150 01180 01210 01230 01250 01270 01280
01340 01400 01470 01490 01510 01530 03130
XXXXXX 3497 04880
ZIPBY 310A 01850 01830
ZIPZY 3199 03000 02980

```

The worst signal abuse takes place during the taping process. The reasonably digital signal is recorded with the poor electronics of a portable tape recorder, and the sharp-edged waves are rounded off by the natural limitations of the tape itself; examine Photo 6-4. Also visible in the photo are residual noise (tape hiss) and the high-frequency recording bias signal.

There is also an unexpected interreaction between the computer's output wiring and some tape recorders that produces a low-pitched hum, called a ground loop. The good data signal can ride on this ground-loop hum to result in sensitive volume settings — too high or low a volume during playback will cause the top or bottom of the digital waveform to travel out of the range of the digital processing.

This digital processing redeems quite a bit of the audio signal, turning it back into a usable digital waveform much of the time. Photo 6-5 shows the results of the reshaping process; it's a fairly good signal that the CPU finally receives and must interpret as data.

Now the names and descriptions of the seven CLOAD culprits:

1. Head misalignment. This is the main cause of bad loads, because misalignment severely cuts the essential high frequencies. The CTR-80 already has a provision for adjusting the playback head. If you have a CTR-41 (or other recorder), drill a hole directly over the playback head adjustment screw (under 1. Head misalignment. This is the main cause of bad loads, because misalignment severely cuts the essential high frequencies. The CTR-80 already has a provision for adjusting the playback head. If you have a CTR-41 (or other recorder), drill a hole directly over the playback head adjustment screw (under the letters ERY in 'battery' on the CTR-41 face plate), and adjust with a small Phillips screwdriver.

2. Speed variations. This one is not obvious, but note that a five percent variation in recorder speed can cause a bad load, especially with long BASIC lines during CLOAD.

3. Bad tape. There is no reason to use low-grade tape, just as there is no reason to buy the best audio tape. Get a good commercial grade, and standardize with it.

4. Dirty head. Clarity and volume are cut down when the head is dirty. Clean it and all parts that contact the tape with isopropyl alcohol.

5. Starting too soon. The beginning of most tape — especially leaderless tape — is often slightly crumpled, and data can be lost right at the start. Count off ten seconds.

6. Magnetized head. Those who depend on cassettes will use the machine often, and the head will build up residual magnetism. Obtain a cassette demagnetizer (degausser) and use it often.

7. Software. Early Level II BASIC ROMs have problems because the timing loops were not written ideally for low-grade audio use. These can be upgraded with new ROMs or the Radio Shack XRX cassette modification. Note, however, that removal of the XRX modification is necessary for use with high/low speed hardware modifications (except the Archbold 1981-82 kit).

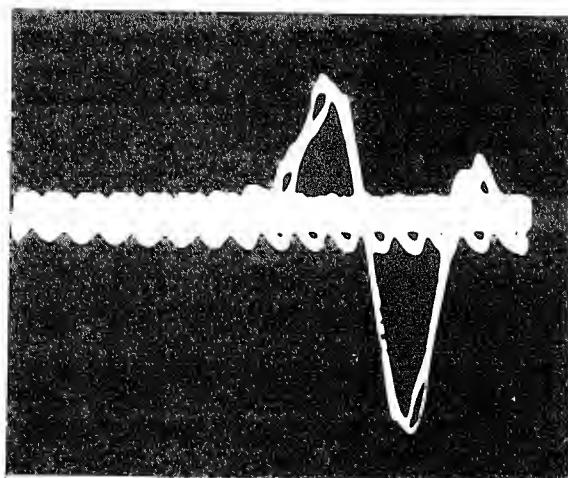


Photo 6-4. Rounded-off digital signal with tape hiss.

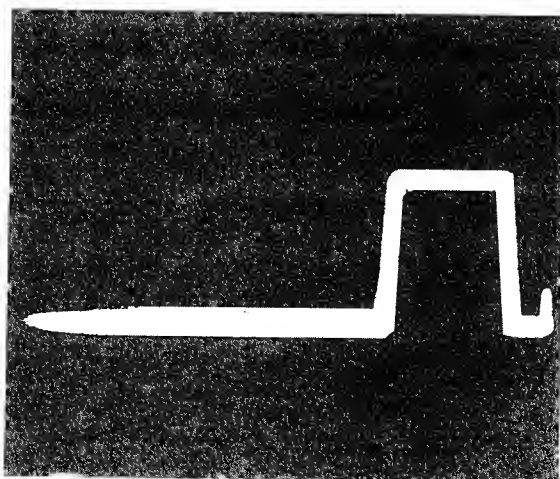
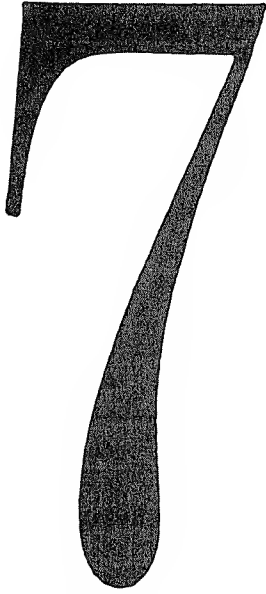


Photo 6-5. Reshaped wave that the CPU finally sees.

NOTES



Controlling the World

The world is out there, waiting to be controlled by your TRS-80. Really. Most of the partly-mechanized daily routine that you find yourself doing by hand might be connected to the '80. Not that it should be, mind you, but it could be. You could successfully operate a clothes washer and drier, a dishwasher, an automatic furnace, an air conditioner, or many other electromechanical devices; turn lights, televisions, and radios on and off, as well as replace your ordinary wake-up alarm; dial the phone, or answer it in a synthetic voice; even act as a digital burglar alarm. But would you really want to?

Admittedly, microprocessors are making their way into more and more appliances, including such diverse items as automobile engines and electric razors. But remember that the TRS-80 is a spectacular *general-purpose* device; although it might do all sorts of control work, it's best suited to applications calling for a large library of different electronic tasks, most with crucial human interaction. Aside from machine-oriented tasks like word processing, calculating, and printing, then, these applications include (generally speaking) measurement, signalling, and cloning.

Measurement involves the evaluation of a real-world occurrence in real time: checking temperature, counting events, comparing relative amounts of light or voltage or sound. Some of these measurements, like counting

events, are inherently digital – an event (say, a person walking through a door) either is taking place or it isn't. A door is closed or it's open. And so on. Other measurements are relative or quantitative, consisting of small increments or continuous change, such as the pressure, the rate of flow, or the quantity of water in a pipe.

Signalling is somewhat the opposite of measurement: a user is informed when some activity has been completed, or when a particular condition in the environment has been reached – including such things as the completion of a mailing list sort or someone breaking and entering.

Cloning is a rather odd phrase, but it's what I would like to call the computer's ability to create precise duplicates of some target object or activity. If that sounds a bit too philosophical, then think of it this way: the computer is fast, which means if it is given a task, it can complete it quickly. The computer is capable of calculating the parameters of its task with enormous accuracy. And finally, the computer works in minute, definable, and identical increments. Simply stated, we can command a computer to do work which, barring glitches and bugs, will be identical every time.

This last concept is the reason microprocessors have become the favored design tools for machining, measurement, and even the creation of music. Jigs wear out, so that tolerances change; but computer programs can be self-correcting. Measuring devices can go out

of alignment: again, computers can be programmed to cross-check and correct these errors. And finally, where electronic design has been sloppy as a result of inaccuracies in the electronic parts themselves (such as in synthesized music), computers can provide the advantages of precise replication where it was not possible before.

The general approach to discrete, digital events is discussed in Chapter (?), where input/output ports are presented. The hardware of interfacing digital inputs and outputs to electrical appliances and other 'real-world' devices is brilliantly described in *TRS-80 Interfacing* (Volumes 1 and 2 — see Bibliography; especially check Volume 2, Chapter 1).

This Chapter will present a few real-world interfacing projects, but will touch only lightly on the theory involved. If you plan to put together your own interfaces, turn to the references cited. Beyond that, here are some rules of thumb before you begin considering interfacing schemes or parts purchases:

Output Interfacing

TTL-level integrated circuits of the type used in your TRS-80 are capable of driving (running) little more than other TTL circuits like themselves or an occasional light-emitting diode (LED). Don't hook your computer expansion card onto any home-made electronic board unless it has TTL inputs; consider the computer's bus to be its most delicate hardware.

For interfacing purposes, integrated circuit peripheral drivers are great. Use them to light small bulbs, turn on miniature relays, and to operate other low-voltage applications needing only limited current. Type 75452 is reliable and cheap, and takes abuse; it's the circuit used by the TRS-80 to run the video dot output and cassette relay. For LEDs, use individual transistors or the inexpensive type 500, 501 and 502 digit and segment drivers.

Isolation of hazardous voltages is essential for external equipment running from your computer. Consider anything plugged into house current to be potentially hazardous, because there's nothing like a power surge or unexpected short circuit to fry your '80 and maybe you too. Use opto-isolators and — or high-current relays for running electric stoves or even AC light bulbs.

Always read the specifications of both the equipment you plan to interface and the device

that's going to do the interfacing. Look for:

- Average voltage and maximum voltage of the interfacing device, and operating voltage of the equipment to be run; the figure is given in volts (V), direct current volts (VDC), or alternating current volts (VAC). The interfacing device must always have a rating higher than the equipment to be run.
- Average current and surge current of the the interfacing device, and operating current of the equipment to be run; the figure is given in amps (A) or milliamps (mA). The interfacing device must always have a rating higher than the equipment to be run.
- Isolation voltage of the interfacing device. This device must have a rating roughly 67 percent higher than the equipment to be run.

Actually, interfacing peripheral equipment is one of the easiest things you can do with your TRS-80. Below is a simple schematic; it shows one latched output line from the computer, and how it might drive:

- (a) an LED
- (b) a relay
- (c) an ordinary house lamp
- (d) a motor
- (e) a high-voltage circuit.

You wouldn't want to use this single line to do all these things, but you might want to combine the LED, house lamp, and motor. That way you could have an indicator near the computer that the motor is on, a bright lamp outside a building to indicate the motor is on, and the motor itself would go on.

The point of this schematic is that it can be tapped at any point on the line of parts shown; use the parts only as far as you need them:

Input Interfacing

Counting events or determining occurrences has an easier group of rules: TTL level signals can interface right to other TTL devices with no work; just don't connect outputs together.

Other on-off signals are interfaced with a device or group of devices which can shape the incoming signal into a neat, tight square wave, at TTL level. Such a device is called a Schmitt trigger, and is available as type 7414 for a few cents.

If the signal is in the range of 4 to 7 volts, the Schmitt trigger will transform it to a fast-moving digital signal. If the signal occurs very often or is erratic or unstable in its rising voltage, then

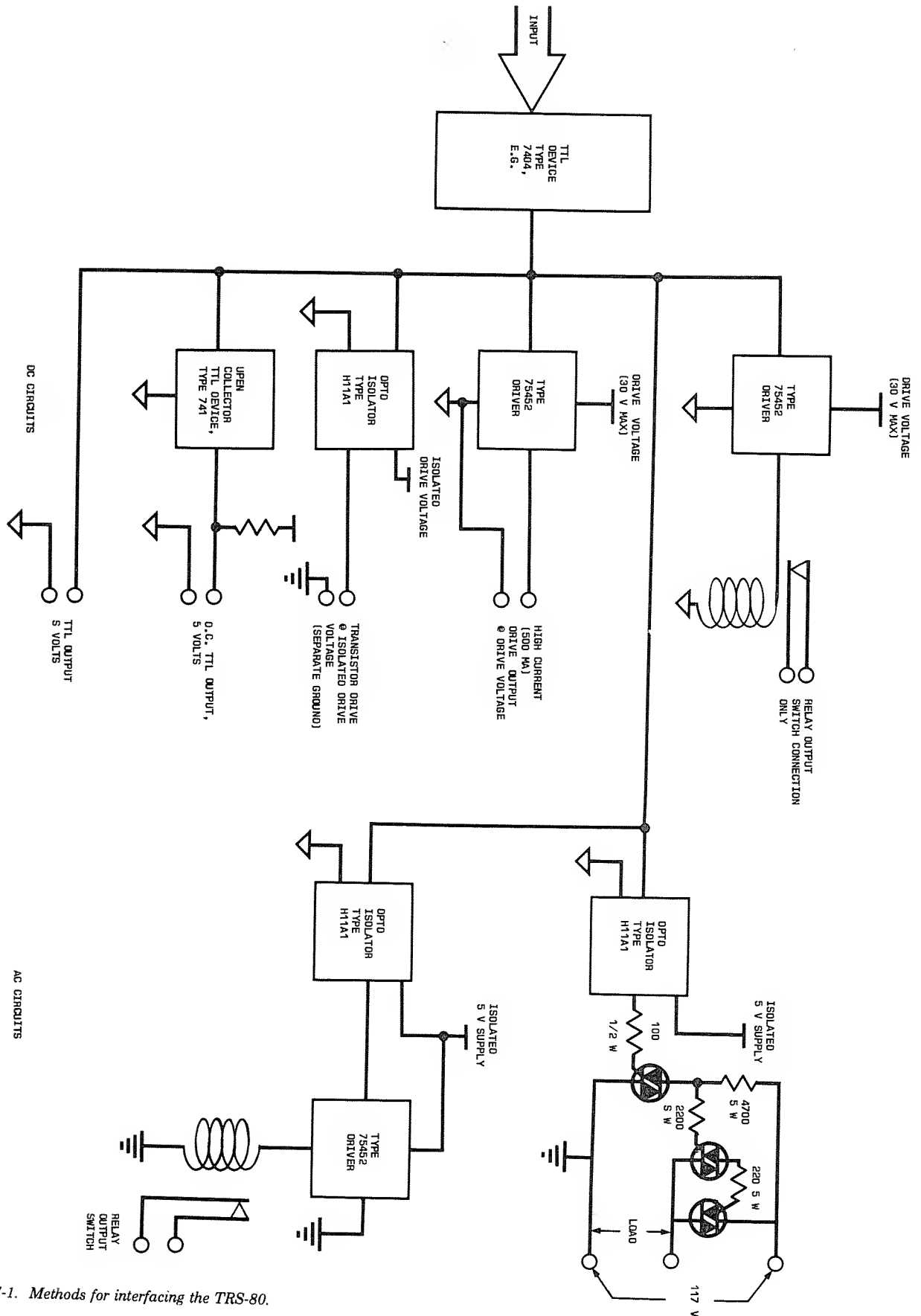


Figure 7-1. Methods for interfacing the TRS-80.

additional hardware or very tolerant software must be used. Use optical isolators to feed the computer or relays for interfacing higher voltages.

For AC input, the signal may be transformed to about 3 to 6 volts, rectified to DC and filtered only if the signal is slow-moving (once a minute or so). Then the signal can be fed to a Schmitt trigger to the computer interface. Otherwise, the AC information can only be transformed to a lower voltage and its pulses (60 per second if it is ordinary house current) counted either by hardware or software. Don't interface AC if you can avoid it, because it can be a genuine pain in the bytes. Instead, have the AC run a fast relay and identify when the relay turns on and off.



A to D, D to A

DAC0808, DAC0807, DAC0806 8-bit D/A converter

general description

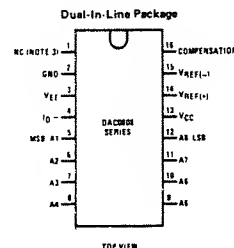
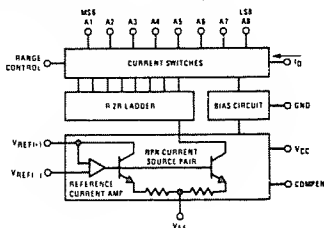
The DAC0808 series is an 8-bit monolithic digital-to-analog converter (DAC) featuring a full scale output current settling time of 150 ns while dissipating only 33 mW with $\pm 5V$ supplies. No reference current (I_{REF}) trimming is required for most applications since the full scale output current is typically ± 1 LSB of $255 I_{REF}/256$. Relative accuracies of better than $\pm 0.19\%$ assure 8-bit monotonicity and linearity while zero level output current of less than $4 \mu A$ provides 8-bit zero accuracy for $I_{REF} \geq 2$ mA. The power supply currents of the DAC0808 series are independent of bit codes, and exhibits essentially constant device characteristics over the entire supply voltage range.

The DAC0808 will interface directly with popular TTL, DTL or CMOS logic levels, and is a direct replacement for the MC1508/MC1408. For higher speed applications, see DAC0800 data sheet.

features

- Relative accuracy: $\pm 0.19\%$ error maximum (DAC0808)
- Full scale current match: ± 1 LSB typ
- 7 and 6-bit accuracy available (DAC0807, DAC0806)
- Fast settling time: 150 ns typ
- Noninverting digital inputs are TTL and CMOS compatible
- High speed multiplying input slew rate: $8 \text{ mA}/\mu\text{s}$
- Power supply voltage range: $\pm 4.5V$ to $\pm 18V$
- Low power consumption: 33 mW @ $\pm 5V$

block and connection diagrams



typical application

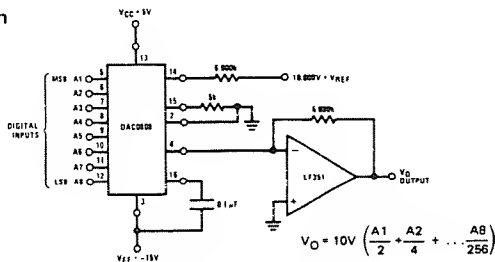


FIGURE 1. $\pm 10V$ Output Digital to Analog Converter

ordering information

ACCURACY	OPERATING TEMPERATURE RANGE	ORDER NUMBERS*			
		D PACKAGE (D16C)	J PACKAGE (J16A)	N PACKAGE (N16A)	
8-bit	$-55^{\circ}\text{C} \leq T_A \leq +125^{\circ}\text{C}$	DAC0808LD	LM1508D-8	DAC0808LJ	LM1508J-8
8-bit	$0^{\circ}\text{C} \leq T_A \leq +75^{\circ}\text{C}$			DAC0808LCJ	LM1408J-8
7-bit	$0^{\circ}\text{C} \leq T_A \leq +75^{\circ}\text{C}$			DAC0807LCJ	LM1408J-7
6-bit	$0^{\circ}\text{C} \leq T_A \leq +75^{\circ}\text{C}$			DAC0806LCJ	LM1408J-6

*Note: Devices may be ordered by using either order number.

Figure 7-2. National Semiconductor data sheet example.

Signals from photocells and similar resistive devices can be fed through Wein bridges or merely fed into an operational amplifier. Again, the references such as *TRS-80 Interfacing* and *Engineer's Notebook* contain plenty of details on interfacing on-off signals.

D-to-A and A-to-D Conversion

Although some human events occur in on-off groups, most of life is pliable, elusive, and relative. It works by image and analogy, not by counting. It is an analog world, and the computer is a digital device. Faced with this dilemma, two important groups of electronic circuits have been developed: the digital-to-analog converter (D/A converter) and the analog-to-digital converter (A/D converter).

The first class accepts a parallel digital input a given number of bits wide, and converts that to individual steps. For N steps, the number of distinct voltages or currents is 2 to the power N. The greater the number of bits, the lesser the relative size of the steps, to the point where the distinction between steps becomes insignificant.

For an ordinary 8-bit data bus such as that on the TRS-80, the available voltage can be divided into 256 parts; for five-volt circuitry, this is 0.0195 volts per step, starting at zero. If greater accuracy is essential, then 12-bit converters can be used (at greater expense), fed by one and one-half bytes of data from the computer. This provides 4,096 steps at 0.0012 volts per step. The most common 8-bit type is the DAC0808 (National Semiconductor) or the MC1408L8 (Motorola); data sheets are found in an Appendix.

With present technology, conversion from analog signals to digital ones is much more

application hints

REFERENCE AMPLIFIER DRIVE AND COMPENSATION

The reference amplifier provides a voltage at pin 14 for converting the reference voltage to a current, and a turn-around circuit or current mirror for feeding the ladder. The reference amplifier input current, I_{QA} , must always flow into pin 14, regardless of the set-up method of reference voltage polarity.

Connections for a positive voltage are shown in Figure 7. The reference voltage source supplies the full current, I_{QA} . For bipolar reference signals, as in the multiplying mode, R_{15} can be tied to a negative voltage corresponding to the minimum input level. It is possible to eliminate R_{15} with only a small sacrifice in accuracy and temperature drift.

The compensation capacitor value must be increased with increases in R_{14} to maintain proper phase margin: for R_{14} values of 1, 2.5 and 5 k Ω , minimum capacitor values are 15, 37 and 75 pF. The capacitor may be tied to either VEE or ground, but using VEE increases negative supply rejection.

absolute maximum ratings ($T_A = 25^\circ\text{C}$ unless otherwise noted)

Power Supply Voltage		Power Dissipation (Package Limitation)	
V _{CC}	5.5 V _{DC}	Cavity Package	1000 mW
V _{EE}	-16.5 V _{DC}	Derate above T _A = 25°C	6.7 mW/°C
Digital Input Voltage, V ₅ –V ₁₂	-10 V _{DC} to +18 V _{DC}	Operating Temperature Range	
Applied Output Voltage V _D	-11 V _{DC} to +18 V _{DC}	DAC0808L	-55°C ≤ T _A ≤ +125°C
Reference Current, I ₁₄	5 mA	QAC0808LC Series	0 ≤ T _A ≤ +75°C
Reference Amplifier Inputs, V ₁₄ V ₁₅	V _{CC} , V _{EE}	Storage Temperature Range	-65°C to +150°C

electrical characteristics

($V_{CC} = 5\text{V}$, $V_{EE} = -15\text{VDC}$, $V_{REF}/R_{14} = 2\text{mA}$, DAC0808L, $T_A = -55^\circ\text{C}$ to $+125^\circ\text{C}$, DAC0808LC, DAC0807LC, DAC0806LC, $T_A = 0^\circ\text{C}$ to $+75^\circ\text{C}$, and all digital inputs at high logic level unless otherwise noted.)

PARAMETER	CONDITIONS	MIN	TYP	MAX	UNITS
E_r Relative Accuracy (Error Relative to Full Scale I_D)	(Figure 4)				%
DAC0808L (LM1508-8), DAC0808LC (LM1408-8)				± 0.19	%
DAC0807LC (LM1408-7), (Note 1)				± 0.39	%
DAC0806LC (LM1408-6), (Note 1)				± 0.78	%
Settling Time to Within 1/2 LSB (Includes t_{PLH})	$T_A = 25^\circ\text{C}$ (Note 2), (Figure 5)		150		ns
t_{PLH} , t_{PHL} Propagation Delay Time	$T_A = 25^\circ\text{C}$, (Figure 5)		30	100	ns
TC_D Output Full Scale Current Drift			± 20		ppm/ $^\circ\text{C}$
M_{SB} Digital Input Logic Levels	(Figure 3)				
V_{IH} High Level, Logic "1"		2			VDC
V_{IL} Low Level, Logic "0"				0.8	VDC
M_{SB} Digital Input Current	(Figure 3)				mA
High Level	$V_{IH} = 5\text{V}$		0	0.040	mA
Low Level	$V_{IL} = 0.8\text{V}$		-0.003	-0.2	mA
I_{15} Reference Input Bias Current	(Figure 3)		-1	-5	μA
Output Current Range	(Figure 3)				mA
$V_{EE} = -5\text{V}$		0	2.0	2.1	mA
$V_{EE} = -15\text{V}$, $T_A = 25^\circ\text{C}$		0	2.0	4.2	mA
I_D Output Current	$V_{REF} = 2.000\text{V}$, $R_{14} = 1000\Omega$, (Figure 3)	1.9	1.99	2.1	mA
Output Current, All Bits Low	(Figure 3)		0	4	μA
Output Voltage Compliance Pin 1 Grounded, V_{EE} Below -10V	$E_r \leq 0.19\%$, $T_A = 25^\circ\text{C}$			-0.55, +0.4 -5.0, +0.4	VDC
SR_{REF} Reference Current Slew Rate	(Figure 6)		8		mA/ μs
Output Current Power Supply Sensitivity	$-5\text{V} \leq V_{EE} \leq -16.5\text{V}$		0.05	2.7	$\mu\text{A/V}$
Power Supply Current (All Bits Low)	(Figure 3)				mA
I_{CC}			2.3	22	mA
I_{EE}			-4.3	-13	mA
Power Supply Voltage Range	$T_A = 25^\circ\text{C}$, (Figure 3)	4.5	5.0	5.5	VDC
V_{CC}		-4.5	-15	-16.5	VDC
Power Dissipation					mW
All Bits Low	$V_{CC} = 5\text{V}$, $V_{EE} = -5\text{V}$		33	170	mW
	$V_{CC} = 5\text{V}$, $V_{EE} = -15\text{V}$		106	305	mW
All Bits High	$V_{CC} = 15\text{V}$, $V_{EE} = -5\text{V}$		90		mW
	$V_{CC} = 15\text{V}$, $V_{EE} = -15\text{V}$		160		mW

Note 1: All current switches are tested to guarantee at least 50% of rated current.

Note 2: All bits switched.

Note 3: Range control is not required.

application hints (Continued)

A negative reference voltage may be used if R_{14} is grounded and the reference voltage is applied to R_{15} as shown in Figure 8. A high input impedance is the main advantage of this method. Compensation involves a capacitor to V_{EE} on pin 16, using the values of the previous paragraph. The negative reference voltage must be at least 4V above the V_{EE} supply. Bipolar input signals may be handled by connecting R_{14} to a positive reference voltage equal to the peak positive input level at pin 15.

When a DC reference voltage is used, capacitive bypass to ground is recommended. The 5V logic supply is not recommended as a reference voltage. If a well regulated 5V supply which drives logic is to be used as the reference, R_{14} should be decoupled by connecting it to 5V through another resistor and bypassing the junction of the 2 resistors with 0.1 μF to ground. For reference

voltages greater than 5V, a clamp diode is recommended between pin 14 and ground.

If pin 14 is driven by a high impedance such as a transistor current source, none of the above compensation methods apply and the amplifier must be heavily compensated, decreasing the overall bandwidth.

OUTPUT VOLTAGE RANGE

The voltage on pin 4 is restricted to a range of -0.6 to 0.5V when $V_{EE} = -5\text{V}$ due to the current switching methods employed in the DAC0808.

The negative output voltage compliance of the DAC0808 is extended to -5V where the negative supply voltage is more negative than -10V. Using a full-scale current of 1.992 mA and load resistor of 2.5 k Ω between pin 4 and ground will yield a voltage output of 256 levels between 0 and -4.980V. Floating pin 1 does not affect

the converter speed or power dissipation. However, the value of the load resistor determines the switching time due to increased voltage swing. Values of R_L up to 500 Ω do not significantly affect performance, but a 2.5 k Ω load increases worst-case settling time to 1.2 μs (when all bits are switched ON). Refer to the subsequent text section on Settling Time for more details on output loading.

OUTPUT CURRENT RANGE

The output current maximum rating of 4.2 mA may be used only for negative supply voltages more negative than -7V due to the increased voltage drop across the resistors in the reference current amplifier.

ACCURACY

Absolute accuracy is the measure of each output current level with respect to its intended value, and is dependent upon relative accuracy and full-scale current drift. Relative accuracy is the measure of each output current level as a fraction of the full-scale current. The relative accuracy of the DAC0808 is essentially constant with temperature due to the excellent temperature tracking of the monolithic resistor ladder. The reference current may drift with temperature, causing a change in the absolute accuracy of output current. However, the DAC0808 has a very low full-scale current drift with temperature.

The DAC0808 series is guaranteed accurate to within ± 1.2 LSB at a full-scale output current of 1.992 mA. This corresponds to a reference amplifier output current drive to the ladder network of 2 mA, with the loss of 1 LSB (8 μA) which is the ladder remainder shunted to ground. The input current to pin 14 has a guaranteed value of between 1.9 and 2.1 mA, allowing some mismatch in the NPN current source pair. The accuracy test circuit is shown in Figure 4. The 12-bit converter is calibrated for a full-scale output current of 1.992 mA. This is an optional step since the DAC0808 accuracy is essentially the same between 1.5 and 2.5 mA. Then the DAC0808 circuit's full-scale current is trimmed to the same value with R_{14} so that a zero value appears at the error amplifier output. The counter is activated and the error band may be displayed on an oscilloscope, detected by comparators, or stored in a peak detector.

Two 8-bit D-to-A converters may not be used to construct a 16-bit accuracy D-to-A converter. 16-bit accuracy implies a total error of $\pm 1/2$ of one part in 65,536, or $\pm 0.00076\%$, which is much more accurate than the $\pm 0.019\%$ specification provided by the DAC0808.

MULTIPLYING ACCURACY

The DAC0808 may be used in the multiplying mode with 8-bit accuracy when the reference current is varied over a range of 256:1. If the reference current in the multiplying mode ranges from 16 μA to 4 mA, the additional error contributions are less than 1.6 μA . This is well within 8-bit accuracy when referred to full-scale.

A monotonic converter is one which supplies an increase in current for each increment in the binary word. Typically, the DAC0808 is monotonic for all values of reference current above 0.5 mA. The recommended range for operation with a DC reference current is 0.5 to 4 mA.

SETTLING TIME

The worst-case switching condition occurs when all bits are switched ON, which corresponds to a low-to-high transition for all bits. This time is typically 150 ns for settling to within $\pm 1/2$ LSB, for 8-bit accuracy, and 100 ns to 1/2 LSB for 7 and 6-bit accuracy. The turn OFF is typically under 100 ns. These times apply when $R_L \leq 500\Omega$ and $C_0 \leq 25\text{pF}$.

Extra care must be taken in board layout since this is usually the dominant factor in satisfactory test results when measuring settling time. Short leads, 100 μF supply bypassing for low frequencies, and minimum scope lead length are all mandatory.

complicated. Most commonly used is a technique which compares the input voltage to known voltages inside the converter. When a match is found, a binary version of the input condition is fed to the computer bus. This process is slow when compared to computer speeds, and such devices usually have start-of-conversion and end-of-conversion signals so the computer will not receive false data by attempting to read the converter's information during the conversion process. The commonly used 8-bit type is ADC0800 (National Semiconductor); see the Appendix.

Music and Sound Effects

There are other ways of producing sound using the TRS-80 which do not involve exclusively software. The software-only approach may be capable of producing sound effects, but in the TRS-80 configuration it cannot provide any

spatial effects (it is monaural only), nor can it offer a large variety of textures. For these, the TRS-80 user must turn to a little extra hardware.

Simplest among the extra hardware is a latched output address. The information written to the address appears at the output of the latch. Write to the address fast enough, and sound is produced, because the latch acts as a kind of electronic window to that single memory location. Feed that digital activity through a few resistors to blend the sounds, and run that to a stereo amplifier. Voila! Spatial sound.

The circuit presented here contains an 8-bit latch (Z1), and address decoder (Z3, Z4 and Z4, mapped to 4FFF), and three-state output buffer (Z2). The computer data is latched into Z1 when Write to 4FFF appears at Z3/4/5. Depending on the data at Z1, any or all of the four buffers in Z2 may be turned on or off.

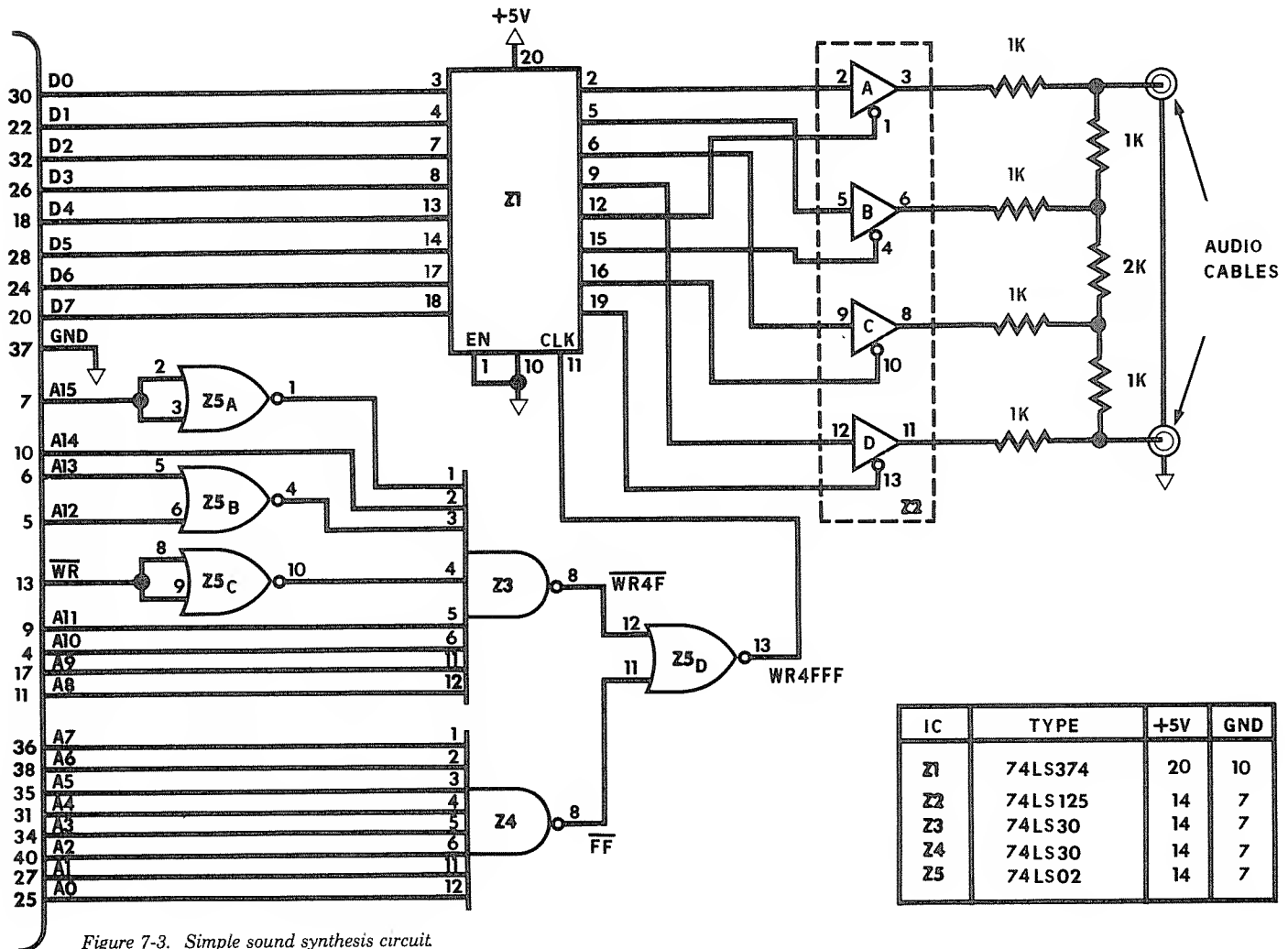


Figure 7-3. Simple sound synthesis circuit.

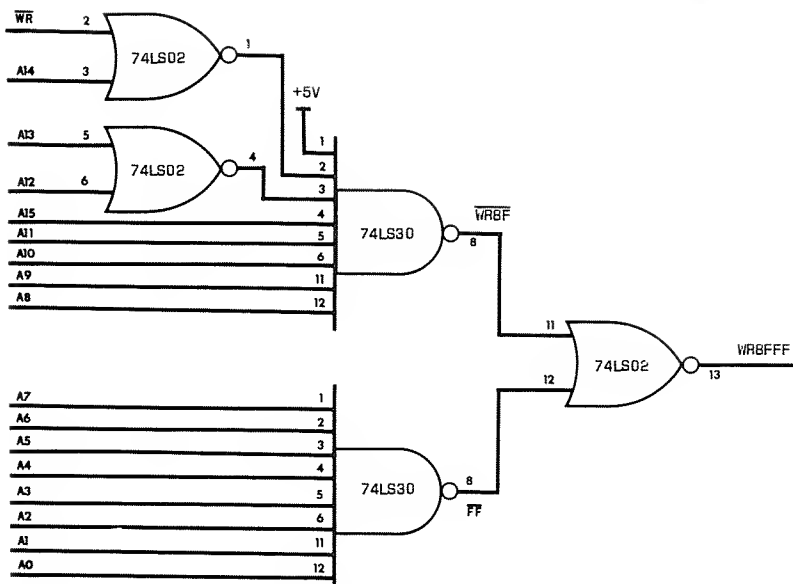


Figure 7-4. Alternate decoding scheme.

An alternate decoding (8FFF) is provided, and is recommended for some of the musical selections presented. The selections include my own arrangement of 'God Rest Ye Merry Gentlemen', first presented in the Christmas 1980 issue of *80 Microcomputing*; and 'Your California', a suite by David Gunn, originally composed for viola duo. The latter group is written for use with the *Exatron Stringy-Floppy*, and will continuously repeat the four parts of the suite.

There are two versions of the software; both are identical, except that the second is a two-voice program intended for pieces such as 'Your California'. Because the TRS-80 is slow (1.77 MHz) and the Z-80 is slow (an enormous number of machine cycles per instruction when compared with other microprocessor families), four-part music will reproduce in the low register. The two-part music is much more satisfying musically.



Figure 7-5. Music notation for people.

The image displays a musical score for a computer, consisting of multiple staves. The notation includes various musical symbols such as notes, rests, and slurs. There are several instances of triplets, indicated by a '3' over a bracket. The score is written in a standard musical notation style, with a key signature of one sharp (F#) and a time signature of 4/4. The music is arranged in a system of staves, with some staves containing multiple measures of music. The notation is clear and legible, with a focus on the rhythmic and melodic elements of the piece.

Figure 7-6. The same music for the computer.


```

10 REM * CLEARING SPACE FOR PROGRAM CHAINING SEQUENCE
20 POKE16634,PEEK(16634)+16 : REM * SIMPLE VARIABLE POINTER
30 CLEAR50 : REM * CLEAR STRING SPACE AND RESET POINTERS
40 @LOAD3 : REM * BRING NEXT PROGRAM IN FROM STRINGY FLOPPY

```

```

10 OUT254,2 : CLS : PRINT"READING.." : REM * OUT254,2 = HISPEED
20 X = 24578 : REM * THIS IS MEM. POSITION OF VOICE NUMBER ONE
30 READ A : IF A = 999 THEN 50 ELSE POKE X,A : REM * EXIT @ 999
40 X = X + 4 : GOTO 30 : REM * VOICE APPEARS EVERY 4 MEM LOC'NS
50 X = 24579 : REM * THIS IS MEM. POSITION OF VOICE NUMBER TWO
60 READ A : IF A = 999 THEN 80 ELSE POKE X,A : REM * PITCH POKE
70 X = X + 4 : GOTO 60 : REM * VOICE APPEARS EVERY 4 MEM LOC'NS
80 X = 24576 : REM * THIS IS THE BEGINNING VALUE FOR DURATIONS
90 READ A : IF A = 999 THEN 110 ELSE POKE X,A*2 : REM * RHYTHMS
100 X = X + 4 : GOTO 90 : REM * NOTICE RHYTHM MULTIPLIER ABOVE
110 RESTORE : REM * THIS GETS THE DATA POINTER BACK TO START
120 X = 24577 : REM * THESE LOCATIONS USED FOR SUBTLE DURATIONS
130 READ A : IF A = 999 THEN 150 ELSE POKE X,1 : REM * RHYTHMS
140 X = X + 4 : GOTO 130 : REM * NO CHANGE IN LSB TIMING ABOVE
150 POKE X,0 : POKE X+4,0 : POKE X+8,0 : REM * MUSIC END CODE
160 OUT254,3 : POKE 16526,96 : POKE 16527,143 : M=USR(0) : CLEAR
170 @LOAD4 : REM * PLAY NORM SPEED, USR CALL, CLEAR, LOAD NEXT
180 DATA0,0,0,0,0,0,0,22,26,30,30,30,30,34,34,34,0,0,0,0,0
190 DATA20,16,13,14,14,14,14,15,15,15,0,0,0,0,0,20,18,18
200 DATA18,18,24,0,24,0,24,0,24,0,24,24,24,0,24,0,24,0,24,0
210 DATA24,24,0,24,0,24,24,24,0,24,0,24,24,0,24,0,24,0,24,0
220 DATA22,22,18,18,16,16,14,14,13,13,12,13,13,13,14,19,27,19
230 DATA13,13,14,19,27,19,13,0,16,0,20,0,22,0,24,0,24,0,24,0
240 DATA0,20,16,13,14,14,20,16,13,0,16,0,16,0,16,12,12,12,13
250 DATA15,15,15,15,18,21,21,21,16,16,0,16,0,16,0,16,0,16,0
260 DATA17,15 : REM * end of page one for voice one
270 DATA20,21,24,26,0,0,0,0,22,26,30,30,30,30,34,34,34,0,0
280 DATA0,24,0,24,24,24,24,0,24,0,24,24,0,24,0,24,0,0,20,18
290 DATA18,18,24,0,24,0,24,0,24,0,21,0,22,0,18,0,30,30,32,0,24,0
300 DATA24,0,24,0,24,0,26,21,16,16,16,16,16,16,16,16,999
310 REM * THIS IS THE END OF THE PITCH SEQUENCE FOR VOICE ONE
320 DATA24,0,24,0,24,0,24,0,24,24,24,0,24,0,24,24,0,24,0,24,0
330 DATA24,0,24,24,24,24,0,24,0,24,24,0,24,0,24,24,0,24,0,26,26
340 DATA0,26,0,27,0,27,0,27,0,40,0,22,26,30,30,30,30,34,34
350 DATA34,0,0,0,0,0,0,20,16,13,14,14,14,14,15,15,15,0,0
360 DATA0,0,27,30,32,36,40,20,18,36,40,20,18,36,40,20,18,18,36
370 DATA36,40,20,18,18,36,36,40,0,30,0,26,0,32,0,27,0,27,0,30
380 DATA0,40,0,22,22,22,26,30,34,34,0,24,0,24,0,24,26,26,22,30,30,30
390 DATA0,24,24,0,24,0,24,24,0,24,0,24,26,26,26,26,22,30,30,30
400 DATA30,32,32,32 : REM * end of page one for voice two
410 DATA32,32,32,0,0,24,0,24,0,24,24,24,0,24,0,24,24,0,24,0,24
420 DATA0,0,20,16,13,14,14,14,14,15,15,15,0,0,0,24,0,26,26
430 DATA0,26,0,27,0,27,0,30,0,30,0,32,32,32,32,32,0,0,34,0
440 DATA34,0,34,0,34,0,32,28,24,24,26,24,26,24,26,28,26,999
450 REM * this is the conclusion of the score for voice two
460 DATA10,2,2,2,10,2,2,2,4,4,2,2,2,8,2,2,2,10,2,2,2,3,3,3
470 DATA2,2,2,2,8,2,2,2,8,2,2,2,8,2,2,2,10,2,2,2,10,2,2,2
480 DATA4,4,2,2,2,8,2,2,2,8,2,2,2,2,3,3,2,2,2,2,8,2,2,2,2
490 DATA10,2,2,2,4,4,4,8,8,8,8,8,8,8,4,4,4,4,8,8,4,4,4
500 DATA4,10,2,2,2,10,2,2,2,10,2,2,2,10,2,2,2,3,3,3,8,8,3,3,3
510 DATA8,10,2,2,2,8,2,2,2,8,2,2,2,8,2,2,2,8,2,2,2,2,2,8
520 DATA2,2,2,2,8,8,8 : REM * end of page one for durations
530 DATA8,8,16,8,8,10,2,2,2,4,4,2,2,2,8,2,2,2,2,10,2,2,2,3,3
540 DATA3,2,2,2,2,8,2,2,2,10,2,2,2,8,2,2,2,10,2,2,2,10,2,2
550 DATA2,10,2,2,2,2,4,8,10,2,2,2,10,2,2,2,4,4,8,8,1,1,1,3
560 DATA6,20,999 : REM * end of score for durations; end score

```

Listing 7-1. Examples of music programs.

More Music

Naturally, there is much more to the creation of music by computer than the mere sounding of tones; the appearance of commercial music-generation peripherals for the TRS-80 attests to that. But though they may be well-designed pieces of electronics, composers and others serious about producing listenable music tend to look for richer sounds and more flexible ways to change that sound.

There are several ways to do this. The simplest is the tone generator with a tempered scale, capable of producing chordal sounds (consisting of three or more pitches sounding simultaneously). The tone color is limited to a fixed palette, and they are suitable for making elementary music. I refer to these generically as 'organs' because their tone colors usually give the illusion of organ stops. Most TRS-80 compatible music boards are simplified variations on the basic electronic organ, using integrated circuit chips like the General Instrument AY-3-8910 and the Texas Instrument SN76489.

A second group of electronic instruments creates music by constructing waveforms from a pre-assigned table of values. Thus, tone color, pitch, and volume can be altered by the selection of parameters in the table. This is digital synthesis in an elementary form, just one step beyond the simple software tone generation presented earlier in this Chapter. The TRS-80 is not capable of operating either with enough speed or electronic flexibility for digital synthesis.

Peripheral to the digital synthesis of tones is the electronic creation of vocal sounds. Texas Instruments, Votrax and other integrated circuits use a complex algorithm called 'linear predictive coding' to select from a known subset of human vocal sounds. Although somewhat convincing voices can be produced this way, mere intelligibility is the least significant criterion in music. Until (and if) predictive devices are developed for a wide array of musical sound, they have no application for producing music.

The most popular electronic music makers for more than a decade have been the analog synthesizers. Traditional oscillators create a sound which can be mutated and transformed until its color is right. These synthesizers were never conceived in computer-compatible terms, but companies such as PAIA have for the past few years offered hybrid analog-digital systems where the computer is used as a super-sequencer,

keeping the notes in order for storage and playback.

There are several reasons why I suggest interfacing analog synthesizers to the TRS-80. First, analog synthesizers are cheap. Small, capable machines can be picked up for the cost of a TRS-80, and kits are sold by PAIA and others for less than \$100. Surplus synthesizers are also available (*Moog, Buchla, PBI, Putney*, and others of early 1970's vintage) at low cost. All of these will take on new musical life when interfaced to the TRS-80.

Second, if you are a performing musician, it's likely you're looking for a musical instrument,

and it's in that area where the analog synthesizer is still champion. They are performance instruments, not electronic widgets. And with a computer interface, they remain stand-alone performance instruments, but with computer assistance where it is wanted.

Figure 7-7 presents the circuit for a 2- to 32-voice analog synthesizer interface. The basic circuit contains a data-line buffer (Z1), a buffer for the most used addresses (Z4), a port decoder and a voice-pair selector (Z2/Z3). In a fully expanded system, sixty-four ports (port 64 to port 127) are used to provide thirty-two voltage outputs, sixty-four envelopes, with sixty-four additional control lines.

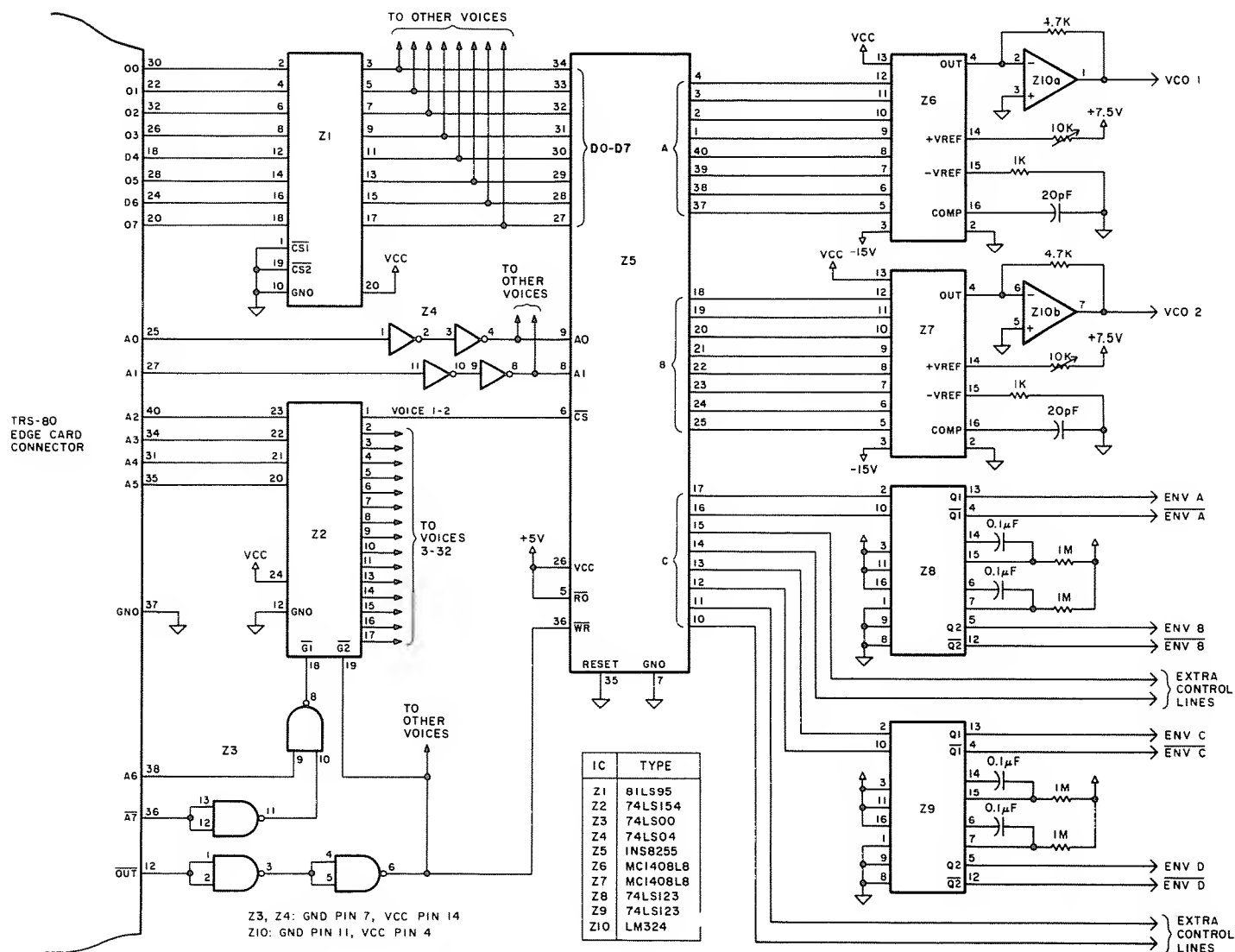


Figure 7-7. TRS-80 to voltage-controlled synthesizer interface.

```

10 OUT 254,2 : CLS : PRINT"READING..." : REM OUT254,2 = HISPEED
20 X = 24578 : REM * THIS IS MEM. POSITION OF VOICE NUMBER ONE
30 READ A : IF A = 999 THEN 50 ELSE POKE X,A : REM * PITCH POKE
40 X = X + 4 : GOTO 30 : REM * VOICE APPEARS EVERY 4 MEM LOC'NS
50 X = 24579 : REM * THIS IS MEM. POSITION OF VOICE NUMBER TWO
60 READ A : IF A = 999 THEN 80 ELSE POKE X,A : REM * PITCH POKE
70 X = X + 4 : GOTO 60 : REM * VOICE APPEARS EVERY 4 MEM LOC'NS
80 X = 24576 : REM * THIS IS THE BEGINNING VALUE FOR DURATIONS
90 READ A : IF A = 999 THEN 110 ELSE POKE X,A*2 : REM * RHYTHMS
100 X = X + 4 : GOTO 90 : REM * NOTICE RHYTHM MULTIPLIER ABOVE
110 RESTORE : REM * THIS GETS THE DATA POINTER BACK TO START
120 X = 24577 : REM * THESE LOCATIONS USED FOR SUBTLE DURATIONS
130 READ A : IF A = 999 THEN 150 ELSE POKE X,1 : REM * RHYTHMS
140 X = X + 4 : GOTO 130 : REM * NO CHANGE IN LSB TIMING ABOVE
150 POKE X,0 : POKE X+4,0 : POKE X+8,0 : REM * MUSIC END CODE
160 OUT254,3 : POKE 16526,96 : POKE 16527,143 : M=USR(0) : CLEAR
170 @LQAO4 : REM * PLAY NORM SPEED, USR CALL, CLEAR, LQAO NEXT
180 DATA0,0,0,0,0,0,0,0,22,26,30,30,30,30,34,34,0,0,0,0,0
190 DATA20,16,13,14,14,14,14,14,15,15,15,0,0,0,0,0,20,18,18
200 DATA18,18,24,0,24,0,24,0,24,0,24,0,24,0,24,0,24,0,24,0
210 DATA24,24,0,24,0,24,24,24,24,0,24,0,24,24,0,24,0,24,0
220 DATA22,22,18,18,16,16,14,14,13,13,12,13,13,14,19,27,19
230 DATA13,13,14,19,27,19,13,0,16,0,20,0,22,0,24,0,24,0
240 DATA24,0,20,16,13,14,14,20,16,13,0,16,0,16,0,16,12,12,12
250 DATA13,15,15,15,18,21,21,21,21,16,16,0,16,0,16,0,16,0
260 DATA16,17,15 : REM * END OF PAGE ONE FOR VOICE ONE
270 DATA20,21,24,26,0,0,0,0,22,26,30,30,30,30,34,34,34,0,0
280 DATA0,24,0,24,24,24,0,24,0,24,24,0,24,0,24,0,24,0,20,18
290 DATA18,18,24,0,24,0,24,0,21,0,22,0,18,0,30,30,32,0,24,0
300 DATA24,0,24,0,24,0,26,21,16,16,16,16,16,16,16,16,999
310 REM END OF PAGE TWO FOR VOICE ONE
320 DATA24,0,24,0,24,0,24,0,24,24,24,0,24,0,24,0,24,0,24,0
330 DATA24,0,24,24,24,24,0,24,0,24,24,0,24,0,24,24,0,24,0,26
340 DATA26,0,26,0,27,0,27,0,27,0,40,0,22,26,30,30,30,30,34
350 DATA34,34,0,0,0,0,0,0,20,16,13,14,14,14,14,15,15,15,15
360 DATA0,0,0,27,30,32,36,40,20,18,36,40,20,18,36,40,20,18,18
370 DATA36,36,40,20,18,18,36,36,40,0,30,0,26,0,32,0,27,0,27,0
380 DATA30,0,40,0,22,22,22,26,30,34,34,34,0,24,0,24,0,24,24,0
390 DATA24,0,24,24,0,24,0,24,24,0,24,0,24,26,26,26,22,30,30
400 DATA30,30,32,32,32 : REM * END OF PAGE ONE FOR VOICE TWO
410 DATA32,32,32,0,24,0,24,0,24,24,24,0,24,0,24,24,0,24,0,24
420 DATA0,0,20,16,13,14,14,14,14,14,15,15,15,0,0,0,24,0,26,26
430 DATA0,26,0,27,0,27,0,30,0,30,0,32,32,32,32,0,0,0,34,0,34
440 DATA0,34,0,34,0,32,28,24,24,26,24,26,24,26,28,26,999
450 REM * END OF PAGE TWO FOR VOICE TWO
460 DATA10,2,2,2,10,2,2,2,4,4,2,2,2,8,2,2,2,2,10,2,2,2,3,3,3
470 DATA2,2,2,2,8,2,2,2,2,8,2,2,2,2,8,2,2,2,2,10,2,2,2,10,2,2,2
480 DATA4,4,2,2,2,8,2,2,2,2,8,2,2,2,2,3,3,3,2,2,2,2,8,2,2,2,2
490 DATA10,2,2,2,4,4,4,4,8,8,8,8,8,8,8,8,8,8,4,4,4,4,8,8,4,4,4
500 DATA4,10,2,2,2,10,2,2,2,10,2,2,2,10,2,2,2,3,3,8,8,3,3,3,3
510 DATA8,10,2,2,2,8,2,2,2,2,8,2,2,2,2,8,2,2,2,2,8,2,2,2,2,8,2
520 DATA2,2,2,8,8,8 : REM * END OF PAGE ONE OF RHYTHMS
530 DATA8,8,16,8,8,10,2,2,2,4,4,2,2,2,8,2,2,2,2,10,2,2,2,3,3
540 DATA3,2,2,2,8,2,2,2,2,10,2,2,2,8,2,2,2,2,10,2,2,2,10,2,2,2
550 DATA2,10,2,2,2,2,2,4,8,10,2,2,2,10,2,2,2,4,4,8,8,1,1,1,1,3
560 DATA6,20,999 : REM * END OF RHYTHMS PAGE TWO AND END PIECE

```

Each pair of voices is managed by programmable peripheral interface Z5, which provides two 8-bit outputs to digital-to-analog converters Z6 and Z7, offering one-part-in-256 accuracy. Remaining from Z5 are eight control lines for envelope, etc. Optionally, six of these remaining control outputs can be used for voltage control of volume or filtering (see Figure 7-7), which can be used with less resolution (one part in 64) than that needed for pitch control.

Figure 7-8 is the power supply, which is similar to others in this book with the exception of the LM340-8, providing eight volts as a reference to the digital-to-analog converters.

The output of the D/A converters Z6 and Z7 is in the range of 0 volts to 4.98 volts, which should be more than adequate to drive most synthesizers across their full range. The envelope triggers provided by Z8 and Z9, however, may need tweaking depending on the type of synthesizer you are using. Some synthos require a negative-going pulse, and others need a positive-going one; likewise, some synthos trigger on the rising or falling *edge* of the wave, while other envelopes will sustain as long as the level of the trigger signal remains high or low. Thus, the resistance and capacitance values given for Z8 and Z9 may have to be changed for sustained envelope triggers, or Z8 and Z9 might be eliminated completely if the envelope is edge-triggered. If you don't have specs on your synthesizer, build the trigger circuit up, then down, until it works for your machine.

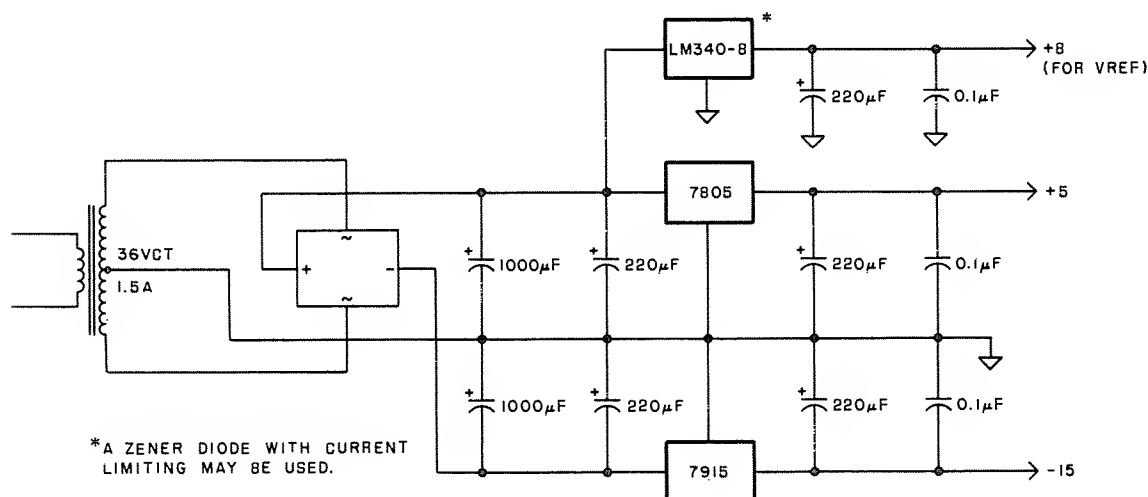


Figure 7-8. Power supply for the synthesizer interface.

```

10 OUT254,2 : CLS : PRINT"READING..." : REM * OUT254,2 = HISPEED
20 X = 24578 : REM * THIS IS MEM. POSITION OF VOICE NUMBER ONE
30 READ A : IF A = 999 THEN 50 ELSE POKE X,A : REM * EXIT @ 999
40 X = X + 4 : GOTO 30 : REM * VOICE APPEARS EVERY 4 MEM LOC'S
50 X = 24579 : REM * THIS IS MEM. POSITION OF VOICE NUMBER TWO
60 READ A : IF A = 999 THEN 80 ELSE POKE X,A : REM * PITCH POKE
70 X = X + 4 : GOTO 60 : REM * VOICE APPEARS EVERY 4 MEM LOC'S
80 X = 24576 : REM * THIS IS THE BEGINNING VALUE FOR DURATIONS
90 READ A : IF A = 999 THEN 110 ELSE POKE X,A*2 : REM * RHYTHMS
100 X = X + 4 : GOTO 90 : REM * NOTICE RHYTHM MULTIPLIER ABOVE
110 RESTORE : REM * THIS GETS THE DATA POINTER BACK TO START
120 X = 24577 : REM * THESE LOCATIONS USED FOR SUBTLE DURATIONS
130 READ A : IF A = 999 THEN 150 ELSE POKE X,1 : REM * RHYTHMS
140 X = X + 4 : GOTO 130 : REM * NO CHANGE IN LSB TIMING ABOVE
150 POKE X,0 : POKE X+4,0 : POKE X+8,0 : REM * MUSIC END CODE
160 OUT254,3 : POKE 16526,96 : POKE 16527,143 : M=USR(0) : CLEAR
170 @LOAD06 : REM * PLAY NORM SPEED, USR CALL, CLEAR, LOAD NEXT
180 DATA0,0,0,0,0,0,0,0,25,29,31,29,3,0,0,0,0,20,16,13,14,14,14
190 DATA14,14,15,15,15,0,0,0,0,0,20,18,18,18,18,24,0,24,0,24
200 DATA0,24,0,24,24,24,0,24,0,24,24,0,24,0,24,0,24,0,24,24
210 DATA24,24,0,24,0,24,24,0,24,0,24,0,24,0,22,22,18,16,16
220 DATA14,14,13,13,12,13,13,13,14,19,27,19,13,13,14,19,27,19
230 DATA13,0,16,0,20,0,22,0,24,0,24,0,24,0,12,12,12,12,13
240 DATA15,15,15,15,18,21,21,21,21,16,16,0,16,0,16,0,16,0,16
250 DATA17,15 : REM * END OF PAGE ONE FOR VOICE ONE
260 DATA3,140,3,22,26,30,30,30,30,34,34,34,0,0,0,24,0,24,24
270 DATA24,24,0,24,0,24,24,0,24,0,24,0,0,0,20,18,18,18,24,0
280 DATA24,0,24,0,21,0,22,0,18,0,30,30,32,0,24,0,24,0,24,0
290 DATA26,21,16,16,16,16,16,16,16,16,999 : REM * END P2V1
300 DATA24,0,24,0,23,3,135,1,1,42,0,24,0,24,0,24,24,24,0,24
310 DATA0,24,24,0,24,0,24,24,0,24,0,26,26,0,27,0,27,0,27,0
320 DATA40,0,22,26,30,30,30,30,34,34,34,0,0,0,0,0,20,16,13
330 DATA14,14,14,14,14,15,15,15,15,0,0,0,0,27,30,32,36,40,20
340 DATA18,36,40,20,18,12,140,140,13,3,40,20,18,18,36,36,40,0
350 DATA30,0,26,0,32,0,27,0,27,0,30,0,40,0,22,22,22,26,30,34,34
360 DATA34,0,24,0,24,0,24,24,0,24,0,24,24,0,24,0,24,24,0,24,0
370 DATA24,26,26,26,26,22,30,30,30,30,32,32,32 : REM * END P1V2
380 DATA32,32,32,0,0,24,0,24,0,24,24,24,0,24,0,24,24,0,24,0,24
390 DATA0,0,0,20,16,13,14,14,14,14,15,15,15,0,0,0,24,0,26,26
400 DATA0,26,0,27,0,27,0,30,0,30,0,32,32,32,32,0,0,0,34,0,34
410 DATA0,34,0,34,0,32,28,24,24,26,24,26,24,26,28,26,999
420 REM * END OF PAGE TWO FOR VOICE TWO
430 DATA10,2,2,2,10,2,2,2,4,4,2,2,2,8,2,2,2,2,10,2,2,2,3,3
440 DATA2,2,2,2,8,2,2,2,8,2,2,2,8,2,2,2,2,10,2,2,2,10,2,2,2
450 DATA4,4,2,2,2,8,2,2,2,8,2,2,2,8,2,2,2,3,3,2,2,2,8,2,2,2
460 DATA10,2,2,2,4,4,4,8,8,8,8,8,8,8,8,4,4,4,4,8,3,11,5,4
470 DATA10,2,2,2,10,2,1,1,10,2,2,2,10,2,2,2,3,3,3,8,8,3,3,8
480 DATA10,2,2,2,8,2,2,2,8,2,2,2,8,2,2,2,8,2,2,2,8,2,2,2,8,2
490 DATA2,2,8,8,8 : REM * END OF PAGE ONE OF RHYTHMS
500 DATA8,8,16,8,8,10,2,2,2,4,4,2,2,2,8,2,2,2,2,10,2,2,2,3,3
510 DATA3,2,2,2,8,2,2,2,2,10,2,2,2,8,2,2,2,2,10,2,2,2,10,2,2
520 DATA2,10,2,2,3,13,5,9,2,2,2,12,10,2,2,2,4,4,8,8,1,1,1,1,3,6
530 DATA20,999 : REM * END PAGE TWO OF RHYTHMS AND END PIECE

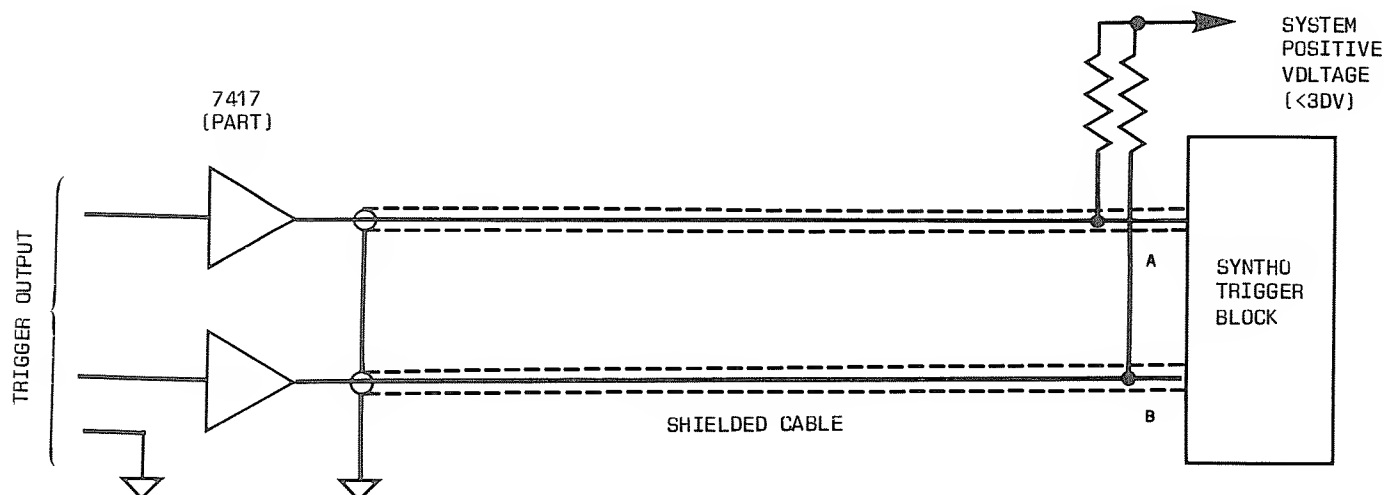
```

Synthesizer Interface Port Addressing

Port Number (Decimal)	Port Number (Hex)	Port Function
64	40	Pitch Control Voice 1
65	41	Pitch Control Voice 2
66	42	Envelopes 1 a/b end 2 a/b Extre Lines 1 a/b end 2 a/b
67	43	Port Control Voices 1 - 2
68	44	Pitch Control Voice 3
69	45	Pitch Control Voice 4
70	46	Envelopes 3 a/b end 4 a/b Extre Lines 3 a/b end 4 a/b
71	47	Port Control Voices 3 - 4
72	48	Pitch Control Voice 5
73	49	Pitch Control Voice 6
74	4A	Envelopes 5 a/b end 6 a/b Extre Lines 5 a/b end 6 a/b
75	4B	Port Control Voices 5 - 6
76	4C	Pitch Control Voice 7
77	4D	Pitch Control Voice 8
78	4E	Envelopes 7 a/b end 8 a/b Extre Lines 7 a/b end 8 a/b
79	4F	Port Control Voices 7 - 8
.	.	.
.	.	.
.	.	.
124	7C	Pitch Control Voice 31
125	7D	Pitch Control Voice 32
126	7E	Envelopes 31 a/b end 32 a/b Extre Lines 31 a/b end 32 a/b
127	7F	Port Control Voices 31 - 32

Table 7-1. Synthesizer interface port addressing.

Using the synthesizer interface is straightforward. Plug the interface into the TRS-80 expansion connector, and run shielded microphone cable from the voltage outputs of the interface to the voltage inputs (marked 'control in', 'voltage in', 'VCO control', 'external in', or something similar) on the synthesizer. Then run either parallel speaker wire or shielded cable from the interface envelope (positive-going edge) to the synthesizer's envelope inputs. The integrated circuits running from the interface are not balanced for long lines; if you plan to use more than a dozen feet of cable, place 7417 open-collector buffers at each control output from the interface, so:



Once the connections are made, run the following few lines, an envelope test:

```

10 OUT254,2:CLS:PRINT"READING..."
12 X=24578
15 READA:IFA=999THEN30ELSEPOKEX,A:X=X+4:GOTO15
30 X=24579
35 READA:IFA=999THEN50ELSEPOKEX,A:X=X+4:GOTO35
50 X=24576
55 READA:IFA=999THEN80ELSEPOKEX,A*2:X=X+4:GOTO55
80 RESTORE:X=24577
90 READA:IFA=999THEN97ELSEPOKEX,1:X=X+4:GOTO90
97 OUT254,3:POKEX+4,0:POKEX+8,0:POKE16526,96:POKE16527,143:N=USR(0):CLEAR
98 @LOAD3
100 DATA20,20,20,20,20,20,18,20,19,17,17,17,17,17,15
110 DATA24,24,21,27,26,27,27,27,30,27,30,27,0,0,20,21,20,19,17
120 DATA16,16,27,24,40,48,48,44,40,36,30,26,26,18,18,14,13
130 DATA11,9,9,12,11,12,10,21,21,27,27,28,27,26,24
140 DATA28,28,27,24,19,17,17,20,13,13,18,16,20,21,21,24,22,22,22,20
150 DATA19,18,36,27,30,30,28,36,19,19,20,13,13,22,22,24,24,17,15,15
160 DATA14,27,24,21,20,21,16,14,13,11,10,13,18,24
170 DATA32,32,32,32,32,32,32,32,32,32,32,32,32,30,27,40,38,38,28,28,26,24,24
180 DATA20,20,20,20,20,20,18,20,19,17,17,17,17,15
190 DATA24,24,21,27,26,27,27,27,30,27,27,24,22,32,32,24,22
200 DATA34,34,32,42,42,38,38,34,30,24,22,20,14,14
210 DATA13,13,14,16,18,18,18,17,16,20,22,26,27,27,27,30,24,18
220 DATA17,17,17,17,18,16,14,13,19,22,26,27,27,34,28,26,22,22,22
230 DATA24,24,20,30,38,38,40,36,42,42,42,44,54,68,68,68,68,48,52,44,40,34,34,36
240 REM END OF PAGE 1
250 DATA38,44,32,32,34,40,28,28,30,36,26,27,32,22,24,30,32,36,38,42,40,36
260 DATA40,40,38,44,32,32,34,40,28,28,38,32,27,27,24,21
270 DATA20,20,20,20,20,20,18,20,19,17,17,17,17,15,24,26,27,30,32,34,36,0
280 DATA40,0,0,0,0,0,0,0,0,999
290 REM END OF PAGE 2
300 DATA60,52,48,40,34,36,42,60,52,48,48,40,34,36,42,42
310 DATA60,52,48,48,40,34,36,42,42,60,52,52,48,40,34,36,36,42
320 DATA54,52,42,36,30,32,38,54,54,52,52,42,36,30,32,38,38
330 DATA54,52,42,36,30,30,32,38,54,52,42,36,30,32,38
340 DATA48,44,38,32,27,27,28,28,34,48,44,44,38,38,32,32,27,28,34,34
350 DATA48,48,44,38,32,32,27,28,34,48,44,44,38,38,32,32,27,27,28,34
360 DATA38,36,36,36,52,52,52,52,54,64,84,68,68,68,68
370 DATA68,80,60,54,42,48,40,44,52,40,54,40,57,72,76,78,44,52,57,64,48,40,54,57
380 DATA60,57,48,40,34,36,42,60,57,48,48,40,34,36,42,42
390 DATA60,57,48,48,40,34,36,42,42,64,68,76,76,48,36,40,40
400 DATA40,57,54,48,60,64,42,52,52,52,48,54,60,68
410 DATA76,72,60,52,42,44,54,54,76,76,72,72,60,52,42,44,44,54
420 DATA40,42,52,60,60,72,72,68,68,57,57,40,42,52,52,60,60,72,68,57
430 DATA38,38,38,44,44,54,54,64,64,60,52,52,36,36,38,44,54,54,64,60,80,52,52
440 REM END OF PAGE 1
450 DATA54,54,72,48,48,48,64,42,42,40,38,38,38,52,52,34,34,34,44,44
460 DATA30,28,27,27,36,24,24,24,32,21,20,19,28,42,68,68
470 DATA60,57,48,40,34,38,42,60,57,48,48,40,34,36,42,42
480 DATA60,57,48,44,42,0,0,0,40,0,0,0,0,0,0,0,0,0,0,0,999
490 REM END OF PAGE 2
500 DATA8,8,8,8,8,8,8,8,4,4,8,8,8,4,4
510 DATA8,8,4,4,8,8,8,8,4,4,8,8,4,4,8,8,4,4,4,8
520 DATA8,8,8,8,8,8,8,8,4,4,4,4,8,8,8,8,4,4,4
530 DATA8,8,8,8,4,4,8,8,8,8,8,8,8,4,4,8,8
540 DATA8,8,8,4,4,4,4,4,4,4,8,8,8,4,4,4,4,4,4,8,8,4,4
550 DATA4,4,4,4,8,8,8,8,8,8,8,8,4,4,4,4,4,4,8,8
560 DATA8,8,8,8,8,8,8,8,8,8,8,8,8,8,8
570 DATA4,4,4,4,4,4,4,4,4,4,4,4,8,8,8,8,4,4,4,4,8,4,4
580 DATA8,8,8,8,8,8,8,8,8,8,8,8,4,4,8,8,8,4,4
590 DATA8,8,4,4,8,8,8,4,4,8,8,4,4,8,8,8,8
600 DATA8,8,8,8,8,8,8,8,8,8,8,8,8,8,8
610 DATA8,8,8,8,8,8,4,4,4,4,4,4,8,8,8,8,4,4,8
620 DATA8,8,8,4,4,4,4,4,4,4,4,4,8,8,8,4,4,4,4,8,8,8
630 DATA8,4,4,4,4,4,4,4,4,4,8,8,4,4,4,4,8,8,4,4,4,4,4,4
640 REM END OF PAGE
650 DATA4,4,8,8,4,4,8,8,8,8,4,4,8,8,4,4,4,4,4,4,4,4
660 DATA8,8,4,4,8,8,8,4,4,8,8,8,8,8,8,8,4,4
670 DATA8,8,8,8,8,8,8,8,8,8,4,4,8,8,8,8,4,4,8,8,8,4,4,8
680 DATA8,8,8,8,8,8,8,8,999

```

```

10 OUT 67,128 : REM SET 8255 PORT
20 A$=INKEY$ : REM SCAN KEYBOARD
30 IF A$="" THEN 20 : REM LOOP IF NO KEY
40 OUT 66,1 : REM ENVELOPE ON SHOT
50 OUT 66,0 : REM ENVELOPE OFF SHOT
60 GOTO 20 : REM LOOP AS NEEDED

```

The envelope should be triggered each time you touch a key on the TRS-80. If the envelope does not trigger, move the interface connection to the negative-going envelope, and try again. If the envelope is still not working, increase the values for C3 or R5, which are found at pins 14 and 15 of Z8. This will lengthen the trigger cycle, and should handle any synthesizer input. Again, try both the positive and negative envelopes.

Now disconnect the envelope trigger, and patch the envelope out of the synthesizer. Try the following program which creates a series of fast-rising whoops if the digital-to-analog converter is working properly:

```

10 OUT 67,128 : REM SET UP 8255 PORT
20 INPUT"TIME DELAY";N : REM GET INTERNOTE DELAY
30 FOR X = 0 TO 255 : REM SET TO RUN ALL NOTES
40 OUT 64,X : REM SEND VOLTS TO SYNTHO
50 FOR Y = 1 TO N : REM SET UP TIMING LOOP
60 NEXT : REM TIME, GET NEXT NOTE
70 GOTO 20 : REM BACK FOR NEXT DELAY

```

Each time you run through the program to the INPUT statement, increase the value for the time delay; the whoops will slow until you can hear a series of discrete pitches. If you have a two-voice synthesizer (or to test the second voice of the interface), connect the voltage output from voice 2, and use OUT 71,128 to set the port and OUT 68,X for the data.

If the pitches fall instead of rise, then your synthesizer responds to a higher voltage as a *lower* pitch, and any music routines you write will have to take this into account. If this is the case, before you run change line 60 to read:

```
60 OUT 66,32 : OUT 66,0 : OUT 64,(255-ML(1,PH))
```

Finally, run Listing 7-4, a rendition of keyboard prelude 23 from Johann Sebastian Bach's Well-Tempered Clavier.

```

10 GOSUB 650 : REM * THE TUNING SECTION IS RETRIEVED FIRST
20 DEFINT V, A, B, C, O, T, X, Y, P, N : REM * INTEGERS ALL
30 V1 = 64 : REM * VOICE ONE USES OUTPUT LOCATION 64
40 V2 = 65 : REM * VOICE TWO USES OUTPUT LOCATION 65
50 EN = 66 : REM * ENVELOPE TRIGGER USES OUTPUT LOCATION 66
60 TN = 32 : REM * TRIGGER TURNED ON WITH VALUE 32
70 TF = 0 : REM * TRIGGER TURNED OFF WITH VALUE 0
80 CT = 67 : REM * 8255 PIA CONTROL REGISTER AT LOCATION 67
90 PT = 128 : REM * PORT OUTPUT ONLY DEFINED BY VALUE 128
100 NN = 200 : REM * NUMBER OF NOTES DEFINED BY VALUE 200
110 DIM V1(2,500) : REM * ARRAY OF NOTES AND THEIR DURATIONS
120 DIM V2(2,500) : REM * ARRAY OF NOTES AND THEIR DURATIONS
130 X=1 : REM * TEMPORARY COUNTER ONLY
140 READ V1(1,X) : REM * READ VOICE ONE
150 READ V1(2,X) : REM * READ VOICE TWO
160 IF V1(2,X) = 0 THEN 170 ELSE X = X + 1 : GOTO 140
170 X=1 : REM * TEMPORARY COUNTER ONLY
180 READ V2(1,X) : REM * READ VOICE TWO
190 READ V2(2,X) : REM * READ VOICE TWO
200 IF V2(2,X) = 0 THEN 210 ELSE X = X + 1 : GOTO 180
210 A=1 : B=1 : C=1 : D=1 : X=1 : Y=1 : Q=0 : R=2
220 OUT CT,PT : OUT V1,255 : OUT V2,255
230 G = V1(R,X) : H = V2(R,Y) : GOSUB 330 : GOSUB 350
240 G = G - A : IF G <= 0 THEN 250 ELSE 270
250 X = X + A : GOSUB 330 : G = V1(R,X)
260 GOTO 280
270 FOR W = 1 TO 14 : NEXT
280 H = H - A : IF H <= 0 THEN 290 ELSE 310
290 Y = Y + A : GOSUB 350 : H = V2(R,Y)
300 GOTO 320
310 FOR W = 1 TO 14 : NEXT
320 IF X > 200 THEN 210 ELSE 240
330 IF V1(A,X) = Q THEN 340 ELSE OUT EN,TN : OUT EN,TF :
OUT V1,V1(A,X) : RETURN
340 FOR W = 1 TO 2 : NEXT : RETURN

```

Listing Continued...

Continued Listing

```

350 IF V2(A,Y) = Q THEN 350 ELSE OUT EN,TN : OUT EN,TF :
    OUT V2,V2(C,Y) : RETURN
360 FOR W = 1 TO 2 : NEXT : RETURN
370 DATA1,4,105,2,102,2,87,4,114,4,102,4,120,2,117,2,132,4,126
380 DATA4,90,4,96,6,93,2,96,4,123,2,117,2,123,4,132,8,138,2,111
390 DATA2,117,2,132,2,129,2,132,2,138,4,144,2,135,2,138,4,138,4
400 DATA108,4,123,4,87,4,117,4,93,4,96,2,90,2,87,2,84,2,87,12
410 DATA108,4,102,12,132,4,129,4,138,4,135,4,129,4,123,4,111,2
420 DATA120,2,123,4,87,4,108,4,75,2,81,2,87,4,108,2,105,2,108
430 DATA2,132,2,138,2,123,2,126,4,111,4,96,4,99,4,87,4,90,8,83
440 DATA8,102,8,105,8,108,4,132,4,93,4,138,4,81,4,105,2,102,2
450 DATA87,4,114,4,102,4,120,2,117,2,132,4,126,4,90,4,87,2,90,2
460 DATA96,2,102,2,96,2,102,2,105,2,87,2,93,2,129,2,132,2,126,2
470 DATA123,2,126,2,120,2,117,2,111,2,105,2,102,2,111,2,108,2
480 DATA111,2,120,2,117,2,111,4,138,2,117,2,123,4,121,2,105,2
490 DATA132,2,126,2,120,2,117,2,111,4,132,2,129,2,132,4,102,4
500 DATA81,4,87,2,90,2,96,2,99,2,96,2,117,2,111,2,108,2,111,4
510 DATA114,4,111,16,0,0,0,36,81,4,105,2,102,2,87,4,114,4,102
520 DATA4,120,2,117,2,132,4,126,4,90,4,96,6,93,2,96,4,123,2,117
530 DATA2,123,4,132,8,138,2,111,2,117,2,132,2,129,2,132,2,138,4
540 DATA144,2,135,2,138,4,138,4,108,4,123,4,87,4,117,4,93,4,96
550 DATA2,80,2,87,2,84,2,87,12,108,4,102,12,132,4,129,4,138,4
560 DATA135,4,129,4,123,4,111,2,120,2,123,4,87,4,105,4,75,2,81
570 DATA2,87,4,108,2,105,2,108,2,132,2,126,2,123,2,126,4,111,4
580 DATA96,4,99,4,87,4,90,8,98,8,102,8,105,8,108,4,132,4,93,4
590 DATA138,4,81,4,105,2,102,2,87,4,114,4,102,4,120,2,117,2,132
600 DATA4,126,4,90,4,87,2,90,2,95,2,102,2,95,2,102,2,105,2,87,2
610 DATA93,2,129,2,132,2,126,2,123,2,126,2,120,2,117,2,111,2
620 DATA105,2,102,2,111,2,108,2,111,2,120,2,117,2,111,4,138,2
630 DATA117,2,123,4,141,2,105,2,132,2,126,2,120,2,117,2,111,4
640 DATA132,2,129,2,132,16,0 : REM * END OF TWO VOICES
850 REM * THE TUNING SECTION PRODUCES OCTAVE PITCHES
660 OUT 67,128 : REM * SET UP THE 8255 PIA FOR ACTION
670 OUT 64,3 : OUT 65,3 : FOR N = 1 TO 1000 : NEXT
680 OUT 64,39 : OUT 65,38 : FOR N = 1 TO 1000 : NEXT
690 OUT 64,75 : OUT 65,75 : FOR N = 1 TO 1000 : NEXT
700 OUT 64,111 : OUT 65,111 : FOR N = 1 TO 1000 : NEXT
710 OUT 64,147 : OUT 65,147 : FOR N = 1 TO 1000 : NEXT
720 A$ = INKEY$ : IF A$ = "" THEN 670 ELSE RETURN

```

```

10 CLS : PRINTCHR$(23) "-" GENERALIZED MUSIC PROGRAM "-"
20 PRINT : PRINT "OENNIS BATHORY KIT52, MAY 1979"
30 FOR X = 1 TO 1000 : NEXT : CLEAR : CLS
40 PRINT"HERE ARE 539 NOTES IN BACH PRELUDE NO. 23" : D = 639
60 GOTO 400 : REM * MOVE TO TUNING ROUTINE BEFORE PLAYING PIECE
80 DIM ML(2,0+50) : REM * EXTRA BYTES SET ASIDE FOR LONGER WORK
70 PRINT : PRINT"NOW READING":0:"NOTES FROM NOTE/RHYTHM DATA."
80 FOR PH = 1 TO D : REM * PH SPECIFIES PITCH & DURATION ARRAY
90 FOR NT = 1 TO 2 : REM * NT SPECIFIES THE TWO ARRAY ELEMENTS
100 READ ML(NT,PH) : REM * EACH PITCH AND RHYTHM FILLS ARRAY
110 NEXT NT,PH : REM * READ ENTIRE ARRAY OF 639 NOTES & RHYTHMS
120 INPUT "PRESS ENTER TO START PIECE":Z : REM * WAIT FOR START
130 FOR PH = 1 TO D : REM * READ ARRAY LOOP TO PLAY BACK MUSIC
140 OUT 55,32 : OUT 68,0 : REM * ENVELOPE TRIGGER ON THEN OFF
150 OUT 64,ML(1,PH) : REM * THEN SEND PITCH TO D/A CONVERTER
160 FOR X = 1 TO ML(2,PH) : NEXT : REM * ARRAY DURATION DELAY
170 NEXT PH : REM * AND CONTINUE UNTIL ALL NOTES ARE PLAYED OUT
180 INPUT "PRESS ENTER TO PLAY AGAIN":Z : REM * OPTIONAL REPEAT
190 GOTO 120 : REM * ANOTHER PIECE MAY BE ADDED AND RUN HERE
200 DATA89,31,80,31,117,31,90,31,86,31,90,31,111,31,90,31,69,
    31,80,31,117,31,90,31,54,31,102,31,128,31,102,31,60,31,
    81,31,105,31,81,31,54,31,81,31,102,31,81,31,60,31,81,31,
    105,31,81,31,45,31,90,31,117,31,90,31
210 DATA48,31,88,31,98,31,69,31,45,31,69,31,90,31,69,31,48,31,
    69,31,86,31,88,31,33,31,81,31,105,31,81,31,39,31,81,31,84,
    31,81,31,33,31,81,31,95,31,81,31,30,31,75,31,90,31,75,31,
    18,31,65,31,84,31,66,31
220 DATA88,31,81,31,90,31,81,31,54,31,69,31,81,31,69,31,45,31,
    54,31,59,31,54,31,33,31,45,31,54,31,45,31,24,31,45,31,51,
    31,57,31,60,31,66,31,69,31,75,31,69,31,60,31,66,31,69,31,
    75,31,81,31,87,31,80,31
230 DATA75,31,87,31,98,31,87,31,60,31,75,31,87,31,75,31,51,31,
    80,31,75,31,60,31,39,31,51,31,60,31,51,31,30,31,51,31,54,
    31,60,31,68,31,72,31,75,31,81,31,75,31,66,31,72,31,75,31,
    81,31,87,31,90,31,86,31
240 DATA81,31,90,31,102,31,90,31,66,31,81,31,90,31,81,31,54,31,
    88,31,81,31,86,31,45,31,54,31,66,31,54,31,33,31,54,31,81,
    31,54,31,24,31,54,31,69,31,54,31,30,31,54,31,75,31,54,31,
    18,31,54,31,84,31,54,31
250 DATA33,31,54,31,81,31,54,31,24,31,54,31,69,31,54,31,30,31,
    54,31,75,31,54,31,18,31,54,31,84,31,54,31,33,31,54,31,81,
    31,54,31,36,31,80,31,90,31,50,31,39,31,50,31,87,31,50,31,
    42,31,68,31,88,31,66,31
260 DATA45,31,86,31,80,31,66,31,9,31,69,31,102,31,69,31,15,31,
    89,31,85,31,58,31,51,31,75,31,105,31,75,31,54,31,75,31,
    102,31,75,31,18,31,81,31,111,31,81,31,24,31,81,31,105,31,
    81,31,80,31,87,31,117,31,87,31
270 DATA88,31,75,31,81,31,87,31,90,31,98,31,102,31,105,31,111,
    31,105,31,102,31,96,31,90,31,84,31,81,31,75,31,81,31,69,
    31,88,31,60,31,54,31,48,31,45,31,39,31,33,31,39,31,45,31,
    38,31,33,31,30,31,24,31,18,31
280 DATA15,31,75,31,105,31,75,31,18,31,75,31,102,31,75,31,9,31,
    88,31,102,31,88,31,15,31,69,31,96,31,69,31,18,31,66,31,96,
    31,66,31,30,31,66,31,80,31,66,31,33,31,60,31,90,31,60,31,
    39,31,50,31,87,31,80,31
290 DATA18,31,75,31,80,31,75,31,66,31,90,31,75,31,66,31,54,31,
    88,31,75,31,68,31,54,31,75,31,56,31,54,31,39,31,54,31,51,
    31,45,31,38,31,33,31,90,31,24,31,18,31,30,31,39,31,54,35,
    68,39,75,42,84,45,93,48
300 DATA8,250,46,15,54,15,89,15,81,15,99,15,105,15,117,15,128,

```

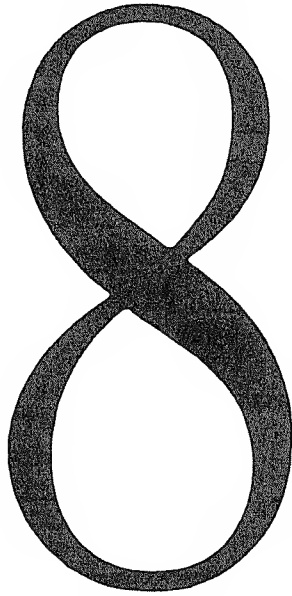
Continued Listing

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200,45,15,54,15,89,15,81,15,99,15,105,15,117,15,126,50,38,
    15,45,15,80,15,85,15,108,15,117,15,126,131,132,31,135,31,
    132,31,126,31,120,31,117,31,120,31
310 DATA128,31,123,31,117,31,111,31,108,31,111,31,117,31,111,
    31,108,31,102,31,86,31,102,31,108,31,102,31,98,31,80,31,
    84,31,81,31,75,31,68,31,83,31,60,31,54,31,48,31,45,31,39,
    31,38,31,30,34,24,37,18,40,12,44,8,48,3,250
320 DATA39,15,48,15,60,15,75,16,98,15,111,15,120,200,39,15,48,
    15,50,15,75,15,98,15,105,15,111,15,120,50,54,15,66,15,75,
    15,90,15,102,15,111,15,120,131,54,31,38,31,45,31,48,31,54,
    31,60,31,86,31,69,31
330 DATA68,31,80,31,54,31,80,31,58,31,89,31,75,31,81,31,75,31,
    69,31,66,31,68,31,75,31,81,31,84,31,80,31,84,31,81,31,75,
    31,81,31,84,31,80,31,88,31,102,31,105,31,111,31,117,31,
    120,34,126,37,132,40,138,44,125,48
340 DATA141,250,33,15,45,15,54,15,89,15,80,15,102,15,105,15,
    117,200,33,15,45,15,54,15,88,15,80,15,102,15,105,15,117,
    50,48,15,80,16,88,16,84,15,98,15,106,15,117,250,51,15,80,
    15,88,15,85,15,106,15,114,200,51,15,60,15,75,15,88,15,106
350 DATA15,111,50,54,15,88,15,75,15,84,15,80,15,102,15,111,250,
    102,15,111,31,102,31,90,31,84,31,80,31,102,31,111,31,102,
    31,111,31,102,31,80,31,84,31,80,31,102,31,111,31,129,31,
    126,31,120,31,114,31,111,31,105,31,102,31,88,31
360 DATA102,31,80,31,102,31,111,31,120,31,128,31,138,31,147,31,
    141,89,111,31,114,31,102,31,105,31,87,31,18,15,38,15,48,
    15,54,15,75,15,80,15,105,15,111,125,81,125,18,15,39,15,48,
    15,54,15,84,16,80,15,102,115,111,031,54,125,80,125
370 DATA33,15,45,15,54,16,63,15,81,15,90,15,105,250,81,31,90,
    31,81,31,88,31,56,31,88,31,81,31,80,31,81,31,80,31,81,31,
    68,31,68,31,68,31,81,31,80,31,108,31,89,31,80,31,81,31,72,
    31,66,31,54,31,81,31
380 DATA60,15,84,15,81,15,84,15,81,16,84,125,81,31,76,31,81,31,
    84,31,81,31,75,31,88,31,66,31,58,31,75,31,89,31,88,31,80,
    31,54,31,60,31,55,31,50,31,54,31,48,31,45,31,48,31,54,31,
    48,31,45,31,39,31,33,31,30,31
390 DATA33,31,45,31,54,31,45,31,45,31,54,31,89,31,54,31,54,31,
    69,31,81,31,89,31,69,31,81,31,80,31,81,31,81,31,90,31,105,
    31,90,31,80,31,105,31,117,33,105,35,105,37,117,39,126,42,
    117,45,141,150
400 REM * * * * * TUNING SECTION FOLLOWS * * * * *
410 PRINT"THIS IS THE TUNING SECTION -- OUTPUT IS 3, 38, 75."
420 PRINT"HOLO SPACE 8AR TO BEGIN READING BACH PRELUDE DATA."
430 OUT 64,3 : FOR X = 1 TO 500 : NEXT : REM * OUTPUT LOW C
440 OUT 64,39 : FOR X = 1 TO 500 : NEXT : REM * NEXT OCTAVE
450 OUT 64,75 : FOR X = 1 TO 500 : NEXT : REM * NEXT OCTAVE
460 OUT 64,111 : FOR X = 1 TO 500 : NEXT : REM * NEXT OCTAVE
470 OUT 64,147 : FOR X = 1 TO 500 : NEXT : REM * TOP OCTAVE
480 A$ = INKEY$ : IF A$ = "" THEN 480 : REM * CHECK KEYBOARD
490 OUT 58,255 : REM * HOPE THAT VERY TOP NOTE IS INAUDIBLE
500 GOTO 60 : REM * NOW RETURN TO MAIN MUSIC READING ROUTINE

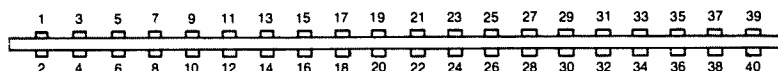
```

Listing 7-4. Johann Sebastian Bach.



P/N	SIGNAL NAME	DESCRIPTION
1	RAS*	Row Address Strobe Output for 16-Pin Dynamic Rams
2	SYSRES*	System Reset Output, Low During Power Up Initialize or Reset Depressed
3	CAS*	Column Address Strobe Output for 16-Pin Dynamic Rams
4	A10	Address Output
5	A12	Address Output
6	A13	Address Output
7	A15	Address Output
8	GND	Signal Ground
9	A11	Address Output
10	A14	Address Output
11	A8	Address Output
12	OUT*	Peripheral Write Strobe Output
13	WR*	Memory Write Strobe Output
14	INTAK*	Interrupt Acknowledge Output
15	RD*	Memory Read Strobe Output
16	MUX	Multiplexor Control Output for 16-Pin Dynamic Rams
17	A9	Address Output
18	D4	Bidirectional Data Bus
19	IN*	Peripheral Read Strobe Output
20	D7	Bidirectional Data Bus
21	INT*	Interrupt Input (Maskable)
22	D1	Bidirectional Data Bus
23	TEST*	A Logic "0" on TEST* Input Tri-States A0-A15, D0-D7, WR*, RD*, IN*, OUT*, RAS*, CAS*, MUX*
24	D6	Bidirectional Data Bus
25	A0	Address Output
26	D3	Bidirectional Data Bus
27	A1	Address Output
28	D5	Bidirectional Data Bus
29	GND	Signal Ground
30	D0	Bidirectional Data Bus
31	A4	Address Bus
32	D2	Bidirectional Data Bus
33	WAIT*	Processor Wait Input, to Allow for Slow Memory
34	A3	Address Output
35	A5	Address Output
36	A7	Address Output
37	GND	Signal Ground
38	A6	Address Output
39	GND	Signal Ground
40	A2	Address Output

NOTE: *means Negative (Logical "0") True Input or Output



Mates with AMP P/N 88103-1 Card
Edge Connector or Equivalent

Figure 8-1. TRS-80 edge card connector.

Adding to the System

In this Chapter we will explore several hardware projects to expand the TRS-80's capabilities :

A parallel printer interface for keyboard units without an expansion interface.

An expansion of system RAM and ROM in a 'reserved' block in the memory map.

Bank selection of RAM and ROM in that 'reserved' blank area, and a memory expansion that includes bank selected RAM.

A programmable input/output port device and a companion interrupt I/O board.

Battery backup and real time clock.

Parallel Printer Interface

Perhaps the simplest addition to the keyboard unit is a printer interface. Radio Shack sells a complete cable for this purpose, but building one is both less expensive and more enlightening.

When an LLIST or LPRINT command is entered, the computer enters a subroutine which plucks each character to be printed, checks its value, and converts it if necessary. Line feeds, for example, are converted to carriage returns, and form feeds are converted to the proper number of line feeds according to the value in the printer's device control block.

The final value is written to an address, almost exactly as if it were being stored in memory. The hardware decodes that memory

address and sends the character along parallel data lines to the printer. The printer then sends a 'busy' signal to the memory address. The LPRINT or LLIST routines read this busy signal, wasting time until the printer is ready. When the busy signal turns off, the next character is plucked, converted and sent to the printer. The process continues until all characters are printed.

There are disadvantages to this system; among them the fact that the printing process 'hangs up' the computer for the duration. This

can easily be avoided in an interrupt-driven system (see Supplement to Chapter 5), and even modified in a less successful way as a background routine to the keyboard scan. But the overall process is simple, requiring neither sophisticated hardware nor complex software.

Figure 8-1 presents the complete circuit for a printer interface board. It consists of three sections :

1. An address decoder for 37E8, the printer location.

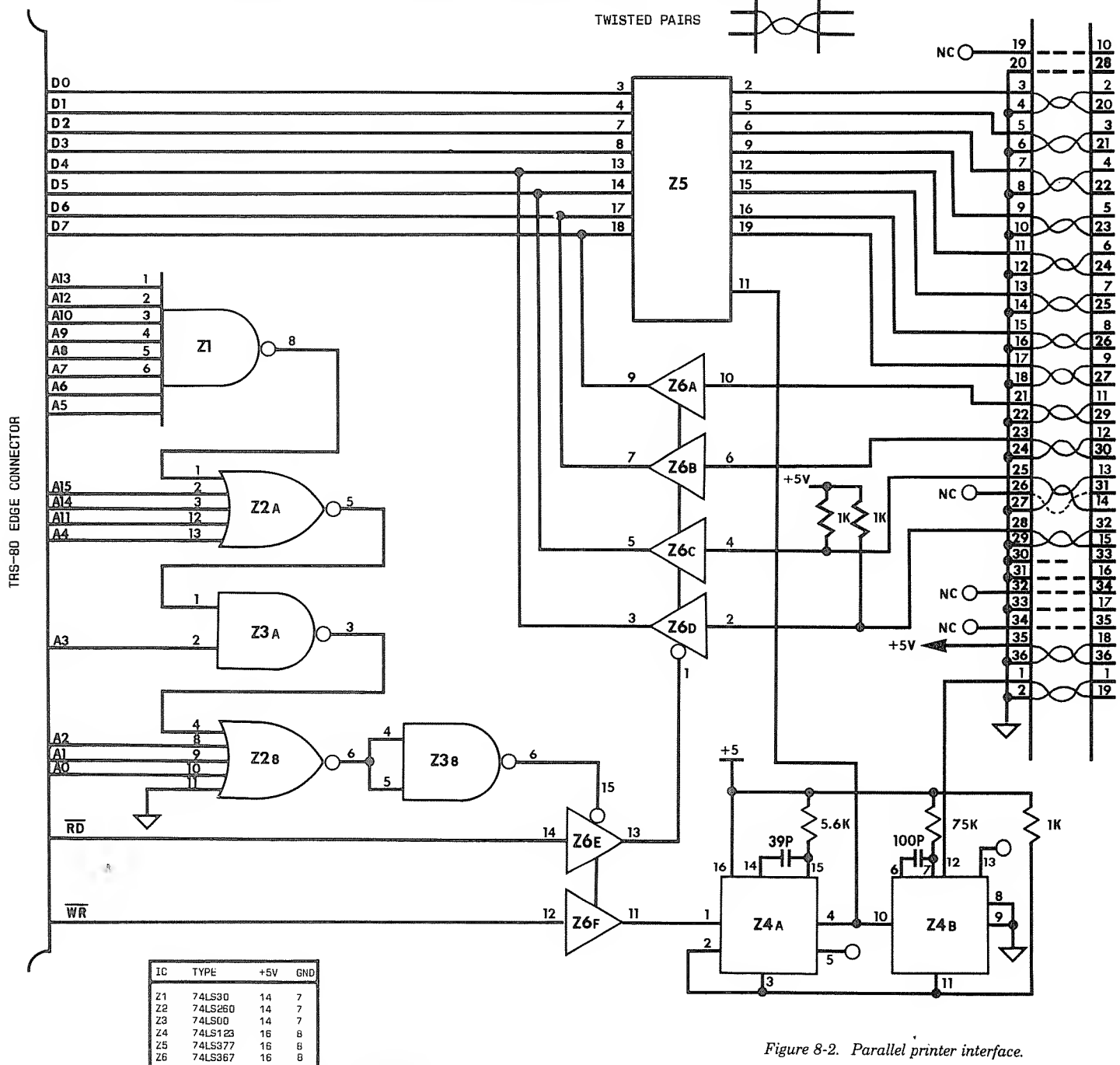


Figure 8-2. Parallel printer interface.

2. An 8-bit output buffer for sending characters to the printer.
3. An input buffer for checking the status of the printer (busy or out of paper).

The address decoder is formed by Z 1 and Z 2 evaluating the address lines in conjunction with RD (read) or WR (write) signals from the computer. If WR is sent, buffer Z 6 is activated, sending a character to the printer. If RD is sent, buffer Z 6 is activated, sending the status of the printer to the CPU.

Power-up the TRS-80, and go through the input sequence. If this process does not act normally, *turn the computer off immediately* and recheck the wiring. If all is normal, it's time for a printer test. Turn the printer on (don't be surprised by a return to MEMORY SIZE? if your printer is an older, electrically noisy one). The simplest way to test the printer interface is to load any program and LLIST it. The printer should spring to life, printing a complete list. If there are problems, they will probably be among the following :

No characters printed at all; computer immediately (or after a short pause) returns to ready.

Characters are being sent to the memory address, but not being received by the printer. Therefore, no 'busy' is being received, and the computer dumps all its characters as fast as it can.

Solution : check wiring of the port address, wiring to the output buffer select line, and see that the board has both power and ground wires connected.

No characters printed at all; computer immediately locks up :

Characters are being sent to the memory address, but if none are printed, they are not being received by the printer. The computer is seeing a constant 'busy' signal, and thus is waiting in a loop.

Solution : check wiring of the port address, wiring to the input buffer from the printer, and see that the printer is enabled (if it has an enable function).

Intermittent but regular characters (every third or fifth or fiftieth character, for example) are being printed.

Characters are being sent through to the

printer, but no handshake ('busy') is being received by the computer.

Solution : check wiring of the input buffer select line, and the wiring from the printer's busy signal. Some printers may not have a busy signal; see box for suggestions..

A single character is printed, then printing stops and the computer locks up.

A constant busy is being received by the computer, and it is waiting in a loop for the busy signal to terminate.

Solution : check for ground shorts in the printer's busy line, or shorts to ground at the input buffer.

Incorrect characters are printed, and none or any of the above symptoms are present.

The data lines are incorrectly wired to the printer or to the board's output buffer. The printer may be wired for complement ASCII.

Solution : check the wiring of the data lines for reversed wires, either at the computer or printer end. If all is well, enter the following short program :

```
10 FOR X = 65 TO 91
20 Y = NOT X AND 255
30 POKE 14312,Y
40 FOR N = 1 TO 100
50 NEXT : NEXT
```

Listing 8-1. Printer interface test routine.

This program produces the complement of the letters from A to Z. If the correct letters are printed this time, replace Z 5 with an inverting buffer, type 81LS96.

Talking with the World – The Computer as Boss

The TRS-80 has remarkable skills for controlling the world around it. Four BASIC commands (POKE, PEEK, INP, and OUT) and their machine language equivalents (LD register, LD memory, IN register, OUT port) are the software conversational tools by which the computer makes its wishes known.

Only one dilemma remains: very little hardware was provided with the TRS-80 to use these powerful features. Just a single port (255) was hard-wired in place, and it is limited to controlling cassette functions and video display size. Memory mapped input/output was left exclusively to the expansion interface. And even then then, only for a printer, dual cassette, RS-232 and disks. In each case, no uncommitted user ports or memory addresses were provided.

Fortunately, creating such input/output (I/O) is not difficult. There are two very effective ways to accomplish it:

1. Using inexpensive, separate logic devices that can be dedicated to their interfacing tasks. The TRS-80 Technical Reference Handbook describes such simple hardware in its 'coffee pot' scenario.
2. Using more costly programmable interface devices (such as the INS8255 peripheral interface adaptor) for handling more flexible, general purpose I/O.

In either case, the input/output device must be identifiable by the computer, which means it must somehow be located. It is assigned a port number or a memory address. But what does this number mean, and how does it work? I have often used the analogy of the key in the lock, because it so well describes the way the electronics can open the doors to the world outside its case.

Figure 8-3 shows a simple-minded lock and key. It is simple minded because there are no fine graduations in the height of the tumblers – the 'pins' either rise to a single height, or are not present at all. In the computer, these pins are really voltages, represented by numbers.

In other words, the higher voltage can be called a 'one', and the lower voltage can be considered a 'zero'. In this way, the key number code shown in Figure 8-3 might look like this :

ON	ON	OFF	ON	ON	OFF	OFF	OFF
1	1	0	1	1	0	0	0

The sample key's code (binary 1101 1000) works out to the hexadecimal value D8, or the decimal equivalent 216 (refer to Chapter 2 for details on binary, decimal, and hexadecimal numbers). The 'key' is the value that the computer will output; the 'tumblers' are the hardware which will unlock when this key is inserted.

Below is the schematic of a general purpose 'tumbler' which can be adjusted to open to any electronic key. As noted in Chapter 2, the triangles are buffers which protect the TRS-80 electronic hardware from overexertion. Once again, the triangle with the 'not' circle at its point is an inverting buffer, which reverses the value of any signal placed at its input.

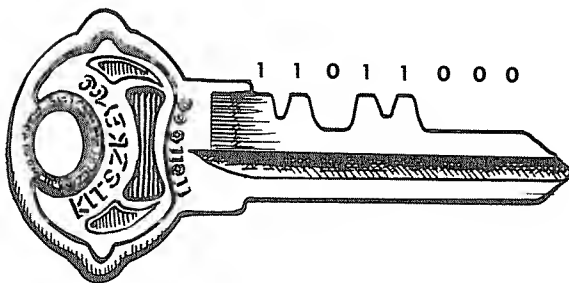


Figure 8-3. Lock-and-key illustration.

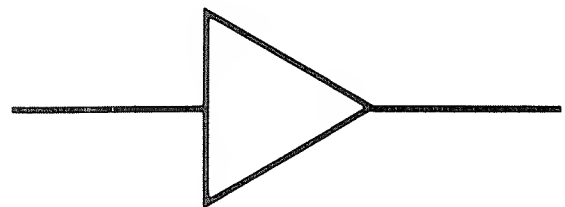


Figure 8-4. Unconnected buffer diagram.

This is very elementary – it will be improved on later – but it is the central scheme of digital operation. Real tumblers must all be lifted to a point so they are in line with the edge of the lock's cylinder. Electronically, the same thing must take place. To turn the electronic cylinder, all the binary input values must be lifted (or depressed) to the same value before the electronic lock will click open.

*****;

Using Exclusive-OR for decoding

Resistor 1K ohms to Ground; Switch to Plus Volts		
Switch Position	Input Value	Output Value
OFF (gate sees 0)	0	0
ON (gate sees 1)	0	1
OFF (gate sees 0)	1	1
ON (gate sees 1)	1	0

Resistor 1K ohms to Plus Volts; Switch to Ground		
Switch Position	Input Value	Output Value
OFF (gate sees 1)	0	1
ON (gate sees 0)	0	0
OFF (gate sees 1)	1	0
ON (gate sees 0)	1	1

Table 8-1. Using exclusive-OR for decoding.

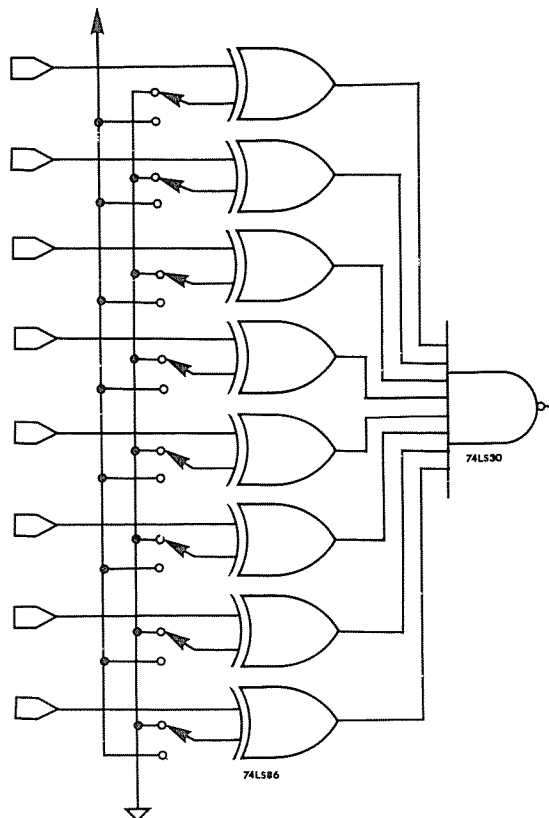


Figure 8-6. Improved port decoding for port addresses.

Figure 8-5 is the schematic for an elementary interface circuit. Z1 and Z2 form the tumblers, and Z3 is the cylinder.

There is another way to produce the same effect as Z1 and Z2, and no rewiring or jumpering is necessary. To do it, you can use one or two 74LS86 exclusive OR gates. Exclusive OR is a remarkable electronic function which states :

If two input signals are alike, the evaluated result will be set to zero. If two input signals are different, the evaluated output will be set to one.

Here's how that might work electronically. One input of an exclusive OR gate is attached to a switch and a resistor. The resistor is attached to ground, with the switch connected to the positive voltage line. When the switch is off, the input looks like a zero. When the switch is on, the lower resistance of the switch makes the input of the gate see a one.

Therefore, with the switch off, a signal coming into the second input of the gate would cause a one output when it is a one. With the switch on, a signal coming into the second input of the gate would cause a zero output when it is a one. Refer to Table 8-1.

By using eight exclusive OR gates (two type 74LS86 circuits) and an 8 position DIP (dual inline package) switch, any one of the 256 possible ports can be selected. The circuit below (Figure 8-6) shows how Z1 and Z2 in Figure 8-5 can be replaced with a switch and the 74LS86's to select the port. Thus, jumpering and soldering from the original Z1 and Z2 can be avoided.

Now, before putting these ports to use, it's time to turn to the other type of input/output device – the programmable interface adaptor. The previous I/O device costs perhaps \$5 to create. The central integrated circuit to the programmable I/O port itself costs about \$8, but it offers some extra features and easier wiring.

The INS8255 is a single integrated circuit capable of providing three complete input/output ports. Each port can act as an input or output, and that condition can be changed via programming. This is how it is done: using a decoder similar to that designed above, a 'chip select' is formed.

When the 8255 receives the chip select, it examines its two address line connections. Two address lines can be configured four ways (00, 01, 10, 11), and so can select one of the three I/O

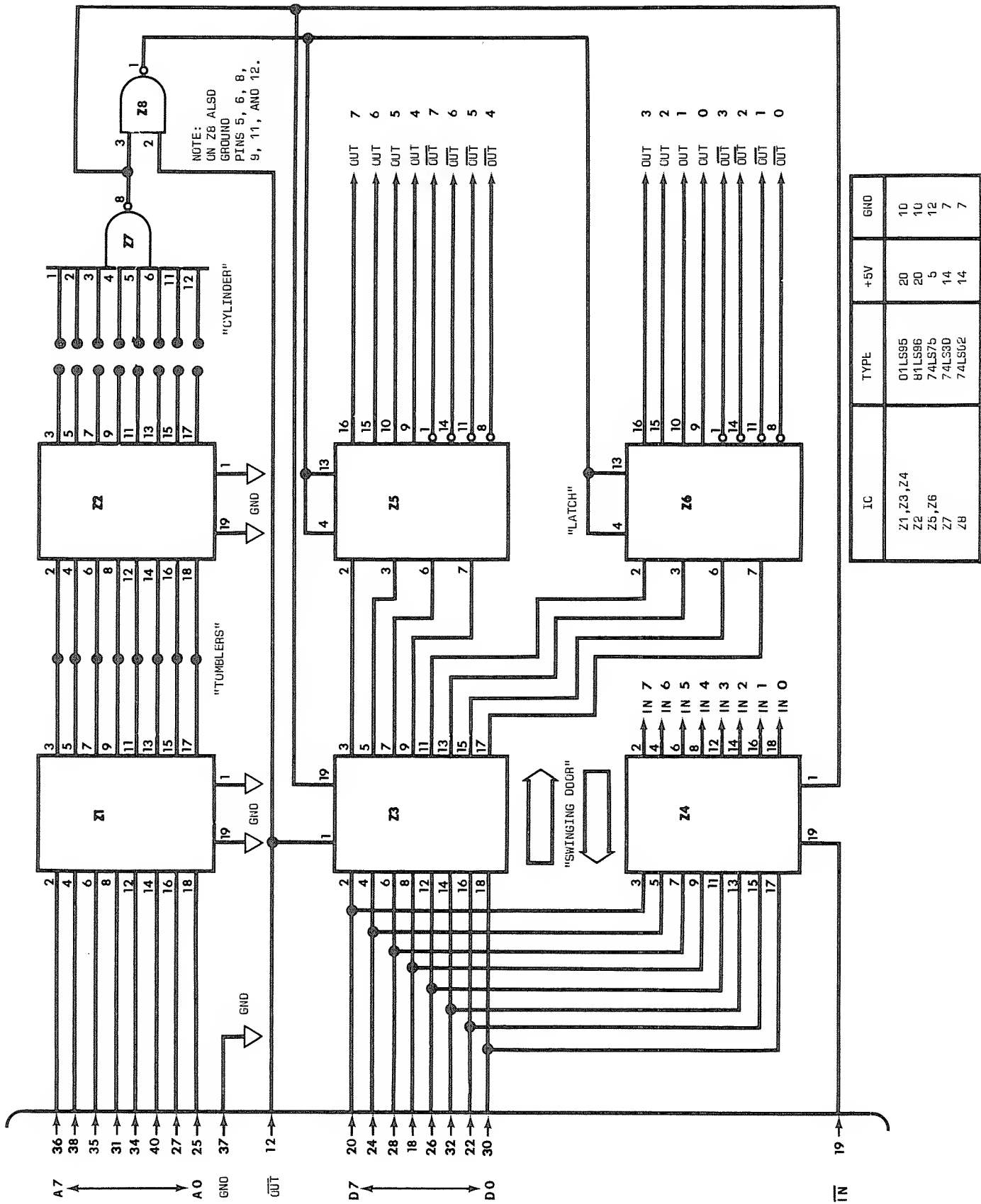


Figure 8-5. Elementary interface.

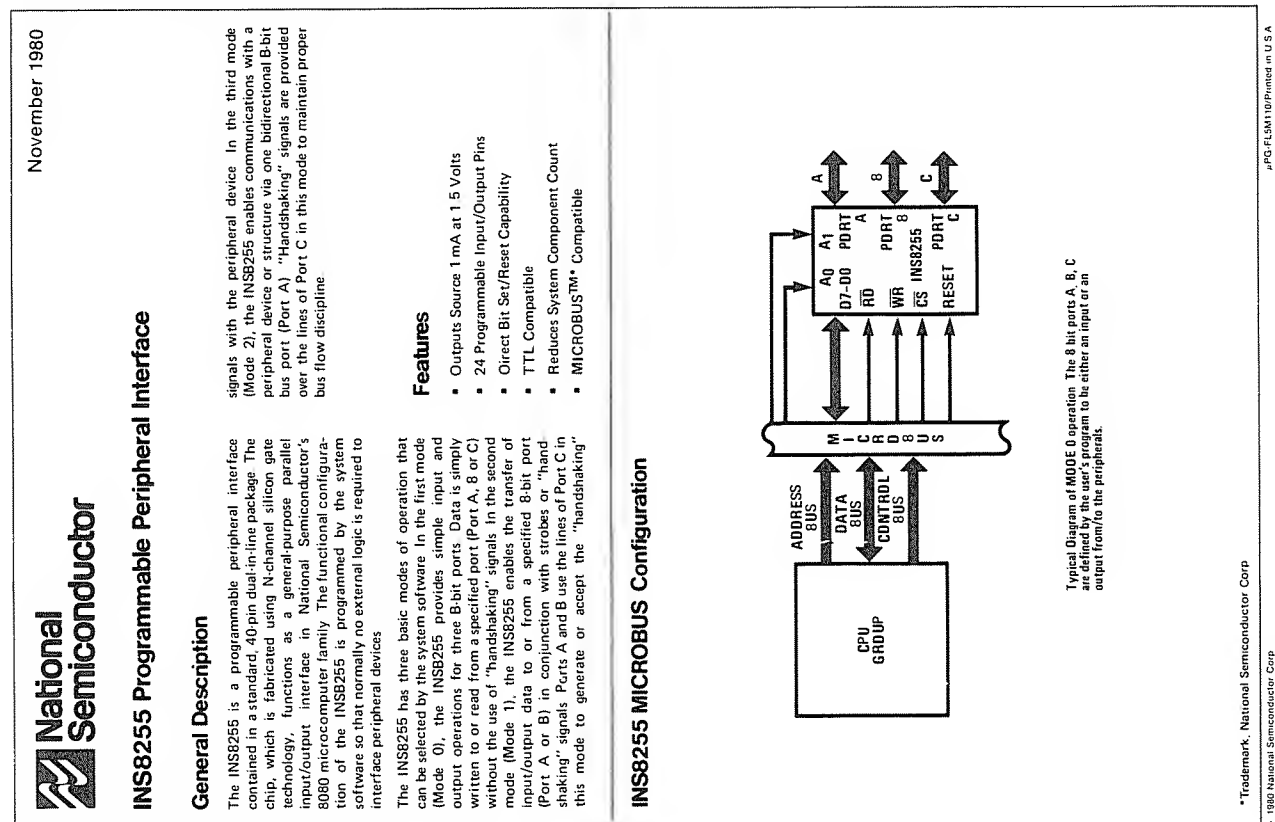


Figure 8-8. 8255 configuration (data).

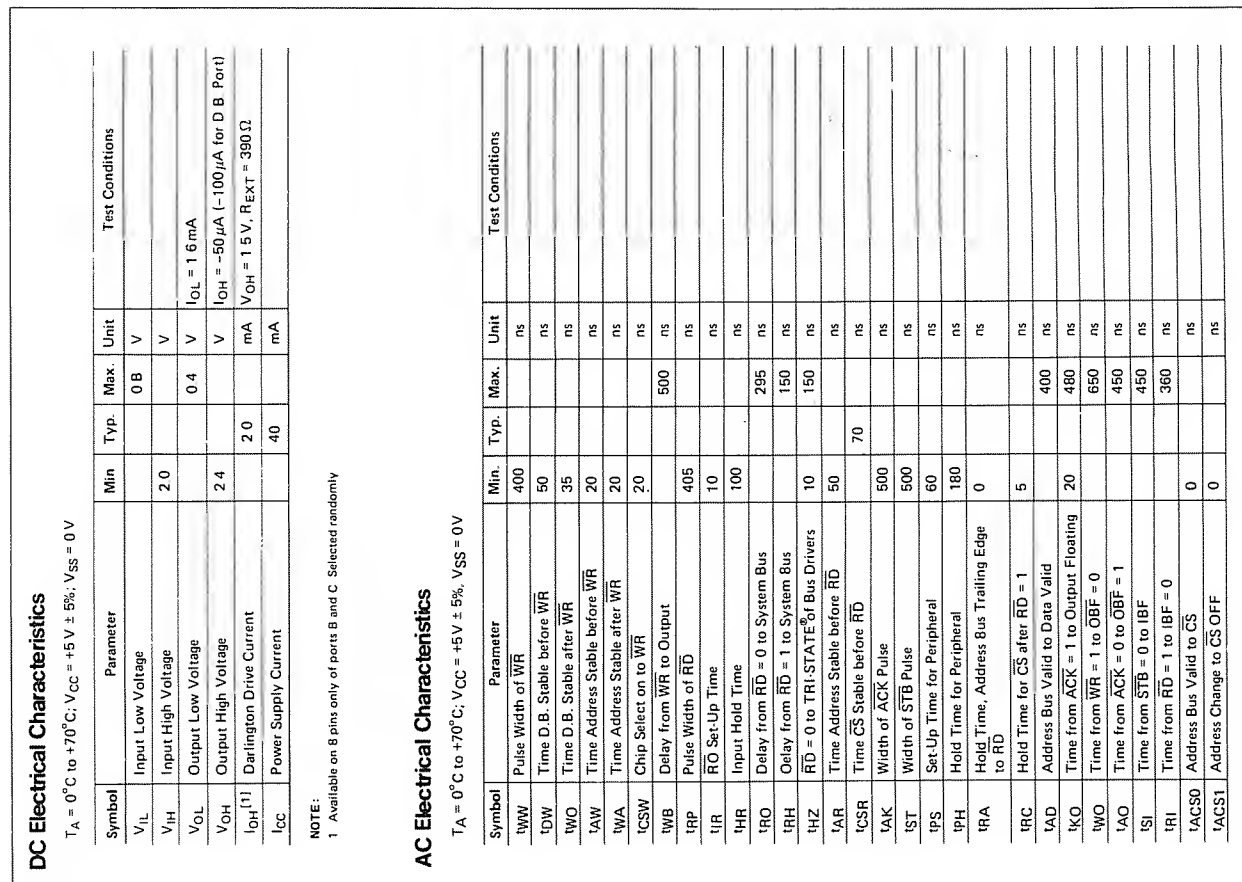


Figure 8-7. 8255 configuration (control).

ports, or – and this is the remarkable part – the 8255's internal *control register*. The 8255 is a smart chip!

When the control register has been selected, the 8255 can be programmed – its ports may be defined as input or output, and other combinations of actions can be selected. The various possibilities are shown in Figure 8-7 and 8-8.

Three Real-Time Clock/Calendars

Telling the time and date is a legitimate concern of computer users, not only for keeping documents in order, but also for observing and controlling experiments. Until recently, the only type of real-time clock available was the one built into the expansion interface, which disk and Level III users could access with the `TIME$` command.

The principle of that clock is simple: forty times each second, a pulse is sent from the expansion box to the TRS-80 keyboard unit, where it is applied to the interrupt (INT) line. The computer responds by temporarily setting aside its other processing activities in order to update the seconds, minutes, and hours. When the user asks for the `TIME$`, the computer checks the area of memory in which this updated information is stored, and sends it to the current display device.

The clocks presented in this section all use the `TIME$` function, but avoid certain problems associated with the expansion box `TIME$` function. First and foremost, it eliminates the need for the expansion box itself, and the special software to use the 25-millisecond interrupt.

Without disk, this special software must be loaded for every session, reloaded if an inadvertent reset should occur, updated should the computer be turned off, and disabled at every `CLOAD` and `CSAVE`. Though convenient in a disk system, this type of approach is more of an annoyance in a Level II system.

There are, however, two distinct advantages to an interrupt-based system for timekeeping: the hardware is simple and cheap to create, and in certain ways the clock pulses generated by the power line are more accurate than crystal-controlled clocks. The power line, because it is linked into a large network of generating systems, must maintain a virtually absolute synchronization over the long term. Short duration lags and leads in the 60-cycle pulses may appear, but the percentage of error over weeks, months, and years is virtually nil.

Time updating from a simple pulse train can be complicated, too, because seconds, minutes and hours, days, weeks, months and years aren't very 'decimal' in their counting. 59 seconds plus one second is one hour and zero seconds. 23 hours and one hour is one day and zero hours. You get the idea. So the software to update a pulse-based calendar needs special math, charts, and tables to keep the time and date in its electronic mind.

REF. OCT. '80
80 MICROCOMP
P. 20

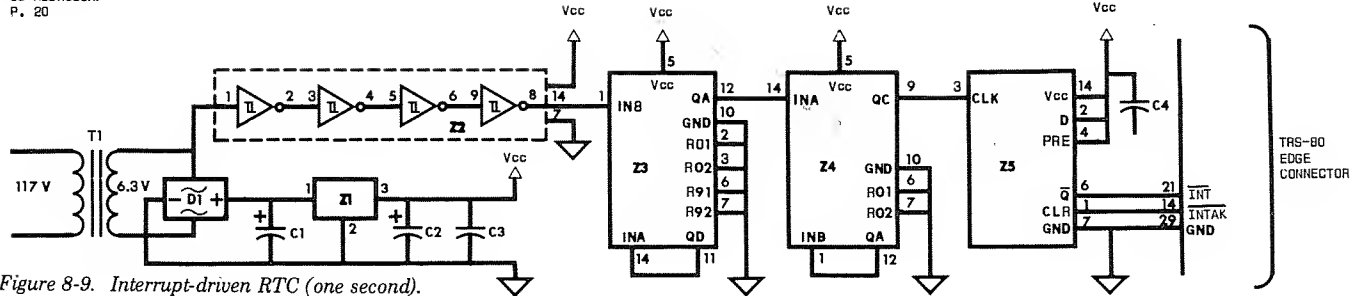


Figure 8-9. Interrupt-driven RTC (one second).

```

00100 ; #####
00110 ; THIS ROUTINE WILL USE THE ONE-SECOND INTERRUPT CREATED
00120 ; BY A HARDWARE ADDITION TO OO TIMEKEEPING. THIS IS
00130 ; TRANSPARENT TO BASIC. NOTE THAT THIS IS NOT INCLUDED
00140 ; IN THE SLASH [/] FUNCTIONS CALLED BY THE CUSTOM
00150 ; INTERPRETER, AND THUS RETURNS ON ITS OWN TO BASIC ROM.
00160 ; #####
00170 ;
7EC0 00180 ORG 7EC0H ; CHANGE TO RELOCATE
00190 ;
00200 ; #####
00210 ; PATCH INTO OOS TIME$ ERROR LOCATION AND CHANGE IT
00220 ; #####
00230 ;
7EC0 F3 00240 ENTRY OI ; DISABLE ACTIVE INTRPTS.
7EC1 210E7E 00250 LO HL,START1 ; ENTRY OF TIMES PROGRAM
7EC4 227741 00260 LO [4177H],HL ; REPLACE ?L3 ERROR MSG.
7EC7 21A07F 00270 LO HL,START2 ; START OF "CMO" PROGRAM
7ECA 227441 00280 LO [4174H],HL ; REPLACE ?L3 ERROR MSG.
7EC0 3EC3 00290 LO A,0C3H ; GET "JUMP" COMMAND
7ECF 321240 00300 LO [4012H],A ; INSERT INT. PATCH POINT
7E02 214C7F 00310 LO HL,SERVE ; INTERRUPT SERV. ROUTINE
7E05 221340 00320 LO [4013H],HL ; INT. PATCH FROM 0038H
7E08 E0S6 00330 IM 1 ; SET INTERRUPT MODE #1
7E0A FB 00340 EI ; ENABLE INTERRUPT LINE
7E08 C3CC06 00350 JP 06CCH ; RETURN TO BASIC "READY"
00360 ;
00370 ; #####
00380 ; PATCH TO INTERCEPT ?L3 ERROR AND CHECK LINE'S SYNTAX
00390 ; #####
00400 ;
7E0E 07 00410 START1 RST 10H ; HOUSEKEEP SPACE, ETC.
7E0F E5 00420 PUSH HL ; SAVE BASIC LINE POINTER
7EE0 3E11 00430 LO A,11H ; LENGTH OF TIMES ITSELF
7EE2 C0S728 00440 CALL 28S7H ; ROM STRING SPACE SETUP
7EE5 2A0440 00450 LO HL,[4004H] ; LOCATION TO FINO TIMES
7EE8 114340 00460 LO OE,SECON0+2 ; POINT OE TO HOURS POS'N
7EE8 C0187F 00470 CALL DISPLY ; CONVERT, PLACE IN TIMES
7EEE 363A 00480 LO [HL],3AH ; PUT COLON INTO TIMES
7EF0 23 00490 INC HL ; BUMP TIMES POINTER
7EF1 18 00500 DEC OE ; SET OE TO MINS. POS'N
7EF2 C0187F 00510 CALL DISPLY ; CONVERT, PLACE IN TIMES
7EFS 363A 00520 LO [HL],3AH ; PUT COLON INTO TIMES
7EF7 23 00530 INC HL ; BUMP TIMES POINTER
7EF8 18 00540 DEC OE ; SET OE TO SECS. POS'N
7EF9 C0187F 00550 CALL DISPLY ; CONVERT, PLACE IN TIMES
7EFC 3620 00560 LO [HL],20H ; PUT SPACE INTO TIMES
7EFE 23 00570 INC HL ; BUMP TIMES POINTER
7EFF 114S40 00580 LO OE,SECON0+4 ; POINT OE TO MON. POS'N
7F02 C0187F 00590 CALL DISPLY ; CONVERT, PLACE IN TIMES
7F05 362F 00600 LO [HL],2FH ; PUT SLASH INTO TIMES
7F07 23 00610 INC HL ; BUMP TIMES POINTER
7F08 18 00620 DEC OE ; SET OE TO DAYS POS'N
7F09 C0187F 00630 CALL DISPLY ; CONVERT, PLACE IN TIMES
7F0C 362F 00640 LO [HL],2FH ; PUT SLASH INTO TIMES
7F0E 23 00650 INC HL ; BUMP TIMES POINTER
7F0F 114640 00660 LO OE,SECON0+5 ; POINT OE TO YEARS POS'N
7F12 C0187F 00670 CALL DISPLY ; CONVERT, PLACE IN TIMES
7F1S C38428 00680 JP 2884H ; FINISH DISPLAY IN ROM
00690 ;
00700 ; #####
00710 ; FIND VALUES IN TIME LOCATIONS AND CONVERT TO ASCII
00720 ; #####
00730 ;
7F18 1A 00740 DISPLY LO A,[OE] ; GET VALUE INTO ACCUM.
7F19 C0407F 00750 CALL N188LE ; SEPARATE INTO 4 BITS
7F1C 47 00760 LO B,A ; VALUE INTO B FOR TEST
7F1D AF 00770 XOR A ; CLEAR A FOR USE IN LOOP
7F1E 04 00780 INC B ; DUMMY INCREMENT ...
7F1F 05 00790 LOOP DEC B ; DECREMENT TO TEST FOR 0
7F20 280S 00800 JR Z,LEAVE ; UPPER N188LE NOW AT 0
7F22 C616 00810 ADD A,16H ; A=A+16 ...HEX-DEC CONV.
7F24 27 00820 OAA ; DEC.ADJ.: 16 BECOMES 10
7F2S 1F8 00830 JR LOOP ; LOOP TILL CONV. DONE
7F27 47 00840 LEAVE LO B,A ; SAVE VALUE BACK IN B
7F28 79 00850 LO A,C ; GET LOW N188LE BACK
7F29 FE0A 00860 CP 0AH ; IS IT GREATER THAN 10?
7F28 3B04 00870 JR C,CLEAN ; NO WORK IF LESS THAN 10

```

Listing 8-2. Interrupt-driven RTC (one second).

Figure 8-9 is a one-pulse-per-second clock that triggers the keyboard unit's interrupt line. Because this clock and the expansion box would compete for that line, this circuit cannot be used in a complete TRS-80 system. But it is a \$5 project, and an inexpensive clock add-on for 16K systems alone.

It consists of five integrated circuits. One regulates the voltage, the second (Z2) takes the sine-wave-shaped signals from the 6.3-volt power transformer and converts them to a sharp-edged digital signal. Z3 and Z4 divide the 60 Hz signal into a one-pulse-per-second signal. Finally, Z5 provides the 1 Hz interrupt signal.

When the interrupting signal is accepted by the TRS-80, an 'interrupt acknowledge' signal is sent back to Z5. There is an important purpose to this action; it turns the interrupt signal off. Why? When the computer receives the interrupt, it sets aside its present software activities to 'service' the interrupt. Unless the interrupting flip-flop is reset, the computer, upon completing the interrupt service routine, will think the previous interrupt is a *new* interrupt, and it will keep updating the time and date.

Software to run this clock is presented in Listing 8-2. It patches into the TIME\$ and CMD locations, and accepts time and date in the following format (use spaces and punctuation exactly as shown):

CMO"09:15:22 05/18/81"

The program checks for correct syntax of the set-time command line, but doesn't verify actual times or dates. So, until the clock is next updated, it will display whatever bizarre time and date you may have set it to!

To print the time and date, enter PRINT TIME\$. You may use TIME\$ just as you would any other strings, including with PRINT, LPRINT, MID\$, LEFT\$, RIGHT\$, concatenation, and other string manipulation. A complete description of these programs is in the supplement to this chapter.

```

7F20 060A      00880      SUB      0AH      ; REDUCE IT TO 0 THRU 5
7F2F C610      00890      A00      A,10H    ; NOW ADD CARRY BIT
7F31 80        00900      CLEAN A00      A,B      ; CREATE A DECIMAL RESULT
7F32 27        00910      0AA      ; DEC. ADJ. THE TOTAL
7F33 C0407F    00920      CALL     NIBBLE   ; SEPARATE INTO 4 BITS
7F36 C630      00930      A00      A,30H    ; CONVERT NIBBLE TO ASCII
7F3B 77        00940      LO       (HL),A   ; PLACE VALUE INTO TIMES
7F39 23        00950      INC      HL       ; BUMP TIMES PTR. BY ONE
7F3A 79        00960      LO       A,C      ; GET VALUE SAVED IN C
7F3B C830      00970      ADD      A,30H    ; CONVERT NIBBLE TO ASCII
7F3D 77        00980      LO       (HL),A   ; PLACE VALUE INTO TIMES
7F3E 23        00990      INC      HL       ; BUMP TIMES PTR. BY ONE
7F3F C9        01000      RET              ; BACK TO 00 PUNCTUATION
01010 ;
01020 ; #####
01030 ; SUBROUTINE TO CONVERT A BYTE AND SAVE IT AS TWO NIBBLES
01040 ; #####
01050 ;
7F40 FS        01060      NIBBLE  PUSH     AF      ; SAVE THE BYTE BRIEFLY
7F41 E60F      01070      ANO      0FH      ; MASK OUT THE HIGH BITS
7F43 4F        01080      LO       C,A      ; SAVE LOW NIBBLE IN C
7F44 F1        01090      POP      AF       ; GET THE WHOLE BYTE BACK
7F45 F1        01100      RRA              ; HOVE THE BYTE RIGHT...
7F46 1F        01110      RRA              ; ... TWO PLACES ...
7F47 1F        01120      RRA              ; ... THREE PLACES ...
7F48 1F        01130      RRA              ; UNTIL MSB BECOMES LSB
7F49 E60F      01140      ANO      0FH      ; MASK OUT THE HIGH BITS
7F4B C8        01150      RET              ; NIBBLES NOW IN A & C
01160 ;
01170 ; #####
01180 ; INTERRUPT SERVICE ROUTINE IS ENTERED AT 1-S CLOCK PULSE
01190 ; #####
01200 ;
4041          01210      SECONO  EQU      4041H    ; LOCATION TO STORE TIMES
7F4C F3        01220      SERVE   OI              ; DON'T BOTHER ME NOW!
7F4D F3        01230      PUSH     AF       ; SAVE ACCUM. & FLAGS
7F4E E5        01240      PUSH     HL       ; SAVE HL REGISTER PAIR
7F4F 05        01250      PUSH     DE       ; SAVE DE REGISTER PAIR
7F50 3A4540    01260      LO       A,(SECONO+4) ; GET CURRENT MONTH VALUE
7F53 5F        01270      LO       E,A      ; SAVE MONTH VALUE IN E
7F54 1600      01280      LO       0,0      ; LET 0=0. REASON FOLLOWS
7F56 214140    01290      LO       HL,SECONO ; START AT SECONOS POS'N.
7F5B 34        01300      INC      (HL)    ; SECONOS = SECONOS + 1
7F5A 7E        01310      LO       A,(HL)    ; GET READY TO COMPARE
7F5B FE3C      01320      CP       600      ; IS IT 60 SECONOS?
7F5D 3824      01330      JR       C,OUT    ; OONE IF NOT 60 SECONOS
7F5F C0B87F    01340      CALL     TICTOC   ; ADVANCE TIME SUBROUTINE
7F62 FE3C      01350      CP       600      ; IS IT 60 MINUTES?
7F64 3B10      01360      JR       C,OUT    ; OONE IF NOT 60 MINUTES
7F66 C0B97F    01370      CALL     TICTOC   ; ADVANCE TIME SUBROUTINE
7F69 FE18      01380      CP       240      ; IS IT 24 HOURS?
7F6B 3B16      01390      JR       C,OUT    ; OONE IF NOT 24 HOURS
7F6D C0B97F    01400      CALL     TICTOC   ; ADVANCE TIME SUBROUTINE
7F70 E5        01410      PUSH     HL       ; SAVE REGISTER BRIEFLY
7F71 21937F    01420      LO       HL,LOOKUP ; DAYS-IN-MONTH TABLE
7F74 19        01430      A00      HL,OE     ; REMEMBER OE? SEE ABOVE
7F75 BE        01440      CP       (HL)     ; IS IT LAST DAY OF MONTH
7F76 E1        01450      POP      HL       ; GET REGISTER BACK NOW
7F77 3B0A      01460      JR       C,OUT    ; OONE IF NOT LAST DAY
7F79 C0B87F    01470      CALL     TIKTOK   ; ADVANCE DATE SUBROUTINE
7F7C FE00      01480      CP       120      ; IS IT 12 MONTHS?
7F7E 3B03      01490      JR       C,OUT    ; OONE IF NOT 12 MONTHS
7F80 C0B87F    01500      CALL     TIKTOK   ; ADVANCE DATE SUBROUTINE
7F83 01        01510      OUT      POP      DE     ; RESTORE DE REGISTERS
7F84 E1        01520      POP      HL       ; RESTORE HL REGISTERS
7F85 F1        01530      POP      AF       ; RESTORE ACCUM. & FLAGS
7F86 FB        01540      EI              ; GET CLOCK TICKING AGAIN
7F87 E040      01550      RETI             ; BACK FROM THE INTERRUPT
01560 ;
01570 ; #####
01580 ; ADVANCE TIME/DATE & RETRIEVE NEW VALUE SUBROUTINES
01590 ; #####
01600 ;
7F8B AF        01610      TICTOC  XOR      A      ; CLEAR ACCUM. TO ZERO
7F8A 77        01620      FINISH  LO       (HL),A   ; HRS, MIN, OR SEC = 0
7F8B 23        01630      INC      HL       ; MOVE TO NEXT POSITION
7F8C 34        01640      INC      (HL)    ; TIME = TIME + 1 (CARRY)
7F8D 7E        01650      LO       A,(HL)   ; SET UP TO TEST VALUE
7F8E C9        01660      RET              ; BACK TO COMPLETE TEST
7F8F 3E01      01670      TIKTOK  LO       A,1      ; A = 1 FOR DAY OR MONTH
7F91 1B7F      01680      JR       FINISH   ; OTHER ROUTINE DOES WORK
01690 ;
01700 ; #####
01710 ; THIS IS THE DAYS-IN-A-MONTH LOOKUP TABLE - NO LEAP YEAR
01720 ; #####
01730 ;
7F93 00        01740      LOOKUP  0EF8     00      ; OOHMY BYTE, BUT THEN...
7F94 20        01750      0EF8     320      ; THIRTY DAYS HATH
7F95 10        01760      0EF8     290      ; SEPTEMBER,
7F96 20        01770      0EF8     320      ; APRIL, JUNE, AND
7F97 1F        01780      0EF8     310      ; NOVEMBER;
7F98 20        01790      0EF8     320      ; ALL THE REST HAVE
7F99 20        01800      0EF8     310      ; THIRTY-ONE,
7F9A 20        01810      0EF8     320      ; 'CEPT FEBRUARY, AND
7F9B 20        01820      0EF8     320      ; THERE'S BEEN ALL
7F9C 1F        01830      0EF8     310      ; TOO MUCH TALK
7F9D 20        01840      0EF8     320      ; ABOUT THE HYRIAQ
7F9E 1F        01850      0EF8     310      ; PRETIDIGITATIONS
7F9F 20        01860      0EF8     320      ; USING THAT MONTH
01870 ;
01880 ; #####
01890 ; "CMD" PATCH CHECKS PARAMETERS, SYNTAX, AND SETS TIME
01900 ; #####

```

The second and third clocks for the TRS-80 are similar in concept, but different in execution. That difference has only to do with manufacturing quality of the specified clock-calendar chip, the MSM5832 (available from *Digi-Key Corporation* and *Hobbyworld Electronics* — see Appendix). This clock-calendar has been designed to interface with microcomputers instead of the familiar red LED readouts. A 32.768 KHz crystal is required for its operation (also sold by the above suppliers).

It provides time in hours (12 or 24), minutes and seconds; month, day, year (leap year as well) and day of the week. It has timing signal outputs for interrupt use, which will not be used in this circuit. A battery backup will keep it in time when the TRS-80 is turned off. Two complete circuits are presented in Figure 8-10 and Figure 8-11.

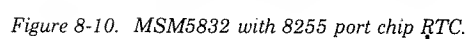
Why two circuits? Because the MSM5832 is a relatively slow electronic circuit, and, depending on the quality of the production run used for the chip you purchase, it may or may not be fast enough to interface directly with your TRS-80!

For a slower chip, you will need to use intermediate logic to latch onto the clock information in its own good time, and feed it to the TRS-80 as the computer's signals speed past. For this job, the INS8255 (as recommended by *OKI*, manufacturers of the MSM5832) is used. As noted above, the 8255 is a peripheral interface adapter, which sets up a private, latched bus between itself and the clock chip. Clock information is sent via Port A. The clock chip's address lines are selected through Port B, and other timekeeping features are selected through Port C.

The 8255 will be placed in the TRS-80 memory map at 37D0 through 37D2, below the cassette/ disk latches, and above the Level II operating system. Z1 and Z2 in Figure 8-10 decode that address group, and pins 8 and 9 of the 8255 (Z3) are used to select the specific address among those.

The last circuit, Figure 8-11, places the MSM5832 directly into the TRS-80 memory map without use of the 8255 interface chip. Because the MSM5832 has four address lines, it is decoded differently from the 8255, but occupies the same general area (37D0 to 37DF).

Wiring all these clocks is a simple procedure because there are few parts. All can be soldered or wire wrapped, though sockets are virtually essential for the 8255 and MSM5832 chips. The latter is a static-sensitive chip, and should be



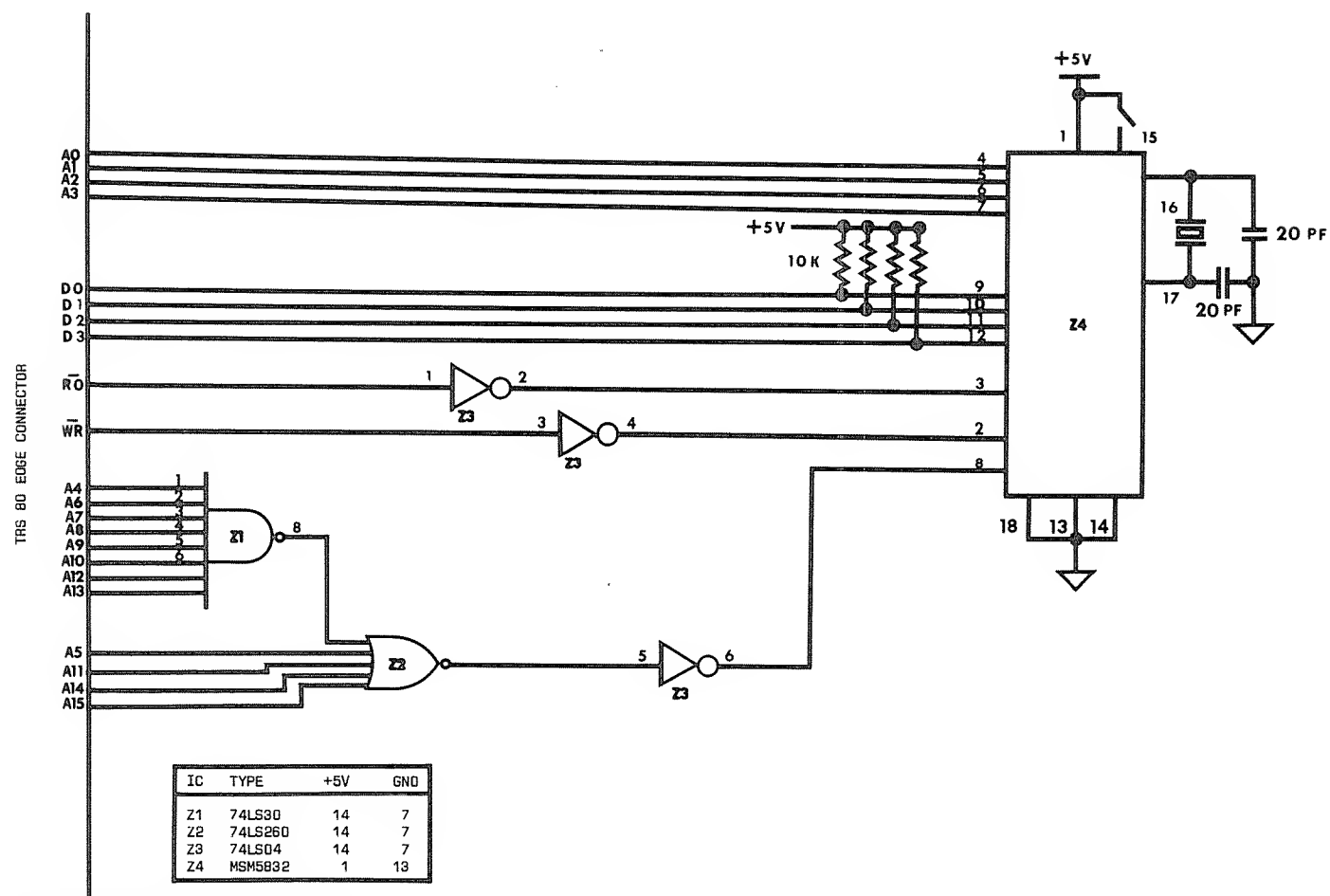
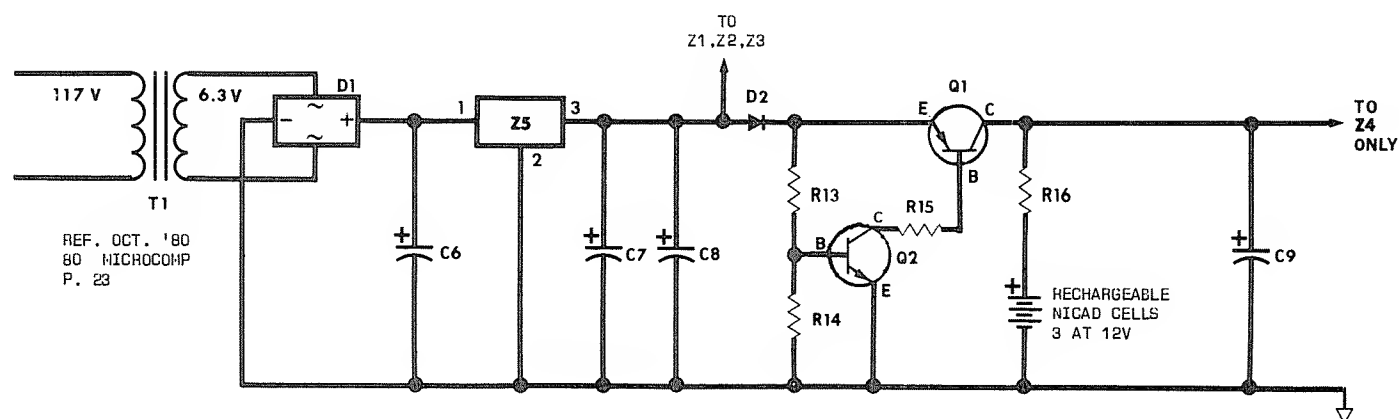


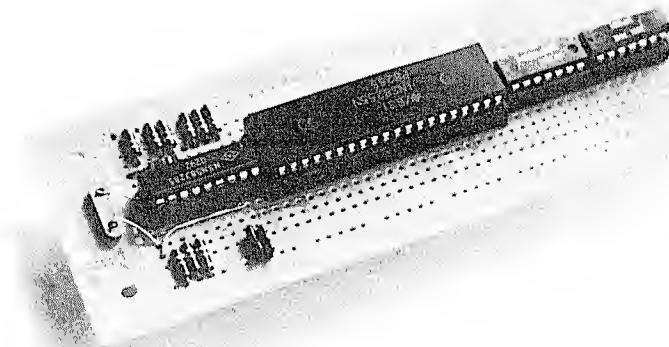
Figure 8-11. MSM5832 direct to bus RTC.



```

01910 ;
7FAD 11434D 0192D START2 LD DE,SECDND+2 ; POINT DE TO HOURS PDS'N
7FA3 7E 0193D LO A,(HL) ; CHAR AT LINE POINTER
7FA4 FE22 0194D CP 22H ; IS IT A QUOTE MARK?
7FA6 204A 0195D JR NZ,OTHERS ; CHECK FOR CMDT OR CMOR
7FAB C00B7F 0196D CALL CONVRT ; READ/CONV. ASCII HR.
7FAB FE3A 0197D CP 3AH ; IS IT A COLON?
7FA0 C29719 0198D SYNERR JP NZ,1997H ; GO TO ?SN ERROR ROUTINE
7FBD C00B7F 0199D CALL CONVRT ; READ/CONV. ASCII MIN.
7FB3 FE3A 0200D CP 3AH ; IS IT A COLON?
7FB5 20F6 0201D JR NZ,SYNERR ; SYNTAX ERRDR IF NOT :
7FB7 C00B7F 0202D CALL CONVRT ; READ/CONV. ASCII SEC.
7FBA FE20 0203D CP 20H ; IS IT A SPACE?
7FBC 20F6 0204D JR NZ,SYNERR ; SYNTAX ERROR IF NOT
7FBE 11454D 0205D LD DE,SECDND+4 ; POINT DE TO MONTH PDS'N
7FC1 C00B7F 0206D CALL CONVRT ; READ/CONV. ASCII MDN.
7FC4 FE2F 0207D CP 2FH ; IS IT A SLASH?
7FC6 20E5 0208D JR NZ,SYNERR ; SYNTAX ERROR IF NOT /
7FCB C00B7F 0209D CALL CONVRT ; READ/CONV. ASCII DAY
7FCB FE2F 0210D CP 2FH ; IS IT A SLASH?
7FC0 200E 0211D JR NZ,SYNERR ; SYNTAX ERRDR IF NOT /
7FCF 11464D 0212D LD DE,SECDND+5 ; POINT DE TO YEARS PDS'N
7FD2 C00B7F 0213D CALL CONVRT ; READ/CONV. ASCII YEAR
7FD5 FE22 0214D CP 22H ; IS IT A QUOTE MARK?
7FD7 20D1 0215D JR NZ,EXIT ; DONE IF A QUOTE MARK
7F09 23 0216D INC HL ; BUMP POINTER PAST QUOTE
7FDA C9 0217D EXIT RET ; BACK TO BASIC
0218D ;
0219D ; #####
0220D ; CONVERT ASCII TO HEX AND PDKE INTO CLOCK TIMES LOCATION
0221D ; #####
0222D ;
7F0B 23 0223D CONVRT INC HL ; BUMP LINE PTR. BY ONE
7F0C 7E 0224D LO A,(HL) ; GET CHARACTER IN LINE
7F00 0630 0225D SUB 30H ; CONVERT ASCII TO HEX
7F0F 3C 0226D INC A ; MAKE A BE AT LEAST 1
7FED 47 0227D LD B,A ; SAVE THAT VALUE IN B
7FE1 3EF6 0228D LD A,0F6H ; A= 1DD HEX MINUS 10 DEC
7FE3 C6DA 0229D MULT ADD A,DAH ; MULTIPLY BY ADDITION
7FE5 10FC 0230D DJNZ MULT ; I.E., A = B TIMES 10
7FE7 47 0231D LO B,A ; SAVE THAT VALUE IN B
7FEB 23 0232D INC HL ; BUMP LINE PTR. BY ONE
7FE9 7E 0233D LD A,(HL) ; GET CHARACTER IN LINE
7FEA D63D 0234D SUB 30H ; CONVERT ASCII TO HEX
7FEC 8D 0235D ADD A,B ; A = (B * 1D) + A
7FE0 12 0236D LO (0E).A ; TIME IS SET, PUT IN 0E
7FEE 1B 0237D DEC DE ; BUMP DE TO NEXT PLACE
7FEF 23 0238D INC HL ; BUMP LINE PTR. BY ONE
7FF0 7E 0239D LD A,(HL) ; GET CHARACTER IN LINE
7FF1 C9 0240D RET ; RETURN FOR FURTHER TEST
7FF2 FE52 0241D OTHERS CP 52H ; IS IT CMDR (CLOCK OFF)?
7FF4 2DD3 0242D CP NZ,NEXT ; NDPE, TRY FOR CMOT
7FF6 F3 0243D OI ; TURN OFF THE CLOCK
7FF7 23 0244D INC HL ; BUMP LINE PTR. BY ONE
7FFB C9 0245D RET ; BACK TO BASIC PRDGRAM
7FF9 FE54 0246D NEXT CP 54H ; IS IT CMOT (CLOCK ON)?
7FFB D080 0247D JR NZ,SYNERR ; NDPE, MUST BE ERRDR
7FFD FB 0248D EI ; TURN ON THE CLOCK
7FFE 23 0249D INC HL ; BUMP LINE PTR. BY ONE
7FFF C9 0250D RET ; BACK TO BASIC PRDGRAM
0251D ;
0252D ; #####
0253D ;
7E0D 0000D TOTAL ERRORS END ENTRY

```



MSM5832 clock board contains only 4 integrated circuits, some resistors, and a crystal. Power supply is external, and includes rechargeable batteries for backup.

handled carefully (see Chapter 4 for details). Use wires as short as possible in the area where the crystal is located, and triple-check all connections before applying the power. The MSM5832 is a delicate circuit, and at \$9 a shot, worth the trouble to check your work. The circuits are connected to the computer via standard edge-card connectors.

Edge Card Connectors: What's Up?

The 40 pin bus of the TRS-80 is a non-standard animal to begin with, and is not particularly logical in its pin assignments. What makes things more frustrating than merely locating the position of the pins, is trying to figure out which way is up.

The edge card of the TRS-80, when viewed from its edge, has pin 1 located at the top left – a logical place. Naturally, because the CPU and expansion unit face each other, the entry to the expansion interface is the mirror image of this design. Yet its top left pin is also called 'pin 1'! So pin 1 (TRS-80) leads to pin 39 (expansion box).

It gets stickier than this, however. When you are building projects, it's likely you won't be doing it with professionally etched edge cards, but rather with headers and wires and cables of various kinds which you assemble yourself. So when you get your 40 pin edge connector to hook onto the TRS-80, notice –

- the numbering of the pins on the connector, which, if industry standards are used, has pin 1 marked on the opposite side, like the expansion box.
- the outside wire, which may either lead to pin 1 or pin 2, depending on the manufacturer.
- the orientation of the connecting header, which may reverse the process one more time.

To be absolutely sure, use a meter to determine the path all the way from the computer edge card to the final connector when it is mounted on the project. It's best to mount the header or connector first, identify its pins with a meter, mark them and *then* begin work wire wrapping or soldering.

```

10 CLS : CLEAR 150 : REM * CRUDE BUT SERVICEABLE CLOCK PROGRAM
20 FOR X = 0 TO 6 : READ OW$(X) : NEXT : REM * ARRAY OF DAYS
30 DATA M O N D A Y, T U E S D A Y, W E D N E S D A Y
40 DATA T H U R S D A Y, F R I D A Y, S A T U R D A Y, S U N D A Y
50 PRINT "ENTER HOURS AND MINUTES, PLUS AM OR PM INDICATION."
60 INPUT "USE FORMAT 0,3,5,8,P (= 3:58 P.M.)":HO,H1,M0,M1,PS
70 INPUT "12-HOUR OR 24-HOUR CLOCK (ANSWER 12 OR 24)":CS
80 IF PS = "P" THEN HO = HO + 4 : REM * BIT 3 INDICATES P.M.
90 IF CS = "24" THEN HO = HO + 8 : REM * BIT 4 FOR 24 HOURS
100 PRINT "DAY OF THE WEEK (ENTER 1 TO 7. MONDAY IS 1.)"
110 INPUT OW : OW = OW - 1 : REM * CLOCK'S MONDAY IS ZERO
120 PRINT "MONTH, DAY AND YEAR IN FORMAT 0,3,3,1,8,0 (3/31/80)"
130 INPUT M2,M3,00,01,Y0,Y1 : REM * LEAP YEAR TEST IN NEXT LINE
140 LY = Y0 + 10 * Y1 : IF LY/4 = FIX (LY/4) THEN O0 = 00 + 4
150 POKE 14291,128 : REM * SET UP 8255 CHIP PORTS
160 POKE 14290,80 : REM * SET UP CLOCK TO READ TIME AND DATE
170 Q = 14289 : REM * THIS IS CLOCK ADDRESS REGISTER
180 POKE Q,2 : POKE Q-1,M1 : POKE Q,3 : POKE Q-1,M0
190 POKE Q,4 : POKE Q-1,H1 : POKE Q,5 : POKE Q-1,H0
200 POKE Q,6 : POKE Q-1,OW : POKE Q,7 : POKE Q-1,O1
210 POKE Q,8 : POKE Q-1,O0 : POKE Q,9 : POKE Q-1,M3
220 POKE Q,10 : POKE Q-1,M2 : POKE Q,11 : POKE Q-1,Y1
230 POKE Q,12 : POKE Q-1,Y0 : REM * TIME AND DATE INFO SET
240 POKE 14291,144 : CLS : REM * DISPLAY SUBROUTINE FOLLOWS
250 PRINT @ 0, "": REM * DISPLAY IS ON TOP LINE OF SCREEN
260 POKE 14290,32 : REM * SET UP CLOCK TO WRITE TIME AND DATE
270 POKE Q,6 : PRINT OW$ (PEEK (Q-1) AND 15), " ";
280 POKE Q,10 : PRINT PEEK (Q-1) AND 15; CHR$(8);
290 POKE Q,9 : PRINT PEEK (Q-1) AND 15; "/";
300 POKE Q,8 : PRINT PEEK (Q-1) AND 3; CHR$(8);
310 POKE Q,7 : PRINT PEEK (Q-1) AND 15; "/";
320 POKE Q,12 : PRINT PEEK (Q-1) AND 15; CHR$(8);
330 POKE Q,11 : PRINT PEEK (Q-1) AND 15; " -- ";
340 POKE Q,5 : PRINT PEEK (Q-1) AND 3; CHR$(8);
350 POKE Q,4 : PRINT PEEK (Q-1) AND 15; " ";
360 POKE Q,3 : PRINT PEEK (Q-1) AND 15; CHR$(8);
370 POKE Q,2 : PRINT PEEK (Q-1) AND 15; " ";
380 POKE Q,1 : PRINT PEEK (Q-1) AND 15; CHR$(8);
390 POKE Q,0 : PRINT PEEK (Q-1) AND 15;
400 POKE Q,5 : IF (PEEK (Q-1) AND 4) = 0 THEN PRINT " A. M.";
410 : IF (PEEK (Q-1) AND 4) = 4 THEN PRINT " P. M.";
420 IF (PEEK(14312) AND 128) = 0 THEN 500 ELSE 250
430 REM : THIS ROUTINE IS MUCH LONGER THAN IT NEED BE,
440 REM : BUT IS SET UP FOR CLARITY, NOT EFFICIENT USE
450 REM : OF MEMORY. IT IS EASIER TO USE THE MACHINE
460 REM : LANGUAGE SUBROUTINE FOR THIS CLOCK CIRCUIT.
500 A$ = "" : FOR X = 15360 TO 15424 : A$ = A$ + CHR$(PEEK(X))
510 NEXT : LPRINT A$ : GOTO 250

```

Listing 8-3. BASIC program for MSM5832/8255.

```

00100 : #####
00110 : MACHINE LANGUAGE CLOCK PROGRAM FOR MSM5832 CLOCK BOARD
00120 : ATTACHED TO THE TRS-80 AT ADDRESS 14288 TO 14300, USING
00130 : THE "TIMES" FUNCTION. THIS ROUTINE IS TRANSPARENT TO
00140 : BUT IS NOT CALLED BY THE CUSTOM INTERPRETER. HENCE,
00150 : IT OPERATES INDEPENDENTLY FROM THE INTERPRETER PATCH.
00160 : #####
00170 :
7E0D 00180 : DRG 7E0DH : CHANGE TO RELOCATE
00190 :
00200 : #####
00210 : PATCH INTO DDS TIME$ ERROR LOCATION AND CHANGE IT
00220 : #####
00230 :
7E0D 21DF7E 00240 ENTRY LD HL,START1 : START OF TIME$ PROGRAM
7E03 227741 00250 LD (4177H),HL : PATCH TIME$ ?L3 ERROR
7E06 218F7E 00260 LD HL,START2 : START OF "CMD" PROGRAM
7E09 227441 00270 LD (4174H),HL : PATCH CMD ?L3 ERROR
7E0C C3CC06 00280 JP 06CCH : BACK TO A BASIC "READY"
00290 :
00300 : #####
00310 : THIS IS THE BEGINNING OF THE "TIME$" PATCH TO READ TIME.
00320 : ROUTINE INTERCEPTS ?L3 ERROR AND CHECKS LINE'S SYNTAX.
00330 : #####
00340 :
7E0F 07 00350 START1 RST 10H : BASIC HOUSEKEEPING
7E10 E5 00360 PUSH HL : SAVE BASIC LINE POINTER
7E11 3E18 00370 LD A,18H : LENGTH OF TIME$
7E13 C05728 00380 CALL 2857H : ROM STRING SPACE SETUP
00390 :
00400 : #####
00410 : SET UP RAM SPACE AND GET CLOCK CHIP READY TO READ TIME
00420 : #####
00430 :
7E16 2A0440 00440 LD HL,(40D4H) : LOCATION TO STORE TIME$
7E19 F0210037 00450 LD IY,37DDH : CLOCK MEMORY ADDRESS
7E1D F0360390 00460 LD (IY+3),90H : SET UP 8255 CHIP PORTS
7E21 C06A7F 00470 CALL DELAY : WAIT FOR SLOW MSM5832
7E24 F0360220 00480 LD (IY+2),20H : SET UP CLOCK TO READ
7E28 C06A7F 00490 CALL DELAY : WAIT FOR SLOW MSM5832
00500 :
00510 : #####
00520 : CLOCK IS READY TO READ ... NOW READ AND CREATE STRING.
00530 : DAY OF THE WEEK IS ALPHABETIC AND WILL BE DONE FIRST.
00540 : #####
00550 :

```

Listing 8-4. Assembly program for MSM5832/8255.

To use the MSM5832 clock with the 8255 interface adaptor, you will need to refer to both chips' programming information. Figure 8-(?), earlier in this chapter, contains the 8255 programming parameters. Figure 8-(?) shows how the clock's registers are set up.

At first, this process may appear confusing. What is being done? Three semi-intelligent electronic devices are being taught to talk to one another. The TRS-80 knows it wants to read and write to memory. The 8255 is that memory. But the 8255 has a mind of its own, and that mind can only be controlled by selecting its control register, telling it what purpose each of its three ports is to serve, and then reading and writing those ports. Finally, the MSM5832 also has a mind of its own. It will neither report nor accept the time until it is told what aspect of the time is needed, and that too is done via a control register.

Figure 8-(?) is a flow chart which describes the process, and Listing 8-(?) is a BASIC program which fairly well describes the steps needed to access the MSM5832 chip, 'way down the chain.

Eliminating the 8255 by using the second circuit means only one set of electronic parts must be taught to speak to each other. The TRS-80 can probe right into the MSM5832 control register to select which aspect of the time and date it wishes.

Each of the machine language programs presented in Listings 8-(?) and 8-(?) use the TIME\$ and CMD commands to set and recall the time. To set the time and date, enter :

CMD"MON 03/14/49 02:29 PM"

Notice that this differs from the interrupt driven clock (Listing 8-(?)) in that the day of the week and morning - afternoon indicators must be given. Remember to use the punctuation and spacing exactly as printed here. As with the interrupt clock, PRINT TIME\$ returns the time and date, and this TIME\$ can be used and manipulated just as any other string.

A complete description of these programs is given in the Supplement to this Chapter.

```

7E28 F0360106 00S60      LO      [IY+1],6      ; POINT TO DAY OF WEEK
7E2F C06A7F 00S70      CALL    DELAY          ; WAIT FOR SLOW MSMS832
7E32 F07E00 00S80      LO      A,[IY+0]        ; GET DUMMY VALUE INTO A
7E35 C06A7F 00S90      CALL    DELAY          ; WAIT FOR SLOW CHIP (1)
7E38 F07E00 00600      LO      A,[IY+0]        ; GET DAY OF WEEK VALUE
7E38 E607 00610      .AND     07H             ; MASK OFF UNUSED BITS
7E30 117S7F 00620      LO      0E,TABLE       ; POINT OE TO DAY TABLE
7E40 3C 00630      INC      A                 ; IT MUST BE AT LEAST 1
7E41 30 00640      LOP1    DEC      A          ; IS ACCUMULATOR ZERO?
7E42 2B07 00650      JR      Z,XLOOP          ; GO OUT OF TABLE LOOP
7E44 0603 00660      LO      B,3              ; NUMBER OF CHARS PER DAY
7E46 13 00670      LOP2    INC      OE         ; MOVE PAST EACH CHAR
7E47 10F0 00680      DJNZ    LOP2P2          ; DO IT TILL AT NEXT DAY
7E49 18F6 00690      JR      LOP1            ; CHECK FOR NEXT DAY
00700 ;
00710 : #####
00720 : VALUE FOR DAY IS FOUND ... NOW TURN IT INTO LETTERS
00730 : #####
00740 :
7E48 0E03 00750      XLOOP   LO      B,3        ; NUMBER OF CHARS TO GET
7E40 1A 00760      YLOOP   LO      A,[OE]      ; CHARACTER TO TRANSFER
7E4E 77 00770      LO      (HL),A            ; XFER DAY NAME TO TIMES
7E4F 23 00780      INC      HL               ; NEXT LOCATION IN TIMES
7E50 13 00790      INC      OE               ; NEXT LOCATION IN TABLE
7E51 10FA 00800      DJNZ    YLOOP          ; LOOP BACK FOR NEXT CHAR
7E53 3620 00810      LO      (HL),20H        ; PUT SPACE AFTER DAY
7E55 23 00820      INC      HL               ; BUMP TIME BUFFER AGAIN
00830 ;
00840 : #####
00850 : DAY OF WEEK IS DONE ... NOW GET MONTH, DAY, AND YEAR
00860 : #####
00870 :
7E56 1E30 00880      LO      E,30H          ; HEX TO ASCII DIFFERENCE
7E58 1608 00890      LO      0,11           ; MONTH HI PORT + 1
7E5A 062F 00900      LO      B,2FH          ; SLASH ("/") CHARACTER
7E5C 0E0F 00910      LO      C,0FH          ; MASK UNUSED PORT BITS
7E5E C0S57F 00920      CALL    FILLER        ; GET MONTH HIGH VALUE
7E61 C0S57F 00930      CALL    FILLER        ; GET MONTH LOW VALUE
7E64 70 00940      LO      (HL),B           ; LOAD SLASH INTO TIMES
7E65 23 00950      INC      HL               ; BUMP TIME BUFFER BY ONE
7E66 0E03 00960      LO      C,3            ; MASK UNUSED CLOCK BITS
7E68 C0S57F 00970      CALL    FILLER        ; GET DAY HIGH VALUE
7E68 0E0F 00980      LO      C,0FH          ; MASK UNUSED CLOCK BITS
7E60 C0S57F 00990      CALL    FILLER        ; GET DAY LOW VALUE
7E70 70 01000      LO      (HL),B           ; PUT SLASH INTO TIMES
7E71 23 01010      INC      HL               ; BUMP TIME BUFFER BY ONE
7E72 16DD 01020      LO      D,13           ; YEAR HIGH VALUE + 1
7E74 C0S57F 01030      CALL    FILLER        ; GET YEAR HIGH VALUE
7E77 C0S57F 01040      CALL    FILLER        ; GET YEAR LOW VALUE
7E7A 3620 01050      LO      (HL),20H        ; VALUE FOR A SPACE
7E7C 23 01060      INC      HL               ; BUMP TIME BUFFER BY ONE
01070 ;
01080 : #####
01090 : MONTH, DAY, YEAR DONE - NOW GET HOURS, MINUTES, SECONDS
01100 : #####
01110 :
7E70 160S 01120      LO      0,S            ; HOURS HIGH VALUE
7E7F F07201 01130      LO      [IY+1],0      ; SET UP CLOCK CHIP PORT
7E82 C06A7F 01140      CALL    DELAY          ; DELAY FOR 8255 CHIP
7E85 F07E00 01150      LO      A,[IY+0]        ; DUMMY VALUE INTO ACC.
7E8B C06A7F 01160      CALL    DELAY          ; DELAY AGAIN FOR CHIP1
7E8B F07E00 01170      LO      A,[IY+0]        ; GET HOURS HIGH VALUE
7E8E F5 01180      PUSH    AF              ; SAVE THIS FOR AM/PM
7E8F 14 01190      INC      0              ; ACCOMMODATE SUBROUTINE
7E90 0E03 01200      LO      C,3            ; MASK UNUSED CLOCK BITS
7E92 C0S57F 01210      CALL    FILLER        ; GET HOURS HIGH VALUE
7E95 0E0F 01220      LO      C,0FH          ; MASK UNUSED CLOCK BITS
7E97 C0S57F 01230      CALL    FILLER        ; GET HOURS LOW VALUE
7E9A 363A 01240      LO      (HL),3AH        ; PUT A COLON IN TIMES
7E9C 23 01250      INC      HL               ; BUMP THE STRING ALONG
7E90 0602 01260      LO      B,2            ; NUMBER MINUTE/SEC LOOPS
7E9F C0S57F 01270      MINSEC CALL FILLER        ; GET, CONVERT, SAVE VALUE
7EA2 C0S57F 01280      CALL    FILLER        ; GET, CONVERT, SAVE VALUE
7EAS 363A 01290      LO      (HL),3AH        ; VALUE FOR A COLON
7EA7 23 01300      INC      HL               ; BUMP TIME BUFFER BY ONE
7EAB 10F5 01310      DJNZ    MINSEC          ; GO BACK FOR MIN/SEC
7EAA 2B 01320      DEC      HL               ; BACK UP TO LAST COLON
7EAB 3620 01330      LO      (HL),20H        ; CHANGE TO STRING END
01340 ;
01350 : #####
01360 : HOURS, MINUTES, SECONDS ARE DONE ... NOW FIGURE AM/PM
01370 : #####
01380 :
7EA0 23 01390      INC      HL               ; BUMP TIME BUFFER BY ONE
7EAE F1 01400      POP     AF              ; GET BACK HOUR HI VALUE
7EAF C857 01410      BIT     2,A            ; CHECK AM/PM INDICATOR
7EB1 2B04 01420      JR      Z,MORNING      ; MORNING IF BIT 2 = 0
7EB3 3650 01430      LO      (HL),50H        ; PUT LETTER "P" IN PLACE
7EB5 1B02 01440      JR      NEXT           ; JUMP PAST LETTER A
7EB7 3641 01450      MORNING LO      (HL),41H      ; PUT LETTER "A" IN PLACE
7EB9 23 01460      NEXT    INC      HL               ; BUMP TIME BUFFER BY ONE
7EBA 3640 01470      LO      (HL),40H        ; PUT LETTER "M" IN PLACE
7EBC C3842B 01480      JP      2B84H          ; BACK TO BASIC ACTIVITY
01490 ;
01500 : #####
01510 : THIS IS THE BEGINNING OF THE "CMO" PATCH TO SET TIME
01520 : CHECK FOR TIME SETTING PARAMETERS AND SYNTAX
01530 : #####
01540 :
7EBF 7E 01550      START2  LO      A,(HL)      ; CHAR AT LINE POINTER
7EC0 FE22 01560      CP      22H            ; IS IT A QUOTE MARK?
7EC2 C29719 01570      JP      NZ,1987H        ; ?SN ERROR IF NO QUOTE

```

Bank Selecting Machine Language in ROM

Warning: Before you begin construction of anything in this section, read the rest of the chapter! You might want to construct the complete ROM/RAM bank select system.

Seriously, one of the most exciting aspects of the TRS-80 is its blank area in the memory map. This has been partially used in this chapter to install a real time clock. It can also be used to select machine language programs or data burned into ROMs. At the time of this book's publication, the cost of a 2K erasable, programmable, read only memory (EPROM) is less than \$8. A year earlier, when these memories were \$27 or more, this project would not have been practical. Now it is.

In quantities of 100, these EPROMs are less than \$5, which means, by using this project, direct access to over 200K of memory is possible for less than \$500. But one or two of such memories are just as valuable.

But first, what is an EPROM? How is it used? An EPROM, as its name suggests, is a memory which the user can program and reprogram as necessary. It is programmed with an EPROM programmer, and erased with ultraviolet light. It maintains its contents with the power off, just like the Level II ROMs themselves. It is used in a way even simpler than the way the TRS-80's RAM memory is used, and in this is found its great advantage: it needs but power, address and data lines to make its data available to the CPU.

By using a decoded output port, one of a bank of these memories may be selected for use. Here is an example; I might want to load a special machine language monitor program. I know that program is located in ROM #26 in my ROM bank. I can command something like -

```

OUT 31,26  <ENTER>
SYSTEM    <ENTER>
/1228B    <ENTER>

```

- and my program will be loaded and active. Only the time for three entries has been spent; not even the time of the disk access. And, beyond that, no RAM memory need be used!

A complete circuit for a ROM select bank is presented in Figure 8-14. Each of the lines marked 'to ROM' will select one of 256 possible ROMs!

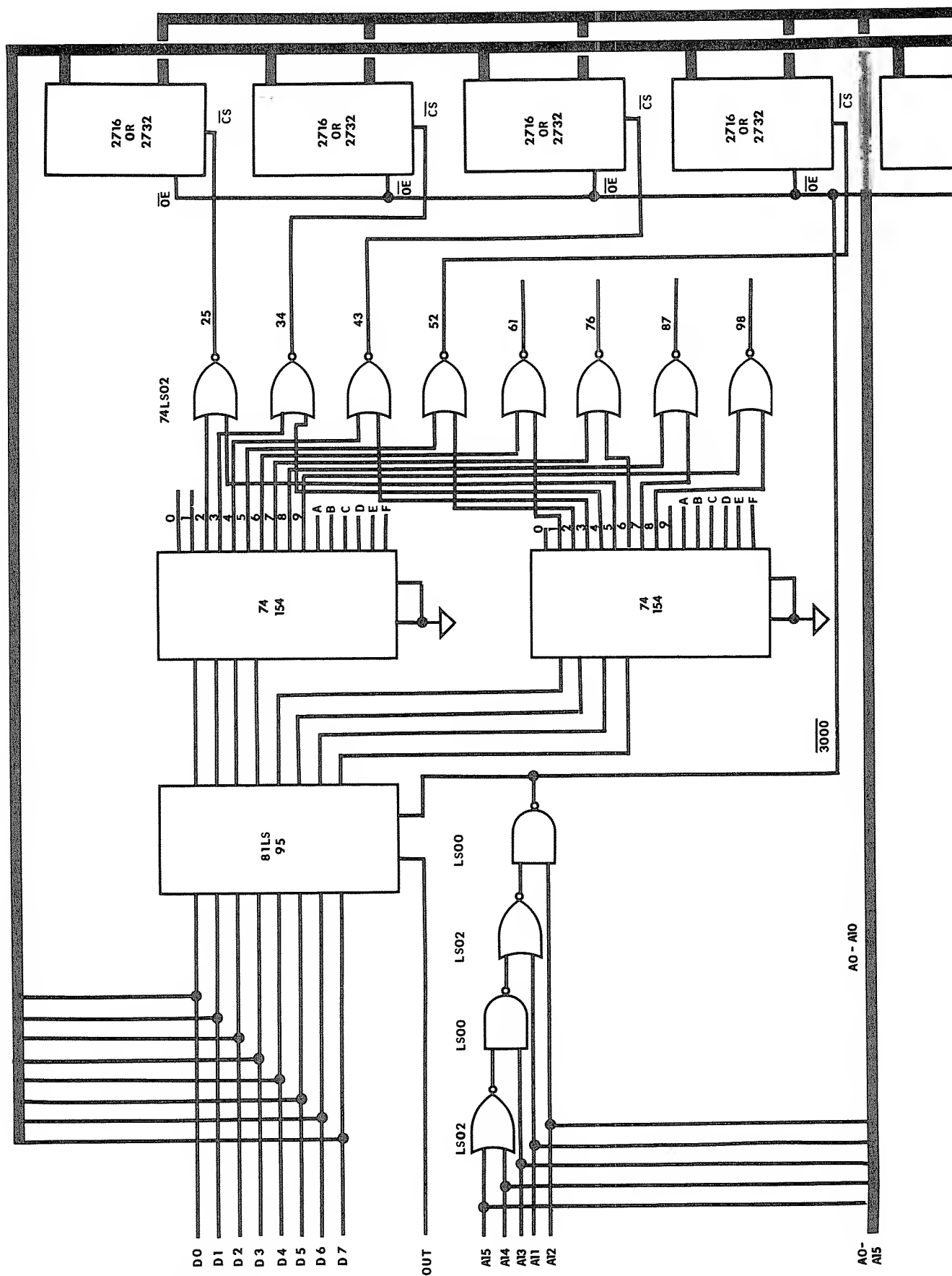


Figure 8-14. Bank-selected ROMs.

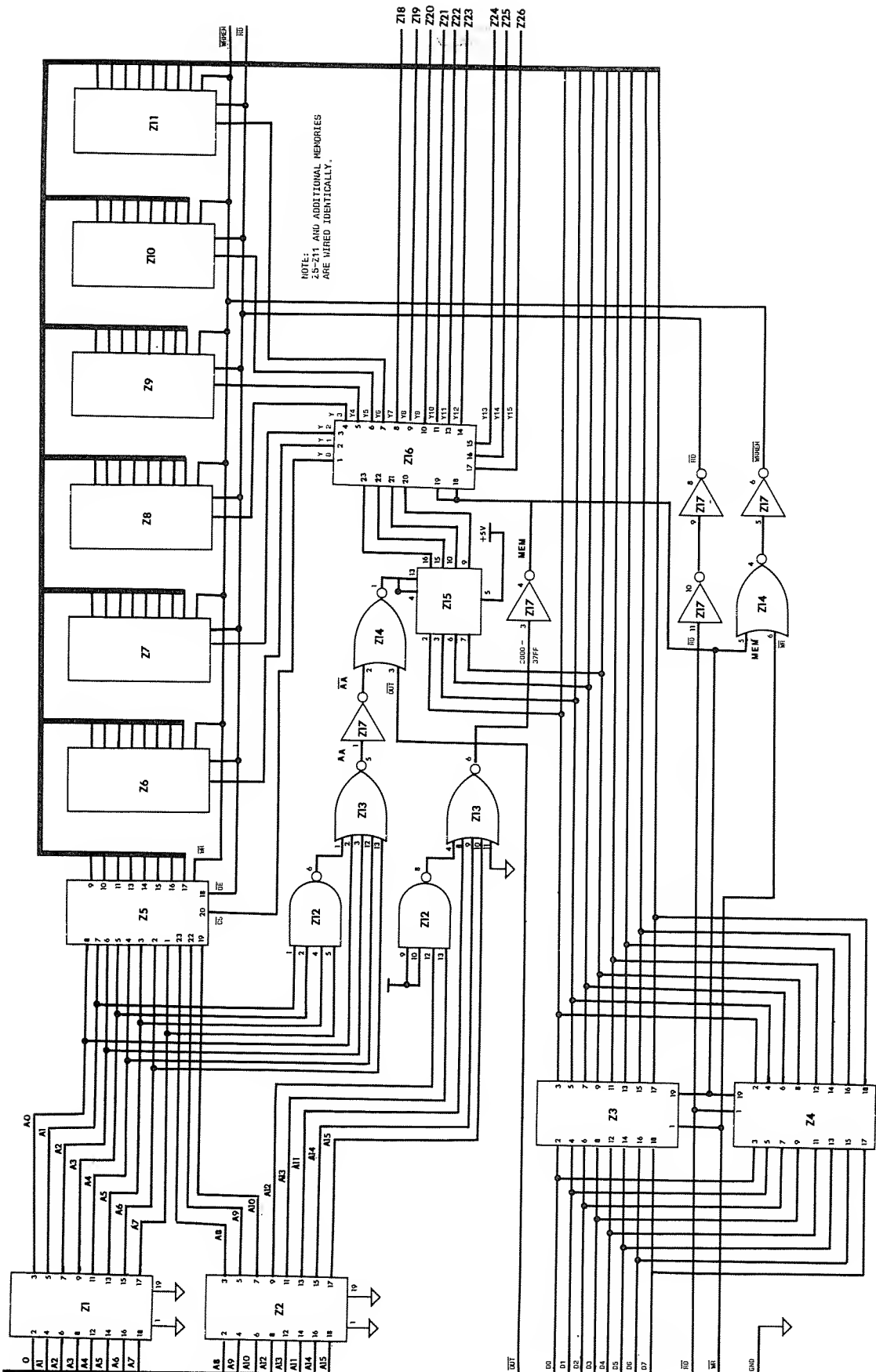


Figure 8-15. Bank-selected ROM/RAM.

Bank Selecting ROMs

```

7ECS 23      01580      INC    HL          ; BUMP LINE PTR. BY ONE
7EC6 ES      01590      PUSH   HL          ; SAVE THE LINE POINTER
7EC7 11757F  01600      LO      OE, TABLE ; GET TABLE OF OAY NAHES
7ECA 0E00    01610      LO      C, 0       ; THIS WILL BE COUNTER
7ECC 0603    01620      OYLOOP POP    B, 3     ; NUMBER OF CHARS IN DAY
7ECE E1      01630      POP     HL          ; GET LINE POINTER BACK
7ECF ES      01640      PUSH   HL          ; SAVE AGAIN FOR LOOP USE
7E00 1A      01650      FINOIT LD      A, {OE}   ; GET 1ST CHAR OF STRING
7E01 A7      01660      AND     A          ; EASY WAY TO SET A FLAG
7E02 2B09    01670      JR      Z, ERROR1    ; VALUE = 0 ... ?SN ERROR
7E04 BE      01680      CP       (HL)        ; CHECK IT AGAINST TABLE
7E05 2B0A    01690      JR      Z, GOTONE    ; GET READY FOR NEXT CHAR
7E07 13      01700      LOOP4  INC     OE     ; RUN PAST VALUES FOR OAY
7E08 10F0    01710      OJNZ   LQDP4       ; BY RUNNING B TO ZERO
7E0A 0C      01720      INC     C         ; NEXT OAY - BUHP COUNTER
7E0B 1BEF    01730      JR      OYLQDP     ; BACK TO NEXT OAY LQDP
7E0D E1      01740      ERROR1 PDP     HL     ; CLEAR STACK OF HL REG.
7E0E C39719  01750      ERROR2 JP      1997H   ; GO TO ?SN ERROR MESSAGE
7EE1 23      01760      GOTONE INC     HL     ; GET NEXT CHAR FROM LINE
7EE2 13      01770      INC     OE     ; BUMP TABLE VALUE ALONG
7EE3 10EB    01780      OJNZ   FINOIT     ; KEEP GOING TILL DONE
              01790      ;
              01800      ; #####
              01810      ; NUHERICAL VALUE FOR OAY IS IN C - PUT IT IN MMSB32
              01820      ; #####
              01830      ;
              01840      ; POP     AF          ; CLEAR STACK OF HL VALUE
7EE5 F1      01850      LD      [IY+3], 8DH ; SET UP B255 TO WRITE
7EE6 F036D3B0 01860      CALL   DELAY    ; THAT 81% SLOW MMSB32
7EEA C06A7F  01870      LD      [IY+2], 5DH ; CLOCK CHIP WRITE VALUE
7EED F036D2S0 01880      CALL   DELAY    ; HOW SLOW IS IT??
7EF1 C06A7F  01890      LD      A, C       ; GET OAY OF WEEK VALUE
7EF4 79      01900      LO      (IY+1), 6 ; READY TO WRITE OAY
7EFS F0360106 01910      CALL   DELAY    ; WAIT (YAWN) TO WRITE
7EF9 C06A7F  01920      LD      [IY+0], A ; WRITE OAY TO CLOCK
7EFC F077D0  01930      ;
              01940      ; #####
              01950      ; DAY IS WRITTEN - FINO HONTH, OAY, YEAR AND WRITE THEH
              01960      ; #####
              01970      ;
7EFF 1608    01980      LO      D, 11      ; VALUE FOR HONTH + 1
7F01 C0377F  01990      CALL   TIHSET   ; WRITE HONTH TO CLOCK
7F04 C0377F  02000      CALL   TIMSET   ; WRITE OAY TO CLOCK
7F07 160D    02010      LO      O, 13      ; VALUE FOR YEAR + 1
7F09 C0377F  02020      CALL   TIHSET   ; WRITE YEAR TO CLOCK
7F0C 1605    02030      LD      D, 5       ; SET TO HOURS HIGH VALUE
7F0E C01A7F  02040      CALL   AHDPRH   ; WRITE HOURS TO CLOCK
7F11 C0377F  02050      CALL   TIHSET   ; WRITE MINUTES TO CLOCK
7F14 0604    02060      LO      B, 4       ; NUMBER OF CHARS LEFT
7F16 23      02070      SNEAK  INC     HL     ; BUMP LINE POINTER
7F17 10F0    02080      OJNZ   SNEAK   ; LOOP PAST "PM" & QUOTES
7F19 C9      02090      RET       ; BACK TO BASIC PROGRAH
              02100      ;
              02110      ; #####
              02120      ; CHECK FOR AH DR PH INDICATION AND WRITE THAT VALUE
              02130      ; #####
              02140      ;
7F1A 23      02150      AMORPM INC     HL     ; BUMP LINE TO NEXT CHAR.
7F1B ES      02160      PUSH   HL     ; SAVE CURRENT LINE PTR.
7F1C 05      02170      PUSH   DE     ; SAVE OTHER VALUES IN OE
7F1D 1106D0  02180      LO      OE, 6     ; HOW MANY SPACES TO MOVE
7F20 19      02190      ADD     HL, OE    ; FINO AH OR PM IN LINE
7F21 7E      02200      LO      A, {HL}   ; GET CHARACTER FROM LINE
7F22 FE41    02210      CP       41H      ; SET FLAG IF CHAR. = "A"
7F24 3E04    02220      CP       A, 4      ; GET PM INDICATOR READY
7F26 2001    02230      JR      NZ, EVENNG ; ZERO FLAG NOT SET IF PM
7F2B AF      02240      XOR     A         ; CLEAR PH INDICATOR
7F29 01      02250      EVENNG POP    OE     ; RESTORE VALUES TO OE
7F2A E1      02260      POP     HL     ; GET ORIGINAL LINE PTR.
7F2B 0602    02270      LO      B, 2      ; SET UP B AS TIHSET LOOP
7F2D 4F      02280      LO      C, A      ; SAVE AM/PH INDICATOR
7F2E 7E      02290      LO      A, {HL}   ; GET VALUE FROM LINE
7F2F 0630    02300      SUB     30H      ; STRIP ASCII HASK
7F31 3BAA    02310      JR      C, ERROR1 ; ERROR IF LESS THAN 0
7F33 B1      02320      ADD     A, C      ; ADD AM/PM BIT TO VALUE
7F34 C3417F  02330      JP      MIDDLE  ; SUBROUTINE FINISHES JOB
              02340      ;
              02350      ; #####
              02360      ; TIME SETTING SUBRDUTINE CHECKS LINE FOR SYNTAX
              02370      ; #####
              02380      ;
7F37 1E3D    02390      TIHSET LO      E, 30H   ; CONVERTS ASCII TO HEX
7F39 D6D2    02400      LO      B, 2     ; LOOP TWICE FOR 2 DIGITS
7F3B 1S      02410      ZLOOP  OEC     D       ; BUMP CLOCK ADDRESS PORT
7F3C 23      02420      INC     HL     ; GET NEXT CHAR FROM LINE
7F3D 7E      02430      LO      A, {HL}   ; MOVE IT TO ACC. TO TEST
7F3E 93      02440      SUB     E         ; STRIP OFF ASCII VALUE
7F3F 369C    02450      JR      C, ERRDR1 ; ERRDR IF LESS THAN 0
7F41 FEDA    02460      HIDDLE CP      0AH     ; CHECK IF GREATER THAN 9
7F43 3D98    02470      JR      NC, ERROR1 ; ERROR IF GREATER THAN 9
7F45 F07201  02480      LO      (IY+1), 0 ; OPEN PORT TO CLOCK
7F48 C06A7F  02490      CALL   DELAY    ; THE USUAL CMOS WAIT
7F4B F077D0  02500      LD      [IY+0], A ; WRITE VALUE TO CLOCK
7F4E C06A7F  02510      CALL   DELAY    ; WAIT! WAIT! WAIT!!!!
7F51 10E8    02520      OJNZ   ZLDOP     ; DO IT FOR 2 DIGITS
7F53 23      02530      INC     HL     ; BUMP PAST / : DR SPACE
7F54 C9      02540      RET       ; BACK TO MAIN PROGRAM
              02550      ;
              02560      ; #####
              02570      ; GET VALUE, CONVERT TO ASCII, AND SAVE IN TIMES BUFFER
              02580      ; #####
              02590      ;

```

Actually, the bank-select ROM idea is only one part of a larger possibility — blocks of ROM and/or RAM placed interchangeably. Ideally the additional RAM would be available in higher memory, but the TRS configuration is so locked into its memory map that expansion in high memory would require major reconfiguring of the memory and refresh circuits. With that in mind, then, all this expansion will be dedicated to the free area located at 3000 to 37E0.

This memory addition is an application of a 'Read-Only-RAM' concept. Simply stated, the memory write (WR) line to read/write memory is disabled by the user in order to outline an area of protected, but not permanently programmed, memory. Machine language programs under development can be emulated with this system. Crucial software can be embedded in crash-proof RAM, and the occasional nuisance of a program gone wild will not affect data in this area.

The inclusion of bank selection in this area permits the use of interchangeable, and constantly on-line, blocks of ROM, RAM, and protected RAM. In Figures 8-16 and 8-17, the ROMs used, as before, are 2716 EPROMs, and the RAMs are of the static variety. Two designs of the circuit are shown: one uses 2114 static RAMs, which are 1K by 4 bits wide; the other uses 4118 static RAMs, each a full 1K by 8 bits wide. The former memories are considerably less expensive, but the latter are easier to wire in a wire-crazy bank select scheme like this one.

As before, output port 31 has been chosen for the memory, and each 2K block is selected by the value output through port 31. Furthermore, RAM data in the unselected blocks remains intact, ready to use whenever a different value is output through port 31.

The bank-selected ROM may be used for the most part as any other memory, with one very interesting exception: routines in *different* ROMs may not call or jump to each others' routines in the normal way. Instead, two special routines must be used.

The jump routine is the simplest: the ROM's position in the bank must be identified, the jump prepared and the OUT statement commanded. The easiest way to do this is shown in Listing 8-7; the AF and HL registers are saved on the stack, then AF is loaded with the location of the routine (which it must obtain from a table of some sort identical in each ROM). Following that, the HL register pair is loaded with the jump address, and the OUT (1F), A command is then

```

7F55 15      02600 FILLER DEC      0      ; BUMP CLOCK PORT ADDRESS5
7F56 F07201 02610      LO      (IY+1),0  ; POINT TO VALUE WANTED
7F59 C06A7F 02620      CALL     DELAY      ; THAT OL' SLOW M5M5832
7F5C F07E00 02630      LD      A,(IY+0)   ; GET DUMMY VALUE INTO A
7F5F C06A7F 02640      CALL     DELAY      ; WAIT AGAIN! SLOW CHIP
7F62 F07E00 02650      LO      A,(IY+0)   ; NOW GET VALIO VALUE
7F65 A1      02660      AND      C          ; MASK UNUSED BITS
7F66 B3      02670      AOR      A,E       ; MAKE IT AN ASCII VALUE
7F67 77      02680      LD      (HL),A     ; PUT VALUE INTO BUFFER
7F6B 23      02690      INC      HL        ; NEXT BUFFER POSITION
7F6B C9      02700      RET              ; BACK TO MAIN PROGRAM
02710 ;
02720 ; #####
02730 ; THIS IS A SETUP WHICH CALLS A DELAY SUBROUTINE IN ROM
02740 ; #####
02750 ;
7F6A C5      02760 DELAY  PUSH      BC          ; SAVE BC REGISTER PAIR
7F6B F5      02770      PUSH     AF          ; SAVE AF REGISTER PAIR
7F6C 01010D 02780      LD      BC,1        ; DELAY FOR MSM5832 CHIP
7F6F C06000 02790      CALL     0060H       ; HERE IS ROUTINE IN ROM
7F72 F1      02800      POP       AF          ; GET AF REGISTER5 BACK
7F73 C1      02810      POP       BC          ; GET BC REGISTER5 BACK
7F74 C9      02820      RET              ; BACK TO MAIN PROGRAM
02830 ;
02840 ; #####
02850 ; THIS IS THE LOOKUP TABLE OF DAYS OF THE WEEK
02860 ; #####
02870 ;
7F75 40      02880 TABLE DEFM      'MON'
02890 ;
7F78 54      02900      DEFM      'TUE'
02910 ;
7F7B 57      02920      DEFM      'WED'
02930 ;
7F7E 54      02940      DEFM      'THU'
02950 ;
7F81 46      02960      DEFM      'FRI'
02970 ;
7F84 53      02980      DEFM      'SAT'
02990 ;
7F87 53      03000      DEFM      'SUN'
03010 ;
7F8A 00      03020      DEF8      0
03030 ;
03040 ; #####
03050 ;
7EDD        03D5D      ENO      ENTRY
00000 TOTAL ERRORS

```

executed, which switches ROMs. The proper jump address is still in HL, so a simple JP (HL) effects the jump.

```

; ROUTINE BEING LEFT MUST PROVIDE THIS
; INFORMATION TO AF AND HL REGISTERS
..ZZZZ      LO      (ZZZZ),A      ; SAVE AF VALUES
..YYYY      LO      (YYYY),HL     ; SAVE HL VALUES
3ENN        LO      A,NN         ; NEW ROM BANK NUMBER
21NNNN      LO      HL,NNNN       ; JUMP ADDRESS
C3WWW        JP      XFER1        ; MAKE ROM TRANSFER
; ALL ROMS CONTAIN THE FOLLOWING IDENTICAL
; BYTES AT THE SAME ADDRESSES IN ROM
031F XFER1   OUT      (1F),A      ; SWITCHES ROMS
E9          JP      (HL)         ; JUMPS TO ROUTINE
..ZZZZ      LD      (ZZZZ),A      ; NOT USED IN JUMP
..XXXX      LO      A,(XXXX)      ; NOT USED IN JUMP
031F        OUT      (1F),A      ; NOT USED IN JUMP
C9          RET                 ; NOT USED IN JUMP
; ALL ROUTINES BEING ENTERED MUST PROVIDE
; THE FOLLOWING RESTORATION CODING
..ZZZZ      LD      HL,(ZZZZ)     ; RESTORE HL VALUES
..YYYY      LO      A,(YYYY)     ; RESTORE AF VALUES

```

Listing 8-6. Accessing multiple ROMs (jumps).

Calling a subroutine in another ROM is more complicated, but still can be done. Listing 8-(?) shows how it might be achieved.

```

; ROUTINE BEING LEFT MUST PROVIDE THIS
; INFORMATION IN AF AND HL REGISTERS
..ZZZZ      LO      (ZZZZ),A      ; SAVE AF DATA
..YYYY      LO      (YYYY),HL     ; SAVE HL DATA
0B1F        IN      A,(1F)        ; GET ROM BANK NUMBER
..XXXX      LO      (XXXX),A      ; SAVE ROM BANK NO.
21NNNN      LO      HL,NNNN       ; GET CALL ADDRESS
C0WWW        CALL     XFER2        ; MAKE ROM TRANSFER
; THE FOLLOWING CALLING ROUTINE MUST BE
; PLACED IN ALL ROMS AT IDENTICAL LOCATION
031F XFER2   OUT      (1F),A      ; SWITCH ROMS
E9          JP      (HL)         ; ENTER ROUTINE
..ZZZZ      LO      (ZZZZ),A      ; SAVE AF DATA
..XXXX      LO      A,(XXXX)      ; GET OLD ROM NO.
031F        OUT      (1F),A      ; TRANSFER BACK
C9          RET                 ; BACK TO CALLER
; THE CALLED ROUTINE MUST MAKE THE
; FOLLOWING IDENTIFICATIONS
..YYYY      LO      HL,(YYYY)     ; GET SAVED DATA
..ZZZZ      LO      A,(ZZZZ)     ; GET SAVED DATA
; EXECUTE SUBROUTINE FOUND HERE
; PERFORM FOLLOWING STEPS WHEN THE
; SUBROUTINE IS COMPLETE AND MUST RETURN
C3 VVVV      JP      BACK        ; GO TO XFER ROUTINE

```

Listing 8-7. Accessing multiple ROMs (calls).

This bank-select method is only one of many options which may be selected to move from one ROM to another; as you can see, three bytes of RAM (marked ZZZZ, YYYY and WWWW in the listings) are needed to store information between transfers. Although this process may initially seem unwieldy, a ROM-resident operating system in each one, together with cautious programming, will provide a remarkably transparent system expansion.

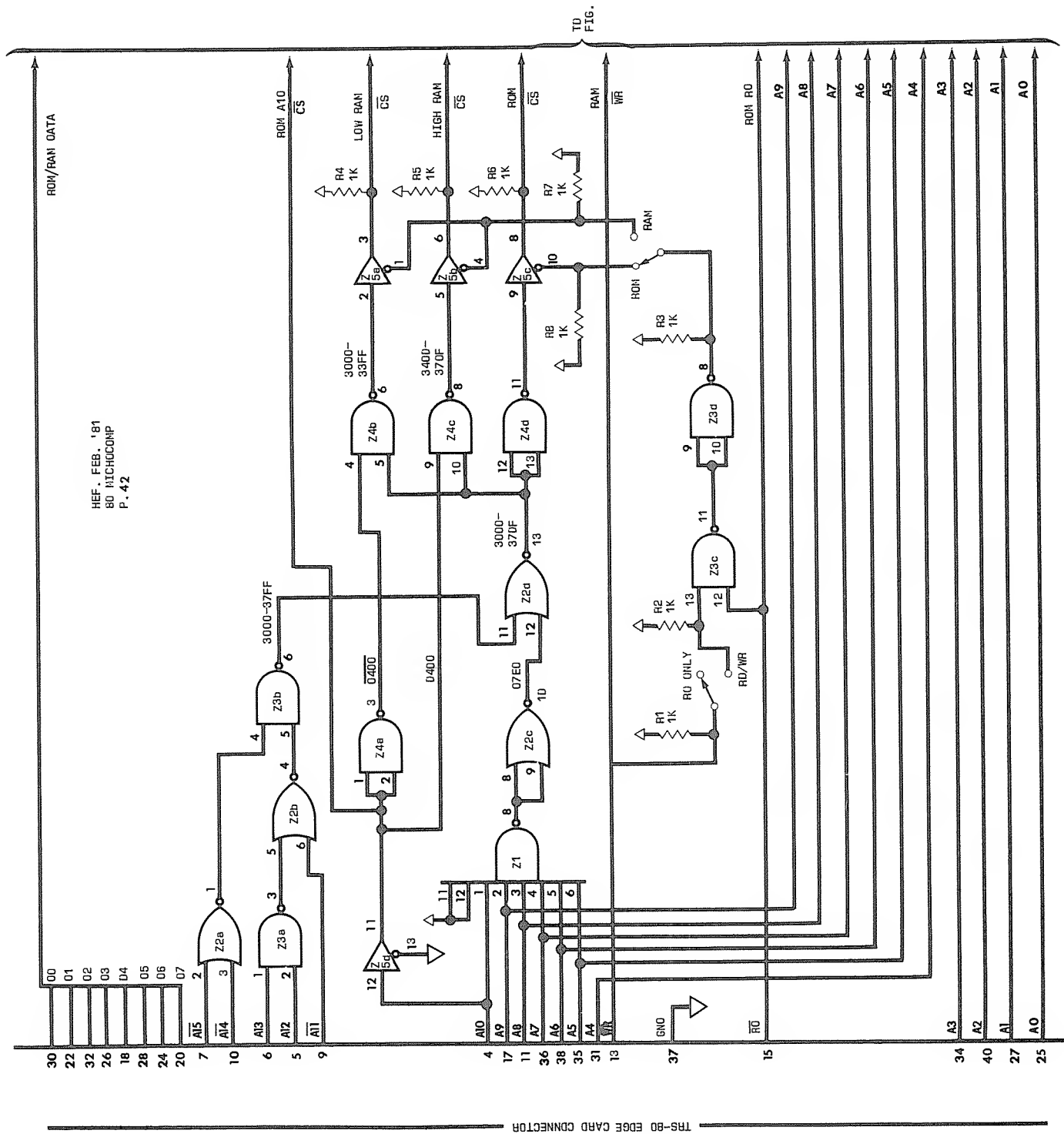


Figure 8-16. Bank-selected 4118 ROM/RAM.

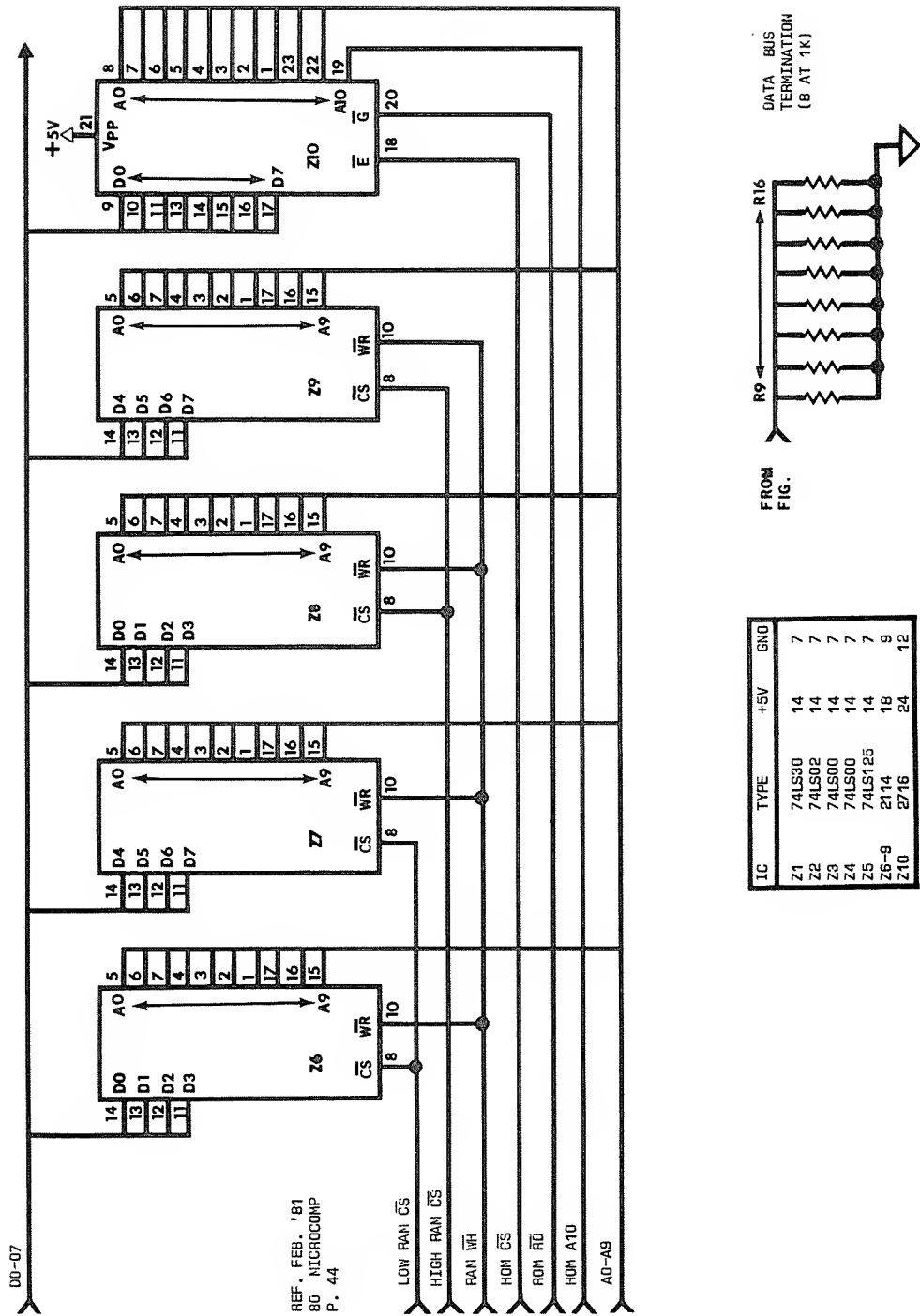
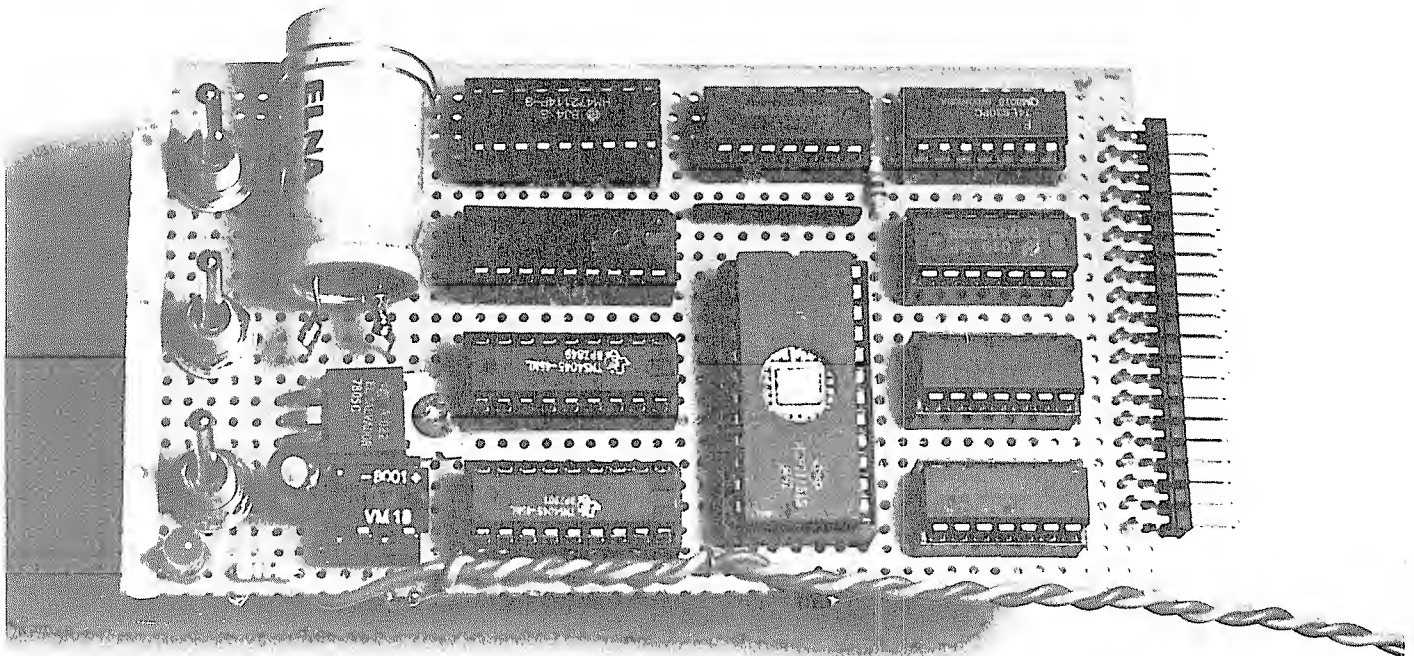
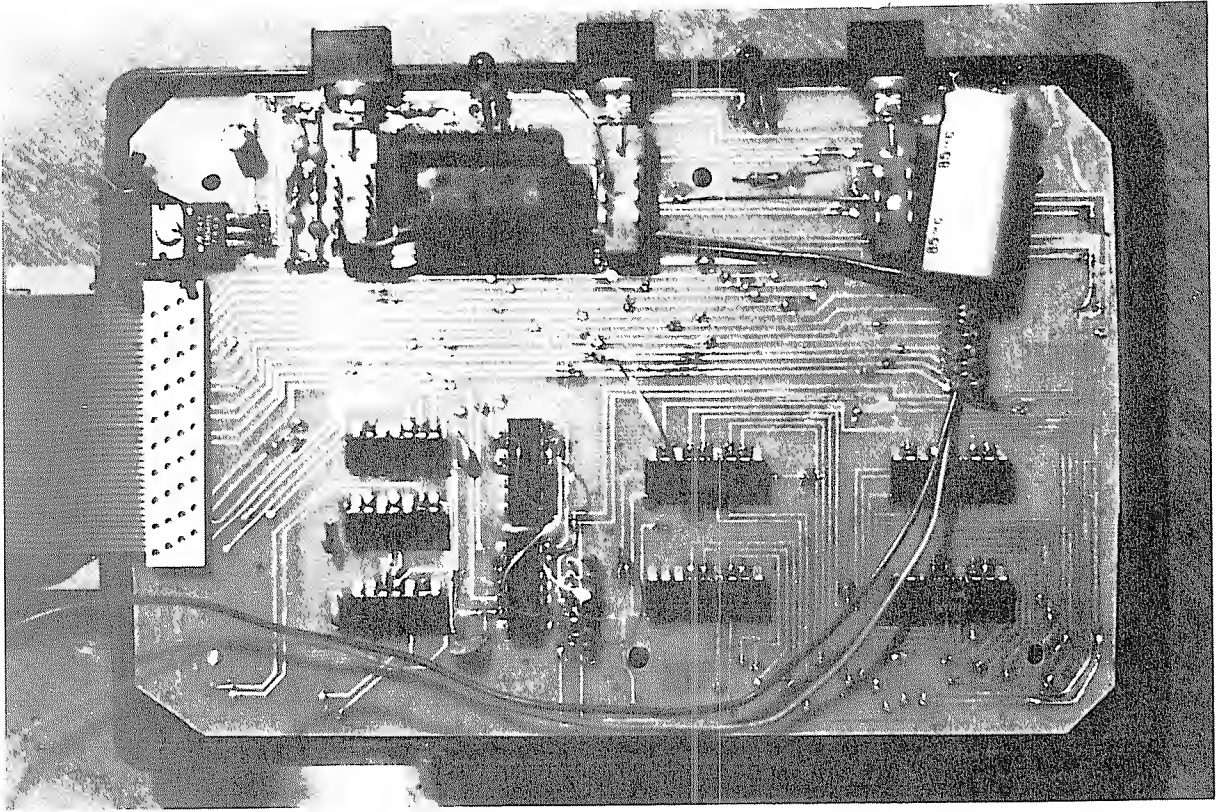


Figure 8-17. ROM/RAM addition without bank-select.



Completed ROM/RAM addition with power supply on board.


```

DD10D ; MEMORY SIDECAR TEST TAPE (C) 1981 D. B. KITSZ
DD11D ;
3CDD DD12D DRG 3CDDM
3CDD 2A DD13D DEFM '** LOADING 2K RAM TEST TAPE '
3C1C 2D DD14D DEFM '— D. B. KITSZ, SEPTEMBER 1981 **'
30DD DD15D DRG 30DDM
30DD 2A DD16D M1 DEFM '*****'
302D 2A DD17D DEFM '*****'
304D 2D DD18D M2 DEFM '----- TESTING "READ-ONLY"'
306D 2D DD19D DEFM '—RAM" MEMORY AREA -----'
308D 54 DD20D M3 DEFM 'THIS SCREEN HAS BEEN LOADED DIRE'
30AD 43 DD21D DEFM 'CTLY FROM THE MEMORY SIDECAR. IF'
30CD 41 DD22D M4 DEFM 'ALL CHARACTERS ARE CORRECT IN TH'
30ED 49 DD23D DEFM 'IS MESSAGE, AND IT HAS LOADED ON'
310D 54 DD24D M5 DEFM 'THE SCREEN NORMALLY, THEN THIS F'
312D 49 DD25D DEFM 'IRST TEST OF YOUR MEMORY SIDECAR'
314D 49 DD26D M6 DEFM 'IS COMPLETED. BELOW IS SHOWN A '
316D 43 DD27D DEFM 'COMPLETE CHARACTER SET AVAILABLE'
318D 54 DD28D M7 DEFM 'TO YOUR TRS-80:'
31AD 2D DD29D DEFM '
31CD 0DD1 DD30D M8 DEFW 010DH
31C2 02D3 DD31D DEFW 0302H
31C4 0405 DD32D DEFW 0504H
31C6 0607 DD33D DEFW 0706H
31C8 0809 DD34D DEFW 0908H
31CA 0AD8 DD35D DEFW 0B0AH
31CC 0CDD DD36D DEFW 0DDCM
31CE 0EDF DD37D DEFW 0FDEH
31DD 1D11 DD38D DEFW 1110M
31D2 1213 DD39D DEFW 1312H
31D4 1415 DD40D DEFW 1514H
31D6 1617 DD41D DEFW 1716H
31D8 1819 DD42D DEFW 1918H
31DA 1A1B DD43D DEFW 1B1AH
31DC 1C1D DD44D DEFW 1D1CH
31DE 1E1F DD45D DEFW 1F1EH
31ED 2D21 DD46D DEFW 212DM
31E2 2223 DD47D DEFW 2322M
31E4 2425 DD48D DEFW 2524H
31E6 2627 DD49D DEFW 2726H
31E8 2829 DD50D DEFW 2928H
31EA 2A2B DD51D DEFW 2B2AM
31EC 2C2D DD52D DEFW 2D2CH
31EE 2E2F DD53D DEFW 2F2EH
31FD 3D31 DD54D DEFW 313DM
31F2 3233 DD55D DEFW 3332M
31F4 3435 DD56D DEFW 3534M
31F6 3637 DD57D DEFW 3736M
31F8 3839 DD58D DEFW 3938H
31FA 3A3B DD59D DEFW 3B3AH
31FC 3C3D DD60D DEFW 3D3CH
31FE 3E3F DD61D DEFW 3F3EH
32DD 4041 DD62D M9 DEFW 4140M
32D2 4243 DD63D DEFW 4342H
32D4 4445 DD64D DEFW 4544M
32D6 4647 DD65D DEFW 4746M
32D8 4849 DD66D DEFW 4948M
32DA 4A4B DD67D DEFW 4B4AH
32DC 4C4D DD68D DEFW 4D4CM
32DE 4E4F DD69D DEFW 4F4EH
321D 5051 DD70D DEFW 515DM
3212 5253 DD71D DEFW 5352H
3214 5455 DD72D DEFW 5554H
3216 5657 DD73D DEFW 5756H
3218 5859 DD74D DEFW 5958H
321A 5A5B DD75D DEFW 5B5AH
321C 5C5D DD76D DEFW 5D5CH
321E 5E5F DD77D DEFW 5F5EH
322D 6061 DD78D DEFW 616DH
3222 6263 DD79D DEFW 6362H
3224 6465 DD80D DEFW 6564H
3226 6667 DD81D DEFW 6766H
3228 6869 DD82D DEFW 6968M
322A 6A6B DD83D DEFW 6B6AH
322C 6C6D DD84D DEFW 6D6CH
322E 6E6F DD85D DEFW 6F6EH
323D 7071 DD86D DEFW 717DH
3232 7273 DD87D DEFW 7372M
3234 7475 DD88D DEFW 7574H
3236 7677 DD89D DEFW 7776M
3238 7879 DD90D DEFW 7978H
323A 7A7B DD91D DEFW 7B7AH
323C 7C7D DD92D DEFW 7D7CM
323E 7E7F DD93D DEFW 7F7EH
324D 8081 DD94D M10 DEFW 818DM
3242 8283 DD95D DEFW 8382M
3244 8485 DD96D DEFW 8584M
3246 8687 DD97D DEFW 8786H
3248 8889 DD98D DEFW 8988H
324A 8A8B DD99D DEFW 8B8AH
324C 8C8D D10DD DEFW 8D8CH
324E 8E8F D1D1D DEFW 8F8EH

```

Remember also that the ROMs do not need to call each other directly. Instead, you may want to establish a lookup and transfer table in RAM, that acts like this :

1. User is in BASIC.
2. OUT 31,0 is entered to select the master ROM.
3. SYSTEM is entered, followed by /12288.
4. The master ROM (ROM #0) may be programmed to reset MEMORY SIZE and relocate a bank of patch points into RAM.
5. The master ROM, having completed its work the way Level II does at power up, returns to a READY in BASIC.
6. The routines are now all ready to use.

In summary, the ROM, RAM and bank-select systems can extend the horizons of your TRS-80 in unique ways. A project for people afflicted with cerebral palsy is developing a system whereby patients, doctors, or nurses need not be concerned if the TRS-80 crashes at some point, or the power is removed. Instead of loading tapes or fumbling with disks that may get damaged through handling or erratic power, these people need only use their bank-selected ROMs to reinstate and activate the machine language programs which drive specially made hardware. This hardware permits them to communicate with and use the computer with great ease.

Without the bank-select feature, however, the 8K control programs would have to be reloaded regularly, at a cost of time and patience, plus the risk of damage, loss or failure.

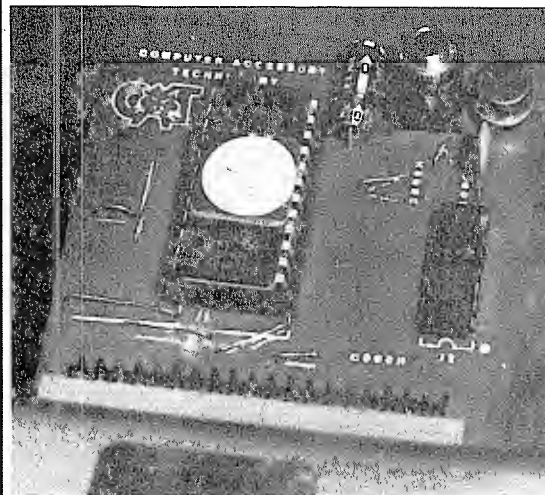
Continued Listing

3250	8091	01020	DEFW	9180H
3252	9293	01030	DEFW	9382H
3254	8485	01040	DEFW	8584H
3256	8687	01050	DEFW	9796H
3258	9889	01060	DEFW	8888H
325A	8A88	01070	DEFW	988AH
325C	8C80	01080	DEFW	808CH
325E	8E8F	01080	DEFW	8F8EH
3260	A0A1	01100	DEFW	0A1A0H
3262	A2A3	01110	DEFW	0A3A2H
3264	A4A5	01120	DEFW	0A5A4H
3266	A6A7	01130	DEFW	0A7ABH
3268	A8A8	01140	DEFW	0A8A8H
326A	AAAB	01150	DEFW	0ABAAH
326C	ACA0	01160	DEFW	0ADACH
326E	AEAF	01170	DEFW	0AFABH
3270	80B1	01180	DEFW	0B180H
3272	82B3	01190	DEFW	0B382H
3274	84B5	01200	DEFW	0B584H
3276	86B7	01210	DEFW	0B786H
3278	88B9	01220	DEFW	0B988H
327A	8AB8	01230	DEFW	0B8BAH
327C	8CB0	01240	DEFW	0B0BCH
327E	8EBF	01250	DEFW	0BF8EH
3280	C0C1	01260 M11	DEFW	0C1C0H
3282	C2C3	01270	DEFW	0C3C2H
3284	C4C5	01280	DEFW	0C5C4H
3286	C6C7	01290	DEFW	0C7C6H
3288	C8C9	01300	DEFW	0C9C8H
328A	CAC8	01310	DEFW	0CBCAH
328C	CEC0	01320	DEFW	0CCCEH
328E	CECF	01330	DEFW	0CFCEH
3290	00D1	01340	DEFW	00100H
3292	02D3	01350	DEFW	00302H
3294	04D5	01360	DEFW	00504H
3296	06D7	01370	DEFW	00706H
3298	08D9	01380	DEFW	00908H
329A	0AD8	01390	DEFW	00BDAH
329C	0C00	01400	DEFW	0000CH
329E	0EDF	01410	DEFW	00F0EH
32A0	0E01	01420	DEFW	0E1E0H
32A2	0E23	01430	DEFW	0E3E2H
32A4	0E45	01440	DEFW	0E5E4H
32A6	0E67	01450	DEFW	0E7E6H
32A8	0E89	01460	DEFW	0E9E8H
32AA	0EAB	01470	DEFW	0EBEAH
32AC	0ECD	01480	DEFW	0EDCEH
32AE	0EEF	01490	DEFW	0EFEBH
32B0	0F01	01500	DEFW	0F1F0H
32B2	0F23	01510	DEFW	0F3F2H
32B4	0F45	01520	DEFW	0F5F4H
32B6	0F67	01530	DEFW	0F7F6H
32B8	0F89	01540	DEFW	0F9F8H
32BA	0FAB	01550	DEFW	0FBFAH
32BC	0FCD	01560	DEFW	0FDFCH
32BE	0FEF	01570	DEFW	0FFFFH
32C0	20	01580 H12	DEFM	
32E0	20	01590	DEFM	
3300	57	01600 M13	DEFM	'WHEN YOU HAVE VERIFIED THAT ALL
3320	43	01610	DEFM	'CHARACTERS HAVE BEEN TRANSFERRED'
3340	43	01620 H14	DEFM	'CORRECTLY TO THE SCREEN, PRESS T
3360	48	01630	DEFM	'HE <ENTER> KEY TO CONTINUE.....'
3380	20	01640 M15	DEFM	
33A0	20	01650	DEFM	
33C0	2A	01660 H16	DEFM	
33E0	2A	01670	DEFM	
3400	2A	01680 M1A	DEFM	
3420	2A	01690	DEFM	
3440	20	01700 H2B	DEFM	
3460	2D	01710	DEFM	
3471	54	01720 M3A	DEFM	
3491	20	01730	DEFM	
34B1	32	01740 M4A	DEFM	
34D1	20	01750	DEFM	
34F1	53	01760 M5A	DEFM	
3511	4F	01770	DEFM	
3531	47	01780 M6A	DEFM	
3551	44	01790	DEFM	
3571	20	01800 M7A	DEFM	
3575	211640	01810 ENTRY	LD	HL,4016H
357B	36E3	01820	LD	(HL),DE3H
357A	23	01830	INC	HL
3578	36D3	01840	LD	(HL),03H
3570	210030	01850 BACKUP	LD	HL,3000H
3580	11003C	01860	LD	0E,3C00H
3583	010004	01870	LD	8C,400H
3586	E0B0	01880	LDIR	
3588	3A403B	01890 L00P1	LD	A,(3840H)
3588	A7	01900	AND	A
368C	28FA	01910	JR	Z,L00P1
368E	CDA436	01920	CALL	0ELAY
3581	2634	01930	LD	H,34H
3583	163C	01940	LD	0,3CH

Romplus and Other ROM Extenders

There are simpler methods of adding ROM to the TRS-80, where some of the work has already been done for you. The *Micro 80 Computer Club of Ontario* (in care of Brian Harron, 67-3691 Albion Road, Ottawa, Ontario K1T 1P2) has produced their Romplus board, capable of handling two 2708 1K ROMs or one 2716 2K ROM. It fits inside the keyboard, and only requires soldering three integrated circuits in place on the board, plus sockets for the ROMs.

Computer Accessory Technology (1307 Bagley Drive, Kokomo, Indiana 46901) has also developed a small board which contains a single 2716 EPROM and an address decoder. It comes caseless, but has a power supply and plugs directly into the TRS-80 edge-card connector. A set of programmed ROMs is available which include utilities of different kinds.



Personal Micro Computers (475 Ellis Street, Mountain View, California 94043) makes the REX-80 ROM extender, which is similar to the C.A.T. board, but comes with case and power supply, plus edge connector and cable. They also provide programmed ROMs. About \$50 for the device, \$25 for the ROMs.

Finally, *The Peripheral People* (P.O. Box 524, Mercer Island, Washington 98040), offers the Memory Sidecar ROM/RAM addition, which I designed, and is identical to the ROM/RAM addition presented in this book, without any bank-select features. The complete unit is \$149, and a blank board is \$25.

Continued Listing

```

3595 0604 01950 LO 8,4H
3597 E080 01960 LOIR
3599 3A4038 01970 LOOP2 LD A,(3840H)
359C A7 01980 AND A
3580 28FA 01990 JR Z,LOOP2
359F C0A435 02000 CALL DELAY
35A2 1809 02010 JR BACKUP
35A4 010040 02020 DELAY LO 8C,4000H
35A7 08 02030 LOOP3 DEC 8C
35A8 78 02040 LD A,B
35A9 81 02050 OR C
35AA 20F8 02060 JR NZ,LOOP3
35AC C9 02070 RET
35A0 20 02080 DEFM ' '
3581 20 02090 M8A DEFM '-----'
3501 20 02100 DEFM '-----'
35F1 20 02110 M9A DEFM ' - WHICH IS THE ACTUAL BLOCK'
3611 20 02120 DEFM ' OF MACHINE CODE THAT IS RUNNING'
3631 54 02130 M10A DEFM ' THIS TEST PROGRAM. TO EXIT THIS'
3651 20 02140 DEFM ' ROUTINE, YOU MUST PRESS RESET.'
3671 54 02150 M11A DEFM ' TO REPEAT THE TEST SEQUENCE, PRE'
3691 53 02160 DEFM ' SS THE <ENTER> KEY.'
3681 20 02170 M12A DEFM ' '
3601 20 02180 DEFM ' '
36F1 54 02190 M13A DEFM ' THE UNUSUAL LINES SEEN HERE ARE'
3711 41 02200 DEFM ' A REPRESENTATION OF MEMORY SPACE'
3731 55 02210 M14A DEFM ' USED BY THE TRS-80 TO CONTROL TH'
3751 45 02220 DEFM ' E DISK DRIVES AND CASSETTE PORT.'
3771 20 02230 DEFM ' '
3791 20 02240 DEFM ' '
37B1 20 02250 DEFM ' '
3C40 02260 ORG 3C40H
3C40 2A 02270 DEFM '** LOADING 2K RAM TEST TAPE COM'
3C5F 50 02280 DEFM ' PLETE. CONTROL TAKEN BY TEST **'
4016 02290 ORG 4016H
4016 7535 02300 DEFW ENTRY
3575 02310 ENO ENTRY
00000 TOTAL ERRORS
29944 TEXT AREA BYTES LEFT

```

```

BACKUP 3570 01850 02010
DELAY 35A4 02020 01920 02000
ENTRY 3575 01810 02300 02310
LOOP1 3588 01890 01910
LOOP2 3599 01970 01990
LOOP3 35A7 02030 02060
M1 3000 00160
M10 3240 00940
M10A 3631 02130
M11 3280 01260
M11A 3671 02150
M12 32C0 01580
M12A 3681 02170
M13 3300 01600
M13A 36F1 02190
M14 3340 01620
M14A 3731 02210
M15 3380 01640
M16 33C0 01660
M1A 3400 01680
M2 3040 00180
M2B 3440 01700
M3 3080 00200
M3A 3471 01720
M4 30C0 00220
M4A 3481 01740
M5 3100 00240
M5A 34F1 01760
M6 3140 00260
M6A 3531 01780
M7 3180 00280
M7A 3571 01800
M8 31C0 00300
M8A 3581 02090
M9 3200 00620
M9A 35F1 02110

```

A Front Panel Monitor

The first personal computers were a hobbyist's dream and a user's nightmare. These large boxes of electronic boards were programmed by hand, one byte at a time. The operation of the processor was stopped, an address was selected, and a byte programmed — all by using nearly 30 switches.

The TRS-80 is a far cry in size, speed, power, and convenience from these early machines (but don't forget that 'early' means 1974!). Yet there was an advantage in these early machines that the TRS-80 and its kin don't have: the front panel display. The front panel not only contained the multiple switches, but also a bank of LEDs so the user could view the contents of memory, registers, etc. (see Photo 8-2.).

The front panel is not an entirely obsolete concept, and can be remarkably valuable when your machine language programs sprint for the exit when you're not looking. Since the front panel visually monitors addresses and data, it provides somewhat of a window opening on the computer's activities.

The front panel can tell you if the machine is caught up in a deadly tight loop, if it is still processing (have you ever waited while the computer sorted string data?), or if it is operating in the area you expect it to. You can follow peripheral accesses like printers and disk, sound output routines and special devices you may have created.

Figure 8-18 presents the circuit for the micro front panel. It is no more than a group of latches which are activated by certain conditions — you may select any combination of input, output, read or write signals to trigger the LED displays. 16 LEDs monitor the address, and eight monitor the data.

The latches are triggered on an upswing of any signal line that is switched into the select gate; the gate may be switched off entirely, leaving the last latched address displayed on the LEDs.

Ideally, the front panel can be created from subminiature, 'grain of wheat' LEDs, and mounted directly in the TRS-80 case and soldered in place. Alternatively, it may be connected when needed by a standard 40 conductor cable. Method of construction is not critical, and can use soldering and wire wrap.

A Front Panel Monitor

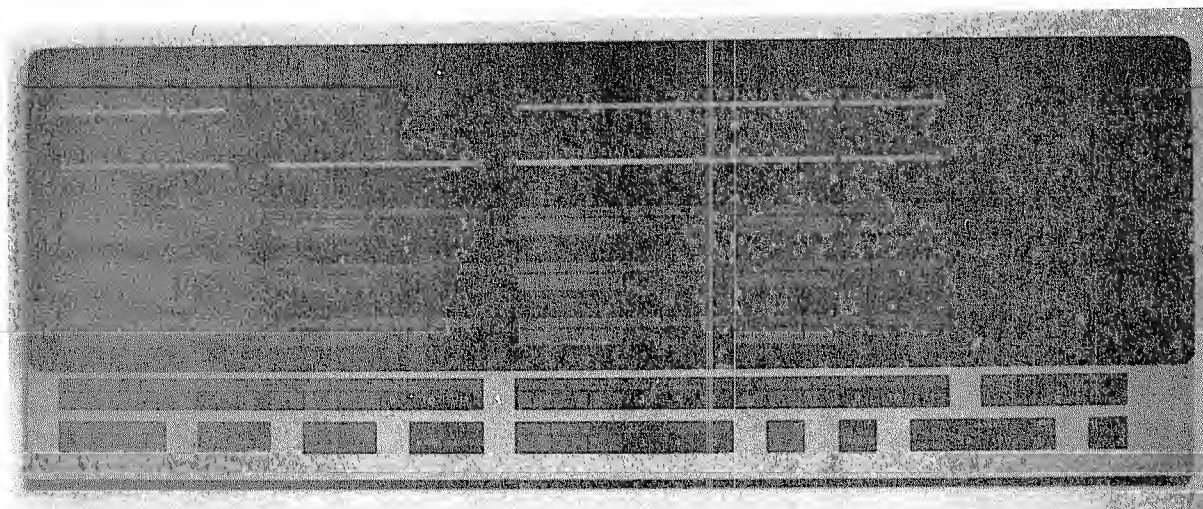


Photo 8-1. DEC front panel.

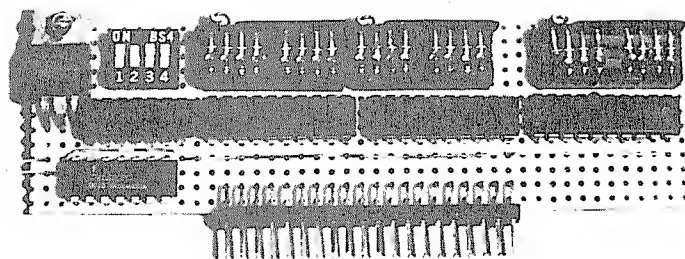
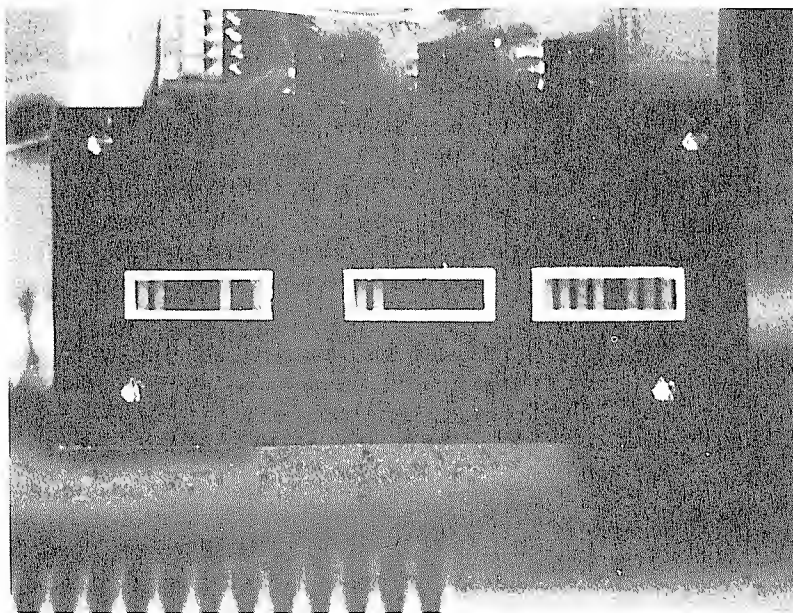


Photo 8-2. Micro front panel.

Micro front panel monitor shows data and address lines from read, write, input, or output conditions.



Micro front panel can be plugged directly into the edge card or put at the end of a connector cable, as shown here.

When using the micro front panel, you'll notice that much of the activity, particularly memory reads, run by very fast. It can't be helped. Because the TRS-80 uses dynamic memory which must be refreshed every few thousandths of a second (see description earlier in this chapter), the instruction clock cannot be stopped. Thus, you'll have to get used to the fast operation, noting from the intensity of the LEDs the frequency with which an area of memory is accessed. The monitor is remarkably useful when running diagnostic routines, because it points out whether there are any obvious signal line flaws. If the diagnostic program is a tight loop (most are – see chapter 10), then the activities of the computer's signal lines will become very obvious on the micro front panel.

In any case, it's easy to build and nifty to watch – very instructive to see how the computer accesses disks, for example, or how one data line pulses during cassette input/output.

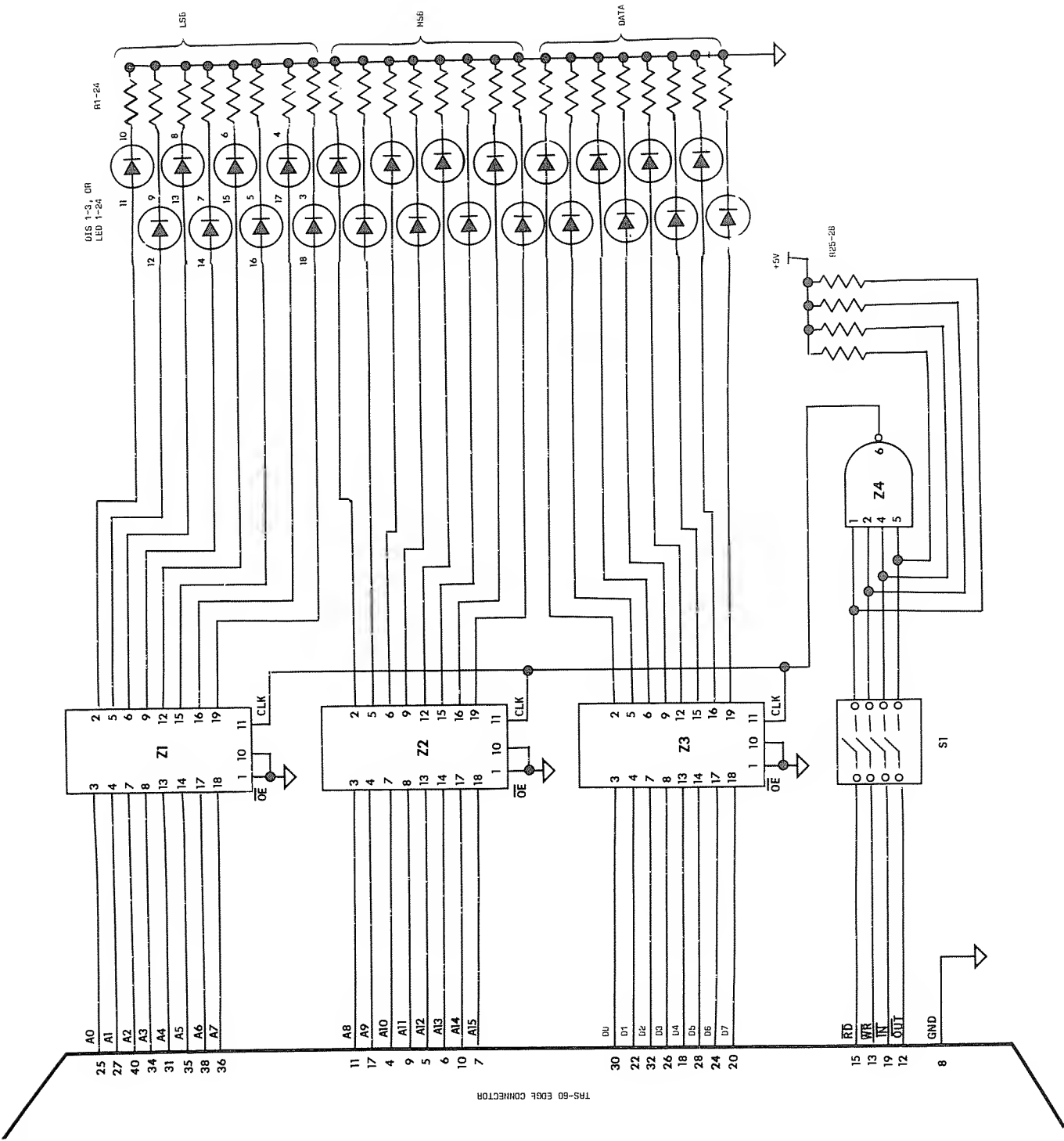
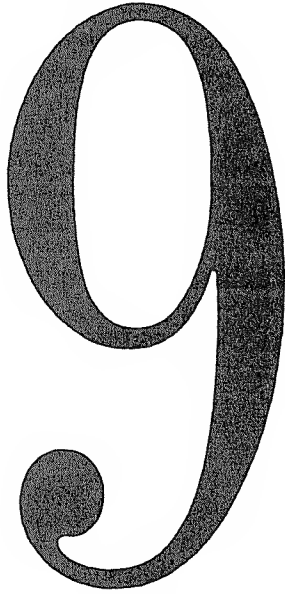


Figure 8-18. Micro front panel.

NOTES



Keeping It Safe: Mass Storage

Mass storage for computer programs and data has evolved from the days of punched paper tape and punch cards to present systems involving magnetic disk storage and bubble memories.

For microcomputer users, mass storage presents a unique problem: slow and apparently unreliable tape systems are inexpensive but frustrating, whereas fast disk-based systems seem transparent, trouble free, and expensive. Both of these generalizations border on myth.

What are the true advantages and disadvantages of the competing systems? This chapter will cover them in some detail, but briefly, here they are:

1. Cassette tape is a slow to medium speed storage medium (500 bits per second is normal for the TRS-80, with systems capable of transferring data at more than 3000 bits per second).
2. Tape is an inexpensive storage medium (3/1000 of a cent per bit), and is available across the counter. The hardware is pre-configured on the TRS-80, and the operating system (such as it is) is in ROM. Total system cost, including several boxes of blank cassettes, is under \$100.
3. Tape is a reliable storage medium, where good materials are used and care is taken. In a properly aligned system using low-noise, high-output cassettes, my own tests showed an average of one loading failure in 2.5 million bits.

4. Higher speed tape loaders add to the cost of the tape electronics, reducing their attractiveness, but increasing their speed and often their reliability where marginally recorded tapes are present.

5. Tape storage is sequential access, except where special digital tape systems are used. A program, once in the process of being loaded, is not tried again and again should a bit failure occur. Data storage problems are compounded, with double-dumping almost a requirement, and sequential conception and operation of programs essential.

6. The tape operating system is an essential part of the language, except where outboard devices (such as *Fastload* and similar systems) are used. Thus, it cannot be changed without adding a non-transparent, RAM-resident patch to the operating system.

7. Disk systems are capable of fast transfer rates (125,000 bits per second is possible), but the orientation towards records, directories, and other internal checking and referencing systems creates an effective transfer rate of 10,000 bits per second or less.

8. Disk systems contain built-in error checking and re-try routines. Thus, although data transfer errors occur as frequently as errors on cassette tape (if not more so), the random-access and retry capabilities make them appear error free

except where disks have been physically damaged, written over in error, or demagnetized. The apparent reliability is very high.

9. Disk systems are comparatively costly in their original hardware configuration, their operating systems, and their medium. However, the storage cost of the medium is comparable to that of cassette tape, at about 2/1000 of a cent per bit. The initial hardware investment is higher, requiring disk control (\$300 for an expansion box or similar controller) and the drive (\$350 or more, with quality increasing with price), and an operating system (\$25 to \$250, again depending on needs). An initial investment in disk storage, including a box of disks, can begin at \$750 and end in the multiple thousands.

10. System flexibility is increased greatly, as the disk's operating system and BASIC language additions overlay each other as needed, and appear almost transparent to system operation. However, the plethora of disk operating systems and approaches limits the interchangeability of information from one TRS-80 to another with a different operating system.

11. Intermediate and hybrid systems are available that encompass some of the features of both standard tapes and disks. Foremost among these is the Exatron Stringy-Floppy endless-loop tape cartridge system. Its operating system is ROM-resident, its transfer rate is 7,200 bits per second.

12. The Stringy-Floppy is probably the most reliable mass storage system under adverse environmental conditions, putting it above tape and far beyond the sensitive (some say temperamental) disk systems. This aspect more than any other probably justifies its consideration as serious mass storage. Based on hi-tech, laboratory models, the ESF is a scaled-down scientific storage system.

13. Like tape, access using the ESF is sequential, but the endless loop makes pseudo-random access possible. With short tapes and programs less than 8K bytes, actual load/save time is faster than disk.

14. Cost of this system is less than disk (\$250 for the hardware, \$3 for the medium,

an endless-loop 'wafer'), and the cost of storage is less than all other complete systems (about 1/1000 of a cent per bit). To be competitive with dropping disk prices, I expect to see the hardware cost drop.

15. Although less fragile than disks (and higher in quality of the magnetic surface), the ESF wafers, because they use thin tape on a tiny, endless-loop hub, can be damaged by the tape binding or pulling from the housing. Unlike the larger cassettes, the tape cannot be successfully reinstalled in the housing, and unlike disks, the undamaged material on the wafer cannot be recovered.

This chapter will present a tour through the available mass storage systems (except tape – see Supplements to Chapters 3(?) and 6(?)), describe the construction of a paper tape reader, and present the construction and operation of a tape storage device using 8-track cartridges.

Disk Drives

Disk drives come in a variety of sizes, shapes and formats. Among those are floppy disks in 5-inch and 8-inch sizes; removable and permanent hard disks; and permanently housed Winchester drives. Miscellaneous variants of all kinds are for sale or under development.

The most popular system for the TRS-80 is the 5-inch floppy disk system. The drive contains a platter which spins the magnetic disk inside its cardboard sleeve, a record/playback head to read the data, and a stepping motor to move the head from concentric track to track. The data is recorded using sharp digital pulses without any sort of audio recording considerations such as high-frequency bias. DC 'trim erase' is used to remove previously recorded data just ahead of the write head.

A single indexing hole near the center hub is used to inform the drive electronics and control software of the disk's position inside the paper sleeve. Other than this hole and a write-protect notch in the top edge of the paper case, there is no other information available to the drive and control software from a blank disk. It must be formatted, which is the process of embedding magnetic information on the disk for later use by the disk operating system.

The information needed for the original Radio Shack specifications included the magnetic outlining of 35 concentric tracks on the disk, and the separation of those 35 tracks into ten discontinuous bands each, called sectors.

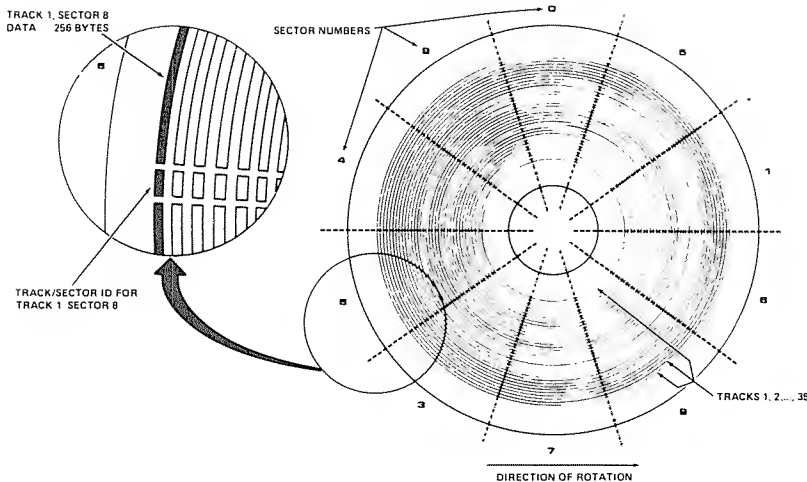


Figure 9-1. Disc and disk system format.

Details of the system by which information is recorded on the disk can be found in the *TRSDOS 9 Disk BASIC Reference Manual*, and *TRS-80 Disk 9 Other Mysteries*. In summary, the original Radio Shack specifications called for a disk containing 35 tracks of 10 sectors of 256 bytes each, for a total of 83,060 available bytes (89,600 are actually recorded, but several thousand are reserved for directory and other information – again, see the references above for details).

Drive manufacturers and software authors immediately began modifying the Radio Shack standards. 40, 77 and 80 tracks could be used, and nearly a dozen different disk operating systems (DOS) made their appearance: the original TRSDOS from Radio Shack (versions through 2.3 now issued); NEWDOS, an updated and corrected version of TRSDOS, sold by Apparat; NEWDOS/80, a complete re-write incompatible with the others; VTOS (now issued through versions 4.0), written by the original author of TRSDOS; MicroDOS and its successor OS/80, a stripped-down, efficient, minimal DOS which is growing larger and hence less attractive; DBLDOS, a Percom entry that operates its double-density (80-track) hardware option; DOSPLUS in single- and double-density versions, a wide-ranging program by Micro Systems Software; CP/M, supposedly a standard DOS for microcomputers, but available only in modified form for the TRS-80; and many special operating systems to load and run protected software.

Choosing a disk operating system is outside the range of this book, to be sure. Incompatibilities ride rampant over the bytes of the TRS-80 computer; the consistent, convenient, accessible Level II ROM gives way under disk control to the

whims of software authors and entrepreneurs, good and bad. It might seem like I'm knocking DOSes; I'm not, because it's truly a customizer's dream. But a problem arises when trying to deal with customization: to begin with, a disk operating system *is* a kind of customization. Hence, it becomes almost impossible to customize one further without being forced to provide versions for every popular DOS. There are too many, and they change quickly . . . to patch them invites problems, ill will, and frustration.

On the other hand, a disk system is ideal for customizing the TRS-80, because, once you have selected the operating system, you may modify and change it, making it your own. In a sense, that is how Apparat created NEWDOS – as a series of patches and improvements to TRSDOS, and in fact it originally required that disk users already own a copy of TRSDOS.

Inside A Disk Drive

The insides of a disk drive seem incredibly simple. In fact, they are. It is only the precision of alignment needed and a few expensive parts which bring the cost so high. The drive consists of an electronic control board which is capable of communicating with the controlling computer, in this case the TRS-80, through a disk controller chip (FDC – see below). There is a motor which spins the disk, and in most inexpensive drives this motor is connected through a drive belt to the disk hub. The motor speed of 300 rpm must be accurate to within five percent; three percent deviation from normal is reasonable, and most drives are capable of a 1.5 percent long-term deviation. The disk is inserted in the housing and held in place by a cone and pressure plate which fit around the hub, and clamp gently but firmly to the disk's center area.

When the door to the drive is closed, the pressure plate moves into place, and the disk is brought into contact with the read/write head. A properly seated disk will present the indexing opening to a light sensor, so that as it spins inside the paper envelope, the index hole will open the light beam as it passes by.

A stepping motor is capable of moving precisely to one position on its axis, on command. This type of motor is used to position a magnetic read/write head on the disk track to be read. The motor is fast and precise, and the assembly which it moves is carefully machined so there is virtually no up-down play in the head. Opposite the head is a pressure pad which forces the disk

to maintain contact with the head, albeit separated from the head by a few microinches.

The head itself is usually of the glass-ferrite variety; it has an extremely smooth, highly polished surface that will not damage the disk, and a very long life that exceeds 20,000 hours of continuous head-to-disk contact. It is capable of handling the high write currents generated by the digital circuitry, and virtually immune to electronic noise in its vicinity.

In the TRS-80 system, the disk drive itself accepts and sends certain pieces of information. They are:

1. A 4-bit drive select indicator sent to the drive; only the drive hard-wired to accept this signal will respond. In the TRS-80 system, this wiring is done in the cable itself.
2. A motor-on signal, which turns on the 300 rpm hub motor.
3. A track-to-track stepping signal and a stepping direction signal.
4. A write-enable signal and a stream of written data.

5. A stream of written data sent by the drive.

6. A write-protect signal to prevent writing to write-protected disks. Most drives will not respond to a write signal, and this write-protect signal is sent so the software can report a write-protected condition.

7. An index pulse to indicate where the disk is currently located in its rotation.

8. A track-zero indicator to identify when the disk head has reached the outermost track on the disk. This is used to locate tracks, relocate tracks or reposition the head, and on initial access to identify the head position.

Except for the drive-select signals, the disk controller chip is responsible for managing all these lines. In the TRS-80, a type 1771 controller is used, manufactured by *Western Digital*. A complete data sheet is provided with the Expansion Interface service manual from Radio Shack.



The author at work.

This controller has eight data input/output lines (DAL0 to DAL7) to the computer, an interrupt output to signal its need for service (INTRQ), read and write enables (RE and WE), and a two-bit address select (A0 and A1). These last are used to select the register that will send or receive data on the DAL lines:

A1	A0	RE	WE
0	0	Status Register	Command Register
0	1	Track Register	Track Register
1	0	Sector Register	Sector Register
1	1	Data Register	Data Register

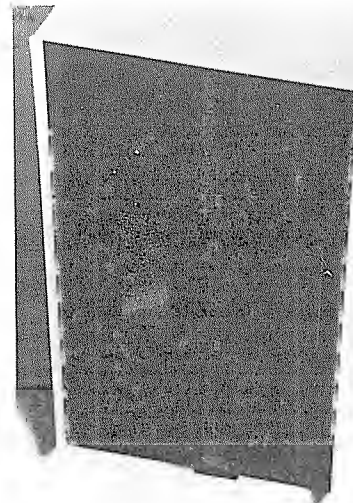
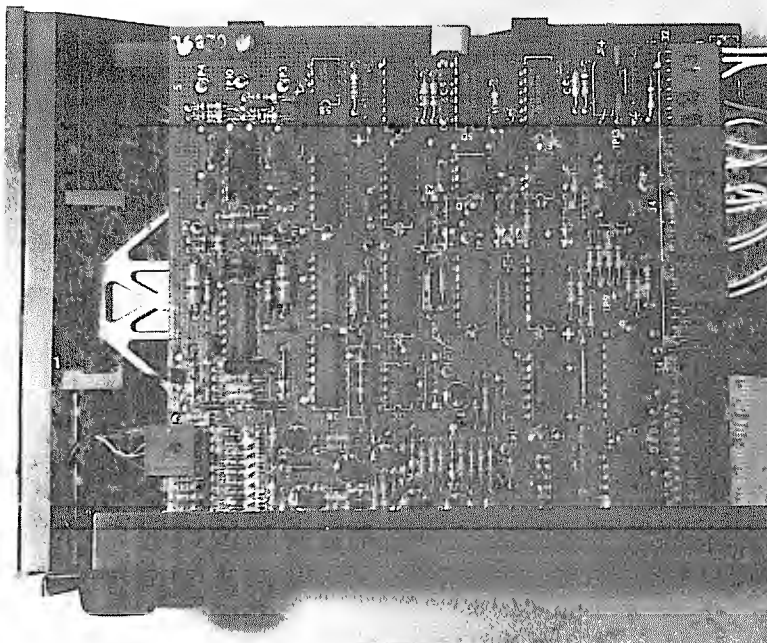


Photo 9-1. Pictorial tour (6) of disk drive.
Front door of disk drive controls mechanism to steady the disk as well as a switch to indicate that the door is closed.



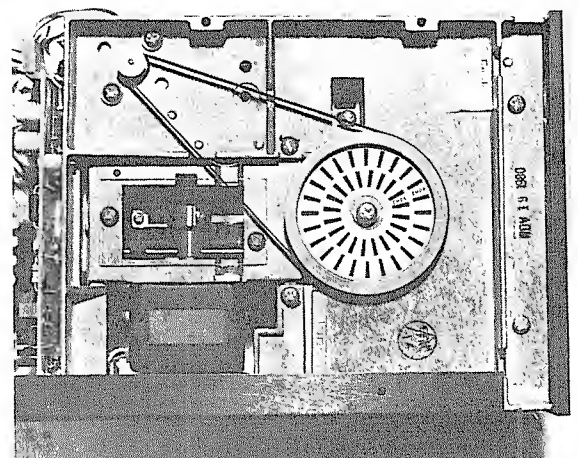
Circuit card contains disk read/write circuitry, motor stepping control, and computer interfacing. Computer connector and termination resistors are seen at lower right.

In the TRS-80, the chip is always selected (CS) because the buffering of data is taken care of by Z33, Z37 and Z38 in the expansion box. The Data Request signal (DRQ) is not used, nor are the three-phase motor signals (PH3 and 3PM), and the track-greater-than-43 signal (TG43). These signals would be present were a more sophisticated disk drive capability intended by Radio Shack in future versions of the TRS-80.

The clock input is provided by the separate oscillator in the expansion interface. Three voltages are needed (+12, +5, and -5), plus ground.

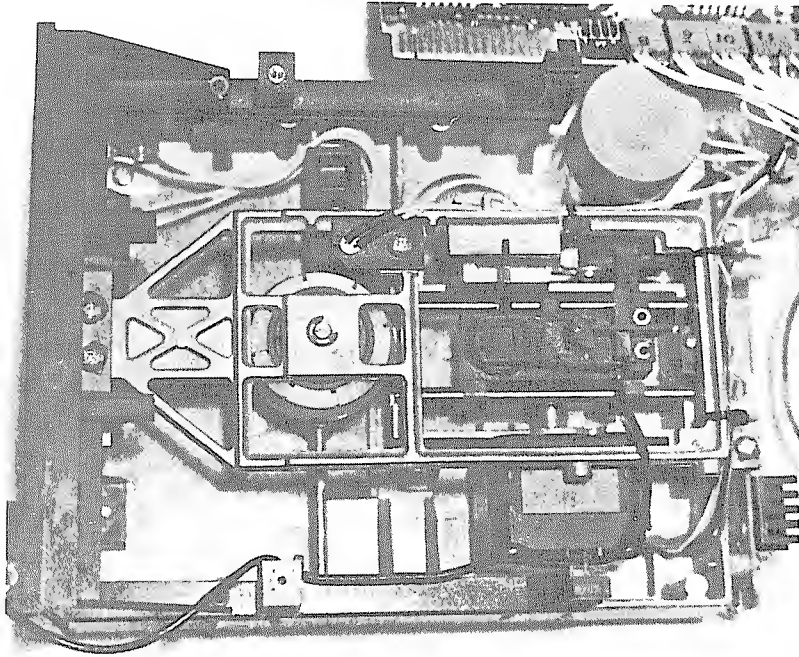
Test, disk initialization (DINT), and write-fault lines (WF) are not used, nor are the infamous external data separation (XTDS) and external data clock (FD CLOCK) lines. Disk data are separated by clock pulses, and high accuracy demands unflinching differentiation between the clock pulses and the data pulses. The Percom data separator plugs into the controller chip socket and makes use of XTDS and FD CLOCK.

The head-load output (which determines whether the read/write head is in place against the disk) is not used, because software controls the timing between motion of the read-write head. Head-load timing (HLT) and READY are both connected via external logic to the TRS-80, where software determines when the drive and read/write head should be ready to read or write. The remaining disk controller signals lead to and from the disk drive itself, and are identical to those listed in the description of the drive signals.

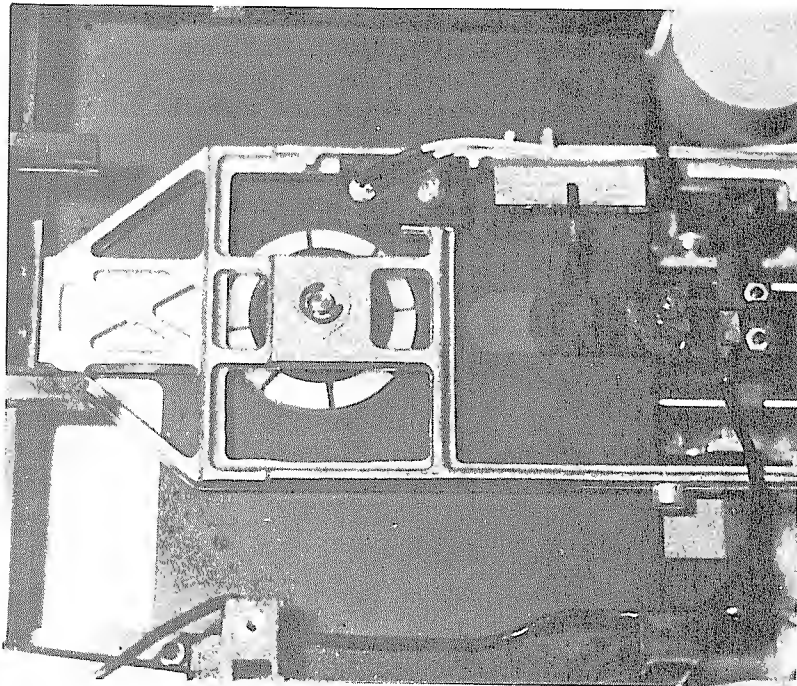


Opposite side from control card is the drive motor with speed strobe disk, and a window revealing the shaft of the head stepping motor.

Inside a Disk Drive



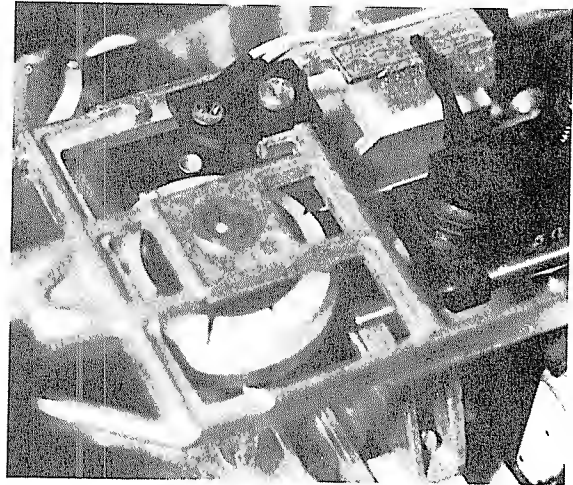
Removing the control circuit card reveals a heavy cast frame to hold the disk in place. In center is the cone that fits through the disk; to the right is a pressure pad.



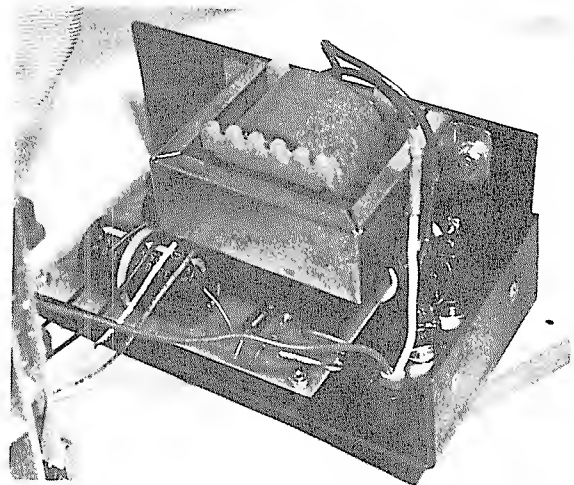
The drive with a disk in place. Cone fits through disk and clamps it in place. Head and index hole light emitter are on the other side of the disk; pressure pad and light sensor are on this side.



Close-up of read/write head and pressure pad. White square is glass ferrite head surface, vertical stripe is the head's read/write surface. Pressure pad moves into place when the disc's front door is closed.



Close-up of positioning sensor. Index hole of a disk allows light to pass from LED emitter to sensor (top center).



Hefty disk drive power supply handles motor drive current and powers the electronics.

A Heavy Dose of DOSes

On the software side, the number of disk operating systems for the Model I continues to increase. Most include 'Level III BASIC', the usual copying and formatting capabilities, but have been expanded to include more than that. For a complete description of the features of competing DOSes, refer to the March 1981 issue of 80 Microcomputing. Among the special features of the most prominent DOSes:

TRSDOS. AUTO, ATTRIB, CLOCK, COPY, DATE, DEVICE, DIR, DUMP, KILL, FREE, LIB, LIST, LOAD, PRINT, PROT, RENAME, TIME, VERIFY. Also BASIC, BASIC2, DEBUG, TRACE. Level III BASIC functions. Details are found in the TRSDOS

manual, and other DOSes include all of these commands in some way or other.

NEWDOS+. Adds COPY, JKL screen print, DIRCHECK directory verification, LMOFFSET tape load offset module, EDTASM with modifications, SUPERZAP for modifying disk contents, LEVEL 1 located in RAM, DISASSEM, LV1DSKSL disk save/load for Level I. CMD"DOS COMMAND NAME" to execute from BASIC, and several commands to re-enter BASIC from DOS.

NEWDOS/80. CHAIN, HIMEM protection from DOS, JKL, MINIDOS don't-distrub-memory DOS, MDBORT for killing MINIDOS, PDRIVE setup for multiple drive types, PURGE for killing file groups, SYSTEM to create special commands, SUPERZAP, LMOFFSET, LEVEL 1, ASPOOL print spooler, DIRCHECK.

VTOS. BOOT software reset, BUILD a group of auto-execute programs, CHAIN to execute them, MEMORY, PURGE, RUN for non-VTOS systems, SYSTEM, XFER disk copy program, PATCH disk modification routine, VTCOMM communications utility, KSR terminal program, ROUTE for changing the destination of data, SET for a user device program, SPOOL, RESET to cancel device setup, LINK for devices, FILTER to be used with device routing, ALLOC setting up disk space in advance.

DOSPLUS. Adds BOOT, BUILD, CLEAR a directory file, DO group of auto-execute programs, FORMS to set up the printer driver, a different variant of FREE, PAUSE for user input in auto-execute routines, RS232 for a report on that status, PURGE, DISKZAP for modifying disk information, CLRFILE for zeroing a file, COPY1 copying utility, CRUNCH space compression for BASIC, TRANSFER program copier.

ULTRADOS. CLEAR for zeroing memory, DEAD for zeroing memory, TOPMEM for setting DOS-protected memory, CMD"C" compression routine, CMD"DOS COMMAND NAME" which executes from BASIC, CMD"O" file buffer allocation, CMD"X" to return to BASIC from DOS, cross referenced listing of variables and line numbers, renumbering from BASIC, shorthand command keystrokes.

A Garden Full of Varieties

The available variety of disk drives is growing. Not only are the capabilities of standard 5-inch drives being stretched (to run 40, 77 and 80 track densities), but 8-inch systems and hard-disk systems are being introduced for the lowly TRS-80. Among them:

Manufacturer	Model/Type	Price	Comments
Minifloppy (5-inch) Disc Systems			
Access	AF0-100	\$315	40 tracks; Tendon?
Aerocomp	40-1	\$350	40 tracks; double-density
	80-1	\$460	80 tracks; double-density
	80-2	\$460	40 tracks; double-sided; double density.
	160-2	\$600	80 tracks; double-sided; double density.
CPU Shop	CCI-100	\$315	40 tracks; Percom?
	CCI-200	\$430	80 tracks; Percom?
MPI	8/51	\$320	40 tracks; 5 ms track-to-track; auto-eject; double-density head.
	8/81	\$425	80 tracks; 3 ms track-to-track; auto-eject; double-density head.
Micropolis	MCP1027	\$300	35 tracks, single head
	MCP1037	\$700	35 tracks, dual head
	MCP1027-2	\$440	77 tracks, single head
	MCP1037-2	\$800	77 tracks, dual head
Percom	TF0-100	\$350+	40 tracks
	TF0-200	\$650+	77 tracks
Pertec	F0200	\$380	40 tracks
Shugert	SA-400	\$330	35 tracks
	SA-410	\$340	40 tracks
Siemens	F00 100-5	\$275	40 tracks, double density
Tendon	—	\$...	40 tracks
TEAC	—	\$300	40 tracks
	—	\$400	80 tracks
Vieta	V-80	\$400	40 tracks
	V-800	\$600	80 tracks
	V-8800	\$775	160 tracks

The Exatron Stringy-Floppy

In between tape systems and disk systems is an unusual device. I only say 'in between' because the name Stringy-Floppy implies that it somehow is a tapelike version of a floppy disk device. In fact, it is not, for two major reasons: it is a unique, high reliability, high endurance storage medium; and it does not (at least in its TRS-80 configuration) contain formatting, sectoring, or record-keeping of the type used for disks.

The *Exatron Stringy-Floppy* (ESF) consists of a small DC motor which uses a tiny plastic belt to drive a capstan, exactly as in a cassette player. The motor's speed, however, is set to 10.5 inches per second, quite a bit faster than cassette, 8-track, or even open-reel. Thus, even though the tape itself is only 1/16 of an inch wide, reasonable data recording can be expected.

In the TRS-80 version, the ESF does not use any standard method of recording. A separate sequence is used to prepare, or 'verify' an endless-loop tape wafer than that which is used to record data. Both methods use a variety of bi-phase recording, in which the bit being read or recorded depends on the polarity of the bit just written or read.

This makes the ESF reliable in spite of a motor speed which varies more than ten percent in either direction of its ideal speed. It also means the ESF can work under harsh conditions which might otherwise throw a standard clock/data system far off.

Programs are recorded on the ESF in a continuous stream, preceded by a leader and followed by information to assist the software in locating the next available blank program space. Programs and data may be read in any order, but must be written in ascending numerical order.

The ESF operating system is contained in a read only memory which resides in the unassigned memory area on the Model I (3000 to 37C0); thus, it is always available and relatively crash-free. It does use a few bytes of RAM, however, and must also patch into the Level II BASIC parameters in low memory.

The wafer itself is available only from *Exatron*, and is based on the wafers used in industrial applications. It consists of a length of high-output digital recording tape, highly polished. This is wound on a small hub, lubricated, and spliced into an endless-loop. A reflective splice is used to determine end-of-tape (EOT/BOT), and is viewed through a window at the top of the wafer.

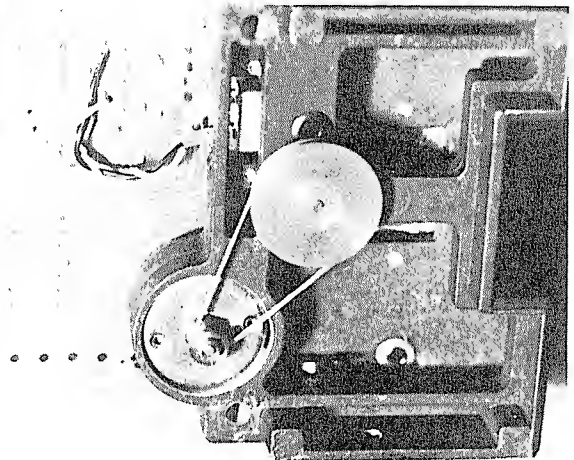
The wafers may be write-protected by affixing a reflective-tape dot on the wafer. This is read by the operating system before the drive is activated. The ESF operating system consists of these commands:

@LOADx Load a program; **x** is the program number, and is optional. An error in the load



Photo 9-2. Pictorial tour (6) of ESF system.

Exatron Stringy-Floppy placed next to Microconnection modem. Both are compact, light devices.



ESF 'rubber band' drive is not as reliable as disk drives, but software is capable of handling greater variations in speed. Small DC motor replaced by a direct-drive type in later versions, say Exatron officials.

terminates the process and the error (parity or checksum) is reported to the user.

@SAVE x Saves a BASIC program; the x is the program number, and is required. The program is verified, and any saving errors are reported.

@SAVE x,r,s,t Saves a block of memory; x is the program number, r is the starting address, and s is the length of the block. These are required. The t is an optional entry address. All numbers are in decimal. The block is verified, and any saving errors are reported.

@NEW x Verifies a wafer. When x is used, the wafer is cleared and verified from that program to the end splice. Available byte count is reported.

Fastload, TC-8, and Other Systems

Most other alternatives to disk systems involve using standard cassettes for high-speed saving and loading. Among these are Fastload, manufactured by *Personal Microcomputers, Inc.*; TC-8, or the 'Poor Man's Floppy', made by *JPC Systems*; and the Beta-80, made by *Meca Technology*.

The advantage to the standard cassette is simple: cost and availability. Unlike disks, they are able to take some measure of abuse, and can be replaced without regret over the expense. And

unlike Exatron wafers (or disks, for that matter), they can be purchased if necessary in the local grocery. Furthermore, cassette tape technology has progressed further than disk technology, providing better surfaces, adhesion, and signal-to-noise ratio. Disks are still in the dark ages of reproduction compared to audio tape.

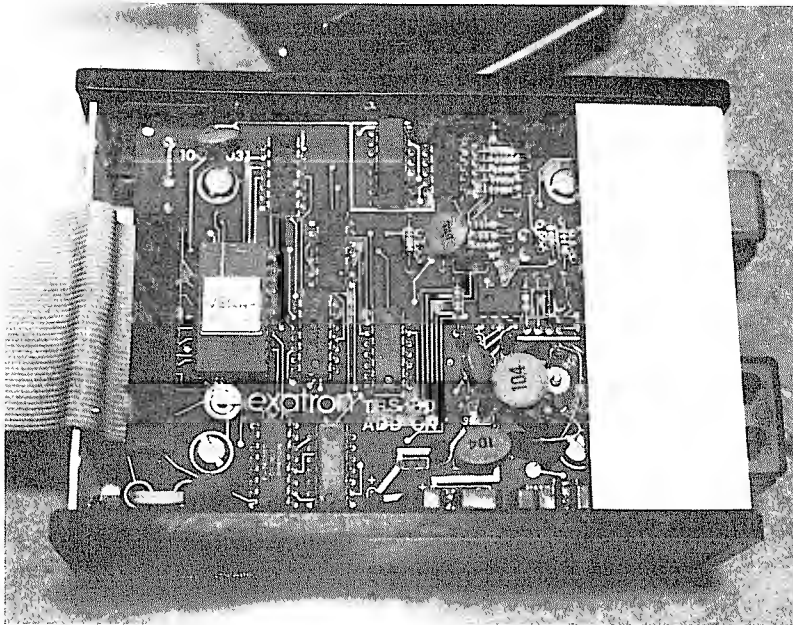
Fastload is a ROM-based operating system combined with a hardware detection and shaping circuit. A modified CTR-41 is used, where the fast-forward button and play button can be locked down. The circuit then can read a standard 500-baud tape in the fast-forward mode quite reliably at about 8000 baud. Debounce, audible beep, and key repeat are included with the operating system.

Although I have not examined the schematics for Fastload, its carefully designed circuit board, with voltage regulator heat-sinked to the case, attest to a cautious, probably over-designed system. The loading system is put into operation by typing `SYSTEM (ENTER), /12288 (ENTER)`; optional debounce/beep/repeat can be added as well.

When in operation, Fastload uses the single `LOAD` command in Level II, and is compatible with disk-based computers by using a `SYSTEM` call. The 500-baud load and save are left undisturbed. Fastload is well-designed, and reasonably reliable. If your original tapes have audible 'bumps' when played at high speed, however, they may not load easily with this system. Commercial tapes, or those recorded by the user on virgin cassettes, load easily and extremely fast. 'Bumpy' tapes, however, do not load well, and have to be re-recorded.

One of the other disturbing features of Fastload is its tendency to give a 'READY' message even when a BASIC program has been loaded incorrectly. Granted, it is almost impossible to be sure a BASIC program has loaded well because, in its normal `CSAVE` format, no checksums have been provided. But the user should be sure to list the program before expecting it to run. `SYSTEM` loads, on the other hand, present checksum error messages on bad loads. Fastload costs \$188 assembled; a modified CTR-41 is \$95.

TC-8, or the Poor Man's Floppy, both saves and records at high speed. It is provided in kit form; assembly can be completed in a few hours. Like Fastload, the TC-8 plugs into the edge connector, but the software is RAM-resident. The standard CTR-80 tape recorder is used, and a two-wire swap is shown to make the CTR-41



Control card of ESF contains drive electronics, tape read/write circuitry, and 2716 EPROM containing the ESF operating system. Power supply is underneath the card.

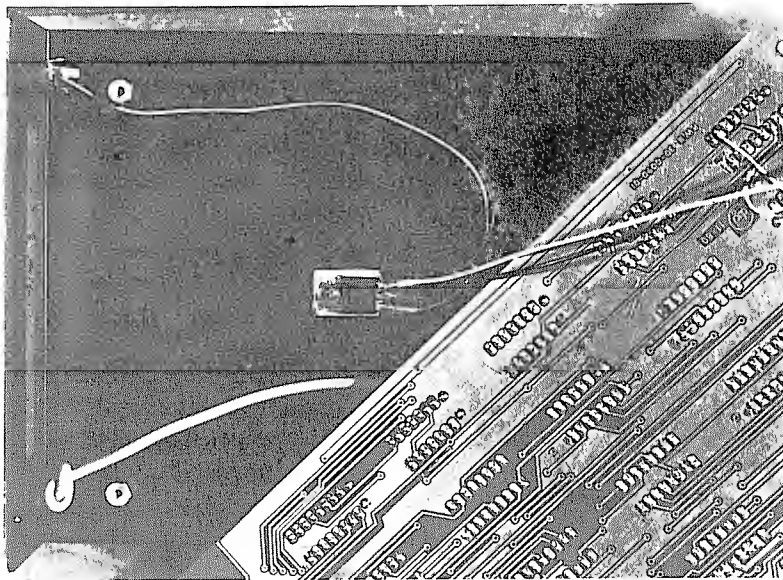


Photo 9-3. Fastload device.

Instead of a small heat sink, Fastload power supply regulator is sunk to the entire metal case. Insulators around regulator show attention to detail.

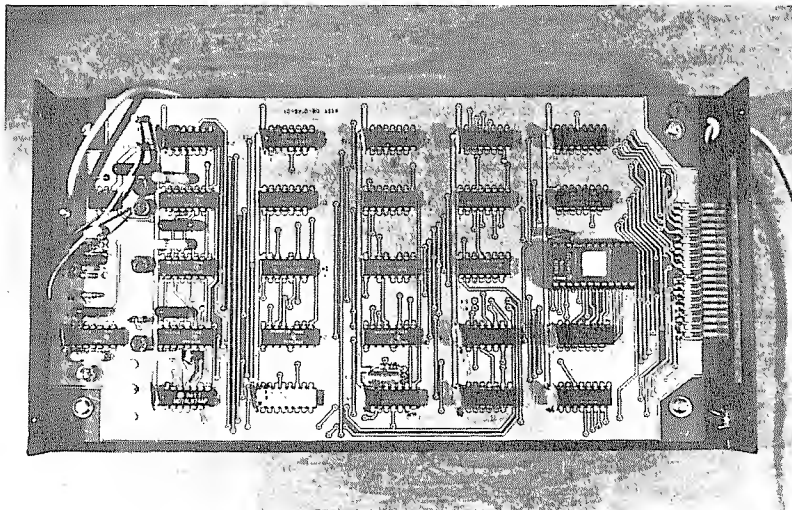


Photo 9-4. Fastload circuit card.

Circuit card for Fastload shows extremely conservative design, with careful attention to grounding and possible noise sources. DIP resistor package is used in place of separate resistors, and the operating system is contained in a 2716 EPROM. Since less than 1K of the EPROM is used for the operating system, the remainder may be programmed by the user with utilities.

run properly (its high current tended to fuse motor relays together on the TRS-80).

Since it is both a record and save device, the TC-8 contains fifteen separate commands, including SAVE (BASIC programs only) and PUT (machine language blocks); LOAD BASIC programs next or by name; LOAD? for verification; LOADN to position the tape; comparable GET commands for machine language programs; RUN as a load-and-go command; RSET for cassette motor on; OPEN, CLOSE, PRINT, and INPUT for file management; and KILL to eliminate file management.

The software can be reconfigured, relocated, and stripped to a memory-saving bootstrap version. The TC-8 is provided with extensive documentation which is clear and literate, providing examples, a sample program using the OPEN, PRINT#, INPUT#, and CLOSE commands, along with recommendations, suggestions, and warnings (plus sympathy!) about the problems of transferring some commercial machine-language software to the TC-8.

The circuit is quite simple, but well designed and accurate; the capabilities of this elegant circuit comes from the software. The construction manual provides soldering suggestions and explanations, recommendations on the soldering iron to purchase, and drawings of properly soldered connections. For convenience, all work is done on the underside of the board; correct placement and orientation of parts is emphasized. Two-color layouts of the board are provided on every page, along with a checklist of each step. These are some of the finest assembly instructions I have seen for a project.

When complete, the unit loads and saves at 2500 baud on off-the-shelf cassettes. The company recommends certain brands, and also sells the brand *JPC products* uses. All signals to the tape player are non-critical, and the loading routine will accept signals over most of the audible, undistorted output spectrum of most cassette recorders. TC-8 costs \$90 in kit form.

The Beta-80 is a tape storage system patterned after professional digital tape devices. Although a standard digital cassette is used, the Beta-80 is configured something like a disk, with directory and fast access. Fast forward and rewind are automatically controlled by the RAM-resident operating system, and the search speed is 100 inches per second.

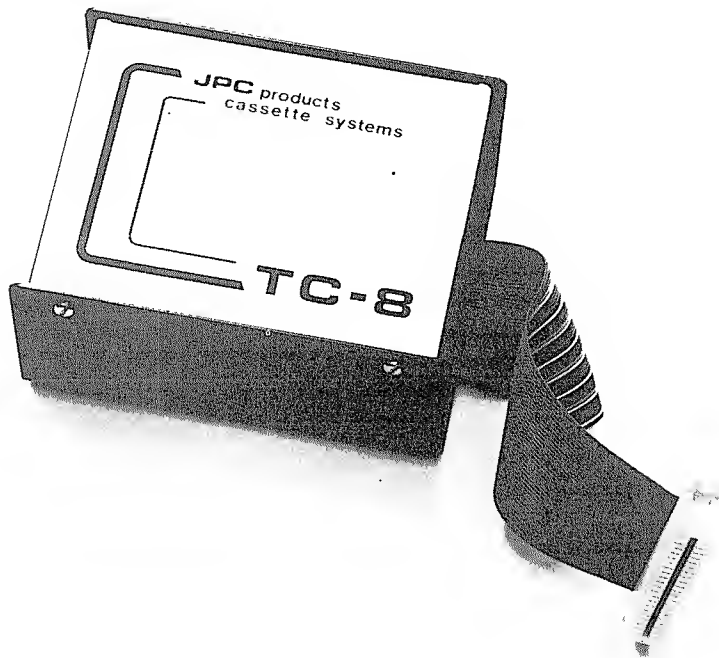


Photo 9-5. TC-8 tape system.

TC-8 tape system attaches to edge-card connector and cassette recorder to provide high-speed record/save of programs. Photo courtesy JPC Products Company.

```

00100 : 
00110 : #####
00120 : HIGH SPEED TAPE LOAD AND SAVE ROUTINE OPERATING AT 1960
00130 : BAUD. THIS CAN BE MADE TRANSPARENT TO BASIC BY USING
00140 : EITHER /SAVE AND /LOAD OR /PUT AND /GET. THE ROUTINE
00150 : PRESENTED HERE SAVES A BLOCK OF MACHINE CODE, BUT MAY
00160 : BE USED FOR SAVING AND LOADING BASIC BY LOCATING START
00170 : AND END OF THE BASIC PROGRAM (FOR WRITE), AND START AND
00180 : END OF BASIC MEMORY (FOR READ). SAMPLE ACCESS ROUTINES
00190 : ARE PRESENTED AT THE END OF THIS LISTING.
00200 : #####
00210 : 
00220 : ORG 7000H ; TOP OF MEMORY - RELOC.
00230 : 
00240 : #####
00250 : WRITE-BYTE SUBROUTINE. BYTE IS ASSUMED IN C REGISTER
00260 : #####
00270 : 
7000 060B 00280 WRITE LO B,B ; NUMBER OF BITS TO WRITE
7002 C5 00290 PUSH BC ; SAVE BITS OF BYTE
7003 0620 00300 LO B,20H ; GET A DELAY VALUE
7005 10FE 00310 OJNZ $ ; AND DELAY 117 USECS
7007 C1 00320 POP BC ; AND RESTORE BITS
7008 C01470 00330 WLOOP CALL TIMEX ; WRITE TIMING BIT
7008 C811 00340 RL C ; ROTATE BIT INTO FLAG
7008 17 00350 RLA ; ROTATE FLAG INTO A
700E C03670 00360 CALL BITEX ; WRITE BIT TO TAPE
7011 10F5 00370 OJNZ WLOOP ; DO IT TOTAL OF 8 TIMES
7013 C9 00380 RET ; BACK FROM WRITE ROUTINE
00390 : 
00400 : #####
00410 : ROUTINE TO WRITE TIMING BIT POSITIVE-NEGATIVE TO TAPE
00420 : #####
00430 : 
7014 C5 00440 TIMEX PUSH BC ; SAVE B ON STACK
7015 3A3040 00450 LO A,(4030H) ; GET SCREEN INFORMATION
701B E6FC 00460 AND OFCH ; MASK OUT LOW TWO BITS
701A F601 00470 OR 1 ; SET LOWEST BIT
701C 03FF 00480 OUT (OFFH),A ; WRITE BIT TO TAPE
701E 0607 00490 LO B,07H ; TIMING VALUE @ 1.77 MHZ
7020 10FE 00500 OJNZ $ ; IS EQUAL TO 67 USECS
7022 E6FC 00510 AND OFCH ; MASK OUT LOW TWO BITS
7024 F602 00520 OR 2 ; SET NEXT LOWEST BIT
7026 03FF 00530 OUT (OFFH),A ; WRITE NEG. BIT TO TAPE
702B 0607 00540 LO B,07H ; TIMING VALUE @ 1.77 MHZ
702A 10FE 00550 OJNZ $ ; IS EQUAL TO 64 USECS
702C E6FC 00560 AND OFCH ; MASK OUT LOW TWO BITS
702E 03FF 00570 OUT (OFFH),A ; WRITE NEUTRAL BIT

```

Listing 9-2. High-speed tape loading routine.

Loading is at 4000 baud (tape running at 5 i.p.s.), twice the speed of the TC-8, but about half that of the Exatron Stringy-Floppy and Fastload, but access time (using the directory) is less than a minute for more than a half megabyte of stored programs.

Commands are LOAD in the form of load, load and run, and load array information; SAVE for programs and arrays; MERGE for append (not a true merge) or append and run; and KILL to delete a program or an array. Other commands operate within the various RAM-resident systems.

The Beta-80 uses a reliable Phi-deck drive, and has impressive features. However, as of this writing, users report little support from Meca for converting and saving machine language programs on the Beta-80. The Beta-80 costs under \$300.

High-Speed Cassette Loading

Speeds higher than 500 baud are achievable entirely through software, and most systems work exceedingly well. Among the most popular are: SPEED, the first such program, loading and saving at 1500 (?) baud; HISPED (Palomar Software, 170 S. Palomar Dr., Redwood City, CA 94062; \$24.95), a 2000-baud system; ZIPLOAD, in the public domain and published in the 80 Encyclopedia (Wayne Green, Inc.); and B-17 (ABS Suppliers, P.O. Box 8297, Ann Arbor, MI 48107; \$25).

Each of these save/load programs patches into the BASIC operating system, and all use the DOS-reserved LOAD and SAVE commands. The last is of particular interest because it is a complete system rather than merely a save/load program.

SAVE, using a six-character name, sends BASIC programs to tape; LOAD? verifies them. LOAD puts the computer into the SYSTEM mode automatically and loads a BASIC program. PUT and GET are used for formatted file arrays. It is also provided with a machine language module. Checksum errors, full memory buffer, continuity errors, and format errors are reported by this module. The entire source code is available for sale if the B-17 program is also purchased.

Aside from B-17's high reliability, it also provides an on-screen prompt indicating saving and loading.

What makes a high-speed loader not only

A Paper Tape Reader

```

7030 0607      00580      LO      B,07M      ; TIMING VALUE @ 1.77 MHZ
7032 10FE      00590      OJNZ     $          ; IS EQUAL TO 113 USECS
7034 C1        00600      POP      BC          ; RESTORE NUMBER OF BITS
7035 C9        00610      RET          ; BACK TO BIT WRITE ROUT.
00620 ;
00630 ; #####
00640 ; ROUTINE TO WRITE INDIVIDUAL BIT TO TAPE (OR NOT IF 0)
00650 ; #####
00660 ;
7036 C5        00670      BITEIX   PUSH    BC          ; SAVE NUMBER OF BITS
7037 E601      00680      AND      1          ; MASK OUT ALL BUT BIT 0
7039 47        00690      LO      B,A        ; SAVE A IN B REGISTER
703A 3A3040    00700      LO      A,{4030H}   ; GET SCREEN STATUS
7030 E6FC      00710      AND      0FCH      ; MASK OUT LOW TWO BITS
703F B0        00720      AOR      A,B        ; SET BIT OR NOT
7040 03FF      00730      OUT      (OFFH),A    ; SEND OUT CASSETTE PORT
7042 C5        00740      PUSH    BC          ; SAVE B REGISTER AGAIN
7043 0605      00750      LO      B,05M      ; TIMING VALUE FOR BIT
7045 10FE      00760      OJNZ     $          ; DELAY 15 65 USECS
7047 C1        00770      POP      BC          ; GET VALUE BACK INTO B
7048 C800      00780      RLC      B          ; ROTATE INTO BIT 1
704A E6FC      00790      AND      0FCM      ; MASK OUT LOW BITS AGAIN
704C B0        00800      AOR      A,B        ; SET BIT 1 OR NOT
7040 03FF      00810      OUT      (OFFH),A    ; SEND OUT CASSETTE PORT
704F C5        00820      PUSH    BC          ; SAVE B REGISTER AGAIN
7050 0606      00830      LO      B,06H      ; DELAY 15 63 USECS
7052 10FE      00840      OJNZ     $          ; DELAY FOR BOTTOM OF BIT
7054 E6FC      00850      AND      0FCH      ; CREATE A NEUTRAL BIT
7056 03FF      00860      OUT      (OFFH),A    ; AND WRITE IT TO TAPE
705B 060B      00870      LO      B,0BH      ; GET DELAY VALUE IN B
705A 10FE      00880      OJNZ     $          ; DELAY 15 110 USECS
705C C1        00890      POP      BC          ; CLEAR STACK OF BIT
7050 C1        00900      POP      BC          ; GET ORIGINAL BITS BACK
705E C9        00910      RET          ; BACK TO SAVING ROUTINE
00920 ;
00930 ; #####
00940 ; ROUTINE TO WRITE LEADER OF 256 BYTES OF FF PLUS SYNC A5
00950 ; #####
00960 ;
705F AF        00970      SYNCMR   XOR      A          ; DEFINE DRIVE NUMBER 0
7060 C01202    00980      CALL     0212H      ; DRIVE MOTOR RUNNING
7063 010000    00990      LO      8C,0        ; NEED A LOOP ABOUT 1 SEC
7066 C06000    01000      CALL     0600H      ; CALL DELAY IN ROM
7069 C5        01010      5YLOOP   PUSH    BC          ; SAVE THIS VALUE (BC=0)
706A E600      01020      LO      C,0        ; WRITE BYTE OF 1111 1111
706C C00070    01030      CALL     WRITE      ; WRITE THAT BYTE
706F 060B      01040      LO      B,8        ; SHORT LOOP BETW. BYTES
7071 10FE      01050      OJNZ     $          ; DELAY JUST 58 USECS
7073 C1        01060      POP      BC          ; RESTORE VALUE FOR USE
7074 10F3      01070      OJNZ     5YLOOP   ; WRITE TOTAL OF FF BYTES
7076 0EA5      01080      LO      C,0A5H    ; SYNCHRONIZATION BYTE
707B C00070    01090      CALL     WRITE      ; WRITE THAT BYTE
7078 C9        01100      RET          ; BACK TO MAIN ROUTINE
01110 ;
01120 ; #####
01130 ; READ BYTE SUBROUTINE: COMPLETED BYTE INTO C REGISTER
01140 ; #####
01150 ;
707C 0608      01160      READ     LO      B,B        ; NUMBER OF BITS TO READ
707E C0BA70    01170      RLOOP    CALL     REDEX      ; READ TIMING BIT
7081 C09C70    01180      CALL     REOBIT   ; READ DATA BIT IF ANY
7084 17        01190      RLA          ; ROTATE BIT INTO CARRY
7085 C811      01200      RL      C          ; ROTATE CARRY INTO C
7087 10F5      01210      OJNZ     RLOOP    ; DO IT TOTAL OF B TIMES
7089 C9        01220      RET          ; BACK FROM READ ROUTINE
01230 ;
01240 ; #####
01250 ; ROUTINE TO READ TIMING BIT FROM TAPE AND DELAY AFTER
01260 ; #####
01270 ;
708A C0C770    01280      REDEX    CALL     INSIG      ; KEEP LOOKING FOR ONE
7080 08FF      01290      REOX2   IN      A,(OFFH)   ; KEEP LOOKING FOR ONE
708F 17        01300      RLA          ; ROTATE INTO CARRY
7090 30F8      01310      JR      NC,REOX2   ; KEEP LOOKING FOR ONE
7092 C5        01320      PUSH    BC          ; SAVE NUMBER OF BITS
7093 060F      01330      LO      B,0FH      ; VALUE TO DELAY
7095 10FE      01340      OJNZ     $          ; AND DELAY FOR 144 USECS
7097 C1        01350      POP      BC          ; RESTORE NUMBER OF BITS
7098 C0C770    01360      CALL     INSIG      ; RESET INSIG FLIP-FLOP
709B C9        01370      RET          ; BACK TO MAIN ROUTINE
01380 ;
01390 ; #####
01400 ; ROUTINE TO READ INDIVIDUAL DATA BIT (IF ANY) FROM TAPE
01410 ; #####
01420 ;
709C C5        01430      REOBIT   PUSH    BC          ; SAVE VALUE IN B AND C
7090 0618      01440      LO      B,1BH      ; VALUE FOR DELAY
709F 10FE      01450      OJNZ     $          ; EQUAL TO 197 USECS
70A1 08FF      01460      IN      A,(OFFH)   ; GET VALUE FROM CASSETTE
70A3 C1        01470      POP      BC          ; RESTORE VALUES TO B & C
70A4 C9        01480      RET          ; RETURN WITH VALUE IN A
01490 ;
01500 ; #####
01510 ; THIS ROUTINE READS THROUGH THE LEADER FOR SYNC BYTE
01520 ; #####
01530 ;
70A5 AF        01540      REOSYN   XOR      A          ; DEFINE DRIVE NUMBER 0
70A6 C01202    01550      CALL     0212H      ; DRIVE IS RUNNING
70A9 C0BA70    01560      CALL     REDEX      ; READ TIMING BIT
70AC C09C70    01570      CALL     REOBIT   ; GET BIT FROM TAPE
70AF 17        01580      RLA          ; ROTATE INTO CARRY FLAG
70B0 30F3      01590      JR      NC,REOSYN   ; KEEP LOOKING FOR A 1

```

possible, but reliable, especially since the TRS-80's 500-baud rate seems to be full of flaws? As noted in the Supplement to Chapter 6 (?), the 500-baud rate is actually fairly reliable, except that a few misconceptions on the part of the ROM designers led to incorrect load timing. The writers of the high speed loaders noted above took considerably more care in designing their input/output schemes, and thus achieved that reliability.

Listing 9-2 presents a pair of high-speed input/output modules that load and save blocks of memory at 2000 (?) baud. These modules may supplant the CALLs to 0235 in all the input/output routines presented elsewhere in this book, allowing faster saving and loading of memory blocks, data, screens, etc. These routines use the normal cassette output. For a high-speed loader using an 8-track system and digital, very reliable recording, see Chapter 9.

A Paper Tape Reader

One of the more common sights in the early days of larger computers was a bank of spinning tape reels for magnetic and punched tapes. Both still exist, but are not common storage modes for microcomputers like the TRS-80.

However, you may have the chance to pick up some terrific programs written for an 8080-based computer, especially those in the National Semiconductor library. These include double-precision mathematical subroutines, text handlers, light-pen readers, etc., and are sometimes available at low prices. But they are stored on rolls of paper tape. Software for devices such as the Computalker Speech Lab is also provided on paper tape.

Furthermore, though paper tape may no longer be the popular program storage medium it once was, for archival storage or communicating among different styles and types of computers, it still has its place. For occasional use, then, paper tape can be used, but an expensive reader won't be a very good investment. For less than \$50, you can interface the TPR-1 paper tape reader and the TRS-80.

The TPR-1 reader is sold by *Raeco*, Box 165, Washington, Maine 04574. The unit consists of a machined, brushed aluminum track for the paper tape, and a circuit board attached to the track. On the board are two integrated circuits, and nine light sensors are on the board under holes in the track; also provided are an LED test light, resistors, and a 14-pin DIP socket. It is sold with a good technical manual for \$32.50; an optional

```

7082 C811      01600      RL      C      ; ROTATE INTO C REGISTER
7084 0607      01610      LD      8,7    ; NUHGER OF 8ITS LEFT
7086 C08A70    01620  SYNCLP CALL  REOEX  ; READ TIMING BIT
7088 C09C70    01630      CALL  REO8IT  ; GET BIT FROM TAPE
708C 17        01640      RLA      ; ROTATE INTO CARRY FLAG
708D C811      01650      RL      C      ; ROTATE INTO C REGISTER
708F 10F5      01660      QJNZ     SYNCLP ; 00 IT 7 TIMES LEFT
70C1 3EA5      01670      LD      A,0A5H ; LOAD SYNC BYTE VALUE
70C3 89        01680      CP      C      ; COMPARE AGAINST C
70C4 C8        01690      RET      Z      ; RETURN IF A MATCH
70C5 180E      01700      JR      REO5YN ; BACK TO KEEP LOOKING
              01710 ;
              01720 ; #####
              01730 ; ROUTINE TO RESET INSIG FLIP-FLOP TO BE ABLE TO READ
              01740 ; #####
              01750 ;
70C7 F5        01760  INSIG  PUSH  AF      ; SAVE VALUE IN AF
70C8 3A3040     01770      LD      A,(4030H) ; GET SCREEN STATUS BYTE
70C8 03FF      01780      OUT     (0FFH),A ; RESET FLIP-FLOP
70C0 F1        01790      POP     AF      ; RESTORE AF STATUS
70CE C9        01800      RET      ; BACK TO READING ROUTINE
              01810 ;
              01820 ; #####
              01830 ; THIS ROUTINE WILL LOAD A BASIC PROGRAM WITH NO FILENAME
              01840 ; #####
              01850 ;
70CF F3        01860  TEMPER  OI      ; GET RID OF BOTHERS
7000 2A8140     01870      LD      HL,(4081H) ; GET TOP OF BASIC MEMORY
7003 E058A440  01880      LD      DE,(40A4H) ; GET BASIC PROGRAM PTR.
7007 05        01890      PUSH    DE      ; SAVE IT A HOHEIT
7008 E052      01900      SBC     HL,DE   ; GET MEHORY AVAILABLE
700A E5        01910      PUSM    HL      ; GET READY FOR TRANSFER
700B C1        01920      POP     8C      ; COUNT OF MEHORY IS IN B
700C E1        01930      POP     HL      ; BEGINNING OF BASIC = HL
700D C5        01940      PUSH    BC      ; SAVE AVAILABLE MEHORY
700E C0A570    01950      CALL    REO5YN ; CALL READ SYNC ROUTINE
70E1 C1        01960      POP     8C      ; RESTORE AVAILABLE HEM.
70E2 C5        01970  TREAD0  PUSH    BC ; SAVE BYTE COUNT
70E3 C07C70    01980      CALL    READ0 ; READ ONE BYTE
70E6 71        01990      LD      (HL),C ; GET VALUE TO VIDEO MEH
70E7 C1        02000      POP     8C      ; GET BYTE COUNT BACK
70E8 08        02010      DEC     8C      ; REDUCE COUNT BY ONE
70E9 78        02020      LD      A,8     ; GET HIGH BYTE OF MEMORY
70EA 81        02030      OR      C      ; AND CHECK AGAINST LOW
70EB CA9719    02040      JP      Z,1997H ; OR ERROR IF TOO MUCH
70EE 7E        02050      LD      A,(ML)  ; GET VALUE IN HL
70EF A7        02060      AND     A      ; TEST IF A ZERO
70F0 C20271    02070      JP      NZ,JUMP3 ; PAST FLASH IF OKAY
70F3 3A3F3C    02080      LD      A,(3C3FH) ; GET PLACE ON SCREEN
70F6 E0A       02090      XOR     0AH    ; AND TOGGLE STAR & SPACE
70F8 323F3C    02100      LD      (3C3FH),A ; AND PUT BACK ON SCREEN
70F8 28        02110      DEC     HL      ; GO BACK SPACE FOR TEST
70FC 7E        02120      LD      A,(HL)  ; GET VALUE THERE
70F0 A7        02130      AND     A      ; TEST IF A ZERO ALSO
70FE CA0571    02140      JP      Z,JUMP4 ; GO TO END ROUTINE IF 0
7101 23        02150      INC     HL      ; BACK TO PROPER BYTE
7102 23        02160  JUHP3   INC     HL ; READY NEXT MEH LOC'N
7103 1800      02170      JR      TREAD0 ; AND THEN GO BACK
7105 23        02180  JUHP4   INC     HL ; GET NEXT MEMORY LOC'N
7106 AF        02190      XOR     A      ; LET A BE EQUAL TO ZERO
7107 77        02200      LD      (HL),A ; AND PUT IT IN PLACE
710B E05BA440  02210      LD      DE,(40A4H) ; GET START OF PRGM PTR
710C 3EFF      02220      LD      A,0FFH ; GET RESETING CODE
710E 12        02230      LD      (0E),A ; PUT AT PROGRAM START
710F C0FC1A    02240      CALL    1AFCH ; RESET ALL LINE NUMBERS
7112 23        02250      INC     HL      ; HL MOVED PAST PROGRAM
7113 22F940    02260      LD      (40F9H),HL ; SIMPLE VARIABLE POINTER
7116 C0C901    02270      CALL    01C9H ; CLEAR THE SCREEN NOW
7119 C06118    02280      CALL    1861H ; CLEAR ALL THE POINTERS
711C C0FE01    02290      CALL    01FEH ; TURN CASSETTE OFF
711F C3CC06    02300      JP      06CCH ; GO TO BASIC "READY"
              02310 ;
              02320 ; #####
              02330 ; THIS ROUTINE IS A GERMINAL ROUTINE TO CSAVE A PROGRAM
              02340 ; WITHOUT A PROGRAM NAME.  FORMAT:  /PUT.  THE 500-8A00
              02350 ; BLOCK SAVE CAN BE USED AS AN EXAPLE OF HOW TO EMPLOY
              02360 ; A PROGRAM NAME IN SAVING A PROGRAM.
              02370 ; #####
              02380 ;
7122 F3        02390  TEMPEX  OI      ; GET RID OF BOTHER
7123 3E2A      02400      LD      A,2AH  ; GET READY AN ASTERISK
7125 323E3C    02410      LD      (3C3EH),A ; PLACE STAR ON SCREEN
712B 323F3C    02420      LD      (3C3FH),A ; PLACE STAR NEXT TO IT
712B 2AA440    02430      LD      HL,(40A4H) ; START OF BASIC PROGRAM
712E E05BF940  02440      LD      DE,(40F9H) ; BOTTOM OF VAR. POINTER
7132 C08F70    02450      CALL    SYNCMR ; WRITE LEADER AND SYNC
7135 7C        02460  L0000P  LD      A,M  ; GET CURRENT MIGH MSB
7136 BA        02470      CP      0      ; SAME AS TARGET MSB?
7137 C23F71    02480      JP      NZ,JUMP01 ; CONTINUE IF NOT SAME
713A 70        02490      LD      A,L    ; SAME HI MSB - READY LOW
713B 8B        02500      CP      E      ; SAME AS TARGET LSB?
713C CA5671    02510      JP      Z,GOOUT ; OONE WITH SAVE IF SO
713F 7E        02520  JUMP01  LD      A,(HL) ; ELSE GET VALUE IN MEH
7140 A7        02530      AND     A      ; TEST IF A ZERO
7141 C24E71    02540      JP      NZ,JUMP02 ; JUST GO ON IF NOT
7144 FS        02550      PUSM    AF      ; SAVE VALUE IN A
7145 3A3F3C    02560      LD      A,(3C3FH) ; GET TOGGLE VALUE TO A
714B E0A       02570      XOR     0AH    ; TOGGLE STAR & SPACE
714A 323F3C    02580      LD      (3C3FH),A ; AND PUT IT ON SCREEN
7140 F1        02590      POP     AF      ; GET VALUE BACK TO A
714E 4F        02600  JUMP02  LD      C,A  ; PUT BYTE IN C REGISTER
714F C00070    02610      CALL    WRITE ; WRITE BYTE TO TAPE

```

case is \$5.00. Photo 1 shows the unit mounted inside the smallest Radio Shack equipment box.

Eight-level (eight bit) paper tape is capable of storing parallel bytes of data by means of holes punched in the tape. A smaller, ninth hole – placed between the third and fourth holes – provides a timing signal for the reading program.

The ninth hole also can be used as a data-ready signal. By the time the light just triggers the circuitry as it passes along the edge of the smaller hole, the larger holes are letting in plenty of light for the data to be stable, ready to read.

The TPR-1 comes ready to hook to a computer bus. Its output is in parallel, and all signals are tri-state. Because it uses only 12 mA, it's possible to run the reader directly from the TRS power supply.

Figure 9-4 presents the diagram of the TPR-1. The low-power CMOS integrated circuits U1 and U2 evaluate the state of the data as seen by the light-sensitive transistors and provide a parallel output. Part of U2 is also used to drive the LED, which blinks on whenever data is stable at the output of the reader.

Figure 9-5 is the TRS-80 interface schematic. Z1 and Z2 decode the port address 3F in order to activate tri-state buffer Z3. This separate port decoding is necessary because the TPR-1 was not designed with the READY line separately activated from the data lines. Were that the case, READY might be tested at all times. That way, data would only be input whenever READY indicated stable data. In its present configuration, however, a separate buffer must be used for the TPR-1 data lines.

Z4 is a flip-flop which produces an interrupt signal and sends it to the TRS-80 INT line; INTAK (interrupt acknowledge) is used to clear the interface flip-flop when data has been read. This configuration is similar to that used for the interrupt based real time clock (see Chapter 8).

The circuit can be wire-wrapped on a small piece of perfboard and mounted inside a case with the TPR-1. A detachable 40-pin cable can also be used to save a few dollars.

Listing 9-2 presents the software to read one page (256 bytes) of data into the TRS-80 and store it in memory. Recall that the interrupt patch point at 4012 is initialized with C9, a RETurn instruction. In its place, then, a patch must be made to one of three interrupt service routines which will read each byte of data as it becomes stable at the output of the TPR-1. Since the reader will not likely be a device used very

A Paper Tape Reader

```

7152 23      02620      INC      HL      ; GET NEXT MEMORY LOC'N
7153 C33571  02630      JP        L0000P ; LOOP AROUND FOREVER
7156 C0FE01  02640      CALL     01FEH   ; TURN OFF CASSETTE
7159 C9      02650      RET              ; BACK TO MAIN ROUTINE
02660 ;
02670 ; #####
02680 ; REMEMBER: ENTRY POINT FOR SAVE ROUTINE IS TEMPEX
02690 ; REMEMBER: ENTRY POINT FOR LOAD ROUTINE IS TEMPER
02700 ; #####
02710 ;
02720      END

0000
BITEX 7036 00670 00360
G00UT 7156 02640 02510
INSIG 70C7 01760 01280 01360
JUMP01 713F 02520 02480
JUMP02 714E 02600 02540
JUMP3 7102 02160 02070
JUMP4 7105 02180 02140
L0000P 7135 02460 02630
READ 707C 01160 01980
RE08IT 709C 01430 01180 01570 01630
RE0EX 708A 01280 01170 01560 01620
RE0SYN 70A5 01540 01590 01700 01950
REDX2 7080 01290 01310
RLOOP 707E 01170 01210
SYLOOP 7069 01010 01070
SYNCHR 705F 00970 02450
SYNCLP 7086 01620 01660
TEMPER 70CF 01860
TEMPEX 7122 02390
TIMEX 7014 00440 00330
TREAD 70E2 01970 02170
WLOOP 7008 00330 00370
WRITE 7000 00280 01030 01090 02610

```

NOTES -
1. OUTPUT DATA IS TRUE.
I.E. A "ONE" (HOLE) ON THE
TAPE GIVES A HIGH (5V)
VOLTAGE LEVEL OUTPUT
2. MAXIMUM LOADING ON ANY
OUTPUT IS TWO STANDARD
TTL LOADS.
3. Q1 THRU Q9 - MRDISO
OPTICALLY MATCHED, DO NOT
SUBSTITUTE.

U1, U2 - 4502

FIG 9-4
9/28/80

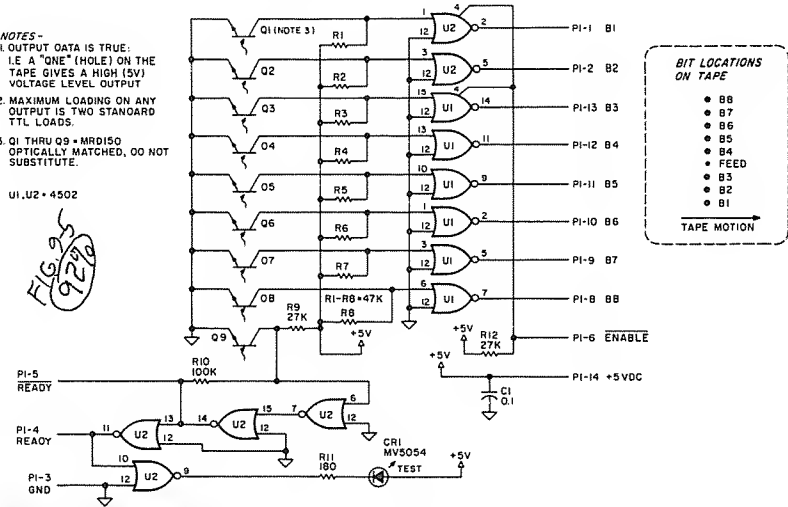


Figure 9-4. TPR-1 schematic diagram.

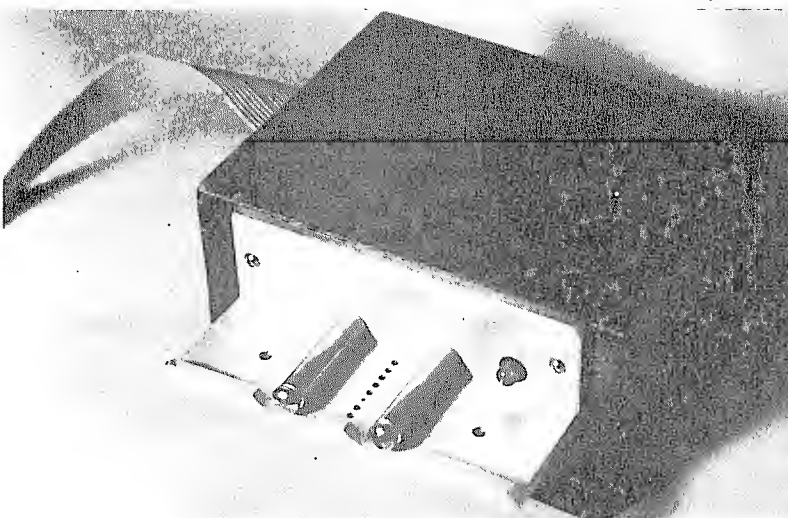
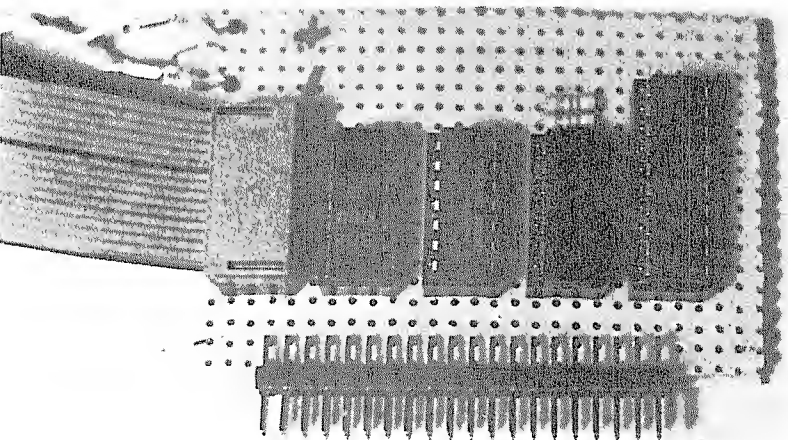


Photo 9-6. Paper tape reader.

TPR-1 tape reader has machined aluminum track and data indicator LED.



Only four integrated circuits form the complete tape reader interface. Power supplies both reader and interface.

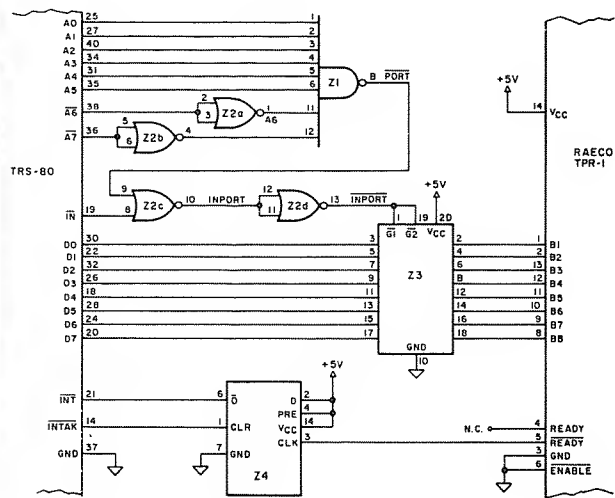


Figure 9-5. TPR-1 interfacing schematic.

The program is entered at line 1160. The screen is cleared, and the user is prompted to enter a base address in hex. This is the address starting at which the tape data is to be loaded into memory. The keyboard is scanned for characters 0 to 9 and A to F; these are displayed, and when ENTER is pressed, the characters are converted to a starting address.

The tape must be threaded before actual data reading is begun, because during threading it's possible to present false information to the TRS-80. The tape reading is begun at line 1740.

Listing 9-3. Paper tape reader program.

```

00100 ; #####
00110 ; TAPE READER AND 256-BYTE LOADER ROUTINE
00120 ; DENNIS BATHORY KITSZ, ROXBURY, VERMONT 05669
00130 ; #####
00140 ;
4013 00150 VECTOR EQU 4013H ;INTERRUPT VECTOR PATCH
06CC 00160 BASIC EQU 06CCH ;RETURN TO BASIC READY
3C00 00170 VIDEO EQU 3C00H ;BEGINNING OF SCREEN
00180 ;
00190 ; #####
00200 ; INTERRUPT VECTOR AT 4012H
00210 ; #####
00220 ;
4012 00230 ; ORG 4012H ;CURRENT VALUE IS C9
4012 C3 00240 0EFB 0C3H ;REPLACE WITH JUMP
7000 00250 ORG 07000H ;ROUTINES BEGIN HERE
00260 ;
00270 ; #####
00280 ; CLEAR SCREEN SUBROUTINE
00290 ; #####
00300 ;
7000 21003C 00310 CLEAR LO HL,VIDEO ;GET START OF VIDEO
7003 11013C 00320 LO 0E,VIDEO+1 ;GET DESTINATION POINT
7006 01FF03 00330 LO BC,03FFH ;GET MEMORY BLOCK SIZE
7009 3620 00340 LO (HL),20H ;WRITE SPACE INTO VIDEO
700B E0B0 00350 LDIR ;BLOCK MOVE CLEAR SCREEN
7000 C9 00360 RET ;BACK FROM SUBROUTINE
00370 ;
00380 ; #####
00390 ; SCAN FOR ENTER SUBROUTINE
00400 ; #####
00410 ;
700E 3A4038 00420 ENTER LO A,(3B40H) ;"ENTER" KEYBOARD ROW
7011 FE02 00430 CP 2 ;"ENTER" KEYBOARD COLUMN
7013 20F9 00440 JR NZ,ENTER ;LOOP UNTIL KEY PRESSED
7015 C9 00450 RET ;BACK FROM SUBROUTINE
00460 ;
00470 ; #####
00480 ; DISPLAY MESSAGE SUBROUTINE
00490 ; #####
00500 ;
7016 7E 00510 DISPLY LO A,(HL) ;GET BEGINNING OF TEXT
7017 A7 00520 AND A ;CHECK IF A NULL
7018 C8 00530 RET Z ;EXIT SUBROUTINE IF NULL
7019 12 00540 LD (DE),A ;DISPLAY CHARACTER IN A
701A 23 00550 INC HL ;GET NEXT MESSAGE LOC'N
701B 13 00560 INC DE ;GET NEXT SCREEN LOC'N
701C 18F8 00570 JR DISPLY ;LOOP FOR CHARACTER TEST
00580 ;
00590 ; #####
00600 ; CONVERT TO ASCII SUBROUTINE
00610 ; #####
00620 ;
701E F5 00630 CONVRT PUSH AF ;SAVE ACCUM. AND FLAGS
701F E6F0 00640 AND 0F0H ;MASK OUT LOW 4 BITS
7021 1F 00650 RRA ;MOVE NIBBLE TO RIGHT...
7022 1F 00660 RRA ; ...SOME MORE
7023 1F 00670 RRA ; ...SOME MORE
7024 1F 00680 RRA ; ...UNTIL ONE
7025 FE0A 00690 CP 0AH ;IS IT TEN OR GREATER?
7027 3004 00700 JR NC,HIBYTE ;MOVE ALONG IF > TEN
7029 C630 00710 ADO A,30H ;ASCII = NUMBER PLUS 30H
702B 1802 00720 JR NEXT ;GO ON TO LOW NIBBLE
7020 C637 00730 HIBYTE ADO A,37H ;ASCII = NUMBER PLUS 37H
702F 77 00740 NEXT LO (HL),A ;DISPLAY FIRST ASCII CHAR
7030 23 00750 INC HL ;GET NEXT SCREEN LOC'N
7031 F1 00760 POP AF ;RESTORE ORIGINAL HEX
7032 E60F 00770 AND 0FH ;MASK OUT HIGH 4 BITS
7034 FE0A 00780 CP 0AH ;IS IT TEN OR GREATER?
7036 3004 00790 JR NC,HIBTE2 ;MOVE ALONG IF > TEN
703B C630 00800 ADO A,30H ;ASCII = NUMBER PLUS 30H
703A 1802 00810 JR NEXT2 ;GO TO DISPLAY & OUT
703C C637 00820 HIBTE2 ADO A,37H ;ASCII = NUMBER PLUS 37H
703E 77 00830 NEXT2 LO (HL),A ;DISPLAY NEXT ASCII VALUE
703F C9 00840 RET ;BACK FROM SUBROUTINE
00850 ;
00860 ; #####
00870 ; CLEAR TAPE READER INTERRUPT ACKNOWLEDGE
00880 ; #####
00890 ;
7040 F3 00900 SERVED OI ;INTERRUPT OFF IN SERVICE
7041 AF 00910 XOR A ;CLEAR ACCUM. & FLAGS
7042 C9 00920 RET ;BACK FROM SUBROUTINE
00930 ;
00940 ; #####
00950 ; PAGE ADDRESS INTERRUPT SERVICE
00960 ; #####
00970 ;
7043 F3 00980 SERVE1 OI ;INTERRUPT OFF IN SERVICE
7044 0B3F 00990 IN A,(3FH) ;GET VALUE FROM READER
7046 C01E70 01000 CALL CONVRT ;CONVERT VALUE TO ASCII
7049 AF 01010 XOR A ;CLEAR ACCUM. & FLAGS
704A C9 01020 RET ;BACK FROM SUBROUTINE
01030 ;
01040 ; #####
01050 ; READ DATA / PLACE ON SCREEN INTERRUPT
01060 ; #####
01070 ;
704B F3 01080 SERVE2 OI ;INTERRUPT OFF IN SERVICE

```

Listing Continued . .

With the software shown, the tape to be read must be in the following format:

1-byte code of information (tape number, address page, etc.), which is displayed but not to be stored in memory.

256 bytes of data.

1-byte simple checksum.

If the tape is not in this format, the program can be easily altered to accommodate any other 256-byte data block format.

Interrupts are then enabled (lines 1810-1820), and a series of short interrupt service routines are activated. The first routine merely waits for the interrupt line to clear, as it may have been set by stray light in the room when the tape is threaded (lines 1820-1870). 256 bytes are then loaded and displayed (lines 1910-2030). The checksum is calculated and displayed (lines 2050-2130), and the checksum is read from tape and displayed (lines 2150-2280). If there is a match, the memory pointer is advanced in order to read the next block of tape; otherwise, it is reset to the beginning of the block, allowing the tape to be read again. (lines 2210-2420). Finally, the option of loading additional blocks or returning to BASIC is presented (lines 2440-2540).

Using the TPR-1, the interface, and this simple software, the wealth of 8080 programs, as well as programs saved in an archival paper tape format, may be read into your TRS-80 and used.

An 8-Track Mass Storage System

Oh, no! Here comes another one! I'd like to join the mass storage fray with another device capable of loading and saving programs at high speed. It's not as slow as a cassette, not as fast as a Stringy-Floppy, but it has one interesting capability: sequential-random access. That's a mythical term for sequential access of more than one track at a time.

Here's how it works: 8-track cartridges play one-quarter their total length on each pass. Then the head switches from the first stereo pair to the second, the second to the third, the third to the fourth, and from the fourth back up to the first. A so-called '40-minute' cartridge is actually 10 minutes long, four passes. The shortest commercially available cartridges are 20 minutes long, five minutes per pass.

An 8-Track Mass Storage System

Continued Listing

```

704C 0B3F 01090 IN A,[3FH] ;GET VALUE FROM READER
704E 77 01100 LO [HL],A ;PUT IT INTO MEMORY
704F B1 01110 AOO A,C ;GET VALUE FROM CHECKSUM
7050 4F 01120 LO C,A ;RESTORE UPDATED CHECKSUM
7051 23 01130 INC HL ;GET NEXT MEMORY LOC'N
7052 AF 01140 XOR A ;CLEAR ACCUM. & FLAGS
7053 C9 01150 RET ;BACK FROM SUBROUTINE
01160 ;
01170 ; #####
01180 ; CHECKSUM INTERRUPT ROUTINE
01190 ; #####
01200 ;
7054 F3 01210 SERVE3 OI ;INTERRUPT OFF IN SERVICE
7055 0B3F 01220 IN A,[3FH] ;GET VALUE FROM READER
7057 47 01230 LO B,A ;SAVE IT IN B REGISTER
7058 AF 01240 XOR A ;CLEAR ACCUM. & FLAGS
7059 C9 01250 RET ;BACK FROM SUBROUTINE
01260 ;
01270 ; #####
01280 ; MESSAGES FOLLOW
01290 ; #####
01300 ;
705A 54 01310 MSGN01 DEFM 'THREAD TAPE AND PRESS CLEAR.'
7076 00 01320 DEFB 00
7077 4C 01330 MSGN02 DEFM 'LOADING PAGE ADDRESS: '
70BE 00 01340 DEFB 00
70BF 42 01350 MSGN03 DEFM 'BYTES LOADING AS FOLLOWS:'
70AB 00 01360 DEFB 00
70A9 43 01370 MSGN04 DEFM 'CALCULATED CHECKSUM IS: '
70C2 00 01380 DEFB 00
70C3 43 01390 MSGN05 DEFM 'CHECKSUM AS READ IS: '
7009 00 01400 DEFB 00
700A 43 01410 MSGN06 DEFM 'CHECKSUM ERROR IN THIS BLOCK.'
70F7 00 01420 DEFB 00
70FB 42 01430 MSGN07 DEFM 'BLOCK LOADED CORRECTLY.'
7E0F 00 01440 DEFB 00
7E10 41 01450 MSGN08 DEFM 'ANOTHER BLOCK? REPLY 1 FOR YES, 2 FOR NO'
7E39 00 01460 DEFB 00
7E3A 50 01470 MSGN09 DEFM 'PRESS CLEAR TO RETURN TO BASIC.'
7E59 00 01480 DEFB 00
01490 ;
01500 ; #####
01510 ; REMEMBER THIS IS ENTRY POINT AND NOT!
01520 ; BEGINNING OF PROGRAM.....
01530 ; CLEAR SCREEN, DISPLAY "THREAD" MESSAGE
01540 ; #####
01550 ;
7E5A C00070 01560 START CALL CLEAR ;OUT TO CLEAR SUBROUTINE
7E50 215A70 01570 LO HL,MSGN01 ;GET MESSAGE #1 LOCATION
7E60 11003C 01580 LO OE,VIDEO ;GET DISPLAY LOCATION
7E63 C01670 01590 CALL DISPLY ;OUT TO DISPLAY SUBROUT.
7E66 C00E70 01600 CALL ENTER ;WAIT FOR ENTER SUBROUT.
01610 ;
01620 ; #####
01630 ; DISPLAY "ADDRESS" MESSAGE & FIND IT
01640 ; #####
01650 ;
7E69 217770 01660 LO HL,MSGN02 ;GET MESSAGE #2 LOCATION
7E6C 11403C 01670 LO OE,VIDEO+40H ;GET DISPLAY LOCATION
7E6F C01670 01680 CALL DISPLY ;OUT TO DISPLAY SUBROUT.
7E72 214070 01690 LO HL,SERVE0 ;GET INT #1 SERVICE ROUT.
7E75 221340 01700 LO [VECTOR],HL ;INSTALL AT INT. VECTOR
7E7B 37 01710 SCF ;CARRY FLAG IS IMPORTANT
7E79 E056 01720 IM 1 ;SET INTERRUPT MODE
7E7B F8 01730 EI ;INTERRUPTS ON & WAITING
7E7C 3BFE 01740 JR C,$ ;SUBROUTINE CLEARS CARRY!
7E7E 214370 01750 LO HL,SERVE1 ;GET INT #2 SERVICE ROUT.
7E81 221340 01760 LO [VECTOR],HL ;INSTALL AT INT. VECTOR
7E84 21573C 01770 LO HL,VIDEO+S7H ;GET DISPLAY LOCATION
7E87 37 01780 SCF ;CARRY DETERMINES LOOP
7E88 F8 01790 EI ;INTERRUPTS ON & WAITING
7E89 3BFE 01800 JR C,$ ;SUBROUTINE CLEARS CARRY!
01810 ;
01820 ; #####
01830 ; DISPLAY "BYTES" MESSAGE & LOAD 256
01840 ; #####
01850 ;
7EBB 21BF70 01860 LO HL,MSGN03 ;GET MESSAGE #3 LOCATION
7E8E 11803C 01870 LO OE,VIDEO+BDH ;GET DISPLAY LOCATION
7E91 C01670 01880 CALL DISPLY ;OUT TO DISPLAY SUBROUT.
7E94 214B70 01890 LO HL,SERVE2 ;GET INT #3 SERVICE ROUT.
7E97 221340 01900 LO [VECTOR],HL ;INSTALL AT INT. VECTOR
7E9A 210030 01910 LO HL,VIDEO+100H ;GET FIRST DISPLAY LOC'N
7E90 AF 01920 XOR A ;CLEAR ACCUM. & FLAGS
7E9E 4F 01930 LO C,A ;CLEAR CHECKSUM REGISTER
7E9F 0600 01940 LO B,00H ;LOAD B REGISTER WITH 256
7EA1 37 01950 LOOP2 SCF ;INSTALL "JR C" LOOP
7EA2 FB 01960 EI ;INTERRUPTS ON & WAITING
7EA3 3BFE 01970 JR C,$ ;SUBROUTINE CLEARS CARRY!
7EAS 10FA 01980 OJNZ LOOP2 ;WRITE ONE PAGE TO MEMORY
01990 ;
02000 ; #####
02010 ; DISPLAY "CHECKSUM CALC" MESSAGE
02020 ; #####
02030 ;
7EA7 215470 02040 LO HL,SERVE3 ;GET INT #4 SERVICE ROUT.
7EAA 221340 02050 LO [VECTOR],HL ;INSTALL AT INT. VECTOR
7EAD 21A970 02060 LO HL,MSGN04 ;GET MESSAGE #4 LOCATION
7E90 11403E 02070 LO OE,VIDEO+240H ;GET DISPLAY LOCATION
7E83 C01670 02080 CALL DISPLY ;OUT TO DISPLAY SUBROUT.
7E86 79 02090 LO A,C ;GET CHECKSUM CALCULATION
7E87 05 02100 PUSH OE ;SAVE DISPLAY INFORMATION

```

Listing Continued . . .

For this system, the shortest 8-track cartridges are used in an 8-track deck with an *electrical* fast-forward mode. There are several loading options:

1. Load the next program on the tape from the current track. The machine fast-forwards to the next leader and loads the program.
2. Load the next program on the tape, with the track specified. The machine moves to the specified track, fast-forwards to the next leader, and loads the program.
3. Load the program specified from the track specified. The machine moves to the appropriate track and reads leaders (in fast-forward mode) until the program is found, and then loads it.
4. Load the program specified; the machine moves ahead and reads the directory immediately following the splice. The program is located and read.

In this way, where the locations of programs are known, they may be loaded immediately. Otherwise, the device is *somewhat* directory organized. I add this reservation because the tape is sequential and programs can't be killed easily unless the tape is re-organized. More on that later.

The advantage of this system is obvious: it provides somewhat faster access and loading than cassettes, and allows fairly fast search and storage. In the fast-forward mode, 20-minute 8-track tapes can be run through completely in less than two minutes. Worst-case program access is then two minutes – when you have just passed the program you want to load. Furthermore, eight individual programs can be stored parallel to each other on the cartridge's tracks.

As noted, to build this device, an electrical fast-forward is necessary. Check the manuals for a two-speed motor; the Craig model H240 playback-only deck is the kind I used. Some modifications are necessary to the tape recorder itself, and alignment is a bit more critical.

Continued Listing

```

7E88 E1      02110      POP      HL      ;TRANSFER IT TO HL PAIR
7E89 C01E70  02120      CALL     CONVRT ;CNVRT. CHECKSUH TO ASCII
02130      ;
02140      ; #####
02150      ; DISPLAY "CHECKSUH REAO" HESSAGE & 00 IT
02160      ; #####
02170      ;
7E8C 21C370  02180      LD      HL,HSGN05 ;GET MESSAGE #S LOCATION
7EBF 11803E  02190      LD      0E,VIDEO+280H ;GET DISPLAY LOCATION
7EC2 C01670  02200      CALL     DISPLY ;OUT TO DISPLAY SUBROUT.
7EC5 F8      02210      EI      ;INTERRUPTS ON & WAITING
7EC6 37      02220      SCF      ;CARRY FLAG LOOP SET
7EC7 3BFE    02230      JR      C,$ ;SUBROUTINE CLEARS CARRY
02240      ;
02250      ; #####
02260      ; DISPLAY CHECKSUH AND CHECK IT
02270      ; #####
02280      ;
7EC9 7B      02290      LD      A,B ;GET REAO CHECKSUM BACK
7ECA 05      02300      PUSH     0E ;STASH 0E REGISTER PAIR
7ECB E1      02310      POP      HL ;TRANSFER TO HL PAIR
7ECC C01E70  02320      CALL     CONVRT ;CNVRT. CHECKSUH TO ASCII
7ECF 78      02330      LD      A,B ;GET VALUE AGAIN FROM B
7E00 89      02340      CP      C ;CHECK WITH CALC CHECKSUH
7E01 2808    02350      JR      Z,CKSHOK ;CHECKSUH OKAY IF A HATCH
02360      ;
02370      ; #####
02380      ; DISPLAY CHECKSUH 8AO HESSAGE
02390      ; #####
02400      ;
7E03 210A70  02410      LD      HL,HSGN06 ;GET MESSAGE #6 LOCATION
7E06 11C03E  02420      LD      0E,VIDEO+2C0H ;GET DISPLAY LOCATION
7E09 C01670  02430      CALL     DISPLY ;OUT TO DISPLAY SUBROUT.
7E0C 1809    02440      JR      LEAVE ;LOAD COMPLETE - GO OUT
02450      ;
02460      ; #####
02470      ; DISPLAY CHECKSUH OKAY HESSAGE
02480      ; #####
02490      ;
7E0E 21F870  02500      CKSHOK LD      HL,HSGN07 ;GET MESSAGE #7 LOCATION
7EE1 11C03E  02510      LD      0E,VIDEO+2C0H ;GET DISPLAY LOCATION
7EE4 C01670  02520      CALL     DISPLY ;OUT TO DISPLAY SUBROUT.
7EE7 21107E  02530      LEAVE LD      HL,HSGN08 ;GET MESSAGE #8 LOCATION
7EEA 11003F  02540      LD      0E,VIDEO+300H ;GET DISPLAY LOCATION
7EE0 C01670  02550      CALL     DISPLY ;OUT TO DISPLAY SUBROUT.
02560      ;
02570      ; #####
02580      ; SCAN KEYBOARD FOR 1 OR 0 & 00 IT
02590      ; #####
02600      ;
7EF0 3A103B  02610      FINDYN LD      A,(3810H) ;GET 0-7 KEYBOARD ROW
7EF3 FE02    02620      CP      2 ;IS IT NUHBER ONE?
7EF5 CASA7E  02630      JP      Z,START ;BACK TO START IF SO
7EF8 FE04    02640      CP      4 ;IS IT NUHBER TWO?
7EFA 2802    02650      JR      Z,00NE ;FINISHED ROUTINE IF SO
7EFC 1BF2    02660      JR      FINDYN ;KEEP LOOKING IF NEITHER
7EFE 213A7E  02670      00NE LD      HL,MSGN09 ;GET MESSAGE #9 LOCATION
7F01 11403F  02680      LD      0E,VIDEO+340H ;GET DISPLAY LOCATION
7F04 C01670  02690      CALL     DISPLY ;OUT TO DISPLAY SUBROUT.
7F07 C0E70  02700      LOOKNG CALL     ENTER ;LOOK FOR ENTER SUBROUT.
7F0A C3CC06  02710      JP      BASIC ;BACK TO BASIC READY
02720      ;
02730      ; #####
7ESA 02740      END      START ;SYSTEM ENTRY POINT
00000 TOTAL ERRORS

```

Eviscerating an 8-Track Cartridge

Chances are that obtaining high quality 8-track cartridges in short lengths won't be easy. It's not hard to make your own in roughly eight minute lengths – two minutes per track, with less than 60 seconds total access time. This will be enough for programs nearly 30,000 bytes in length.

Purchase a few cheap 8-tracks to experiment with. All you will need is a piece of wood, a pair of scissors, and a package of silver foil sensing tape. This is available at Radio Shack (catalog number 44-1155, under \$2).

Most 8-tracks are fastened together with plastic tabs recessed in holes on top or bottom of the case. There are usually five: one on each back corner, one under the label in the center, one through the capstan, and one to secure the remaining corner. First, insert the cartridge into a player and run it ahead to its splice.

Make sure you are holding the cartridge with the label down. Slip a thin piece of hobby 'half-round' – a round stick split lengthwise down the center – into each hole and push the tab back, splitting the plastic case gently with your hand. The dull end of a small crocheting hook also works well. When all the tabs are released, split the cartridge apart to about one-quarter inch. Now *turn it over*. Gradually lift the cartridge apart, being careful to note the tape path and *exactly* how much slack is present in the loop.

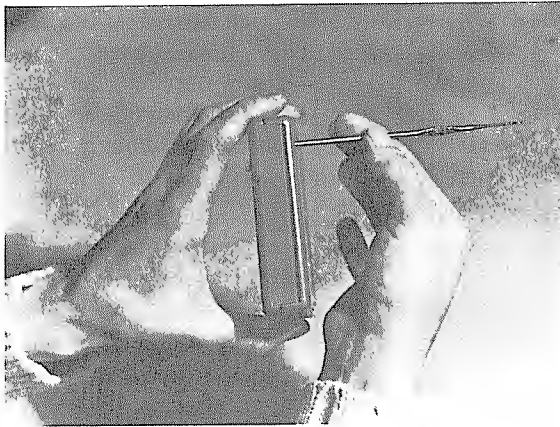
Snip the splice out first clipping out a length of tape that runs to within an inch of the center hub. Take the loose end of the tape that winds around the outside of the tape 'pancake', and begin to pull it out. Let the hub spin freely as you do this, so that the tension on the wound tape does not change. Measure the tape removed (to produce the 37.5 feet for an 8-minute cartridge) using this table:

Total Cartridge Time	Total Feet	Remove to Create 8-minute Cartridge
20 minutes	94	56
22 minutes	103	65
30 minutes	141	103
40 minutes	188	150
44 minutes	206	169
45 minutes	211	173
60 minutes	281	244

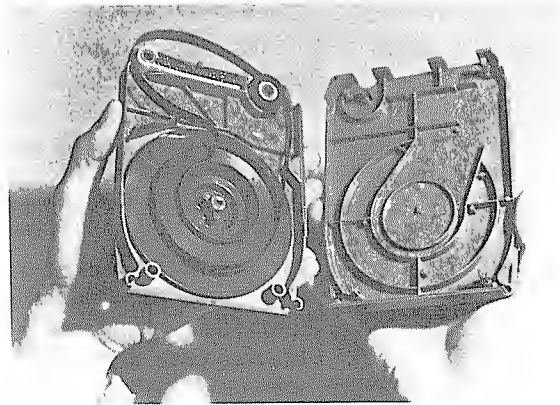
Note: Do not use cartridges over 60 minutes long.

Sound time consuming? It can be if you measure a foot at a time. Instead, anchor the cartridge on a table, and pull the end of the tape across the room (or down the hall). If you've purchased a 40-minute cartridge, that's only ten trips.

Eviscerating an 8-Track Cartridge (5 photos).



Blunt tool is inserted in slots at bottom of the cartridge as the sides are separated and held steady.

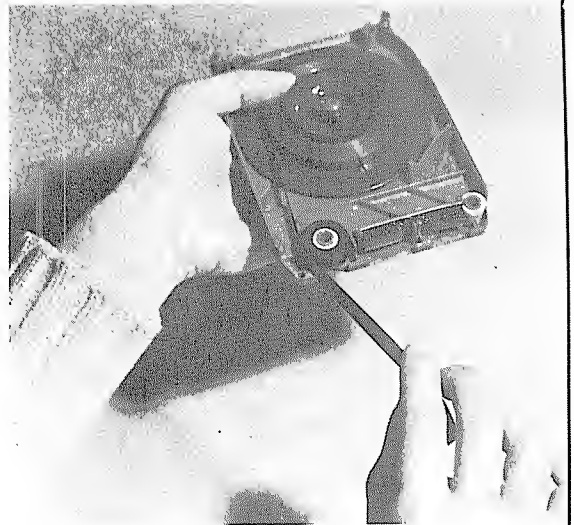


Cartridge is flipped right side up and top of case is removed.

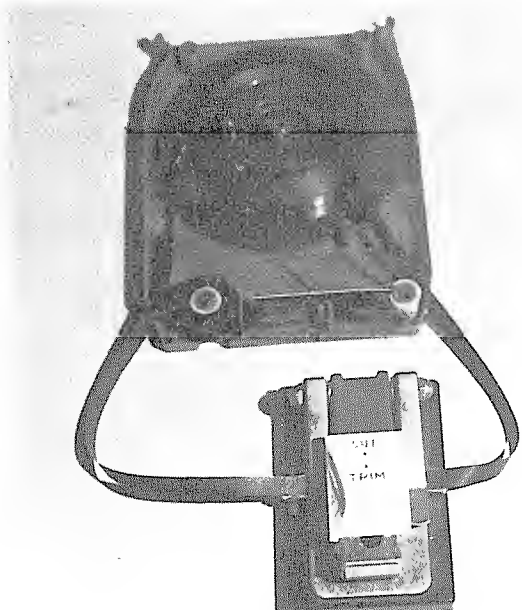


Tape is removed from its path and cut at splice.

When you've removed all the tape you need to, leave the same amount of slack as before and splice the two ends together. The oxide (playing) side of the tape is spliced with metal foil sensing tape. Make sure you splice the same sides to each other! Now reroute the tape along the original path, holding it gently in place; use thin cloth gloves if your hands tend to perspire.



Excess tape is pulled gently from outside of the hub, letting hub and platter rotate freely. Note that end near center of hub has been cut fairly short to eliminate the possibility of tangling with the tape being removed.



When enough tape has been removed, ends are spliced together with silver sensing foil. The tape is then re-routed and the case snapped back together.

Without bending the tape or getting it caught, slip the cartridge top back in place. Note the front of the cartridge as you do this, making sure the tape does not slip on the outside of the 'window frames'. You now have an 8-minute cartridge, ready to use.

If you end up with some slack tape hanging out of the cartridge, put on some thin cloth gloves. Now pull the tape out of the case gently from the end opposite the rubber puck. The tape will pull into the other end faster than it is being pulled out. When most of the slack is taken up, give a gentle tug with a bit of 'snap'; the momentum will spool the tape into the cartridge housing.

If you are using the Craig H-240, the directions below will apply directly; since 8-track playback decks are similar, you can put some of them to use in remodeling your 8-track. It is also possible to purchase drives without electronics from surplus houses. These are sold for under \$10, and come complete with head assembly, motor, and track-change solenoid. Be sure that you get one with a two-speed motor if you want to implement the fast forward features of the system; an excellent single-speed drive is sold for only \$8 by BNF Enterprises.

First of all, you should know that the deck won't be usable as an audio deck when you're done; so don't split up your hi-fi system in hopes that this 8-track storage system will serve double duty.

1. Remove the case. Two screws hold the back panel, and screws through the four feet keep the frame attached to the housing. Set the back aside, and slide the electronics out through the front of the case.

2. Pull off the buttons, and set them aside. Remove the four screws which hold the frame to the face plate, and set the face plate aside.

3. Heat the soldering iron. You will be removing these wires:

- The two wires connected to the on-off switch; contact is made when a cartridge is inserted in the deck.

- The two wires from the foil sensor pickup. This is located about an inch from the on-off switch.

- All wires running to the three switches underneath the deck.

4. Remove the two front screws holding the three switches onto the front plate. Set this switch block aside; also remove all loose wires (those desoldered at both ends), and desolder the far ends of remaining wires which had run to the switches.

5. Unscrew the electronics control board (two screws), which is found to the front of the transformer. Desolder all wires leading to this board, and discard the board, scavenge it for parts, or keep it. It is a legitimate 8-track preamplifier, and can still be used if you need such an animal.

6. The following parts are still intact:

- The motor and the three wires leading from it. These wires are still attached to a terminal strip.

- The head assembly. This will be modified later. At present, it contains a shielded, three-wire cable leading from the playback head itself, and a five-wire assembly from the track-select switch.

- The capstan, drive belt, and track-change mechanics. These remain intact.

- The transformer. Three wires run from it; the center tap (black wire), won't be used, so cover it with tape or a wire nut.

- The terminal strip and two audio output jacks. These will be used.

By removing these parts, you have returned the drive to its 'naked' state. If you are using a surplus 8-track drive, this is the condition in which it will be shipped.

To use this in a digital system, several important conditions have to be met:

1. The recording and playback must be done in a digital format.

2. The track, splice, and tape-in-place status must be readable by the computer.

3. Speed must be controlled by the computer.

Figure 9-6. presents the complete circuitry to convert the Craig H-240 to a digital record/playback system. Incoming data is latched by Z12 on the occurrence of the command OUT (0AAH), A and is buffered by Z2a/b. It is fed to a symmetrical pair of output-coupled buffers (Z1 and Z6).

Since the output of these CMOS buffers is capable of rising very high (within a few millivolts of the 12-volt supply voltage), this provides a fast rising pulse to the recording head. The data is recorded in a bipolar manner: that is, when buffers Z1a/b rise, Z1c/d fall, and vice versa.

This information is recorded directly on the tape. During playback, the raw waveform is fed to Z4a, an LF353 FET operational amplifier (a plug-in replacement for the more commonly available LM747, which can be substituted with some signal degradation). This amplifier is set up in an inverted configuration with high gain; it produces a strong waveform which is then fed into Z4b, configured as a high-gain ‘clipping’ or ‘squaring’ amplifier (contributed by diodes D1 and D2).

This output is stabilized and buffered by comparator Z5a, and fed to Z5b (Z5 is a simple LM 339 comparator), arranged as a TTL-level driver. Z3a, a three-state inverter, is connected directly to a TRS-80 data line.

Figure 9-7 provides drive status information. Z10e is hooked to the former cartridge on-off switch, informing the computer of the presence of a cartridge in the drive. The track select switches feed their 12-volt signals to Z10a-b, which report the track pair in use. The foil sensor triggers a flip-flop made up of Z8a/b, which latches the fact that the foil has passed, until the computer resets it via the RESET SPLICE line.

The unit is turned on from the front, and power is always applied to the electronics. When the program reads or writes to tape, the motor is turned on via data bit 5; for fast read search and write, the motor is activated by data bit 6. The splice status is reset via data bit 3, and tracks are changed via data bit 4. Writing to the deck is enabled by data bits 1 and 2 (for the upper and lower of the track pairs, respectively), and the actual data writing is done through data bit 0.

Addressing of port AA is provided by Z14a-d and Z13, and this signal is combined with the computer's OUT and IN signals by Z8c/d. All the input data is latched by Z12.

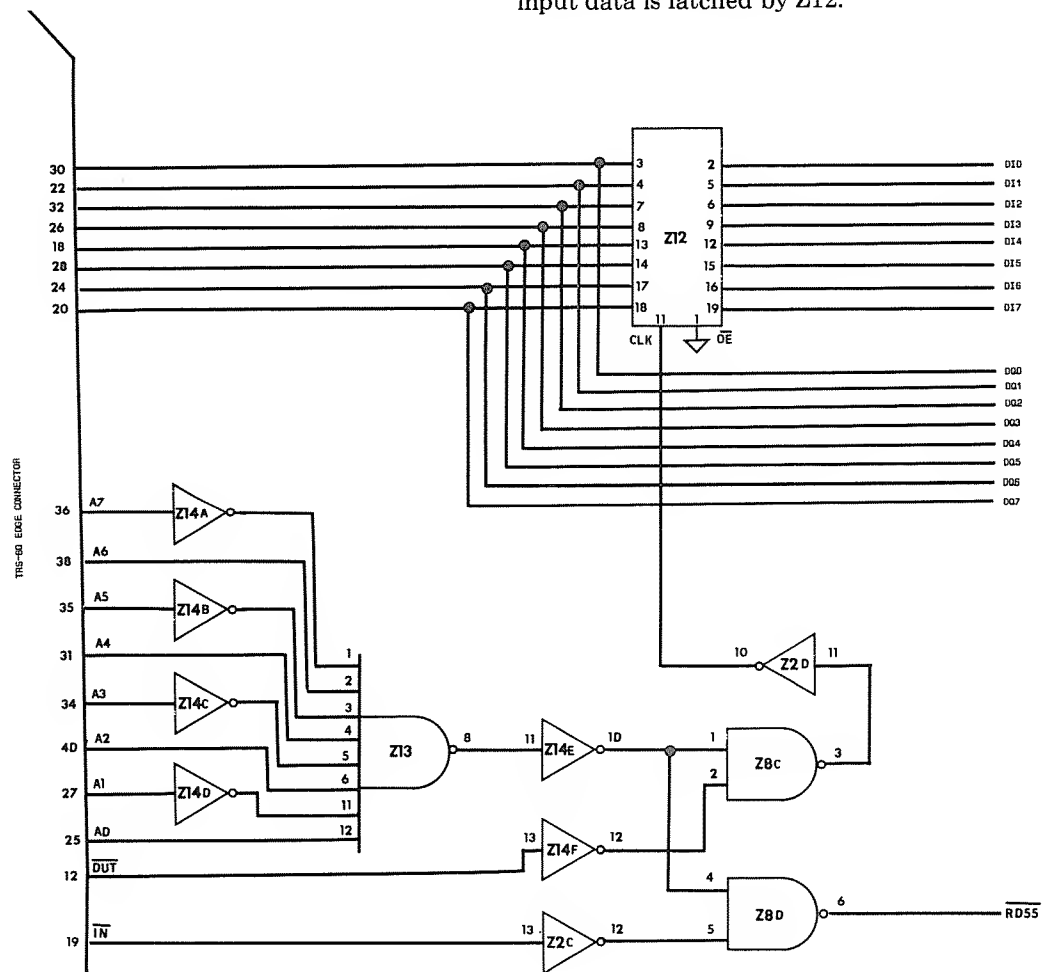


Figure 9-6. Schematic for 8-track storage system.

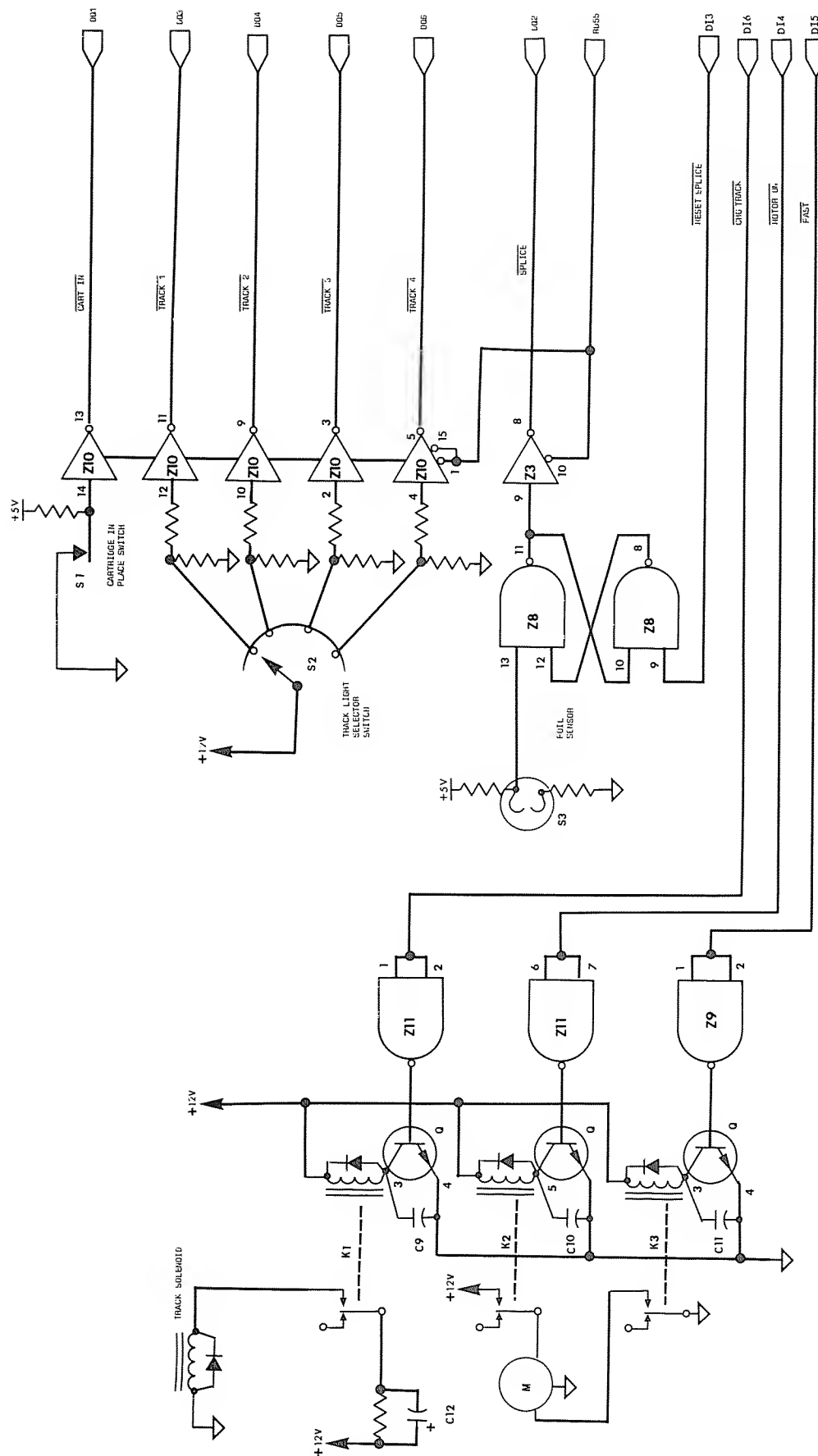
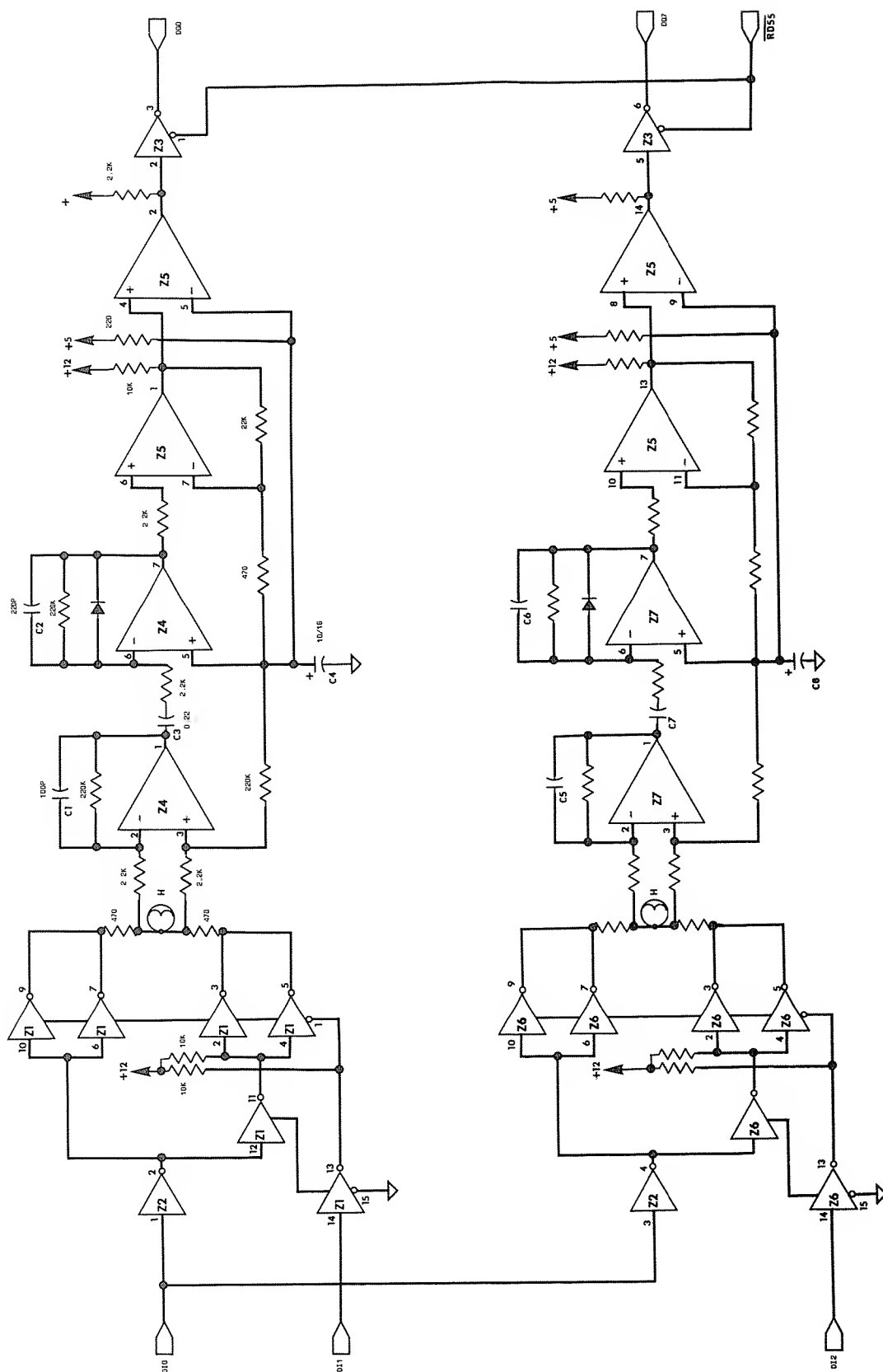


Figure 9-7. Schematic of tape player modification.



Listing 9-4. 8-track load and save cartridge system.

```

20 CLS : REM * THIS PROGRAM PRINTS CURRENT 8-TRACK DRIVE STATUS
20 OUT85,16 : REM * START 8-TRACK DRIVE BY SETTING OUTPUT BIT 4
30 PRINT@0,"DATA 1:          DATA 2:" : REM * L & R TRACKS
40 PRINT@128,"TRACK NO.:" : REM * CURRENT HEAD POSITION PROMPT
50 A=INP(85) : REM * GET DATA FROM 8-TRACK DRIVE AT PORT 55 HEX
60 PRINT@66,(A AND 128)/128 : REM * TEST FOR HIGH DATA BIT 7
70 PRINT@86,(A AND 1) : REM * TEST FOR HIGH DATA BIT 0 & PRINT
80 B=(A AND 120) : REM * 120 = 78HEX = 01110000 TRACK CONDITION
90 IF8=112THENPRINT@140,"1" : REM * 01110000 = LO BIT 3 = TK1
100 IF8=104THENPRINT@140,"2" : REM * 01101000 = LOW BIT 4 =TK2
110 IF8=88THENPRINT@140,"3" : REM * 01011000 = LOW BIT 5 = TK3
120 IF8=56THENPRINT@140,"4" : REM * 00111000 = LOW BIT 6 = TK4
130 IF(A AND 4)=4 THEN PRINT @ 256,"SPICE DETECTED" : REM * GOTD190
140 IF(A AND 2)=2 THEN PRINT @ 320,"CARTRIDGE IS IN " ELSE
    PRINT @ 320,"INSERT CARTRIDGE " : REM *
150 Q = RND(120) : REM * RANDOM TRACK SWITCH FOR TESTING ONLY
160 IF B=Q THEN OUT 85,80 ELSE 30 : REM * SWITCH TRACKS HERE
170 FOR N = 1 TO 100 : NEXT : OUT 85, 16 : REM * RESUME NORMAL
180 GOTD30 : REM * RANDOM TRACK SWITCH COMPLETE; BACK TO START
190 OUT85,0 : AS=INKEY$ : IFAS=""THEN190 : REM * OFF, THEN TEST
200 OUT85,16 : GOTD30 : REM * TURN BACK ON WHEN ANY KEY PRESSED
210 OUT254,2 : REM * HIGH-SPEED SELECT FOR TESTING PURPOSES ONLY
20 CLS : REM * THIS ROUTINE CHECKS TOTAL CURRENT 8-TRACK STATUS
30 Q=148 : OUT85,16 : REM * SET Q VALUE AND TURN ON TAPE DRIVE
40 A=INP(85) : REM * GET CURRENT STATUS OF TAPE DRIVE FROM PORT
50 B=(A AND 128)/128 : REM * GET VALUE AT DATA TRACK 1 (BIT 7)
60 C=(A AND 64)/64 : REM * GET VALUE AT TRACK POSN 4 (BIT 6)
70 D=(A AND 32)/32 : REM * GET VALUE AT TRACK POSN 3 (BIT 5)
80 E=(A AND 16)/16 : REM * GET VALUE AT TRACK POSN 2 (BIT 4)
90 F=(A AND 8)/8 : REM * GET VALUE AT TRACK POSN 1 (BIT 3)
100 G=(A AND 4)/4 : REM * GET VALUE OF SPICE CONDITION (BIT 2)
110 IF G=1 THEN OUT 85,0 : GOSUB 220 : REM * TURN OFF IF SPICE
120 H=(A AND 2)/2 : REM * GET VALUE IF CARTRIDGE IS IN (BIT 1)
130 I=A AND 1 : REM * GET VALUE AT DATA TRACK POSN TWO (BIT 0)
140 PRINT@Q,"07 06 05 06 03 02 01 00" : REM * PRINT DATA HEAD
150 PRINT@Q+64,8;C;D;E;F;G;H;I : REM * PRINT VALUES CALCULATED
160 PRINT@Q+192,"0 T T T T T S C 0 " : REM * PRINT SOME
170 PRINT@Q+256,"A R R R R P A A" : REM * PRETTY
180 PRINT@Q+320,"T K K K K L R T" : REM * WORDS
190 PRINT@Q+384,"A # # # # C T A" : REM * FOLLOW
200 PRINT@Q+448,"1 4 3 2 1 E N 0" : REM * STATUS.
210 GOTD40 : REM * AND REPEAT THE PROCESS AS THE TAPE CONTINUES
220 OUT 255,255 : OUT 255, 0 : REM * A LITTLE SCREEN SHAKING
230 AS=INKEY$ : IF AS="" THEN 220 : REM * TEST, LOOP IF NO CHAR
240 OUT85,24 : RETURN : REM * TAPE BACK ON, START PROCESS OVER

00100 : #####
00110 : SHORT 8-TRACK OPERATING SYSTEM FOR CONVERTED 8-TRACK
00120 : TAPE DECK, FORMAT:
00130 : ON, FAST FORWARD, FIND SPICE, WRITE ENABLE, WRITE
00140 : 256 BYTES OF DIRECTORY ENTRY, LOOP UNTIL 10 ENTRIES ARE
00150 : WRITTEN OR NUMBER TRANSFERRED FROM BASIC IS ACHIEVED.
00160 : FILL OUT TRACK WITH FF'S, VERIFY WRITE (IF VERIFY FLAG
00170 : IS ON), FIND SPICE, CHANGE TRACKS, REPEAT PRYESS FOR
00180 : ALL FOUR TRACKS.
00190 : DIRECTORY READ AT HIGH SPEED, ENTRIES ENTERED AT LOW
00200 : SPEED IN 256-BYTE PACKETS. EACH PACKET HAS A NUMBER
00210 : WHICH IS STORED IN THE DIRECTORY. NO DEFAULT ON TRACK
00220 : OR DIRECTORY ENTRY IS ALLOWED. DIRECTORY ENTRY #0 IS
00230 : RESERVED FOR FREE PACKETS ON EACH TRACK AND SECTOR.
00240 : *** NOTE: THIS LISTING IS NOT A COMPLETE OPERATING
00250 : SYSTEM. IT CONTAINS A SAMPLE DIRECTORY FORMATTING AREA
00260 : AND SAMPLE HIGH-SPEED TAPE READ/WRITE ROUTINES. A FULL
00270 : 8-TRACK OPERATING SYSTEM IS AVAILABLE; WRITE TO:
00280 : MS8 ELECTRONICS, MR#1 DRAWER 766, BARRE, VERMONT 05641
00290 : OR DETAILS. THE INFORMATION PRESENTED IN THIS LISTING
00300 : IS FOR DEMONSTRATION, ALTHOUGH IT MAY BE ORGANIZED,
00310 : TOGETHER WITH OTHER STANDARD SAVE/LOAD PARAMETERS, INTO
00320 : A FULL OPERATING SYSTEM. ***
00330 : #####
00340 :
7000 00350 : ORG 7000H
00360 :
00370 : ROBYTE EQU $ : ***** REASSIGN TO READ BYTE
00380 : WRBYTE EQU $ : ***** REASSIGN TO WRITE BYTE
00390 : LEADER EQU $ : ***** REASSIGN TO LEADER WRITE
00400 : LEADRO EQU $ : ***** REASSIGN TO LEADER READ
00410 : DELAY EQU 0060H : ; DELAY VALUE IN ROM
00420 : HOWMNY DEF8 100 : ; 10 UNLESS BASIC XFER
00430 : VERFLG DEF8 00H : ; TRANSFERRED FROM BASIC
00440 : VERFON EQU 01H : ; STATUS REQUIRED BY TOS
00450 : PORT EQU 55H : ; DEFINED AT B5 DECIMAL
00460 : WRITEN EQU 01H : ; 010 = HIGH FOR WR ENABL
00470 : WRITEA EQU 02H : ; 011 = HI/LO FOR WRITE A
00480 : WRITEB EQU 04H : ; 012 = HI/LO FOR WRITE B
00490 : RESSPL EQU 08H : ; 013 = HI TO RESET SPICE
00500 : START EQU 10H : ; 014 = HIGH FOR MOTOR ON
00510 : FSTFWO EQU 20H : ; 015 = HIGH FOR FAST FWD
00520 : CHANGE EQU 40H : ; 016 = HIGH TO CHNGE TRK
00530 : READA EQU 01H : ; 020 GOES HI/LO DATA A
00540 : CARTIN EQU 02H : ; 021 GOES HIGH AT CARTIN
00550 : SPLICE EQU 04H : ; 022 GOES HIGH AT SPLICE
00560 : TRACK0 EQU 08H : ; 023 GOES HIGH AT TRK 0
00570 : TRACK1 EQU 10H : ; 024 GOES HIGH A TRK 1
00580 : TRACK2 EQU 20H : ; 025 GOES HIGH AT TRK 2
00590 : TRACK3 EQU 40H : ; 026 GOES HIGH AT TRK 3
00600 : READB EQU 80H : ; 027 GOES HI/LO DATA B
7002 00 00610 : STATUS DEF8 00H : ; COMPOSITE OF ABOVE
00620 :
00630 : BASIC PATCH HERE — USE CUSTOM INTERPRETER
00640 :
00650 : #####
00660 : START 8-TRACK DECK MOTOR RUNNING (GIVE IT ENOUGH TIME)
00670 : #####
00680 :
7003 0E55 00690 : AAAAAA LO C,PORT : ; START VALUE
7005 3E10 00700 : LO A,START : ; BITS SET TO START

```

Listing Continued . . .

Figure 9- 8 presents the optional decoding of a ROM to contain the 8-track operating system. Notice that the decoding of the addresses is incomplete, so that data from 37C0 to 37FF must not be entered into the ROM to avoid bus conflict. The ROM should remain in its erased (all one's) condition in that memory area.

Listing 9-4 is an operating system for the 8-track storage system. It is made up of four major sections:

1. **Initialization.** Patching the operating system into the BASIC interpreter.
2. **Formatting.** The directory is set up past the splice on track zero in fast-forward mode. Each track is then written with program #1 headers. Because the tape is sequential, this is not a true disk-style directory. Instead, the directory stores the order and track number for each program, so that the correct track may be searched at high speed for its leader.
3. **Load Module.** This accepts the command, checks for correct syntax and program type, activates the tape deck and searches for the program.
4. **Save Module.** This module also accepts the command, checks syntax and program type, and writes the program to tape. The directory is updated.

The most interesting aspect of the software is the method used for recording the data. After a start level, a low is written to tape. Each subsequent bit changes the level either once (a zero) or twice (a one). Clock bits are not used in this scheme.

The operating system presented here is very basic, allowing only the elementary program save and load functions. However, data of all types may be stored in the system with additions to the software via additional commands.

Several problems may initially be encountered in using this system:

1. System does not respond. Make sure power is on; check for power to track lights. One should be lit. Find a cartridge which does not have its foil splice visible, insert the cartridge, switch to track 1 and PRINT INP (170). The value returned (in binary) should be x001000x. The x's are either one or zero, depending on the status of the data read outputs. Remove the cartridge. The value will change to x101000x. Switch to track 2, and replace the cartridge. Now it should read x000100x.

An 8-Track Mass Storage System

Continued Listing

```

7007 F620 00710 OR FSTFWD ; GET INTO HIGH GEAR
7009 320270 00720 LO (STATUS),A ; PUT IT INTO PLACE
700C 06FF 00730 LO 8,OFFH ; GET LONG LOOPS
700E E079 00740 OUT (C),A ; START 8 TRACK
7010 10FC 00750 QJNZ GOING ; MANY TIMES FOR SURE
00760 ;
00770 ; #####
00780 ; MOVE TO TRACK ZERO, CHECKING STATUS OF TRACK EACH TIME
00790 ; #####
00800 ;
7012 1600 00810 LO 0,0 ; TRACK VALUE TO CHECK
7014 E078 00820 IN A,(C) ; CHECK TRACK VALUE
7016 8A 00830 CP 0 ; AGAINST VALUE WANTED
7017 2800 00840 JR Z,JUMP1 ; GO IF TRACK = 0
7019 3E40 00850 LO A,CHANGE ; TRACK CHANGE VALUE
7018 2A0270 00860 LO HL,(STATUS) ; GET CURRENT STATUS
701E 84 00870 OR H ; CHECK PREVIOUS INFO
701F 320270 00880 LO (STATUS),A ; AND PUT INTO PLACE
7022 E078 00890 OUT (C),A ; CHANGE TRACK VALUE
7024 18EE 00900 JR LOOP1 ; GO BACK UNTIL TRK = 0
00910 ;
00920 ; #####
00930 ; FIND END-OF-TAPE SPLICE AND DELAY TO GET PAST SPLICE
00940 ; #####
00950 ;
7026 E078 00960 JUMP1 IN A,(C) ; GET VALUE FROM PORT
7028 E604 00970 ANO SPLICE ; CLEAR ALL OTHER VALUES
702A FE04 00980 CP SPLICE ; CHECK IF SPLICE DET'D
702C 20F8 00990 JR NZ,JUMP1 ; BACK UNTIL FOUND
702E E078 01000 JUMP1A IN A,(C) ; GET VALUE FROM PORT
7030 E604 01010 ANO SPLICE ; CLEAR ALL OTHER VALUES
7032 FE04 01020 CP SPLICE ; CHECK IF STILL SPLICE
7034 28F8 01030 JR Z,JUMP1A ; BACK UNTIL SPLICE OVER
7036 010010 01040 LO 8C,1000H ; GET DELAY VALUE
7039 C06000 01050 CALL DELAY ; DELAY TILL GOOD TAPE
01060 ;
01070 ; #####
01080 ; WRITE LEADER (NOTE: LEADER WRITE ROUTINE NOT INCLUDED)
01090 ; #####
01100 ;
703C C00070 01110 CALL LEADER ; WRITE LEADER
01120 ;
01130 ; #####
01140 ; FORMAT DIRECTORY (NOTE: THIS METHOD IS FOR EXAMPLE)
01150 ; #####
01160 ;
703F 210070 01170 LO HL,HOWMNY ; HOW MANY DIR ENTRIES?
7042 AF 01180 XOR A ; START WITH ENTRY #0
7043 06FF 01190 LO 8,OFFH ; 256 BYTES TO WRITE
7045 C00070 01200 CALL WRBYTE ; WRITE THEM IN PLACE
7048 F5 01210 LOOP2 PUSH AF ; SAVE ACCUM TEMPORARILY
7049 E078 01220 IN A,(C) ; GET VALUE FROM DEVICE
7048 E604 01230 ANO SPLICE ; CLEAR ALL OTHER VALUES
7040 FE04 01240 CP SPLICE ; SEE IF SPLICE HIT
704F 2802 01250 JR Z,JMPX ; GO IF SPLICE IS HIT
7051 10F2 01260 QJNZ LOOP2 ; ELSE ON TO NEXT BYTE
7053 F1 01270 JMPX POP AF ; RESTORE ENTRY VALUE
7054 3C 01280 INC A ; UP TO NEXT ENTRY
7055 8E 01290 CP (HL) ; LAST ENTRY COMPLETED?
7056 30E8 01300 JR NC,LOOP3 ; OUT IF NOT DONE
01310 ;
01320 ; #####
01330 ; FILL OUT TRACK WITH A SINGLE VALUE (AGAIN, OPTIONAL)
01340 ; #####
01350 ;
7058 3EFF 01360 LOOP4 LO A,OFFH ; WHEN DONE GET ALL FF'S
705A C00070 01370 CALL WRBYTE ; WRITE FF BYTE TO TAPE
7050 E078 01380 IN A,(C) ; CHECK CONTENTS OF PORT
705F E604 01390 ANO SPLICE ; MASK OUT OTHER BITS
7061 FE04 01400 CP SPLICE ; CHECK IF TO SPLICE
7063 20F3 01410 JR NZ,LOOP4 ; IF NOT KEEP WRITING
01420 ;
01430 ; #####
01440 ; GET NEXT TRACK VALUE AND GET TRACK FORMATTING READY
01450 ; #####
01460 ;
7065 14 01470 INC 0 ; GET NEXT TRACK VALUE
7066 7A 01480 LO A,0 ; CHECK CURRENT VALUE
7067 FE04 01490 CP 4 ; IS IT FOURTH TRACK?
7069 2819 01500 JR Z,JUMP2 ; IF NOT THEN BACK
01510 ;
01520 ; #####
01530 ; FORMAT NEXT TRACK, AND CONTINUE UNTIL PROCESS COMPLETE
01540 ; #####
01550 ;
7068 3E40 01560 LO A,CHANGE ; GET CHANGE TRACK VALUE
7060 2A0270 01570 LO HL,(STATUS) ; GET CURRENT STATUS
7070 84 01580 OR H ; STATUS PLUS CHANGE
7071 E078 01590 OUT (C),A ; AND CHANGE THE TRACK
7073 3E40 01600 LO A,CHANGE ; GET CHANGE TRACK AGAIN
7075 2F 01610 CPL ; SWITCH THE BITS
7076 2A0270 01620 LO HL,(STATUS) ; GET THE STATUS
7079 A4 01630 ANO H ; STATUS PLUS CHANGE
707A 010002 01640 LO 8C,200H ; GET DELAY VALUE
7070 C06000 01650 CALL DELAY ; AND INVOKE THE DELAY
7080 E079 01660 OUT (C),A ; AND TURN IT OFF
7082 18A2 01670 JR JUMP1 ; AND GO BACK FOR MORE
01680 ;
01690 ; #####
01700 ; CHECK IF VERIFY FLAG IS ON (NOTE: ADD AS BASIC COMMAND)
01710 ; #####
01720 ;
7084 3A0170 01730 JUMP2 LO A,(VERFLG) ; STATUS OF VERIFY FLAG
7087 FE01 01740 CP VERFON ; SEE IF FLAG IS ON
7089 2036 01750 JR NZ,JUMP3 ; IF NOT THEN SKIP PAST
01760 ;
01770 ; #####
01780 ; MOVE TO THE FIRST TRACK, CHECK EACH AS IT PROGRESSES
01790 ; #####
01800 ;
7088 1600 01810 LO 0,0 ; ST COUNTER TO ZERO
7080 E078 01820 JUMP4 IN A,(C) ; GET STATUS FROM DECK
708F 8A 01830 CP 0 ; CHECK AGAINST TRK 0
7090 2819 01840 JR Z,JUMP5 ; IF AT ZERO, THEN GO
7082 3E40 01850 LO A,CHANGE ; ELSE BEGIN TO CHANGE
7084 2A0270 01860 LO HL,(STATUS) ; GET VALUE FROM STATUS

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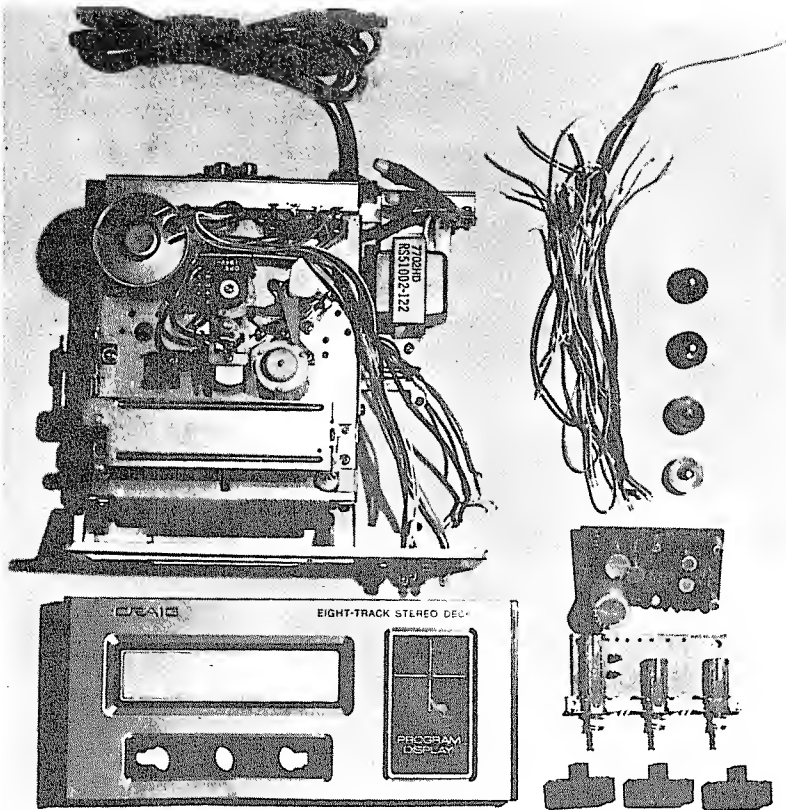
Listing Continued . . .

2. Tracks do not switch. Check wiring to the solenoid, and that the 75452 is wired correctly. Listen to hear if the solenoid attempts to react (a light click or start). Remove the 75452 from its socket and short the free lead of the solenoid to ground. It should switch. Replace the 75452 if necessary.

3. Programs do not load. If programs do not load at all, check the cartridge on an audio deck to see if something has been written. If not, go on to #4 below. If so, listen for occasional changes in pitch as the machine switches from fast forward to normal. Lengthen the speed change wait period in the program if you can hear the pitch slide as it restarts at a new level. If a loading message is displayed, but an error is detected, try to read from another track. Tracks 1a and 4b are at the edge of the tape, and lower-quality tapes may drop out occasionally in this area. The head may be badly misaligned and not make good contact with the tape. This can be heard as shifting or slewing in the sound. Adjust the Phillips alignment screw on the head to match a prerecorded commercial tape of good quality.

4. Programs do not save. Begin the program-saving process, and place the signal lead of a small amp against one lead of the recording head. If the signal is present, the program should be saving. If not, check the wiring of the buffer IC's, which may not be letting the signal through. Also check that the software is entered correctly, and that a signal is actually being sent to the device (correct connection of the write line, and proper wiring of the address decoding and data latch). If the motor turns on and switches tracks properly, the signal is probably being held up by incorrectly wired buffers to the recording head.

5. Motor speed does not change. Make sure that the third lead from the motor is being switched to ground, not positive voltage. This lead reacts best when switched below ground, and ground potential is its minimum position. If you have substituted another solid-state switch for the one shown, make sure it goes to full ground potential when switched in place.



Craig playback deck disassembled. Front panel, main frame and feet are maintained, but switch panel, playback electronics, wires, and buttons are discarded. Front buttons may be kept for appearance.

Continued Listing

```

7097 B4      01870      OR      H      ; AND CHECK STATUS VAL
7098 E079    01880      OUT     (C),A    ; AND SEND OUT CHANGE
709A 3E40    01890      LO      A,CHANGE ; GET VALUE AGAIN
709C 2F      01900      CPL      ; REVERSE THE BITS
709D 2A0270  01910      LO      HL,(STATUS) ; GET VALUE FROM STATUS
70A0 A4      01920      ANO      H      ; GET VALUE FROM STATUS
70A1 010002  01930      LO      BC,200H   ; GET DELAY VALUE
70A4 C06000  01940      CALL    DELAY    ; CALL DELAY IN ROM
70A7 E079    01950      OUT     (C),A    ; AND SEND OUT CHANGE
70A9 1B2     01960      JR      JUMP4    ; AND GO TO THE REST
              01970
              01980
              01990      ; *****
02000      ; VERIFY THE TAPE FORMATTING AND QUALITY IF FLAG IS ON
02010      ; *****
02020 JUMP5   LO      E,0      ; PACKET NUMBER TO START
02030 LOOP7   LO      B,OFFH   ; BYTE TO VERIFY WITH
02040      CALL  LEA0R0       ; GET SYNC AND READ LEAOR
02050      CALL  ROBYTE      ; BEGIN READING THE BYTES
02060      CP      E          ; AND CHECK FORMAT INFO
02070      JR      NZ,VERERR   ; IF NOT CORRECT, ERROR
02080      DJNZ   LOOP6       ; AND GO BACK UNTIL DONE
02090      INC      E          ; GO TO NEXT PACKET
020A 1C      02100      LO      A,(HOWMNY) ; CHECK IF FORMAT DONE
020B 3A0070  02110      CP      E          ; BY GETTING PACKET #
020E 8B      02120      JR      NZ,LOOP7   ; BACK IF NOT DONE
020F 20EC    02130
              02140
              02150      ; *****
02160      ; STOP THE TAPE MACHINE AND CLEAR OUT ALL STATUS INFO
02170      ; *****
02180 JUMP3   XOR      A          ; CLEAR ACCUM TO ZERO
02190      LO      (STATUS),A    ; PUT IN STATUS LOCATION
021A 320270  02200      OUT     (C),A    ; AND SEND IT OUT PORT
021B E079    02210
              02220
              02230      ; *****
02240      ; GET AND DISPLAY FORMAT COMPLETE MESSAGE (OR ERROR MSG)
02250      ; *****
02260      LO      HL,MSG01      ; GET THE OKAY MESSAGE
02270      CALL  2BA7H          ; LEVEL II DISPLAY ROUT.
02280      JP      107BH        ; BACK TO INTERPRETER
02290 MSG01  DEFM  'FORMATTING COMPLETE'
022A 00      02300      VERERR  LO      HL,MSG02      ; VERIFY ERROR MESSAGE
022B 21EC70  02310      CALL  2BA7H          ; DISPLAY SUBROUTINE
022C C0A72B  02320      JP      06CCH        ; BASIC COMMAND LEVEL
022D C3CC06  02330 MSG02  DEFM  'FORMAT ERROR'
022E 46      02340
              02350
              02360      ; *****
02370      ; THESE ROUTINES MAY BE CALLED WITH THE /LOAD COMMANDS OR
02380      ; THE /SAVE COMMANDS, AS WELL AS THE FORMATTING ABOVE.
              ; NOTE: THESE ROUTINES USE AN EXCLUSIVE-OR METHOD OF

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Listing Continued . . .

Construction and Checkout

Since I recommend wire-wrapping the projects in this book, there are some different considerations when wrapping all the resistors and capacitors necessary for this project. Wire-wrapping is done best when there are sharp, square pins on which to wrap; since resistors and capacitors have round leads, it is best to insert these parts into sockets. Not only will the wrapping be more firmly attached, but the parts will not fall out when you are turning the board over during construction.

Before beginning construction, note that the transformer installed in the Craig recorder is marginal; ideally, it should be replaced with a 12 volt, 1.5 amp transformer of the type sold by Radio Shack. This will fit, though not comfortably, in the present transformer's location. The center tap is not used.

Because there are mechanical and inductive parts in this device, a great deal of electromagnetic noise can be produced. Be sure, then, to do the following:

1. Use bypass capacitors (.01 mF or so) between power and ground on every integrated circuit.
2. Use bypass capacitors (.01 mF) on the 75452 peripheral driver chips as noted, as well as diodes (hobbyist 1N914's are fine) across the three relay coils. Leave the diode already present on the track change solenoid. *Don't* skip these parts; they make the difference between a working system and a continuous return to MEMORY SIZE.
3. Use a central grounding point for all wires on the circuit board, and another central point for all wires on the drive. Don't weave wires from one IC to another.
4. The resistors at the inputs of the CMOS chips (the 80C96 and 80C98s) are vital for a reliable signal, and also prevent damage to the chips themselves.
5. The wires leading to the recording head should be twisted or shielded to prevent picking up extraneous signals. Remember also that the connections of the top and bottom pairs of pins on the head should be the same.

Construction and Checkout

Continued Listing

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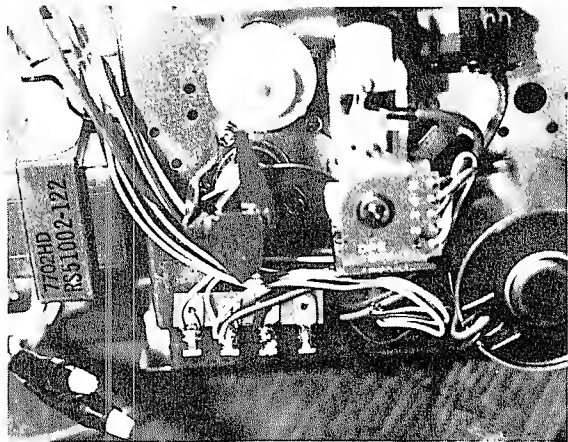
02390 ; RECORDING, WHICH WILL FUNCTION WITH A DIRECT DIGITAL
02400 ; RECORDING SITUATION SUCH AS THE B-TRACK DEVICE. FOR
02410 ; FURTHER INFORMATION ON ITS OPERATION, ASK FOR THE
02420 ; OHIO SCIENTIFIC TAPE WAFER INTERFACE, SOLD BY MS8
02430 ; ELECTRONICS (ADDRESS ABOVE), AND EXAMINE THE CODE.
02440 ; #####
02450 ;
40F9 02460 ; BOTTOM EQU 40F9H ; BOTTOM OF BASIC PROGRAM
06CC 02470 ; BASIC EQU 06CCH ; RETURN TO READY MODE
02480 ;
70F8 F3 02490 ; ENTER OI ; KILL THEM SUCKERS
70F9 2AF940 02500 ; LO HL,(BOTTOM) ; GET FIRST PROGRAM LOC'N
70FC 3A3040 02510 ; LO A,(4030H) ; VALUE IN THE MASK
70FF E6FC 02520 ; AND OFCH ; SET BIT TO ZERO
7101 323040 02530 ; LO (4030H),A ; PUT BACK INTO LATCH
7104 5F 02540 ; LO E,A ; WILL BE USED A LOT
02550 ;
7105 7E 02560 ; NEXT LO A,(HL) ; SPECIAL TEST (EXAMPLE)
7106 57 02570 ; LO 0,A ; THIS IS THE BIG BUMPER
7107 0608 02580 ; LO 8,8 ; NUMBER OF BITS TO WRITE
7109 3ED1 02590 ; LO A,1 ; THIS IS THE START BIT
7108 83 02600 ; OR E ; GET PROPER LATCH MASK
710C 03FF 02610 ; OUT (OFFH),A ; WRITE OUT START LEVEL
710E C04371 02620 ; CALL DELAY1 ; WRITE A START DELAY
7111 AF 02630 ; XOR A ; SET UP A WITH 0 I/O BIT
7112 B3 02640 ; OR E ; GET PROPER LATCH MASK
7113 03FF 02650 ; OUT (OFFH),A ; WRITE OUT STARTING EDGE
02660 ;
02670 ; #####
02680 ; ADDING 66 T-STATES HERE TO TOTAL 95 US
02690 ; #####
02700 ;
7115 00E5 02710 ; PUSH IX ; 15 T-STATES
7117 00E1 02720 ; POP IX ; 14 T-STATES
7119 00E5 02730 ; PUSH IX ; 15 T-STATES
7118 00E1 02740 ; POP IX ; 14 T-STATES
711D 00 02750 ; NOP ; 4 T-STATES
711E 00 02760 ; NOP ; 4 T-STATES
02770 ;
02780 ; #####
02790 ;
02800 ; THIS IS SET UP FOR 3300 BAUD
02810 ; TOTAL LOOP TIME PER BIT SHOULD BE 300 US
02820 ; TOTAL T-STATES FOR EACH BIT LOOP = 169 = 95 US
02830 ; TOTAL T-STATES FOR BIT DELAY LOOPS AT 0060 NOTED BELOW
02840 ; #####
02850 ;
711F C04E71 02860 ; LOOP CALL DELAY2 ; WRITE A NORMAL DELAY
02870 ;
7122 AA 02880 ; XOR 0 ; TEST THE FIRST BIT
7123 E601 02890 ; AND 1 ; MASK OUT OTHER 8 BITS
7125 F5 02900 ; PUSH AF ; SAVE THE VALUE
7126 83 02910 ; OR E ; VALUE OF LATCH MASK
7127 03FF 02920 ; OUT (OFFH),A ; WRITE IT
02930 ;
7129 F1 02940 ; POP AF ; ORIGINAL VALUE BACK
712A 2F 02950 ; CPL ; REVERSE THE BIT
712B E601 02960 ; AND 1 ; MASK OUT ALL BUT ONE
712C C04E71 02970 ; CALL DELAY2 ; WRITE A NORMAL DELAY
7130 F5 02980 ; PUSH AF ; SAVE THE PROPER BIT
7131 B3 02990 ; OR E ; AGAIN GET THE MASK
7132 220000 03000 ; LO (0000),HL ; NEED DELAY TIME
7135 00 03010 ; NOP ; BIT MORE DELAY TIME
7136 03FF 03080 ; OUT (OFFH),A ; WRITE IT
03090 ;
7138 F1 03100 ; POP AF ; GET ORIGINAL BITS BACK
7139 CB0A 03110 ; RRC 0 ; ORIENT TO NEXT BIT
713B 10E2 03120 ; ORJNZ LOOP ; WRITE OUT 8 BITS
03130 ;
03140 ; XOR A ; CLEAR ACC. TO ZERO
03150 ; OR E ; GET MASK FROM 403D
03160 ; OUT (OFFH),A ; SEND OUT ZERO BIT
7130 C04E71 03170 ; CALL DELAY2 ; SPACE OUT LAST BIT TOO
03180 ;
7140 23 03190 ; INC HL ; GET NEXT MEMORY LOC'N
03200 ;
03210 ; #####
03220 ; PUT TESTING FOR MENTOP HERE
03230 ; #####
03240 ;
03250 ; LO A,H ; CURRENT MEMORY STATUS
03260 ; CP 40H ; END? (EXAMPLE ONLY!!)
03270 ; JR NZ,NEXT ; GO BACK IF NOT DONE
03280 ; JP 06CC ; READY (USE RETURN!)
03290 ;
7143 F5 03300 ; DELAY1 PUSH AF ; SAVE AF REGISTERS
7144 C5 03310 ; PUSH BC ; SAVE BC REGISTERS
7145 012800 03320 ; LO 8C,28H ; GET DELAY VALUE
03330 ;
03340 ; #####
03350 ; LINE ABOVE COMPLETES A TOTAL 1000 US DELAY LOOP
03360 ; #####
03370 ;
7148 C06000 03380 ; CALL 0060H ; MAKE A DELAY
714B C1 03390 ; POP 8C ; RESTORE BC REGISTERS
714C F1 03400 ; POP AF ; RESTORE AF REGISTERS
714D C9 03410 ; RET ; BACK TO MAIN ROUTINE
03420 ;
714E F5 03430 ; DELAY2 PUSH AF ; SAVE AF REGISTERS
714F C5 03440 ; PUSH BC ; SAVE BC REGISTERS
7150 010300 03450 ; LO BC,03H ; GET SHORTER DELAY VALUE
03460 ;
03470 ; #####
03480 ; LINE ABOVE COMPLETES AT TOTAL 205 US DELAY LOOP
03490 ; #####
03500 ;
03510 ; CALL 0060H ; MAKE A DELAY
7153 C06000 03520 ; POP 8C ; RESTORE BC REGISTERS
7156 C1 03530 ; POP AF ; RESTORE AF REGISTERS
7157 F1 03540 ; RET ; BACK TO MAIN ROUTINE
7158 C9 03550 ;
03560 ; #####
03570 ; ENO ENTER
00000 TOTAL ERRORS
21490 TEXT AREA BYTES LEFT

```

Position sockets as close together as possible so the final board will fit into the tape drive's case. Fill the board with sockets and parts, and test its size before beginning the wire-wrapping.

Wire-wrap all connections completely before installing the 80C96 and 80C98 ICs. Because they are static-sensitive CMOS, they can be damaged by improper handling or application of power to partly-connected ICs.

When construction is complete, install the ICs, connect the unit to the TRS-80, and apply power.



Detail view of Craig mechanics. Track change mechanism is operated by a spinning cog on the drive capstan (lower right). Four track change lights are illuminated by switch contacts to the rear of the playback head.

```

AAAAA 7003 00690
BASIC 06CC 02470
BOTTOM 40F9 02460
CARTIN 0002 00540
CHANGE 0040 00520
DELAY 0060 00410
DELAY1 7143 03300
DELAY2 714E 03430
ENTER 70F8 02490
FSTFW0 0020 00510
GOING 700E 00740
HOWMNY 7000 00420
JMPX 7053 01270
JUMP1 7026 00960
JUMP1A 702E 01000
JUMP2 70B4 01730
JUMP3 70C1 02180
JUMP4 70B0 01820
JUMP5 70AB 02020
LEADER 7000 00390
LEADER 7000 00400
LOOP 711F 02860
LOOP1 7014 00820
LOOP2 7045 01200
LOOP3 7043 01190
LOOP4 705B 01360
LOOP6 70B2 02050
LOOP7 70A0 02030
MSG01 7000 02290
MSG02 70EC 02330
NEXT 7105 02560
PORT 0055 00450
ROBYTE 7000 00370
READA 0001 00530
READB 0080 00600
RSPICE 0008 00490
SPICE 0004 00550
START 0010 00500
STATUS 7002 00610
TRACK0 0008 00560
TRACK1 0010 00570
TRACK2 0020 00580
TRACK3 0040 00590
VERERR 70E3 02300
VERFLG 7001 00430
VERFON 0001 00440
WRBYTE 7000 00380
WRITEA 0002 00470
WRITEB 0004 00480
WRITEN 0001 00460
02500
00850 01560 01600 01850 01890
01050 01650 01340
02620
02860 02970 03170
03570
00710
00750
01170 02100
01250
00840 00990 01670
01030
01500
01750
01960
01840
01110
02040
03120
00900
01260
01300
01410
02000
02120
02260
02300
03270
00690
02050
01400
00700
00720 00860 00880 01570 01620 01860 01910
02190
00560
00570
00580
00590
02070
01730
01740
01200 01370

```

10

And Now It's Broken

It is not inevitable that your TRS-80 will fail during your lifetime, but there's always that chance. And if it happens, there's no reason to truck the computer down to the nearby Cost-a-Buck repair center. Do it yourself. This chapter will present the most likely failures or dilemmas you may encounter with your TRS-80, including:

Setting up a reliable, crash-free environment in a typically casual home.

Curing memory crashes in the CPU or the expansion interface, and replacing failed memory.

Solving the garbage-on-screen power-up failure.

Discovering the many sources of mysterious program crashes and keyboard lockup, and how to cure them.

Aligning your video display to cure images off-screen, tearing or jitter.

How program bugs can look like hardware failures, and vice versa.

'Routine' maintenance – the hidden cure for many failures.

Handling the computer and its peripherals.

Overview of difficulties in disk drivers, cables, cassette devices, printers, RS-232 boards, and other add-ons.

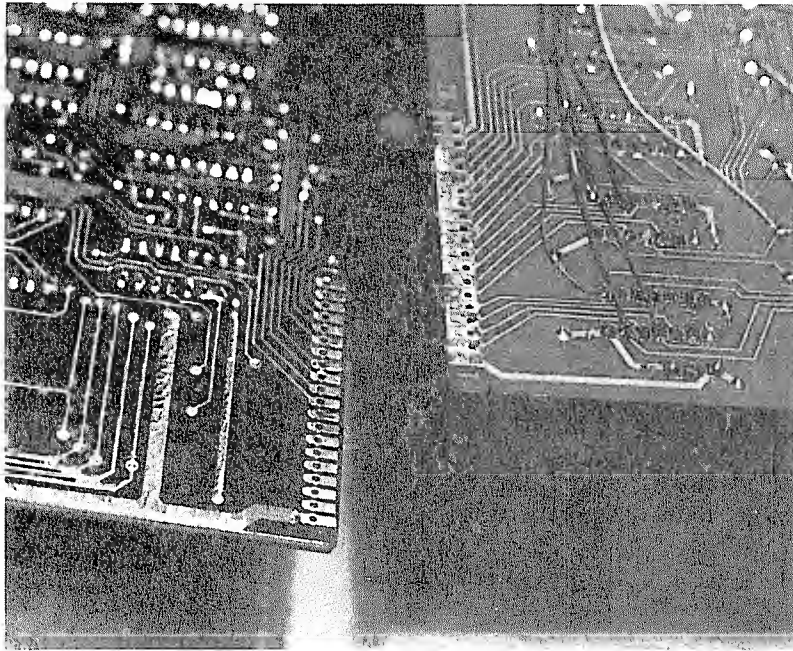
A New Keyboard Cable

Virtually every modification to the TRS-80 contains a warning like this: 'Carefully open the case, and carefully take out the unit then carefully spread it out. Work carefully so as not to bust the keyboard cable'. So what's with the cable that makes it so fragile? The connections to the computer boards seem secure, but if you look carefully at the cable itself, you will see that it is made up of flat copper bands inside an insulating strip. The bands themselves are strong, and the connections to the computer are strong.

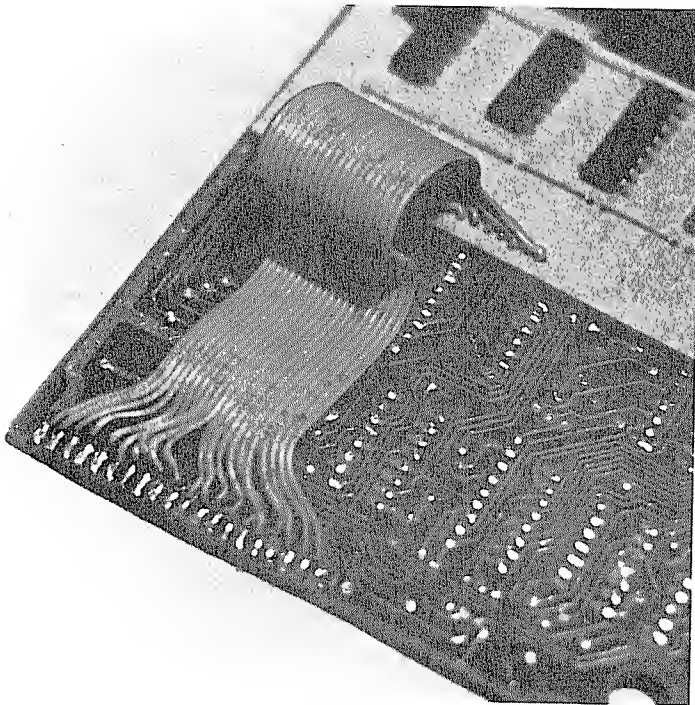
The problem occurs at the point where the copper bands are clamped to the connectors that attach to the circuit board. Hairline cracks develop in the copper bands, separating them almost invisibly from the connecting pins. The only evidence of these cracks comes when odd combinations of letters appear sporadically while you are typing.

If you plan to open the machine more than a half dozen times, you should replace the connection, with a 20-wire multiconductor cable or very flexible single conductor phono cartridge wire.

First, cut the present cable off with shears, cutting at the point where the copper bands meet the connectors. Now, with generous amounts of solder-wick, remove the 40



Replacing the keyboard cable: Solder-wick and flux cleaner are used to remove all traces of the old connection cable.



Replacing the keyboard cable: A new flexible cable is attached to both boards. A removable cable and connector can be used, but the permanent flexible cable serves well enough once most modifications have been made.

connectors on the keyboard and CPU board. Make sure all the copper is out of the holes, and clean what is surely to be a mess with flux remover, then buff the solder so new wires will slip easily through the holes.

The new cable can be six inches long or more without affecting the operation of the unit. My own added keyboard, which is wired into place in the same manner as the original, has a 20 foot cable and suffers no ill effects or program crashes. Cut and strip no more than one quarter inch of insulation off each end, and tin all 40 ends.

Tinning is the process of running some hot solder onto the wire to prepare it for easy soldering to the board. The wire will remain the same diameter, the solder should not lump up, but will become smooth, hard and shiny with a solder coating. This will allow the wire to pass easily through the holes without random strands sticking out sideways and shorting against the neighboring connection. If the insulation creeps back from the heat, finish tinning all the ends and then clip them back to one quarter inch.

Pass the multiconductor through all the holes in advance, and secure it temporarily with masking tape or other tape not affected by warmth. Getting all the wires through can be tricky, but accept the absurdity of the process in advance; outsiders tend to burst into laughter at the sight of a grown adult trying to thread 20 needles simultaneously, so if you can't take a joke, do this with the door locked.

When all the wires are attached to either the keyboard or the CPU board (I recommend the keyboard first because it is lighter), check for shorts, then thread the cable through the other board. Solder, then look carefully for shorts before applying the power; both five volts and ground run through the cable, so be sure everything is well.

Apply power. All characters should work properly, and no odd combinations of letters should be produced. If the power LED does not come on, either a connection was left off or you have a short. Power off immediately, and check again. If letters are missing, a connection was probably not made. If odd combinations of letters appear, you probably have a broken wire. If incorrect letters appear, you have switched wires. Once it's all working, you can ignore those 'carefully . . .' warnings and concentrate on the modifications to be made.

Home and Small Business Environments

Computers are designed in laboratories, tested in laboratories, examined by engineers and run by programmers. There could hardly be a more unlikely manner of producing an appliance-type product like the TRS-80. When you unpack the home computer and plug it in, you begin a torture test unimagined by the professionals in their sleek, air conditioned factories.

Below is a computer environment quick quiz; if you answer 'yes' to any question, your TRS could be in trouble. The more 'yesses', the more potentials for disaster:

- Your computer and peripherals are plugged into extension cords or cube-taps.
- A refrigerator, toaster, water pump, washer, dryer, or other large appliance is on the same fuse or breaker circuit as your computer.
- You have an electric mixer, blender, or food processor which is used when the computer is used.
- You have an electric drill, jigsaw, or table saw which is used when the computer is used.
- Your computer table desk lamp is a push-on and hold fluorescent type.
- You use your computer on or near a television set while the tv is turned on.
- Your computer is used in the kitchen or shop, you or a family member smokes, or you heat with wood or coal.
- You have fur-bearing animals in the house near your computer, especially cats.
- Your power is supplied by a rural cooperative, or is generated by a local power company especially with low-head hydropower systems.
- You live in an industrial area where heavy electrical equipment (winches, cranes) is used.
- Your electrical thermostat is located near the computer.
- Your home is especially dry, and you do not use a humidifier.
- You live in a coastal area, or by a salt-water lake.
- Your computer area is located near a railroad or by a highway traveled by heavy trucks.
- You move your computer while it is on.

Chances are you've got at least one check mark on the list above. Here are some solutions to these environmental problems:

1. Plug your computer directly into a wall outlet, or use a commercial 'power strip'.
2. Make sure that the computer is on a single-appliance circuit, even though it uses very little power.
3. Brush-motor electrical appliances like mixers and drills create an enormous amount of electrical noise. Reach a compromise with the culinary artist or the shop craftsperson to use those tools at some other time, or at least far away from the computer.
4. Use incandescent lamps, which send out virtually no electrical hash, and get rid of the fluorescent ones. Don't use a television for a computer table, because it creates heavy electromagnetic fields. And move the computer or the thermostat; there's a tremendous electrical noise jolt transmitted when that thermostat turns on.
5. Keep smoke and grease of all kinds out of your computer's atmosphere, as well as animal (and people) hair. Tape and disk drives hate the stuff, and cables, connector, and keys build up greasy mudpiles because of them.
6. If your power is unsteady because of an inept or unconcerned power company, or because local industry unexpectedly drains a large amount of power, you will have to install some sort of power regulator. Types such as Solatron, Mayday, and Topaz provide different qualities of regulation, at corresponding costs (see Appendix).
7. Use a humidifier, or place pans of water on radiators or stoves if you house is especially dry in winter, because static electricity is quite powerful. If you live in a salt-air atmosphere, air-condition your computer area during warm, humid days. Salt air corrodes cables and connections.
8. If railroads or trucks are nearby, cushion the computer. Vibration can cause noise in cables and especially the expansion box, and make disk reading and writing very failure-prone. Likewise, moving the computer (or even just pulling or straining the expansion or peripheral cables) can create bursts of electrical noise through the computer. Don't do it. Interestingly, even

pushing away a hard chair that vibrates on a linoleum or hardwood floor can cause disk read/write errors. And don't forget the dummy plug for CTR-41 tape recorders when saving programs!

The home and small business environment is not often conducive to these suggestions, and some of them may not be necessary, depending on other factors. If you maintain your cables, edge card connectors, keyboards, etc., and keep your computer cushioned and seated on a conductive tabletop, you've gone a long way to increasing reliability.

Furthermore, merely thinking of your micromputer in the same terms that used to be reserved for larger computers in the past (COMPUTER ROOM! NO SMOKING! . . . CAUTION! SENSITIVE ELECTRONIC EQUIPMENT! . . . NO FOOD OR DRINK IN THIS ROOM!), then you have the right idea. Electronically, your TRS-80 can take a great deal of abuse and still function. But this abuse cannot take place *while the computer is running*. And that is the clue: treat your operating computer as if you were paying \$50 an hour in time-sharing charges.

When the Memory Crashes

In early TRS-80's, memory crashes were the most prominent sort of failure. The type of memories used were the culprit, partly because of expansion box design problems, but especially because an unusual condition called the 'soft error' had not then been diagnosed.

The 'soft error' was the tendency of a perfectly good program to crash with some error message when no such error was present. A simple CONTINUE command would restore the system. These errors were caused by the internal structure of the memory chips themselves, which, because they are 'dynamic', use a peculiar and surprising principle for their operation.

Memories maintain information. That is their job. 'Static' memories retain information so long as the power is applied to the computer. 'Dynamic' memories, the type used in the TRS-80, retain their information for only a few thousandths of a second, requiring a electronic prodding, called *refresh*, to remember their data. They depend on their internal capacitance, acting much like a leaky tire.

In the early days, this odd way of maintaining memory resulted in occasional erratic behavior, sometimes because the chip itself was fluke, and sometimes because normal low-level radioactive

alpha-decay, *present right on the base of the chip*, could knock a memory bit from one state to another. This radioactivity is so delicate that a single sheet of paper can stop it. So this memory failure would occur only when the radioactive alpha particle actually struck a junction, and only when that junction was struck at precisely the right billionth of a second. Newer memories use a 'cool' base which does not emit alpha radioactivity, and so this rather bizarre problem has finally disappeared.

But memory crashes still occur, and they come in a few major forms:

Temporary crashes due to electrical noise in the vicinity of the computer.

Temporary crashes in the expansion interface due to a badly attached or otherwise noisy set of refresh lines.

Temporary crashes due to improperly seated or corroding memory chips.

Permanent crashes due to bad memory chips.

Repairable crashes due to a damage to one of the three lines responsible for memory refresh.

Electrical storms. I'll digress just a moment on this one. During a summer meeting of the Vermont Computer Guild, an electrical storm approached rapidly. We began to engage in a sly but nervous game of electronic chicken, leaving our micros not only attached to the power, but running! Naturally, that was too much of an invitation to Mom Nature, who zapped our very power line with a basketful of megavolts. My TRS-80 kept working; another hung, but reset. An Apple winked and sighed, and its memory cleared. A KIM turned tail completely, taking with it video routines and all. And the expensive DEC PDP/11 was broiled, losing more than a handful of expensive chips. Moral: You know it.

Also, there are crashes which appear to be memory crashes when they are in fact otherwise. Among these:

A **wayward** program containing errors in PEEK or POKE statements, or machine language subroutines.

A **damaged** CPU chip or blown buffer chip.

A **cracked** circuit board trace, solder ball, or solder splash.

Cleaning the Edge Connectors

The expansion connector on the back of the CPU, and the various ports around the expansion interface, were all manufactured using a solder coating instead of gold plating. Because solder is lead, and lead tends to corrode badly, these connections will inevitably get electrically noisy.

Cleaning them is quite easy, and should be done, depending on your environment, from as often as weekly to a minimum of monthly.

Turn the equipment off and remove all cables. Check the edges of the cables for internal bend pins, hairs, or other damage or obstructions. Next, remove the top and bottom of the keyboard and expansion interface cases; if you don't wish to go inside your computer, this process can still be done, although it is

awkward.

Using a new dollar bill, a piece of the finest grade of emery paper, or a fine talc buffing wheel, bring the contacts to a bright shine. Brush off any remaining particles, and spray with contact cleaner (sometimes called tv tuner cleaner). Repeat this for all the edge connectors, and reinstall the cables.

If this process is repeated regularly, interconnect noise problems will be completely eliminated. However, if cables slide along the contacts when the computer is in use, noise may still be produced. To cure this, place no strain on the cables, and replace the keyboard expansion cable with a longer, more flexible unit. If your TRS-80 uses the buffered cable, instead of replacing it, merely add to it with one of the cable extenders sold by *Exatron* and others (see Appendix).

There are two software tools which can be used to eke out memory problems: the memory test printed in Chapter 3, and MEMORY SIZE? itself. Running the memory test will describe which memory locations and chips may be bad. When the MEMORY SIZE? question is answered with only a carriage return, a Level II subroutine begins testing each memory location until it finds a bad one; that final non-memory or bad memory location is found by typing:

```
PRINT PEEK (16561) + 256 * PEEK (16562) + 2
```

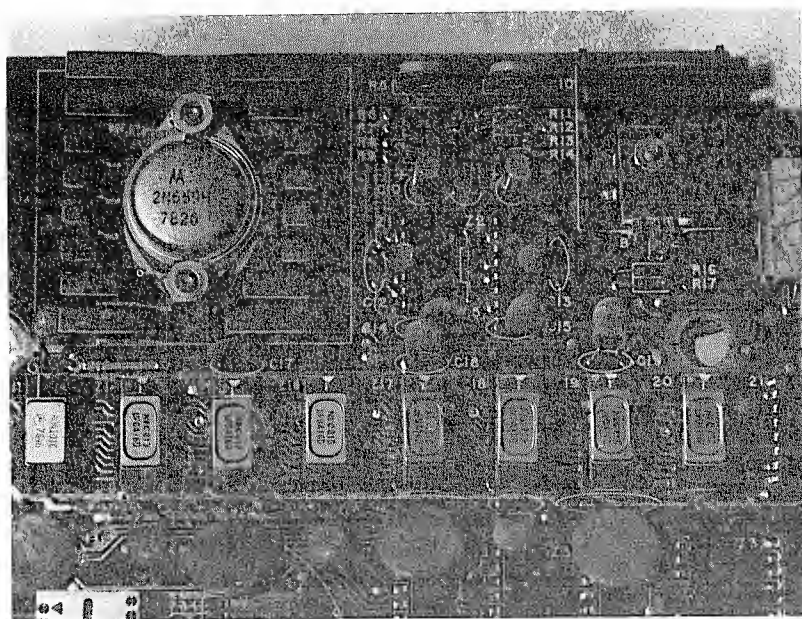


Photo 10-1. Power supply transistors.

In a 16K machine, the value returned should be 32768. If it is less, then the value returned is a bad memory location. Since the MEMORY SIZE? test is simplistic, it will not identify all possible memory errors, and if you suspect one, run the memory test.

If the bad memory location seems to move about, then perhaps it is not a true memory problem at all, but rather due to bad connections or electrical noise. Be sure noisy electrical equipment is not located near the computer. Clean cable and edge card connectors (see Box), and see if the problem recurs. If it does, open the case of your keyboard unit (and expansion box), and remove and reseat each memory chip. Look for corrosion, especially on newer chips without gold-plated leads.

If the problem is still not cured, the difficulty may be in the power supply. The location to suspect first is the large transistor (Q 4 - see

Photo 10-1), which is screwed down to the circuit board. Corrosion can build up between this transistor and the solder-plated circuit board. Loosen and remove the screws that attach the transistor, but do not attempt to remove the transistor itself. Slide fine emery paper - definitely *not* steel wool or sandpaper - face down between the transistor and the board, and clean out any corrosion.

If a lockwasher was not used between the transistor and the board, insert one, and reinstall and tighten the screws. This should stabilize the power supply inside the case.

Using an Oscilloscope

The oscilloscope is a very sophisticated tool, but in this book it is used for an elementary purpose: merely to see if a signal is 'there' or not, and if it looks pretty good. Almost all TRS-80 circuitry failures can be traced this way, and it requires no previous experience using an oscilloscope, and no special training in reading waveform timings.

The height of an oscilloscope screen trace changes in proportion to the signal present at its input. To see this, plug in the scope, turn it on, and adjust the intensity until a flat line is visible in the center of the display. Attach its cables to the vertical input. Leaving the ground (black) wiring hanging, hold the signal (red) wire in one hand. The flat line displayed on the screen should go wild.

Adjust the vertical calibration (or voltage range) until you can see the trace. Adjust the sweep, sweep vernier and synchronization controls until the trace stabilizes, and looks like this:

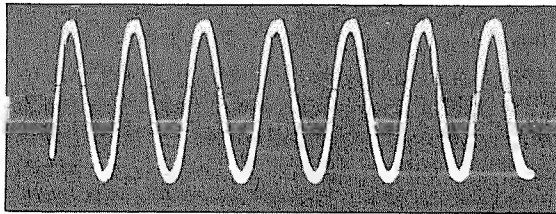


Photo 10-3. Photo of sine wave.

Adjust the vertical and horizontal centering until the trace is in the center of the screen. What you see is your own body acting as a kind

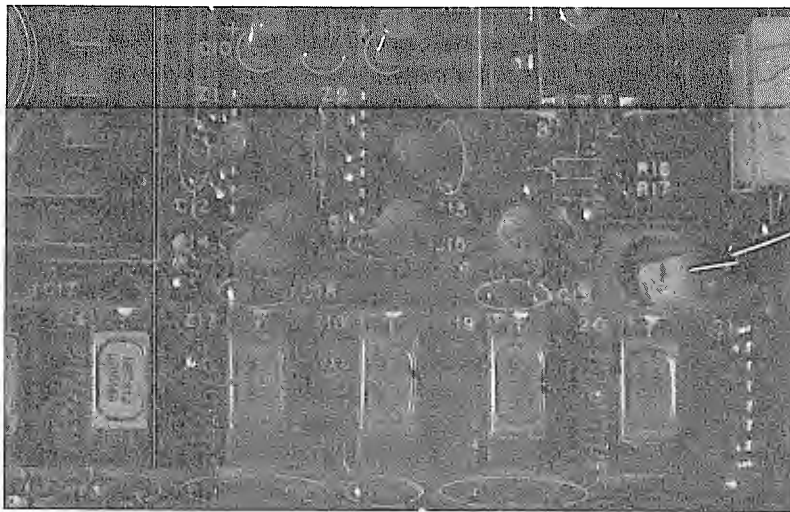


Photo 10-4. Photo of TRS-80 ground point for scope.

of transformer, soaking up all the electrical signals from the wiring and equipment all around you.

Now hold the black ground wire in your other hand. The wild trace should now flatten out considerably, as you become both signal and ground, essentially 'shorting out' the bulk of the signal received by the oscilloscope.

Next, power down your TRS-80 and open it. Using clips or by soldering a wire temporarily into place, attach the black ground wire to the point shown in the photo below:

Hold onto the red cable, and power-up the computer. Now find pin (?) of the Z-80 microprocessor (Z40 on the circuit board), and hold the red probe to it. Adjust the screen so you can see all of the trace, and it consists of a regular series of pulses.

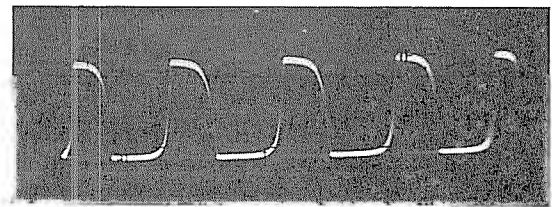
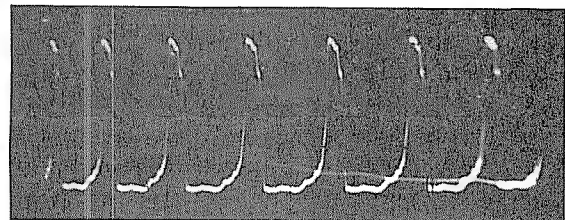


Photo 10-5. 1.77 MHz clock pulses.

What you are looking at is the actual 1,774,000 pulse per second master clock of the TRS-80. All the signals you see throughout the computer should be of this quality – sharp-edged and clean. They will not have the same regularity, because different lines are turned on and off only as needed. But the signals should always be of the same crispness.



For practice, place the red probe at various points on the computer, being very careful not to short two points together with the tip of the probe. You will see the many kinds of signals present throughout the computer. If your computer has a problem, you will be looking for signal lines which have failed: they will be rounded, or angled instead of vertical, or will be much lower in level than the ones you see now. This is a clue to the area of the failure.

The next step is to switch the external black power supply box with a new one, because your present supply may have a damaged, highly resistive, or intermittent fuse wire. Later supplies can be opened to check this; there are screws under the feet. If the new supply works, have the old one repaired, or open it yourself and replace the fine fuse wire with a new one.

Memory problems caused by damaged refresh lines are subtle, but can be discovered. Load a BASIC program, LIST it, and re-LIST it. If gradually the program shows changes, suspect that the refresh lines may be weak. To diagnose this job fully, you will need an oscilloscope.

Photo 10-2 shows how cleanly a good signal is represented on the scope screen. Anything less than a sharp-edged, rectangular pulse should be suspect. Furthermore, a signal that is distinctly lower than others can be suspected.

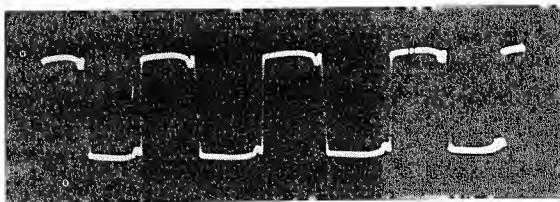


Photo 10-2. Digital scope photo.

Aligning the Video

Video jitter in most cases is due to incorrect settings of either the vertical or horizontal controls on the back of the monitor itself. To align these, enter the following short program:

```
10 CLS
20 FOR X = 15360 TO 16383
30 POKE X,191
40 NEXT
50 GOTO 50
```

Listing 10-1. Video centering routine.

This will paint the screen white. Any bulging on either side of the screen can be cured by adjusting the horizontal control on the back of the monitor. Turn it until the bulging is reduced as much as possible. Instability or rolling is most easily cured by adjusting the vertical control on the back of the monitor. 'Zero' it in by making the screen first roll forward, then back, then forward, and back again until you have what feels like the central position of the control.

Open the TRS-80, and set the program above in motion again. The horizontal and vertical positioning controls are located at the far right of the circuit board as shown in Photo 10-1. Adjust these until the image is properly centered, and drop a very small amount of glue or nail

polish on the control's plastic handle to hold it in that position against the control's body.

Any remaining instability in the screen display can usually be attributed to low voltage present on your house current line, or problems with the TRS-80 power supply (see above).

Routine Maintenance

Microcomputers like the TRS-80 are generally thought of as maintenance free, except for peripherals with moving parts. But in truth, the keyboard unit and expansion box always need some maintenance, all of which is covered throughout this book. In summary:

1. Clean edge card connectors and check cables regularly.
2. Remove keytops (only on units with the older keyboards susceptible to keybounce) and blow out dust and clean contacts regularly.
3. Blow out dust from the unit itself, and give it a shake, to remove any solder balls and splashes that may come loose as the unit gets older.
4. Keep the unit covered when not in use.

Care of Peripherals

Cassette players. Using ordinary rubbing alcohol and cotton swabs, clean the head, capstan and pinch roller to remove all traces of brown tape oxide. Demagnetize the head using a commercial demagnetization cassette, or a head degausser (available at Radio Shack); follow the directions carefully so you don't make the problem worse.

Always handle the player's buttons gently, as snapping them vigorously will not only shorten their life, but create enough vibration to throw the head out of alignment. If the head does get misaligned, or you suspect it is, turn to the details of cassette player maintenance in Chapter 6.

In summary: clean all parts that come in contact with tape, and treat the unit's controls gently.

Printers. Dot-matrix types like the Radio Shack line printers are workhorses, but there are a few cautions. First and foremost, never move or remove paper when the printer is in motion. The printed characters are formed by fine wires striking against the paper, and these wires can be bent if the paper is pulled against them.

Keep ribbons in good condition, because a ribbon with a slight tear can also catch on a head

and damage the dot-matrix wires. Open the printer and blow out dust and dirt regularly; if you are skilled in handling small machinery, lubricate the printer according to the Radio Shack maintenance manual. If you are a klutz, don't try; alignment of printers has to be accurate to 1/100th of an inch.

In summary: keep printers free of dust and ribbons fresh.

Disk Drives. These horrors are terribly sensitive. Like tape machines, the heads need cleaning, but use only head cleaning diskettes especially made for the purpose. If you have regular data read/write problems, there are several solutions: purchase a better disk operating system (DOS), and/or purchase and install an external data separator.

Keep out of the drive's insides as much as possible. Alignment is sometimes vital, but again, unless you are very good with precision machines, don't try it yourself. And even if you are a born watchmaker, be overly careful with double-sided or double-density drives.

In summary: beyond a regular cleaning, leave them alone. Oh yes, keep the drive cables well away from the monitor or any AC power cables.

Exatron Stringy-Floppy. Despite the manufacturer's warnings, this unit does not come well aligned from the factory, and many units were shipped without head alignment cement (normally Lok-Tite, a grey compound) to hold heads firmly in place. Remove the cover of the unit and check for this gop. If it is not there, turn to Chapter 9 and follow the alignment directions.

If the gop is not there, you'll need an oscilloscope. Turn the scope on, and adjust the vertical calibration for full scale (maximum). Hook the scope's ground (black) lead to a convenient ground point on the ESF or the TRS-80, and place the hot (red) probe on Z6 — pin 7. Start the ESF reading a prerecorded tape from the manufacturer. If the wave won't stay in synchronization on the screen, or if it is outside the vertical bounds of the screen, adjust the scope until it is visible and stable. With a small screwdriver, adjust the playback head's alignment until the amplitude (height) of the waveform is at its maximum. Try several prerecorded tapes to make sure of this position, and put some lok-tite on it. If you have your own tapes recorded on the ESF, you may have to move back and forth to and from the correct position, reading in the original position, and re-recording in the corrected one.

Clean the heads and capstan with a very sparing and gentle application of rubbing alcohol on the end of a lint-free swab. Don't use Q-tips or their equivalent, but rather wrap surgical gauze around a lollipop stick. Blow dust out of the unit regularly. Remember never to operate the ESF in any position but on its feet.

In summary: blow out dust, clean heads, and check for the presence of 'Lok-Tite' on the head screws.

RS-232 Board. In another piece of dream engineering, the connection of this board to the expansion box was done via a flukey, solder-coated connector board pressed against a plastic bridge with plain metal contacts. Clean those contacts regularly with a fine cloth, and brush the solder-coated board connectors vigorously with a buffing cloth or extra-fine emery paper. Spray both sparingly with tuner cleaner, and quickly reattach them, constantly rocking the board to assure a firm connection until you have tightened the screws down.

In summary: clean both sets of contacts occasionally, especially if you notice more than the usual transmission/reception errors.

Diagnostic Programs and Loops

Loading large-scale programs for diagnosis is not only time-consuming, but as often as not impossible when your computer is not working well. If BASIC can be brought up at all, then simple machine language diagnostics can be POKEd into place from command level. This section will present several of those loops, their use in trapping some of the possible difficulties, and how they can be used to examine the operation of the system. All these routines can be POKEd anywhere in your machine's memory; these examples all start at 5000 (20480 decimal).

1. Checking the Write Circuits.

This routine merely writes to a location and jumpsback to itself:

```
Coding:
AF          XOR      A
32 00 3C    LO      (3C00),A
3C          INC      A
1B FA      JR      $-4
```

Listing 10-2. Write circuit diagnostic/machine.

```
From BASIC:
X=20480:POKE X,175:POKE X+1,50:POKE X+2,0:POKE X+3,60:
POKE X+4,60:POKE X+5,24:POKE X+6,250 <ENTER>
SYSTEM <ENTER>
/20480 <ENTER>
```

Listing 10-3. Write circuit diagnostic/BASIC.

This routine will write the value in A, which is incremented from 00 to FF (0 to 255) each time the loop is passed through. The write line will pulse each time the LDd (3C00),A instruction is commanded. The first position on the video screen will flicker, as this is the memory location being written to. Press the Reset button to return from this routine.

2. Checking the Read Circuits.

Since each instruction must be fetched, this is only a simple loop:

```
Coding:
18 FE      JR      $
```

Listing 10-4. Read circuit diagnostic/machine.

```
From BASIC:
POKE20480,24:POKE20481,254      <ENTER>
SYSTEM                          <ENTER>
/20480                          <ENTER>
```

Listing 10-5. Read circuit diagnostic/BASIC.

The Read line will pulse four times for every loop through this routine: twice for each instruction fetch, and twice for each refresh action. Press the Reset button to return from this routine.

Listing 10-6. Output circuit diagnostic/machine.

3. Checking the Output Circuits.

This diagnostic is very much like that for examining the Write line, except that it triggers the Out line.

```
Coding:
AF      XOR      A
03 FF      OUT    (FF),A
3C      INC      A
18 F8      JR      $-3
```

```
From BASIC:
X=20480:POKEX,176:POKEX+1,211:POKEX+2,255:POKEX+3,60:
POKEX+4,24:POKEX+5,251      <ENTER>
SYSTEM                      <ENTER>
/20480                      <ENTER>
```

Listing 10-7. Output circuit diagnostic/BASIC.

Each time the loop is passed through, the OUT line will be pulsed once, and the data present on the data lines will be incremented. This routine is also very useful because the cassette relay, the cassette data output, and the video screen will all demonstrate activity as the routine is looped through. Press the Reset button to return from this routine.

4. Checking the Input Circuits.

This routine is similar to the write routine, but does not include any accumulator changes.

```
Coding:
08 00      IN      A,(00)
18 FC      JR      $-2
```

Listing 10-8. Input circuit diagnostic/machine.

```
From BASIC:
X=20480:POKEX,219:POKEX+1,0:POKEX+2,24:
POKEX+3,252      <ENTER>
SYSTEM          <ENTER>
/20480          <ENTER>
```

Listing 10-9. Input circuit diagnostic/BASIC.

Press the Reset button to return from this routine.

5. Checking the HALT Line.

Since the HALT line is gated together with the Reset button, executing HALT should return to ready (or reboot in a disk system).

```
Coding:
76      HALT
```

Listing 10-10. HALT line diagnostic/machine.

```
From BASIC:
POKE20480,118      <ENTER>
SYSTEM            <ENTER>
/20480            <ENTER>
```

Listing 10-11. HALT line diagnostic/BASIC.

Since the machine should return to READY, the Reset button need not be used.

6. Checking the Video.

This routine presents a screen full of characters, all identical, and increments through all 256 of the possible characters. With no lower case modification the screen will display three sets of upper case characters, followed by two sets of graphics characters. With a lower case modification, two upper case sets and one lower case, or one of each case plus a set of control characters will be displayed.

```
Coding:
AF      XOR      A
21 00 3C  LD      HL,3C00
11 01 3C  LD      DE,3C01
01 FF 03  LD      BC,03FF
77      LD      (HL),A
E0 80      LOIR
3C      INC      A
F5      PUSH    AF
01 00 C0  LD      BC,C000
C0 60 00  CALL    0060
F1      POP     AF
18 E9      JR      $-210
```

Listing 10-12. Video routine diagnostic/machine.


```

From BASIC:
X=20480:POKE X,175:POKE X+1,33:POKE X+2,0:POKE X+3,60:
POKE X+4,17:POKE X+5,1:POKE X+6,60:POKE X+7,1:
POKE X+8,255:POKE X+9,3:POKE X+10,119:POKE X+11,237:
POKE X+12,176:POKE X+13,60 <ENTER>
POKE X+14,245:POKE X+15,1:POKE X+16,0:POKE X+17,192
POKE X+18,205:POKE X+19,96:POKE X+20,0:POKE X+21,241
POKE X+22,24:POKE X+23,233 <ENTER>
SYSTEM <ENTER>
/20480 <ENTER>

```

Listing 10-13. Video routine diagnostic/BASIC.

Also, the VID* line (noted on the schematic in the Technical Reference Handbook) should pulse in very noticeable groups as each screen is printed. Use the Reset button to exit this routine.

7. Checking the Cassette Output.

This routine calls the byte-output routine, and should write an FF (equal to eight timing pulses and eight data pulses) to port FF (the cassette output).

```

Coding:
3E FF      LO      A,OFF
C0 35 02   CALL   0264
1B F9      JR      5-5

```

Listing 10-14. Cassette output diagnostic/machine.

```

From BASIC:
X=20480:POKE X,62:POKE X+1,255:POKE X+2,205:POKE X+3,100:
POKE X+4,2:POKE X+5,24:POKE X+6,249 <ENTER>
SYSTEM <ENTER>
/20480 <ENTER>

```

Listing 10-15. Cassette output diagnostic/BASIC.

You should be able to measure (or hear) a constant group of pulses output to the cassette player. Note that the cassette player must be in record position, and must be running (the small plug must be removed) because this routine does not turn the cassette machine on. Use the Reset button to exit this routine.

8. Checking the Printer.

By writing data to address 37E8, the printer should react by printing that character. The following two routines output the letter 'A' to the printer; the first loops through a delay, outputting about five characters per second. The second waits for a printer handshaking signal.

Print-and-Delay Routine

```

Coding:
3E 41      LO      A,41
32 E8 37   LO      (37E8),A
01 00 30   LO      8C,3000
C0 60 00   CALL   0060
1B F3      JR      5-110

```

Listing 10-16. Printer diagnostic I/machine.

```

From BASIC:
X=20480:POKE X,62:POKE X+1,65:POKE X+2,50:POKE X+3,232:
POKE X+4,55:POKE X+5,1:POKE X+6,0:POKE X+7,48:
POKE X+8,205:POKE X+9,96:POKE X+10,0:POKE X+11,24:
POKE X+12,243 <ENTER>
SYSTEM <ENTER>
/20480 <ENTER>

```

Listing 10-17. Printer diagnostic I/BASIC.

Print-and-Wait Routine

```

Coding:
3E 41      LD      A,41
21 E8 37   LD      HL,37E8
77         LD      (HL),A
C8 76      BIT     6,(HL)
20 FE      JR      NZ,$
1B F4      JR      5-100

```

Listing 10-18. Printer diagnostic II/machine.

```

From BASIC:
X=20480:POKE X,62:POKE X+1,65:POKE X+2,33:POKE X+3,232:
POKE X+4,55:POKE X+5,119:POKE X+6,203:POKE X+7,118:
POKE X+8,32:POKE X+9,254:POKE X+10,24:POKE X+11,244 <ENTER>
SYSTEM <ENTER>
/20480 <ENTER>

```

Listing 10-19. Printer diagnostic II/BASIC.

9. Checking the Interrupt Line

For this, of course, interrupt hardware must be attached to the system. In the case of the expansion box, the interrupt flip-flop must be cleared (see Supplement to Chapter 4). This routine performs that function. It is identical to the first service routine presented in the Supplement to Chapter 4, and so only the BASIC coding is presented.

From BASIC, attempt to use this program instead of direct POKEs, because there are so many values:

```

10 FOR X = 20480 TO 20926 : READ A : POKE X,A : NEXT
20 STOP
30 DATA 243,62,195,50,18,64,33,20,50,34,19,64
40 DATA 33,25,26,229,237,86,251,205
50 DATA 243,245,229,213,197,58,236,55
60 DATA 58,224,55,33,17,1,17,37,62,1,26,0
70 DATA 237,176,193,209,225,251,201

```

```

SYSTEM <ENTER>
/20480 <ENTER>

```

Listing 10-20. Interrupt line diagnostic/BASIC.

This routine, as described, returns to BASIC command level and displays a RADIO SHACK LEVEL II BASIC message continuously on the screen.

10. Checking Speed Modifications

Among the possible high speed failures are miswirings, ROMs which are too slow, and RAMs which are too slow. Combinations of these can make diagnosing a locked-up computer very frustrating.

The two routines below test, respectively, RAM calling ROM, and RAM alone. The RAM routine may be moved to higher memory for testing that area; a version at 5000 and B000 are provided.

Coding:

RAM-Resident Version – Origin is at 5000

```

32 60 30      LO      (3060),A
3C            INC     A
F5            PUSH    AF
AF            XOR     A
03 FE        OUT     (FE),A
F1            POP     AF
32 61 30      LO      (3061),A
F5            PUSH    AF
3E FF        LO      A,FF
03 FE        OUT     (FE),A
01 00 10      LO      BC,1000
* 0B          OEC     BC
* 7B          LO      A,B
* B1          OR      C
* 20 FB       DJNZ    FB
F1            POP     AF
% C3 00 50    JP      5000

```

ROM-Resident Version

Remove lines marked (*) and replace with:

```

C0 60 00      CALL    0060

```

High RAM-Resident Version

Origin is at B000; replace line marked (%) with:

```

C3 00 B0      JP      B000

```

For the ROM-RAM version, replace the last group of POKEs with:

```

POKE+21,205:POKE+22,96:POKE+23,0:POKE+24,241:
POKE+25,195:POKE+26,0:POKE+27,B0      <ENTER>

```

For the high RAM version, let X = -20480, and replace the last value POKEd (in either configuration) with 176.

These routines will present, just below the center of the screen, a pair of characters. The first of these is printed in the machine's normal configuration (if OUT 254,0 sets your machine to normal). The second character is printed with the effect of OUT 254,255, which should turn all modifications on. Any failure in these routines should point to the specific area of breakdown.

Listing 10-21. High-speed modification diagnostic machine.

RAM/ROM Version

Remove lines marked (*) and replace with:

High RAM-Resident Version

Origin is at B000; replace line marked with percent sign with:

Since speed changes can affect the operation of BASIC, a version with POKE may be useless; if not, however, here is the RAM-resident version:

```

X=20480:POKE,50:POKE+1,96:POKE+2,61:POKE+3,60:
POKE+4,245:POKE+5,175:POKE+6,211:POKE+7,254:
POKE+8,241:POKE+9,50:POKE+10,97:POKE+11,61:
POKE+12,60:POKE+13,245:POKE+14,62:POKE+15,255:
POKE+16,211:POKE+17,254:POKE+18,1:POKE+19,0:
POKE+20,16      <ENTER>
POKE+21,11:POKE+22,120:POKE+23,177:POKE+24,32:
POKE+25,251:POKE+26,241:POKE+27,195:POKE+28,0:
POKE+29,80      <ENTER>

```

Listing 10-22. High-speed modification diagnostic BASIC.

NOTES

11

111 CURES FOR THE COMMON CRASH

Because TRS-80 computers are sold at the 'appliance' level, salespeople and users often forget that they are indeed computers, prone to all the problems that larger computers face. As a user of a personal computer, you become your own service representative, your own diagnostician, and your own repair agency. This is even more true if you have a custom TRS-80.

This chapter is dedicated to you, the user, as service rep, diagnostic engineer, and repair person.

I. SOFTWARE

Symptoms of software problems: all error messages, system lockup and crash, just about everything that doesn't snap, crackle or pop.

1. Correct the program – the greatest cause of crashes.

By far the largest cause of computer 'crashes' is the software. Mishandling of POKE, PEEK and VARPTR are most common as well as improper input/output routines, and insufficient error-checking.

If crashes seem to occur only in specific programs, check these for errors before turning to hardware cures, although a 'stuck bit' may always turn up a consistent error.

2. Set memory size correctly for machine language or hybrid programs.

Memory size is a boundary above which BASIC is prohibited. BASIC is not merely the source program, but consists of variable and string storage and the BASIC stack. Some of this uses high memory, and when memory size is not set as specified, these variables and stack information can run into the high-memory machine language program.

3. Wait out or re-write long string searches and sorts.

The memory allocation method for strings has made 'garbage collection' techniques necessary. When sorting is complex and the number of strings is large, this process can take from a few minutes to an hour, during which the computer seems locked up. One cure is clearing as much string space as possible; when all variables have been defined, break and PRINT MEM. To the program, add a line such as CLEAR MEM-N, where N is about 100 bytes more than the difference between the

program with all variables defined and your total memory; subtract about 20 bytes from the CLEAR statement for each nested FOR-NEXT loop and GOSUB. Refer to 'BASIC Faster and Better' and other books for techniques of rewriting string handling using variable pointers, and otherwise using memory economically.

4. Read the manuals.

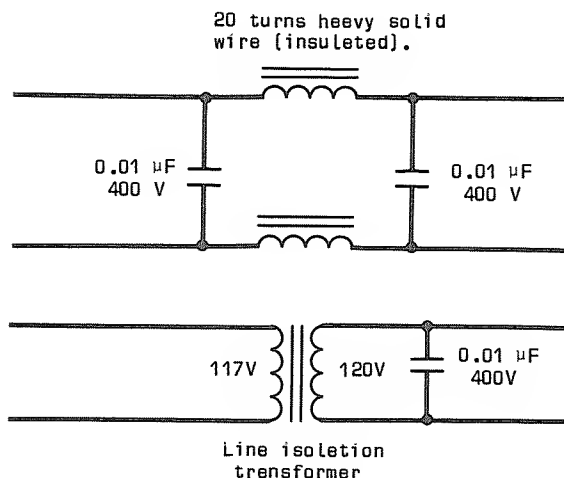
Most bugs are called 'features' by program authors. Discover these features in the program documentation.

I. POWER

Problems Caused by Poor Power: system reboot, unexpected syntax errors, speckles on screen, shrinking screen, unexpected disk or Stringy-Floppy startup, total lockup of system.

5. Add a line filter or power supply monitor.

'Clean' power enhances system reliability. The simplest help is a line filter such as that sold by Radio Shack. More severe power problems can be ameliorated by using power supply monitors such as Topax or Mayday. Isolating other appliances (power tools, washers, refrigerators, toasters, water pumps, fluorescent lamps, etc.) from your computer's outlet will enhance reliable operation. You can even build some very simple filters yourself:



6. Use 3-wire grounded circuitry for all equipment.

3-wire circuitry is electrically quieter and safer. Some equipment, particularly printers, have damaged computers when used in an ungrounded or improperly-grounded environment. If you have an older, 2-wire system, test for floating voltages using a neon test light, available from hardware stores. Reverse the plugs in the socket until all voltages disappear.

7. Stop strong power drains (granite sheds, arc welders).

Strong power drains are visible as brief but distinct shrinking and dimming of the video. Cassette and disk I/O problems are most likely, but power drains can sometimes result in program failure. A high quality power supply monitor can help. Also, if you don't live in an industrial area, your power company or zoning commission might be able to assist.

8. Fuse in the power supply might be going or gone.

All separate TRS-80 power supplies contain fuses in order to receive an Underwriters Laboratories approval. Earlier supplies are sealed, but the seal can be broken by forcing a screwdriver around the joint at the bottom of the supply and gradually popping off the top section. Later supplies have screws under the rubber feet. Fuses are cheap. Don't test the fuse, replace it. It will need to be resoldered since these fuses are not in sockets.

9. Replace or resolder faulty power cables at the connector.

The power connector is not molded in place, and after a year of abuse, the soldered connections can break inside. Lift the tabs on the plastic sheath and pull it back. The connector can be resoldered.

If the power remains intermittent, the problem may be at the point where a band of metal is crimped around the cable. Loosen it, cut the cable shorter, and resolder. The 5-pin DIN plug is sold by Radio Shack if you want to replace it completely.

10. Install a lightning arrestor in spite of the phone company.

If you have a direct-connect modem, lightning strikes may be damaging your system, or at least affecting its performance. For best protection, always unplug the system from the phone and power lines when a storm is in the area. You can also obtain a TII lightning arrestor from Datadyne, 450 Seventh Avenue, NY 10001.

II. EDGE CONNECTORS

Symptoms of edge connector problems: system reboot, unexpected syntax and line number errors, loss of the end of a program or text file, system lockup, return to READY before program end.

11. Clean edge connectors by erasing them.

Since the solder-coated edge connectors are prone to corrosion, they must be cleaned regularly. The simplest methods are vigorous rubbing with a piece of coarse paper like a dollar bill, or erasing with pink pearl (good) or white plastic (best) erasers.

12. Spray edge connectors with contact cleaner and swab with cotton.

Once you have cleaned the edge connectors thoroughly, a weekly application of a small amount of spray contact cleaner (such as Radio Shack 64-2320), followed by rubbing with a cotton swab, will keep the contacts in good shape.

13. Emery-paper edge connectors for really bad corrosion.

For really bad corrosion or scoring from continuous insertion and removal of connectors, the finest grade of emery paper or cloth can be used. Use wet-or-dry paper soaked in contact cleaner or diluted isopropyl alcohol, and rub smoothly along the connectors. Follow up with more contact cleaner and cotton swabbing. Rub just enough to bring up a shine, and do not use coarse emery paper. Never sand down to the copper traces.

14. Silver solder the edge connectors with Silver-It.

Corrosion can be reduced to a minimum while maintaining the original physical size and shape of the edge connectors by obtaining a Silver-It kit from Fuller Products, Grand Prairie, Texas. This is a process that must be completed with great care, but will result in connectors that (except in chemically violent atmospheres in major cities) need virtually no cleaning. For Radio Shack repairs, this process looks like no modifications have been made.

15. Gold plate the edge connectors.

Most difficult of the solutions is gold plating, which is also a poisonous process. However, when it is complete, the result will be connectors of the original physical size with a corrosion-free gold coating. Refer to 80 Microcomputing, December 1981, for full details.

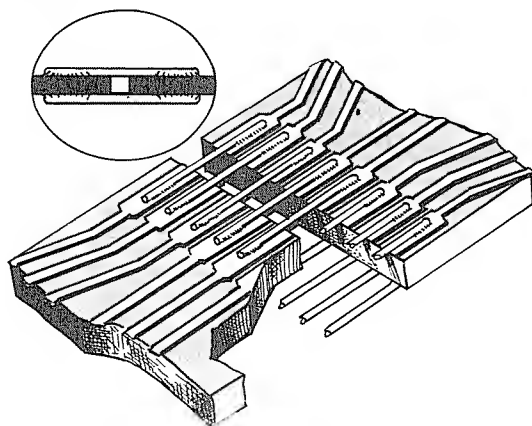
16. Replace the connectors with gold plugs from EAP.

Easier than silvering or gold plating and less annoying than regular cleaning are replacement gold edge connectors (Gold Plug 80) sold by EAP Products, Box 14, Keller, Texas. Because the connectors are made by Kel/Am, they do not mate properly with AP Products connectors. Peri-

peripherals using AP connectors can be easily changed to use the T&B/Ansley connector sold by Radio Shack. Also, the Gold Plug 80 connectors will protrude from the keyboard unit and expansion box somewhat less than an inch.

17. Solder all the connectors into one box and one board.

Model III cases can be purchased as replacement parts through Radio Shack's National Parts distribution system. If you wish to have a one-piece system, this can also cure the edge connector problem. Heavy pieces of copper wire can be soldered to the keyboard and expansion edge connectors, effectively creating a single large board. Don't solder a cable in place and then put one board on top of the other, because electronic noise will be a problem.



18. Both unbuffered and buffered cables are inserted only one way.

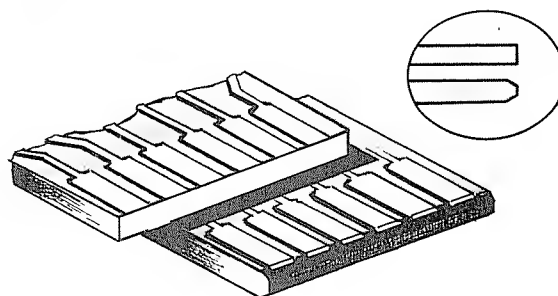
The buffered cable is marked clearly as to which end connects to the keyboard and which end to the expansion interface. However, the unbuffered cable is also directional because it contains a shield which should be connected to ground. If the cable is reversed, the large shield will be connected to a system reset signal — a sure troublemaker. Examine the cable for a fine copper wire sliding out from between the grey plastic 'sandwich'. It should be on the left side of the cable when the cable is attached, with the smooth face of the cable toward the top.

19. Cables not inserted all the way or crooked on connector.

Interconnect cables, particularly AP Products connectors, need to be fully inserted to make electrical contact. Although the expansion cable is not usually inserted partly or crooked, it is easier to do with the printer and disk cables, which are somewhat out of sight. Also, Kel/Am connectors have a limited life — about 50 insertions and removals before the contacts bend inside the connector. Examine the end — all contacts should be even.

20. File carefully to cure a tight fit on the edge connector.

Most Model I TRS-80's contain printed circuit board slightly thicker than that normally used with edge connectors. Although the cable will go on, it has to be forced. To avoid this, smooth the very edge of the circuit board with a fine file, to round the edge. File back not more than 1/8 inch:



III. MEMORY

Symptoms of memory problems: erratic operation, unexpected syntax, unidentified-line, next-without-for and subscript-out-of-range errors, incorrectly reported memory on power-up, program lock-up.

21. Replace slow memory with faster chips, but check speeds.

All dynamic memory has a window during which data is valid for the CPU. Refer to Chapter 4 for a table of memory speeds for keyboard unit and expansion interface. Also, note that memories can deteriorate with age, and poorer access time is one of the first signs of age.

22. Replace bad memory with good memory, using a memory test.

A memory test will reveal memory chips which are genuinely bad, identifying both their address and the bit (chip) involved. Remember that addresses **4000-7FFF** are in the keyboard, and the expansion interface contains two banks, **8000-BFFF** and **C000-FFFF**.

23. Replace older memory with newer (post-1979) memory.

Early dynamic memories were plagued with the occasional 'soft error' — a program crash not caused by any permanent hardware affliction. The error was finally identified as a slow emission of alpha particles from the substrate (base) of the memory circuitry inside the chip. These particles would strike a memory cell, changing it. The occurrence was unpredictable. The best cure for this dilemma is the replacement of all such memories. Date codes are stamped on the chip in year/week format, such as 7851 (end of December 1978); use memories with an 8001 or later date code.

24. Increase access with a modification to Z69.

If the problem seems to be slow or aging memory, the Z69 modification can help. Cut trace from Z69 pin 5. Attach Z69 pin 12 either to Z69 pin 10 or to Z69 pin 13.

25. Install a buffered cable according to modified Radio Shack instructions.

When memory failure seems to occur as more items are added to the system via the expansion interface, then (in earlier units) a buffered cable is called for. It provides additional fan-out (drive) capabilities to the keyboard unit. This item is available from Radio Shack at no charge

to owners of early expansion interfaces. You can identify the earlier interfaces because the 32K memory runs from back to front; in newer boxes it runs along the back. Refer to Chapter 4 for installation instructions.

26. Install the twisted-pair DIN plug modification.

In addition to the buffered cable, some earlier interfaces may require the 'twisted pair' modification. This is a 6-wire DIN cable combination to carry memory select signals RAS, CAS and MUX. Refer to Chapter 4 for installation instructions.

27. Modify twisted pair resistors upward to 470/680 ohm pairs.

If the installation of the twisted pair mod deteriorates rather than ameliorates operation, change the resistor values specified by Radio Shack to 470 ohms to ground, and 680 ohms to 5 volts. Both are originally 220 ohms.

28. Straighten bent pins under socketed ICs.

Memory chips and the CPU were the only chips socketed in earlier machines. However, various runs of the computer had the character generator, line buffers, and so forth, in sockets. Evidence of bent pins is a machine which acts up when given a light physical shock — such as hard typing. Lift each IC and check for bent pins; be sure not to bend any pins when putting it back! If your machine has had a recent trip to salt-water environments, the problem will be more severe; see #104.

IV. FIRMWARE

Symptoms of firmware problems: incorrectly read data, equality failures in IF-THEN statements.

29. POKE 16553, 255 to correct READ/DATA error.

In Version 1 of the Level II ROM, under certain conditions the same data would be read over and over, with the data pointer not being stepped through memory. The most common fix is POKE 16553,255 added as the first line of any program. Also, any INPUT statement before the reading of data begins will take care of properly stepping the pointer.

30. There are floating point accuracy errors, such as X-Y<> X-Y%.

Because of the way real numbers are handled digitally, there is a very small numeric residue left after some types of calculations. Where IF-THEN statements don't seem to work when you know they should, break into the program and print the offending variables. Chances are you will see that a value, instead of being 20 as you expected, is actually 20.00001 as a result of residual binary information. Use integers where you can for such tests; see the Level II manual for details.

V. RS-232

Symptoms of RS-232 problems: incorrectly received or transmitted characters, system crashing or lockup when not using RS-232, electronic failure.

31. Place bar across RS-232 to keep it in place.

Heat buildup and general stress in the expansion box will cause a warping of the RS-232 board, lifting some contacts above the connection pins, or making the unit vibration sensitive. By using longer mounting screws and a heavy insulated metal bar, the warp can be prevented. You may have one made at a local metal shop, then cover it with 'heat-shrink' tubing sold by Radio Shack. Recently, some sources have been making such bars available; look for ads.

32. Clean RS-232 contacts vigorously with cotton cloth.

These contacts too are solder-plated, and prone to corrosion. But they are also very thin, and only a clean cotton cloth - together with vigorous rubbing (and perhaps a little contact cleaner) should be used. Make the contacts shiny.

33. Reseat the RS232 board, checking for bent pins.

The RS-232 board contacts are 1/20 inch from center to center, exactly half the size of the edge connectors. Bent pins on the expansion interface RS-232 connector can cause shorts which will affect not only the operation of the RS-232 system, but the computer as a whole. Reseat the board, checking under strong light that all the connecting pins are straight and contact properly.

34. COM/TERM on RS-232 must be in correct position.

Recheck the RS-232 manual, and make sure that the Communications/Terminal switch is in the correct position for the software you are using.

VI. DISK

Symptoms of disk problems: lost data, hang-up during disk access, rattling and banging of disk mechanism.

35. Use better disks, the best you can get if data is critical.

If you've spent an hour entering data, you've paid for the best diskette you can buy. Purchasing cheap disks is false economy.

36. Install a data separator, doubler, or other separation.

Data is stored on disk as a continuous stream of clock pulses separating data pulses. As a matter of economy, Radio

Shack chose to use 'internal' separation — that is, using the disc controller integrated circuit to distinguish between clock and data pulses. The disc controller's manufacturer does not recommend this method; therefore, a piggyback data separator (sold by several sources) reinstates the proper electronic design. Similarly, a double-density controller contains the essential data separation. If you obtain a 1771B-01 data sheet from Western Digital Corporation, you can build your own data separator.

37. Align the disk drive read/write head professionally.

When a single-drive system is in use, misalignment of a drive head might not be noticed. However, in a multi-drive system, one drive may produce an unusual number of re-seeks or error messages. If you suspect misalignment of the head, don't attempt to service it yourself. You can obtain an alignment diskette (a good one is manufactured by Dysan) to confirm your suspicions, but professional service is called for in this case.

38. Clean the disk drive read/write head with cleaning disks.

When a large number of disks are in use, especially inexpensive ones, tiny bits of oxide are shed onto the drive read/write head. Since this is not immediately visible, it's good practice to obtain a disk cleaning kit (Scotch and others) and use it weekly.

39. Erase-clean connections to the disk drive.

The edge connector to the disk drive is afflicted with the same corrosion problem as other edge connectors. The solutions are the same: cleaning, plating, and replacement with gold plugs. See under edge connectors (#11 above) for more information.

40. Replace 74LS38/LS16 clock generator ICs and socket them.

The cable to disk drives is long and must be driven by higher-current integrated circuits, type 74LS38 (74LS16 in the newer expansion box). These can often break down under continuous use and in situations where the drives are often plugged in incorrectly. The most evident symptom is failure during formatting or backup, because the stepping signal (low to move outward, high to move inward) deteriorates and missteps the read/write head. Remove these chips, and solder in sockets. Then keep a small stock handy so they can be replaced when these symptoms show up.

41. Lubricate disk drive and rails sparingly with silicones.

Many disk drives squeal and clatter. There is no need to risk mechanical failure. Remove the cover and lubricate the motor bearings and the guide rails; use only a very light grade silicone lubricant, and wipe any extra lubricant off once it has spread across the area that needs it. Refer to Exclusive Oracle, 80 Microcomputing, January 1982, for details.

42. Keep the disk door closed till it stops to protect loaded heads.

Your disk drive may contain 'loaded' heads, which means that the door mechanics do not lift the head away from the disk when the door is opened. If you remove or insert a disk before the select light goes out, you may damage the head assembly. You can tell if your drive has loaded heads from the documentation, or by listening for a 'clack' when the drive is selected, and a second 'clack' when the select light goes out. In any case, it is good practice to leave the disk in place in any drive while the motor is spinning.

43. Put disks in correct keyed cable position.

Disk drive selection (drive 0 to 3) can be made in two ways: the drive itself may be programmed with internal jumpers, or the cable may have teeth pulled to eliminate the unwanted select signals. If your drive programming and missing teeth do not match, then the drive will appear dead when selection is attempted. Make sure your drives are marked 0, 1, 2 and 3, and that you match them correctly to the cable.

44. Disk cable must be right side up.

If your disk drive keeps running and doesn't otherwise work, the drive is probably plugged into the cable upside down. Try reversing the connection. Since many drives have cables protruding out the back, and the physical position of these cables is not standard, the plug-in may look correct with respect to the other drives, but be backwards.

45. Disk cable must be fully installed inside case.

The cable protruding from the back of some brands of disk drives is usually not hard-wired, but simply an extension cable from an internal edge card. Under the weight of the long multi-drive cable, or simply from regular moving and unplugging, this internal connector can come loose, resulting in erratic operation. Remove the case top and re-insert the extension cable. A piece of strong plastic packing tape can keep this cable from shifting.

46. Insert the disk correctly into the drive.

This is not as obvious as it seems. If you mix different brand drives on your system, you may discover some in which the disk must be inserted with the write-protect notch pointing down rather than up.

Usually the write-protect notch points toward the side of the drive where the select light is mounted, but even that is not standard. Be especially careful about this when using an unfamiliar system, as a hangup during disk access can be fatal to data in memory.

47. Remove or pad sources of vibration nearby disks.

If you have lived or worked in an area for some time, you have probably blocked out sources of vibration, such as heavy equipment, trains, etc. However, this kind of vibration can ruin a disk write or read. The simplest solution is padding: a layer of heavy cloth, a layer of cork, a layer of wood or metal. If you have occasional inexplicable read/write errors, tune your senses to the environment.

48. Update the DOS, especially if it's an early version of TRSDOS.

Most of the bugs and inconveniences of early versions of TRSDOS have been corrected either by Radio Shack, or by other software houses who have created new disk operating systems. To avoid frustrating errors, update your DOS.

49. Keep disks clean, unbent, and store them straight up in cases.

Though this may seem obvious, there are hidden causes of dirt: smoking, heavy dust or other airborne particles, air freshener sprays, animals, etc. Bending can be caused by storing disks sideways, keeping them in a car window, or inserting them hastily into the drive. Don't drink soda nearby — unless you put your disks right back in their sleeves (which you should do); the bursting bubbles of carbonation can carry sugar residue to the disk surface, damaging the recorded data and abrading the disk head.

VII. TAPE

Symptoms of tape problems: loading impossible, hangup during load, tape won't go on or off, tape won't record.

50. Align the tape recorder head by drilling a hole.

Misalignment is singly the largest tape problem. Put the recorder in playback mode with no cassette in place, and shine a bright light so you can see the Phillips alignment screw to the left of the tape head. Drill a hole directly above it. Align the head by popping in a good commercial music tape, and turning the screw until the sound is at its brightest. Use this for standard recording and playback, but readjust for any commercial tapes that don't sound 'bright'.

51. Clean the tape recorder head with isopropyl alcohol only.

An oxide buildup is common on all tape recorders. Using head cleaner or isopropyl alcohol (not acetone!), swab the tape head and other metal parts which exhibit brown oxide caking. This will prevent scratching or scraping of the tape surface, as well as ensure good contact with the tape head.

52. Demagnetize the tape recorder head with proper devices.

The high frequencies are the most crucial element in good tape loading, and a magnetized tape head erases some of these high frequencies every time the tape is played. Pick up a cassette tape head demagnetizer, either a plug-in type or the kind packaged in a cassette case, and demagnetize the cassette player at least monthly.

53. Use better tape but not the very best audio stuff.

Good tape will always give good loads. Avoid inexpensive tapes like Certron, Concertape, and questionable depart-

ment store house brands. Radio Shack red-label Realistic tape is just fine, as is most any good audio tape. Very high quality tapes (chromium dioxide, metal, etc.) are not necessary, except for archival backups. Digital tapes from Microsette are only about \$.65 and are sold in handy lengths (C-10 and C-20).

54. Modify the CTR-80 with a diode to prevent head glitches.

If the stop button was pressed during loading, a head field collapse in early CTR-80 tape recorders would put glitches on program tapes. Radio Shack provides a free modification for this problem. If you wish to do it yourself, a small silicon diode (such as type 1N4148) can be connected across the tape head contacts.

55. Replace cassette relay to prevent sticking.

The current drain of the CTR-41 tape recorder, and many non-Radio Shack recorders, is too high for the relay contacts in the keyboard unit. This causes it to stick closed, keeping the tape recorder running when it should not. There are two options: change tape recorders, or change relays. The relay change is permanent, and a new unit can be obtained from Lab Service, Inc., in Hustisford, WI.

56. Use CTR-80 to prevent cassette relay burnout.

If you don't want the trouble of replacing a relay, use a CTR-80 tape recorder. It also prevents wear and tear on the cables, because rewind and fast-forward can be used without pulling the motor-control plug.

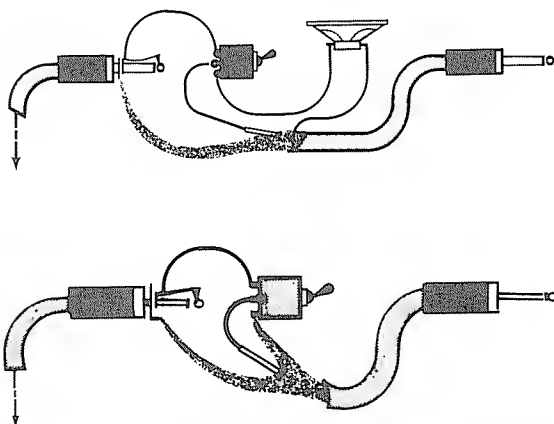
57. Clean the capstan and pinch roller with isopropyl alcohol.

As the tape recorder sees a lot of use, two things will happen to the rubber pinch roller: oxide buildup and glazing. Oxide buildup is similar to that on the tape head,

but on the roller it can cause tape to slip out of place, creasing it. Glazing is a shininess of the rubber surface, also causing tape slip and speed variations. Both can be cured with the liberal application of isopropyl alcohol, acetone, or tape roller cleaner. Hard-glazed rollers need to be brought down with fine emery paper.

58. Add a switch box to the cable assembly for hand operation and sound.

A very useful addition to your tape system is a small switch box containing a speaker; headphone speakers are best. With a simple two-switch box you can listen to programs being saved, keep the motor running, and stop it when you wish:



59. Get the XRX-II cassette modification.

Radio Shack created a special cassette system modification, which is available at no charge. It improves 500-baud loading, but consists of a 500-baud window — which means using cassettes running at any other baud rate (hardware or software based) is impossible. However, the modification can be switched out if you wish to have the advantages of good 500-baud loading and not lose the option of other speeds:

60. Ground loop hum in cassette system is cured by breaking the loop.

Some loading problems can be attributed to hum during program saves. This is due to a 'ground loop' created by the input and

output cables. Pull back the plastic sheath on the computer end of the cassette cable, and cut either one — but not both — of the two shields going to ground (pin 2, the center pin). This will kill ground loop hum.

61. Replace or resolder a faulty cassette cable at the connector.

If you often pull the cassette cables out of the recorder, you may break the internal wiring. This will be evident as a crackling sound during program saves, or partial good loads from tapes which once loaded perfectly. You can replace the connectors with mini and submini plugs, or you can obtain a complete new cable.

62. Replace the dual cassette relay or driver IC.

If the dual-cassette system from the expansion interface ceases to function, either the relay or its associated driver integrated circuit (Z41) may be bad. You can test the operation with:

```
1 A=14308:POKEA,0:FORX=1TO9:NEXT
:POKEA,1:FORX=1TO9:NEXT:GOTO1
```

The relay in the expansion box should clack rapidly.

63. Cassette cables must go to correct expansion box positions.

If one cassette seems to load and the other doesn't, then check that the cassette cables are properly installed. The order (looking from the back) is not particularly logical: Common Cable, Cassette #1, Cassette #0.

VIII. HARDWARE

Symptoms of hardware problems: loss of power, changing of memory contents, keyboard lockup, no apparent expansion interface memory, miscellaneous woes.

64. Tighten the power transistor screw and resolder screw head.

The keyboard unit's power transistor (Q2) was screwed into place, with the screw acting as an electrical connection. This screw can corrode, resulting in erratic operation and frequent complete system crashes. Remove, clean and tighten this screw.

65. Terminate data lines either up or down, but only once.

The eight data lines are flying free in the TRS-80 system. For reliable operation, they should be terminated with resistors — but only once. Since some peripherals contain terminating resistors, the best idea is to obtain an edge connector, and solder resistors to it. Eight resistors (680 ohms, 1/4 watt) go from data lines to ground, eight resistors (470 ohms, 1/4 watt) go from data lines to 5 volts.

66. Add two capacitors to RAM bank the expansion interface.

The absence of two capacitors from the expansion interface has never been explained. There was a place for them, and their inclusion improves operation. Install two electrolytic capacitors (10 mF, 16V) in the positions marked C54 and C55 in the two rows of memory in the expansion interface.

67. Replace the CPU with a Z80A for high speed, or a Z80B for best operation.

If you are running your computer at high speed (100% increase or greater), you are exceeding the formal specifications for the Z80 processor. Although many Z80s can run at that higher speed, certain complex operations (stack operations, for example) can fail. Replacing the chip with a Z80A (4MHz version) or Z80B (6MHz version) will enhance reliability.

68. Replace memory for high speed, but observe DDU wiring changes.

Replacing expansion interface memory for high speed operation may not always help run the machine reliably. Newer expansion boxes contain a digital delay unit for memory selection, which has a fixed access time. Refer to Chapter 4 for wiring changes to the DDU for high speed.

69. Check power supply voltages with a calibrated meter.

Since the miniature voltage adjustment controls were not lacquered, vibration can cause them to move, resulting in the voltages being out of calibration. Units showing as little as 3.2 volts on the 5-volt line have been seen. Obtain a calibrated meter (not a \$10 off-the-shelf number) and adjust the voltages — 12 volts first, then 5 volts. These are R10 (12 volts) and R5 (5 volts) in the keyboard unit.

70. Increase C48 in the expansion interface to keep the disk going.

A discouraging aspect of some disk operating systems is that they cannot recover from a disk drive which has 'timed out' — the motor has stopped. Some DOSes contain a Shift/Break option, but this does not always function, depending on when the disk timed out. A better solution is to increase the value of C48 (C62 in the newer expansion box) from 33 mF to 47 mF or 68 mF. Use a bead tantalum capacitor, not an aluminum electrolytic, if possible.

71. Add or remove disk termination resistors.

In order for disk selection to take place properly, termination resistors are required. Normally, these should only be installed in the furthest drive on the cable. However, as people trade or sell units, or purchase them from a variety of manufac-

turers, the number of termination resistors may vary. These can be found by removing the disk drive cover; they are a red, white or blue integrated-circuit sized package, usually near the edge connector. Make sure only one set of resistors is in place, no matter how many drives are in the system.

72. Make the reset modification to the expansion interface.

It seems the option to turn off a disk system should have been provided with the expansion interface. Ironically, other manufacturers have followed the original unresettable design. You can get out of many program hangups by pressing reset in a Level II keyboard only — so why not with an expansion box attached? Refer to Chapter 4 for details.

73. Change LNW termination resistors to 470/680 pairs.

LNW expansion systems have placed very low termination resistor values in their boards. If too many peripherals (or any peripherals with their own termination) are added to the system, computer lockup will probably occur. For reliability, these values should be changed to 680 ohms to ground, and 470 ohms to 5 volts.

74. Check the DDU in the newer expansion interface.

Sudden failure of memory in new expansion interfaces can almost always be traced to failure in the digital delay unit (DDU), marked Z37. This unit should be slightly warm — neither hot nor cold — to the touch. Since replacement DDU's are about \$20, get some help if you're not sure whether the DDU is bad.

75. Make sure the ROM cable is okay, neither pulled from its sockets, nor with bent pins.

If the computer crashes with an occasional screenful of garbage, or patterns of @9@9,

@A@A, etc., or it fills the screen from the bottom left with A A A , then suspect a fluke Level II ROM cable, or a cracked DIP shunt (see #7 below). Make sure the cable is fully installed. No pins should be bent or broken on the cable; replace it with another 24-pin jumper cable if any pins are bent or broken. (If they're bent, chances are they'll be broken when you try to bend them back).

76. Cold solder joints, splashes, balls, etc., can be anywhere.

This is the worst problem. TRS-80 computer boards are soldered mechanically, and residual solder bits are cleaned away. However, a few balls, splashes or hairs of solder may remain, breaking loose after the vibration of a year of use to cause trouble. Remove all cables and shake out the computer; small solder bits may drop out of the case. Also, broken traces can occur, particularly where any scratches in the green solder mask have occurred. This might be a 'professional help' category.

77. Check DIP shunts for correct or accidental breaks.

Another cause of @9@9, @A@A, @S@S, or the moving A A A A are DIP shunts. Only certain pins (see Chapter 4) should be broken in the DIP shunts (Z3 and Z72 in the keyboard unit). DIP shunts may have hairline cracks; remove them and check with an ohmmeter to be sure. You can replace DIP shunts with DIP switches or ordinary staples.

78. Reset problems are a bad CPU and related capacitor, no disk system in use, or are program-caused.

Users who recently install an expansion interface often forget that the reset button no longer works as it used to, and this is particularly a problem if no disk drives are attached (see #72 above). However, there are three other reasons for reset failure, even if disks are in place: the reset

pin of the Z80 CPU is bad (meaning replace it), or the associated charging capacitor and bleeder resistor (R47 and C42) are bad (replace them), or the program attempts to use the machine language HALT instruction (see Chapter 3).

IX. KEYBOARD

Symptoms of keyboard problems: key-bounce, keys not working, multiple different characters, continuous scrolling?SN or ?S errors on screen, constant repeating memory size question followed by odd characters.

79. Clean keyboard contacts — old one only — not ALPS.

The most straightforward cure for keybounce is cleaning the keyboard. On the bounce-prone keyboards (see photos in Chapter 4), the keycaps lift off. The keys can be cleaned with dry air, spray cleaner (works well but requires cleaning more often), or filled with silicone grease (some people say this improves the feel of the keyboard). The latter technique reduces bounce from vibration as well.

80. Use KBFIX and other software instead of a new keyboard.

A debounce program — automatic on most disk systems — can be loaded at the beginning of each session. Radio Shack's KBFIX is clumsy to get loaded, but works. The debounce routine presented in Chapter 3 works well. Note that new ROMs ('R/S L2 BASIC') contain a debounce routine, and a second routine added to this can result in very slow key response.

81. Get ALPS keyboard, which is the debounced keyboard.

If you can afford the \$75 replacement cost, a Radio Shack Hall-effect ('ALPS') keyboard will permanently take care of bounce problems. If you switch often between BASIC and machine language

programs, play games, or have a variety of disk operating systems or text editing programs, the ALPS keyboard may be the best solution.

82. Replace 74LS05s in the keyboard if multiple or unusual characters appear.

If multiple or unusual combinations of characters appear unexpectedly on the screen, or if characters begin to repeat by themselves, then the 74LS05s on the keyboard baseplate are probably bad. Replace them both, in sockets.

83. Replace or repair the keyboard cable if odd stuff shows up.

Another cause of odd character combinations — usually this like 'FIAQ9' or some such when a single or pair of letters is pressed — is a cracked or intermittent keyboard cable. Replace it; see Chapter 10.

X. VIDEO

Symptoms of video problems: no video, dim video, screen tearing or twitching, blurred screen, screen glitches.

84. Add a buffer stage or resistor to cure video tear.

Lowering the value of R14 in the video monitor will reduce the 'tearing' present when large blocks of graphics are displayed, especially using reverse video. See Chapter 4 for details. The other option is to obtain a video buffer circuit from Archbold Electronics.

85. Add a deglitch modification for a prettier screen.

The video 'hash' created when graphics are being drawn is caused by a conflict between the relatively independent video display circuitry and the need of the CPU to access video memory. Add the deglitch modification for a clearer screen, presented in 80 Microcomputing, Feb. 1982.

86. Get a new character generator, with descenders.

Upper case characters on the TRS-80 have always been consistent, but lower case characters can be displayed either with flying letters (g, p, q and y) or with descenders. The presence of a 'flying a' is normal on early computers. The new character generator with descenders can be obtained as a 'word processing character generator' from Radio Shack.

87. Horizontal and vertical image adjustment can be done.

If the image is not centered on the screen (use

```
10 FORX=15360TO16383:POKEX,191:NEXT
20 GOTO20
```

as a test program), adjust variable resistors R20 and R21 in the keyboard unit.

88. Adjust horizontal and vertical sync on the monitor.

If the image flickers badly, tears sideways, or rolls, the video monitor may be out of adjustment. Turn the white horizontal and vertical adjustment controls on the back of the monitor, just as with an ordinary television that exhibits the same symptoms.

89. Add a capacitor to cure video twitch.

A continuous, annoying screen twitch is the result of oscillations present at the video output. The insertion of a small capacitor (47 to 220 pF) between Z50 pin 3 and ground will eliminate the twitch.

90. Dimmers off! Get rid of the little runners on screen.

Similar to a twitch is the 'runner', a shaking horizontal line that works its way up or down through the screen, rocking one line of letters back and forth. This is a

kind of reverse RFI — not caused by the computer for a change — which can be cured by turning off light dimmers, faulty fluorescent or neon lights, or similar interference producers. If you live in an apartment building, neighboring dimmers should not (but may) affect your computer.

91. Adjust the monitor for blurred characters.

There is no high-voltage adjustment for blurred characters; however, there are a few kludges. First, a higher line voltage (a full 120 volts) will increase the sharpness. Adjusting the internal vertical height control may reduce the image to the more in-focus (center) area of the screen.

Also, replacing the power transistor on the bottom of the chassis and the high-voltage rectifier (they both may exhibit undesirable characteristics) can improve the image. Unfortunately, the monitor is a very basic video display, and has few adjustments.

92. Correct the lowercase software for LDOS.

The 'universal' (alas, a dangerous word) lowercase modification presented in this book is not universal at bootup for one disk operating system — LDOS. However, the LDOS driver can be invoked separately and will work; refer to the documentation.

93. Check the 5 volts or the optoisolator in the monitor.

A hardly existent or dim picture can be caused by insufficient 5 volts into the video cable from the computer (check for a bad connection), or a weak or dying optical isolator on the plug-in card inside the video monitor. First, try another monitor. If the system works with that monitor, resolder the cable connection if necessary; if that does not work, replace the optical isolator.

XI. RFI

Symptoms of RFI (Radio Frequency Interference): herringbone across television screens, complete loss of TV stations, whistling on radio (AM or FM), disconnection of wireless telephones, blackout of shortwave transmissions.

94. Shield the entire system for RFI.

The TRS-80 system is a broad-band interference generator; that is, what it sends out affects all bands of radio and television reception. The interference can be lowered by shielding: spray the inside of the case with aluminum paint and hook that to signal ground; use shielded multi-conductor cable for all peripherals (it's expensive); and create a 'Faraday cage' — if your decor will allow it — by shielding the room in which the computer is used with fine mesh.

XII. HEAT

Symptoms of heat problems: loss or lockup of program consistently after the machine has been on several hours.

95. Ventilate case or add fan, latter if it's an all-in-one system.

Normal ventilation using the slots cut into the TRS-80 and expansion box is adequate. However, if speed modifications, internal memory additions, etc., have been made, the power supply is asked to do extra work. A hard desk with good circulation around the computer is essential, and a small 'Sprite' type fan can be used occasionally to cool the system. A quiet fan may be added for continuous use, and is a necessity for an all-in-one-case system.

96. Remove power supplies from the expansion interface.

The two power supplies in the expansion interface generate a great deal of heat. With the video monitor stacked above,

this can create an undesirably hot environment for memory. Remove the two supplies from the expansion box; the system won't look as compact, but it will have a longer life.

XIII. PRINTER

Symptoms of printer problems: will not backspace or underline, will not move up a line at a time during program listings.

97. Printer may need a line feed with every carriage return.

Many printers require not only a carriage return, but a line feed as well; examples are Teletypes and some Centronics printers. If you have this problem, first check with the manufacturer for a modification kit or instructions. If none is available, use a disk operating system with a CR/LF option, a printer driver patch, or a text-editing system with the CR/LF option. A hardware addition can be made to most printer interfaces to generate a linefeed when a carriage return is received.

XIV. USER PROCEDURE

98. Hide furry animals, keep away from wood stoves, etc.

Insignificant as it may seem, smoke is a severe abrasive to disks. Smoking, wood stoves, unskilled kitchen use (ahem), animal hair, and so forth, can result in airborne particles that affect disks, whose rotation acts as a static vortex to pull in those particles. Similar cautions apply to the piles of hair and grit that can gather in the keyboard. If you wonder just how much dirt is in the air, open the video monitor (with the power off, of course), and have a look. All around the high voltage will be piles of grit pulled in because of the electrical attraction.

99. Don't pull the cables while you're using the system.

This might seem obvious, but realize that pressing the reset button jostles the cable, just as does moving the keyboard to make it comfortable. Dropping a pencil, a cassette, or a diskette case on the cable may generate noise; see Chapter 7 for vibration protection.

100. Wait before power-up after power-down to protect memory.

When the manual cautions to wait ten seconds before repowering the system, it is not merely because the system doesn't always fully reset during that time, but also because the application of power to the memory must be done in this order: -5 volts, 5 volts, 12 volts. The -5 volt line can be lost with too hasty repowering of the system, and memory will be physically damaged.

101. Neither steel wool, nor metal filings, nor metal bits should be nearby.

Using the computer in a home 'shop' is dangerous. Wood bits may cause key-bounce and disk damage, but metal filings may cause the entire system to fail. A buildup of metal dust will decrease local resistance levels to short-circuits, or create unexpected current drains and intermittent operation.

102. No water or drinks nearby or open windows.

Drinking cups can damage disks, water glasses can fall over, carbonated soda can bubble tiny sugar globs onto tapes and disks, and open windows can invite water-carrying breezes and even rainstorms.

103. No magnets nearby, especially unobtrusive refrigerator types.

Refrigerator or bulletin board magnets are handy things, but because they are so

ubiquitous, their danger to magnetic media is easy to forget. If there are any in the house, keep track of them . . . magnets shaped like daisies or fruits, bars, magnetic kiddie letters, even magnetic screwdrivers and scissors.

104. Watch out for salt air areas.

There's nothing like salt air for corroding metal, and metal is the heart of the computer's interconnections. Integrated circuit pins can corrode, cables can corrode, even screw connections can corrode. In boards previously reliable, salt air can corrode unseen bent pins, making the system flakey.

105. Keep telephone bells away, which are bulk erasers.

More than 50 volts shoot into a telephone bell's electromagnet. It's a strong magnetic field, and can act as a bulk eraser for disks sitting under them. Keep phones 'way back on the desk.

106. Keep out of Xray at airports; hand check disks and tape.

Xrays are damaging to magnetic materials. When going through an airport check, keep all your disks and tapes under your arm. Don't check them with your luggage, and don't let them go through the carry-on luggage conveyor belt Xray device. Insist on a hand check for those items.

107. Attach cables correctly.

Make sure all cables are correctly in place. Mark disk cables (the worst culprit) with 0, 1, 2 and 3, as well as 'top' and 'bottom'. Mark the expansion cable the same way (the metal wire or blue stripe is to the left, smooth side up). Mark all peripherals, particularly those which require power from the expansion box. Also, be sure the cassette cables are incorrectly (see #63 above).

108. Attach cables with power off, no matter who says what.

Occasionally, a manual may say something about turning the computer on, then attaching the cables. This is dead wrong. Correctly designed peripherals are always connected with the power off. If it is a construction project, avoid it; there's something wrong with the author's judgment. If it is a commercial product, return it; it shows poor design sense.

109. Disconnect power especially during summer.

High transient voltages can be present over the power and telephone lines during electrical storms. If your area is prone to electrical storms, keep your system unplugged — not just turned off. Also, disconnect direct-connect modems which can carry high voltages straight through to the rest of the system.

110. Obtain a static-free mat for computer and peripherals.

In dry climates, static buildup is common. A static zap can: change memory contents; damage memory; glitch a disk or cassette write; glitch a load; crash a program; reboot the system; or simulate just about any crash your system might be sensitive to. Obtain a static-free mat, or line your table with aluminum foil, and discharge to it — don't touch your computer first.

111. Keep temperature 55 to 80 degrees, relative humidity 50 to 80 percent.

This is simply good practice; although my computer is used in 40 degrees or below, disk reliability is lessened. Dry weather (low humidity) encourages static (see #110), and wet weather encourages corrosion (see #104).

112. Test homebrew devices before installation.

There is very little more to say; always

'proofread' your circuit with a second person, no matter how exhausting it may seem. It will prolong the life of your computer, and the homebrew device might even work the first time.

What are the chances that your problem will be one of these? Probably, the difficulty will be a combination. This list is derived from work on hundreds of TRS-80 system combinations and relatives. Virtually every suggestion has been made and every cure implemented. Sometimes the problems were multiple: one unit suffered from salt air, bent pins, corroded connectors, unbuffered cables, slow memory, and a bum program. Another was the victim of modifications made with a Boy Scout woodburning kit instead of a soldering iron.

Even if your problem is not included here, these suggestions should give you a clue about where and how to begin looking. Use the Radio Shack technical manual diagnostic chart — but don't believe the 'ROM is bad' section, because I've never seen a bad ROM. Instead, suspect a bus driver difficulty (Z22, 38, 39, 55, 75 and 76 in the keyboard). Otherwise, the manual will give you a good start on the tough stuff.

Last Thoughts

The solons at Radio Shack have done something very impressive: they have created a popular personal computer. From it have come the Models II and III, the Pocket Computers I and II, the Color Computers, the Model 16 is on the way, and a host of engineers across the world have been encouraged to come up with TRS-80 compatible hardware and software, and even full computer improvements like the LNW-80.

But they have done something unwittingly even more special: they have, through strict and strange corporate policies, challenge users to create the Custom TRS-80. Because computers become appliances more and more each day, there will only be one Custom TRS-80 . . . the humble Model I.

NOTES

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* indicates discussion of topic
indicates a program listing
\$ indicates a schematic diagram

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Appendix I

Parts Suppliers

The Four Stars

Digi-Key Corporation, Hiway 32 South, P.O. Box 677, Thief River Falls, MN 56701. 800 346-5144. *COD, check, money order, credit cards. Volume discounts over \$100; shipping, insurance prepaid.*

This company is in my opinion the hobbyist's best. Shipping is fast (five days from ordering to my door in Vermont), and most items are in stock. Their catalog is monthly, and items not stocked are not listed. All merchandise is prime; no bubble packs.

QT Computer Systems, 15335 S. Hawthorne Blvd., Lawndale, CA 90260. 800 421-5150. *COD under \$100, check, money order; credit cards preferred. Quantity discounts, no insurance.*

A good hobbyist catalog similar to Digi-Key, with competitive prices. This is a new company, but they are already beginning to make a mark for promptness, exceeding courtesy, and prime parts. Their catalog is very complete and quite up-to-date.

Advanced Computer Products, 1310 E. Edinger, Santa Ana, CA 92705. 800 854-8230. *COD, check, money order, credit cards. Volume discounts, no insurance.*

One of the best catalogs in the business, prompt, but be wary of substitutions in orders. Specify voltages of devices and check upon receipt. Expect harassment from Customer Service. Otherwise, they have what you can't get anywhere else.

Jameco Electronics, 1355 Shoreway Road, Belmont, CA 94002. 415 592-8097. *COD, check, money order, no credit cards. No discounts, no insurance.*

One almost wonders why to put Jameco in with the four best, but their selection is contemporary and their response prompt. They have items others don't have in stock for the popular computer hobbyist. Bubble pack stuff on retail store racks at Lafayette Radio and others. Highest prices in the business.

And Others

Jade Computer Products, 4901 West Rosecrans Ave., Hawthorne, CA 90250. 800 421-5500. *No CODs; checks, money order; credit cards preferred. Quantity discounts; insurance under 50 lbs.*

This company works hard at immediate hobbyist needs and some unusual items. Get their catalog, but consult monthly ads in electronics magazines for hot items.

Priority One Electronics, 16723C Roscoe Blvd., Sepulveda, CA 91343. 800 423-5633. *No CODs; check, money order, credit cards. Quantity discounts, insurance.*

Priority deals for the most part in larger items for computer hobbyists, with only a token selection of small parts. This company concentrates on boards and naked disk drives, and heavier hardware.

Hobbyworld Electronics, 19511 Business Center Dr., Northridge, CA 91324. 800 423-5387, (800 382-3651 in CA). *COD (\$1.25 extra), check, credit cards. No discounts, no insurance.*

Hobbyworld is the computer hobbyist's pop culture. It stocks all the hot items with a high turnover. Look to them for low prices on items you need right away.

Electrolabs, P.O. Box 6721, Stanford, CA 94305.

The best part is always their funny and schizophrenic catalog with an honest selection and a wealth of good information.

For example . . .

"Save yourselves \$6.75 and use a 25 cent transistor the next time your looking for a temperature probe." Also, TTL Family rules of Incest are great.) The shipping was always prompt and the merchandise prime.

Not Recommended

Active Electronic Sales, P.O. Box 1035, Framingham, MA 01701. 617 879-0077. *Minimum \$10, handling \$2, check (wait to clear), money order. No discounts, no insurance.*

This group claims to be "The World's Largest International Semiconductor Distributor", which implies lots of stock, in stock. No way. All my orders have been returned 25 percent filled, with 50 percent errors.

Appendix II. Bibliography.

Reference Books

Manuals, Guides, and Data Books

Radio Shack, Tandy Corporation. Fort Worth, Texas.

TRSDOS & Disk BASIC Reference Manual, 1979.

TRS-80 Micro Computer Technical Reference Handbook, 1978.

Printer Interface Cable, Service Manual.

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National Semiconductor, 2900 Semiconductor Drive, Santa Clara, California 95051.

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BYTE. 70 Main St., Peterborough, New Hampshire 03458. \$19 per year.

onComputing. 70 Main St., Peterborough, New Hampshire 03458. \$12 per year.

80-U.S. Journal. 3838 South Warner Street, Tacoma, Washington 98409. \$24 per year.

Softside. 6 South Street, Milford, New Hampshire 03055. \$24 per year.

Popular Electronics. One Park Avenue, New York, New York 10016. \$14 per year.

The Alternate Source. 1806 Ada Street, Lansing, Michigan 48910. \$12 per year.

Radio Electronics. 200 Park Avenue South, New York, New York 10003. \$9.98 per year.

CLOAD. P.O. Box 1267, Goleta, California 93017. \$42 per year, on cassette.

Fairfield County Users Group, *Voice of the 80*. C/O Alan Abrahamson, 10 Richlee Road, Norwalk, Connecticut 06851.

Marin County TRS-80 Users Group Newsletter. P.O. Box 895, Novato, California 94948.

Computer Base Lubbock. C/O Roger Smith, 2601 Nonesuch Dr., Lot 1802, Abilene, Texas 79606.

TCS, Club Project of the Tidewater TRS-80 Users Group. P.O. Box 10281, Norfolk, Virginia 23513.

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NOTES

APPENDIX 3

Byte Values and Their Equivalents

Code		PDKE displays:			PRINT:	Edtesm	Stendrd
Hex Val	Dec Val	Without L/Case	With Dld LC	With New LC	w/ LC Driver	Z8D Dpcode	ASCII Meaning
00	0	@	gaming	@	none	NDP	NUL
01	1	A	gaming	A	none	LD BC,NN	SOH
02	2	B	gaming	B	none	LD (BC),A	STX
03	3	C	gaming	C	none	INC BC	ETX
04	4	D	gaming	D	none	INC B	EDT
05	5	E	gaming	E	none	DEC B	ENQ
06	6	F	gaming	F	none	LD B,N	ACK
07	7	G	gaming	G	none	RLCA	BEL
08	8	H	gaming	H	bkspace	EX AF,AF'	BS
09	9	I	gaming	I	bkspace lin	ADD HL,BC	HT
0A	10	J	gaming	J	linefeed	LD A,(BC)	LF
0B	11	K	gaming	K	none	DEC BC	VT
0C	12	L	gaming	L	formfeed	INC C	FF
0D	13	M	gaming	M	car retn	DEC C	CR
0E	14	N	gaming	N	crsr on	LD C,N	SD
0F	15	O	gaming	D	crsr off	RRCA	SI
10	16	P	gaming	P	none	DJNZ d	DLE
11	17	Q	gaming	Q	none	LD DE,NN	DC1
12	18	R	gaming	R	none	LD (DE),A	DC2
13	19	S	gaming	S	none	INC DE	DC3
14	20	T	gaming	T	none	INC D	DC4
15	21	U	gaming	U	none	DEC D	NAK
16	22	V	gaming	V	none	LD D,N	SYN
17	23	W	gaming	W	widemode	RLA	ETB
18	24	X	gaming	X	bkspace crsr	JR d	CAN
19	25	Y	gaming	Y	adv crsr	ADD HL,DE	EM
1A	26	Z	gaming	Z	dn linfd	LD A,(DE)	SUB
1B	27	l brace	gaming	up err	up linfd	DEC DE	ESC
1C	28	dn arr	gaming	dn err	home crsr	INC E	FS
1D	29	r brace	gaming	left err	strt crsr	DEC E	GS
1E	30	rt arr	gaming	rt err	erese lin	LD E,N	RS
1F	31		gaming		clr frame	RRA	US
20	32	space	space	space	space	JR NZ,d	space
21	33	!	!	!	!	LD HL,NN	!
22	34	"	"	"	"	LD (NN),HL	"
23	35	#	#	#	#	INC HL	#
24	36	\$	\$	\$	\$	INC H	\$
25	37	%	%	%	%	DEC H	%
26	38	&	&	&	&	LD H,N	&
27	39	'	'	'	'	DAA	'
28	40	((((JR Z,d	(
29	41))))	ADD HL,HL)
2A	42	*	*	*	*	LD HL,(NN)	*
2B	43	+	+	+	+	DEC HL	+
2C	44	,	,	,	,	INC L	,
2D	45	-	-	-	-	DEC L	-
2E	46	LD L,N	.
2F	47	/	/	/	/	CPL	/
30	48	0	D	D	D	JR NC,d	D

Code		POKE displays:			PRINT:	Edtasm	Standrd
Hex Val	Dec Val	Without L/Case	With Old LC	With New LC	w/ LC Driver	Z8D Opcode	ASCII Meaning
31	49	1	1	1	1	LD SP,NN	1
32	50	2	2	2	2	LD (NN),A	2
33	51	3	3	3	3	INC SP	3
34	52	4	4	4	4	INC (HL)	4
35	53	5	5	5	5	DEC (HL)	5
36	54	6	6	6	6	LD (HL),N	6
37	55	7	7	7	7	SCF	7
38	56	8	8	8	8	JR C,d	8
39	57	9	9	9	9	ADD HL,SP	9
3A	58	:	:	:	:	LD A,(NN)	:
3B	59	;	;	;	;	DEC SP	;
3C	60	<	<	<	<	INC A	<
3D	61	=	=	=	=	DEC A	=
3E	62	>	>	>	>	LD A,N	>
3F	63	?	?	?	?	CCF	?
40	64	@	open quote	@	@	LD B,B	@
41	65	A	A	A	A	LD B,C	A
42	66	B	B	B	B	LD B,D	B
43	67	C	C	C	C	LD B,E	C
44	68	D	D	D	D	LD B,H	D
45	69	E	E	E	E	LD B,L	E
46	70	F	F	F	F	LD B,(HL)	F
47	71	G	G	G	G	LD B,A	G
48	72	H	H	H	H	LD C,B	H
49	73	I	I	I	I	LD C,C	I
4A	74	J	J	J	J	LD C,D	J
4B	75	K	K	K	K	LD C,E	K
4C	76	L	L	L	L	LD C,H	L
4D	77	M	M	M	M	LD C,L	M
4E	78	N	N	N	N	LD C,(HL)	N
4F	79	D	O	D	D	LD C,A	O
50	80	P	P	P	P	LD D,B	P
51	81	Q	Q	Q	Q	LD D,C	Q
52	82	R	R	R	R	LD D,D	R
53	83	S	S	S	S	LD D,E	S
54	84	T	T	T	T	LD D,H	T
55	85	U	U	U	U	LD D,L	U
56	86	V	V	V	V	LD D,(HL)	V
57	87	W	W	W	W	LD D,A	W
58	88	X	X	X	X	LD E,B	X
59	89	Y	Y	Y	Y	LD E,C	Y
5A	90	Z	Z	Z	Z	LD E,D	Z
5B	91	l.brkt.	l.brkt.	up arr.	up arr.	LD E,E	l.brkt.
5C	92					LD E,H	slant
5D	93					LD E,L	r.brkt.
5E	94					LD E,(HL)	carat
5F	95					LD E,A	l.arr.
60	96	@	open quote	pound	@	LD H,B	undef.
61	97	A	a	a	a	LD H,C	a
62	98	B	b	b	b	LD H,D	b
63	99	C	c	c	c	LD H,E	c
64	100	D	d	d	d	LD H,H	d
65	101	E	e	e	e	LD H,L	e

Code		PDKE displays:			PRINT:	Edtasm	Standrd
Hex Val	Dec Val	Without L/Case	With Dld LC	With New LC	w/ LC Driver	ZBD Dpcode	ASCII Meaning
66	102	F	f	f	f	LD H,(HL)	f
67	103	G	g	g	g	LD H,A	g
68	104	H	h	h	h	LD L,B	h
69	105	I	i	i	i	LD L,C	i
6A	106	J	j	j	j	LD L,D	j
6B	107	K	k	k	k	LD L,E	k
6C	108	L	l	l	l	LD L,H	l
6D	109	M	m	m	m	LD L,L	m
6E	110	N	n	n	n	LD L,(HL)	n
6F	111	D	o	o	o	LD L,A	o
7D	112	P	p	p	p	LD (HL),8	p
71	113	Q	q	q	q	LD (HL),C	q
72	114	R	r	r	r	LD (HL),D	r
73	115	S	s	s	s	LD (HL),E	s
74	116	T	t	t	t	LD (HL),H	t
75	117	U	u	u	u	LD (HL),L	u
76	118	V	v	v	v	HALT	v
77	119	W	w	w	w	LD (HL),A	w
78	120	X	x	x	x	LD A,B	x
79	121	Y	y	y	y	LD A,C	y
7A	122	Z	z	z	z	LD A,D	z
78	123					LD A,E	L,brace
7C	124					LD A,H	separator
7D	125					LD A;L	r,brace
7E	126					LD A,(HL)	wave
7F	127					LD A,A	delete

Upper Character Set

HEX Val	DEC Val	POKE Display:	BASIC Keyword:	PRINT Display:	ZBD Dpcode:
80	128	/G/	END	/G/	ADD A,B
81	129	/G/	FDR	/G/	ADD A,C
82	130	/G/	RESET	/G/	ADD A,D
83	131	/G/	SET	/G/	ADD A,E
84	132	/G/	CLS	/G/	ADD A,H
85	133	/G/	CMD	/G/	ADD A,L
86	134	/G/	RANDDM	/G/	ADD A,(HL)
87	135	/G/	NEXT	/G/	ADD A,A
88	136	/G/	DATA	/G/	ADC A,8
89	137	/G/	INPUT	/G/	ADC A,C
8A	138	/G/	DIM	/G/	ADC A,D
8B	139	/G/	READ	/G/	ADC A,E
8C	140	/G/	LSET	/G/	ADC A,H
8D	141	/G/	GDTD	/G/	ADC A,L
8E	142	/G/	RUN	/G/	ADC A,(HL)
8F	143	/G/	IF	/G/	ADC A,A
90	144	/G/	RESTORE	/G/	SUB B
91	145	/G/	GOSUB	/G/	SUB C
92	146	/G/	RETURN	/G/	SUB D
93	147	/G/	REM	/G/	SUB E

Upper Character Set

HEX	OEC	POKE	BASIC	PRINT	ZBO
Val	Val	Display:	Keyword:	Display:	Opcode:
94	14B	/G/	STOP	/G/	SUB H
95	149	/G/	ELSE	/G/	SUB L
96	150	/G/	TRON	/G/	SUB (HL)
97	151	/G/	TROFF	/G/	SUB A
9B	152	/G/	DEFSTR	/G/	SBC A,B
99	153	/G/	DEFINT	/G/	SBC A,C
9A	154	/G/	DEFSNG	/G/	SBC A,D
9B	155	/G/	DEFOBL	/G/	SBC A,E
9C	156	/G/	LINE	/G/	SBC A,H
90	157	/G/	EOIT	/G/	SBC A,L
9E	15B	/G/	ERROR	/G/	SBC A,(HL)
9F	159	/G/	RESUME	/G/	SBC A,A
A0	160	/G/	OUT	/G/	AND B
A1	161	/G/	ON	/G/	AND C
A2	162	/G/	OPEN	/G/	AND O
A3	163	/G/	FIEL0	/G/	AND E
A4	164	/G/	GET	/G/	AND H
A5	165	/G/	PUT	/G/	AND L
A6	166	/G/	CLOSE	/G/	AND (HL)
A7	167	/G/	LOAO	/G/	AND A
AB	16B	/G/	MERGE	/G/	XOR B
A9	169	/G/	NAME	/G/	XOR C
AA	170	/G/	KILL	/G/	XOR D
AB	171	/G/	LSET	/G/	XOR E
AC	172	/G/	RSET	/G/	XOR H
AO	173	/G/	SAVE	/G/	XOR L
AE	174	/G/	SYSTEM	/G/	XOR (HL)
AF	175	/G/	LPRINT	/G/	XOR A
B0	176	/G/	DEF	/G/	OR B
B1	177	/G/	POKE	/G/	OR C
B2	17B	/G/	PRINT	/G/	OR O
B3	179	/G/	CONT	/G/	OR E
B4	1B0	/G/	LIST	/G/	OR H
B5	1B1	/G/	LLIST	/G/	OR L
B6	1B2	/G/	ODELETE	/G/	OR (HL)
B7	1B3	/G/	AUTO	/G/	OR A
BB	1B4	/G/	CLEAR	/G/	CP B
B9	1B5	/G/	CLOAO	/G/	CP C
BA	1B6	/G/	CSAVE	/G/	CP D
BB	1B7	/G/	NEW	/G/	CP E
BC	1B8	/G/	TAB(/G/	CP H
BO	1B9	/G/	TO	/G/	CP L
BE	190	/G/	FN	/G/	CP (HL)
BF	191	/G/	USING	/G/	CP A
CO	192	/G/	VARPTR	/G/	RET NZ
C1	193	/G/	USR	TAB+01	POP BC
C2	194	/G/	ERL	TAB+02	JP NZ,NN
C3	195	/G/	ERR	TAB+03	JP NN
C4	196	/G/	STRING\$	TAB+04	CALL NZ,NN
C5	197	/G/	INSTR	TAB+05	PUSH BC
C6	19B	/G/	POINT	TAB+06	AOO A,N
C7	199	/G/	TIME\$	TAB+07	RST 00H

Upper Character Set

HEX Val	OEC Val	POKE Display:	BASIC Keyword:	PRINT Display:	ZBO Dpcode:
CB	200	/G/	MEM	TAB+DB	RET Z
C9	201	/G/	INKEY\$	TAB+D9	RET
CA	202	/G/	THEN	TAB+1D	JP Z, NN
CB	203	/G/	NOT	TAB+11	<Note 1>
CC	204	/G/	STEP	TAB+12	CALL Z, NN
CD	205	/G/	+	TAB+13	CALL NN
CE	206	/G/	-	TAB+14	AOC A, N
CF	207	/G/	*	TAB+15	RST OBH
00	208	/G/	/	TAB+16	RET NC
01	209	/G/	**	TAB+17	POP OE
02	210	/G/	AND	TAB+18	JP NC, NN
03	211	/G/	OR	TAB+19	OUT (N), A
04	212	/G/	>	TAB+20	CALL NC, NN
05	213	/G/	=	TAB+21	PUSH DE
06	214	/G/	<	TAB+22	SUB N
D7	215	/G/	SGN	TAB+23	RST 10H
0B	216	/G/	INT	TAB+24	RET C
D9	217	/G/	ABS	TAB+25	EXX
0A	218	/G/	FRE	TAB+26	JP C, NN
0B	219	/G/	INP	TAB+27	IN A, (N)
0C	220	/G/	POS	TAB+2B	CALL C, NN
0D	221	/G/	SQR	TAB+29	<Note 2>
DE	222	/G/	RND	TAB+3D	SBC A, N
0F	223	/G/	LOG	TAB+31	RST 1BH
E0	224	/G/	EXP	TAB+32	RET PD
E1	225	/G/	COS	TAB+33	PDP HL
E2	226	/G/	SIN	TAB+34	JP PD, NN
E3	227	/G/	TAN	TAB+35	EX (SP), HL
E4	22B	/G/	ATN	TAB+36	CALL PO, NN
E5	229	/G/	PEEK	TAB+37	PUSH HL
E6	230	/G/	CVI	TAB+3B	AND N
E7	231	/G/	CVS	TAB+39	RST 2DH
E8	232	/G/	CVD	TAB+40	RET PE
E9	233	/G/	EOF	TAB+41	JP (HL)
EA	234	/G/	LOC	TAB+42	JP PE, NN
EB	235	/G/	LOF	TAB+43	EX DE, HL
EC	236	/G/	MKI\$	TAB+44	CALL PE, NN
ED	237	/G/	MKS\$	TAB+45	<Note 3>
EE	23B	/G/	MKO\$	TAB+46	XDR N
EF	239	/G/	CINT	TAB+47	RST 2BH
F0	240	/G/	CSNG	TAB+48	RET P
F1	241	/G/	COBL	TAB+49	PDP AF
F2	242	/G/	FIX	TAB+50	JP P, NN
F3	243	/G/	LEN	TAB+51	DI
F4	244	/G/	STR\$	TAB+52	CALL P, NN
F5	245	/G/	VAL	TAB+53	PUSH AF
F6	246	/G/	ASC	TAB+54	DR N
F7	247	/G/	CHR\$	TAB+55	RST 3DH
FB	248	/G/	LEFT\$	TAB+56	RET M
F9	249	/G/	RIGHT\$	TAB+57	LD SP, HL
FA	250	/G/	MIO\$	TAB+5B	JP M, NN
FB	251	/G/	_____	TAB+59	EI

Upper Character Set					
HEX	DEC	PDKE	BASIC	PRINT	ZBD
Val	Val	Display:	Keyword:	Display:	Dpcode:
FC	252	/G/	----	TAB+6D	CALL M,NN
FD	253	/G/	----	TAB+61	<Note 4>
FE	254	/G/	----	TAB+62	CP N
FF	255	/G/	----	TAB+63	RST 38H

Note 1: CB includes RLC, RRC, RL, RR, SLA, SRL, BIT, RES, SET.

Note 2: DD includes IX register manipulations.

Note 3: ED includes miscellaneous IN, OUT, LD, and ADC functions as well as block move and compare commands, and interrupt mode setting and handling.

Note 4: FD includes IY register manipulations.

Graphics Note: In the table above /G/ means a graphics character will be displayed. The table below shows the graphics characters and their ASCII code in decimal.

128	129	130	131
132	133	134	135
136	137	138	139
140	141	142	143
144	145	146	147
148	149	150	151
152	153	154	155
156	157	158	159
160	161	162	163
164	165	166	167
168	169	170	171
172	173	174	175
176	177	178	179
180	181	182	183
184	185	186	187
188	189	190	191

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00100 ;
00110 ;
00120 ; ***** K E E P I T 3 . 2 *****
00130 ;
00140 ; >>>> SYSTEM /12345 TO ENTER <<<<
00150 ;
00160 ; >>>> USE *READ FOR MENU <<<<
00170 ;
00180 ;
00190 ; COPYRIGHT (C) 1980, 1981 BY DENNIS BATHORY KITSZ
00200 ; ALL RIGHTS RESERVED. NO PART OF THIS PROGRAM MAY
00210 ; BE REPRODUCED BY ANY MEANS, ELECTRONIC, ELECTRO-
00220 ; MECHANICAL, OR IN PRINT, WHETHER BY METHODS IN
00230 ; USE OR CONCEIVED IN THE FUTURE, WITHOUT SPECIFIC
00240 ; WRITTEN PERMISSION OF THE AUTHOR.
00250 ;
4099 00260 INKEYS EQU 4099H ;INKEY$ BYTE STORAGE AREA
03FB 00270 KYSCAN EQU 03FBH ;KEYBOARD SCAN ROUTINE
4030 00280 PORTFF EQU 4030H ;CASSETTE OUTPUT PORT
401A 00290 KPLACE EQU 401AH ;1-BYTE KEYSTROKE STORE
4019 00300 SHIFTR EQU 4019H ;STORAGE FOR LC DRIVER
4004 00310 RESTRT EQU 4004H ;BASIC INTERP PATCH POINT
40A0 00320 STGEN0 EQU 40A0H ;END OF STRINGS POINTER
40E8 00330 STACKR EQU 40E8H ;BASIC STACK POINTER
40F0 00340 VARENO EQU 40F0H ;END OF VARIABLE POINTER
06CC 00350 READY EQU 06CCH ;RETURN TO READY INTACT
40B1 00360 MENTOP EQU 40B1H ;TOP OF BASIC MEMORY
4000 00370 SETPTS EQU 4000H ;SETPOINT RESTART ADDRESS
0264 00380 WRB8YT EQU 0264H ;ROM BYTE WRITE ROUTINE
1078 00390 BYTE EQU 1078H ;ROM READ KEYS & TOKENIZE
3C00 00400 VIDEO EQU 3C00H ;FIRST SCREEN LOCATION
00410 ;
3039 00420 ORG 3039H
00430 ; PREPARE RAM AREAS FOR USE
3039 2AB140 00440 START LD HL,(40B1H)
303C E058A040 00450 LD OE,(40A0H)
3040 0607 00460 LD B,7
3042 2B 00470 DEC HL
3043 18 00480 DEC DE
3044 10FC 00490 DJNZ $-2
3046 22B140 00500 LO (40B1H),HL
3049 ED53AD40 00510 LD (40A0H),DE
3040 23 00520 INC HL
304E E5 00530 PUSH HL
304F FDE1 00540 PDP IY
00550 ; READY INTERPRETER VECTOR PATCH
00560 ; GET CURRENT CONTENTS OF 4004H
3051 3EC3 00570 LD A,0C3H
3053 F07700 00580 LD (IY+0),A
3056 ED580440 00590 LD DE,(RESTRT)
305A F073D1 00600 LD (IY+1),E
3050 F07202 00610 LD (IY+2),0
3060 F077D3 00620 LD (IY+3),A
3063 117631 00630 LD OE,SKPSTP
3066 F07304 00640 LD (IY+4),E
3069 F072D5 00650 LD (IY+5),D
306C 23 00660 GRINS INC HL
306D 23 00670 INC HL
306E 23 00680 INC HL
306F 220440 00690 LD (RESTRT),HL
3072 AF 00700 XOR A
3073 321940 00710 LD (SHIFTR),A
00720 ; CALL CLEAR AND RETURN TO BASIC
3076 C0C901 00730 RSLVII CALL 01C9H
3079 C0611B 00740 CALL 1661H
307C 211101 00750 LD HL,0111H
307F C0A72B 00760 CALL 2BA7H
3082 C3CC06 00770 JP READY
00780 ; START KEYBOARD SCAN
3085 213640 00790 KBPFIX LO HL,4036H
308B 01013B 00800 LD BC,3B01H
308B 1600 00810 LD D,D
00820 ; CHECK EACH ROW OF KEYS
308D 0A 00830 KEYPRS LD A,(BC)
308E 5F 00840 LD E,A
308F A3 00850 AND E
3090 2D1B 00860 JR NZ,STROKE
3092 77 00870 LD (HL),A
00880 ; INC AND ROTATE TO CHECK NEXT ROW
3093 14 00890 RECHK INC 0
3094 2C 00900 INC L
309S C801 00910 RLC C
3097 79 00920 LD A,C
00930 ; CHECK IF LAST ROW [NOT INCL. SHIFT]
309B 06B0 00940 SUB 80H
309A 20F1 00950 JR NZ,KEYPRS
00960 ; CHECK IF KEYBOARD CLEAR
309C 0607 00970 LD B,7
309E 20 00980 CLRMEM DEC L
309F B6 00990 ADO A,(HL)
30AD 10FC 01000 DJNZ CLRMEM
30A2 A7 01010 AND A

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30A3 3E00 01020 LD A,0
30AS C0 01030 RET NZ
01040 ; RESET AUTOREPEAT DELAY TO 0
30A6 321A40 01050 LD (KPLACE),A
30A9 C9 01060 RET
01070 ; IF KEYSTROKE FOUND CHECK AUTOREPEAT LOOP
30AA A6 01080 STROKE AND (HL)
30AB 2B1E 01090 JR Z,FOUND
30AD 3A9940 01100 LD A,(INKEYS)
30B0 A7 01110 AND A
30B1 20E0 01120 JR NZ,RECHK
30B3 3A1A40 01130 LD A,(KPLACE)
30B6 3C 01140 INC A
30B7 321A40 01150 LD (KPLACE),A
30BA FEFF 01160 CP OFFH
30BC 2B0B 01170 JR Z,OECA
30BE CS 01180 PUSH BC
30BF 06FF 01190 LD B,OFFH
30C1 10FE 01200 THWSTE DJNZ THWSTE
30C3 C1 01210 POP BC
30C4 1BC0 01220 JR RECHK
30C6 30 01230 OECA DEC A
30C7 321A40 01240 LD (KPLACE),A
01250 ; GET AND SAVE FOUND KEYBOARD BYTE
30CA 78 01260 LD A,E
30CB 73 01270 FOUND LD (HL),E
30CC 7A 01280 LD A,0
30CD 07 01290 RLCA
30CE 07 01300 RLCA
30CF 07 01310 RLCA
30D0 57 01320 LD
30D1 0E01 01330 LD C,1
30D3 79 01340 BACKUP LD A,C
30D4 A3 01350 ANU E
30D5 2005 01360 JR NZ,AROUND
30D7 14 01370 INC 0
30D8 C801 01380 RLC C
30DA 1BF7 01390 JR BACKUP
30DC 3AB038 01400 AROUND LD A,(38B0H)
30DF 47 01410 LD B,A
30E0 7A 01420 LD A,D
30E1 C640 01430 ADD A,40H
30E3 FE60 01440 CP 60H
30E5 3016 01450 JR NC,ZD429H
30E7 57 01460 LD 0,A
30E8 3A4038 01470 LD A,(3840H)
30EB E610 01480 AND 10H
30ED 2009 01490 JR NZ,CNTRDL
30EF 7A 01500 LD A,0
30F0 C808 01510 RRC B
30F2 383D 01520 JR C,GOAWAY
30F4 C62D 01530 ADD A,20H
30F6 1839 01540 JR GOAWAY
30F8 7A 01550 CNTRDL LD A,D
30F9 0640 01560 SUB 40H
30FB 1834 01570 JR GOAWAY
30FD 067D 01580 ZD429H SUB 70H
30FF 3D10 01590 JR NC,ZD430H
3101 C640 01600 ADO A,40H
3103 FE3C 01610 CP 3CH
3105 3B02 01620 JR C,ZD435H
3107 EE1D 01630 XDR 10H
3109 C808 01640 ZD435H RRC B
310B 3024 01650 JR NC,GOAWAY
310D EE10 01660 XOR 1DH
310F 1820 01670 JR GOAWAY
3111 07 01680 ZD430H RLCA
3112 C808 01690 RRC B
3114 3001 01700 JR NC,ZD443H
3116 3C 01710 INC A
3117 212131 01720 ZD443H LD HL,TABLET
311A 4F 01730 LD C,A
311B 0600 01740 LD B,D
311D 09 01750 ADO HL,BC
311E 7E 01760 LD A,(HL)
311F 1B10 01770 JR GOAWAY
3121 000D 01780 TABLET OEFW 0D00H
3123 1F1F 01790 DEFW 1F1FH
3125 D101 01800 DEFW 0101H
3127 5B1B 01810 DEFW 1B5BH
3129 0A00 01820 DEFW 000AH
312B 0B18 01830 DEFW 1B0BH
312D 0919 01840 OEFW 1909H
312F 2020 01850 OEFW 2020H
3131 57 01860 GOAWAY LD D,A
3132 3A103B 01870 BEEEEEP LD A,(3B1DH)
313S FE01 01880 CP 1
3137 2016 01890 JR NZ,BLEEPP
3139 3AB038 01900 LD A,(38B0H)
313C FE01 01910 CP 1
313E 20UF 01920 JR NZ,BLEEPP
3140 3A1940 01930 LD A,(SHIFTR)
3143 EE01 01940 XOR 1

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3145 321940 01950 LD (SHIFTR),A
3148 010005 01960 LD BC,500H
3148 C060DD 01970 CALL 006DH
314E C9 01980 RET
01990 ; OEBOUNCE
314F 018001 02000 BLEEP LO BC,180H
3152 C060DD 02010 CALL DD6DH
3155 7A 02020 LO A,0
02030 ; BEEP ROUTINE ON FOUND KEYSTROKE
3156 C5 02040 PUSH BC
3157 F5 02050 PUSH AF
3158 0640 02060 LD B,40H
315A 3A3040 02070 LO A,(PORTFF)
3150 E6FD 02080 ANO DFOH
315F 67 02090 LO H,A
3160 F6D2 02100 OR 2
3162 6F 02110 LO L,A
3163 7D 02120 BEEPER LO A,L
3164 D3FF 02130 JUT [OFFH],A
3166 7C 02140 LD A,H
3167 D3FF 02150 JUT [OFFH],A
3169 C5 02160 PUSH BC
316A 0640 02170 LD B,40H
316C 10FE 02180 FREQCY QJNZ FREQCY
316E C1 02190 POP BC
316F 10F2 02200 QJNZ BEEPER
3171 F1 02210 POP AF
3172 C1 02220 POP BC
3173 C35204 02230 JP 0452H
02240 ; CHECK FOR STATUS OF BASIC STACK
3176 02250 SKP5TP EQU $
3176 E3 02260 BEGIN EX (SP),HL
3177 7D 02270 LD A,L
3178 FE58 02280 CP 58H
317A 2003 02290 JR NZ,NOTROY
317C 7C 02300 LD A,H
317D FE10 02310 CP 10H
317F E3 02320 NOTROY EX (SP),HL
3180 C27B10 02330 JP NZ,BYTE
02340 ; CHECK TO SEE IF SPECIAL STAR (*) COMMAND
3183 C07810 02350 CALL BYTE
3186 F6CF 02360 CP 0CFH
318B 2B03 02370 JR Z,DKSTAR
318A 28 02380 DEC HL
3188 F0E9 02390 JP (IY)
318D C07810 02400 DKSTAR CALL BYTE
3190 CABE31 02410 JP Z,SYNERR
02420 ; CHECK STATUS OF SAVE COMMAND
3193 FEA0 02430 SAVE CP 0A0H
3195 287F 02440 JR Z,SAVER
3197 FE88 02450 CP 088H
3199 2866 02460 JR Z,RENEW
3198 FEA2 02470 CP 0A2H
3190 CAE432 02480 JP Z,OPENER
31A0 FECC 02490 CP 0CCH
31A2 CAB034 02500 JP Z,STP5ET
31A5 FECB 02510 CP 0CBH
31A7 CACE34 02520 JP Z,MEM5ET
31AA FEF2 02530 CP 0F2H
31AC 284A 02540 JR Z,INBEEP
31AE FEB2 02550 CP 0B2H
31B0 2829 02560 JR Z,LOWCA5
31B2 FEA0 02570 CP 0A0H
31B4 283A 02580 JR Z,BIPOFF
31B6 FEA4 02590 CP 0AAH
31B8 282E 02600 JR Z,UPPPER
31BA FE88 02610 CP 088H
31BC 2803 02620 JR Z,MENU
31BE C39719 02630 SYNERR JP 1997H
02640 ; THIS IS THE MENU
31C1 213B35 02650 MENU LO HL,INTRO1
31C4 C0A72B 02660 CALL 28A7H
31C7 21A335 02670 LO HL,INTRO2
31CA C0A728 02680 CALL 28A7H
31C0 214236 02690 LD HL,INTRO3
31D0 C0A72B 02700 CALL 28A7H
3103 210B36 02710 LO HL,INTRO4
31D6 C0A728 02720 CALL 28A7H
3109 1B23 02730 JR READYX
02740 ; THIS PUTS LOWER CASE IN PLACE
310B 3E01 02750 LOWCA5 LO A,1
3100 321940 02760 LO (SHIFTR),A
31E0 211035 02770 LO HL,LOWER
31E3 221E40 02780 LO (4016H),HL
31E6 1816 02790 JR READYX
02800 ; THIS REMOVES LOWER CASE DISPLAY
31EB 215B04 02810 UPPPER LO HL,0458H
31EB 221E40 02820 LD (4016H),HL
31EE 180E 02830 JR READYX
02840 ; THIS REMOVES KBPFIX ROUTINE
31F0 21E303 02850 BIPOFF LO HL,03E3H
31F3 221640 02860 LO (4016H),HL
31F6 1806 02870 JR READYX
02880 ; THIS PUTS KBPFIX IN PLACE
31FB 21B530 02890 INBEEP LO HL,KBPFIX
31FB 221640 02900 LO (4016H),HL
31FE C3CC06 02910 READYX JP READY
02920 ; THIS IS BEGINNING OF RENEW SEQUENCE
3201 E05BA440 02930 RENEW LO DE,(40A4H)
3205 3EFF 02940 LD A,0FFH
3207 12 02950 LD (0E),A
3208 C0FC1A 02960 CALL 1AFCH
3208 23 02970 INC HL
320C 22F940 02980 LD (40F9H),HL
320F E078EB40 02990 LD SP,(40EBH)
3213 C37630 03000 FGHJ JP RSLVII
03010 ; THIS CHECKS NEXT BYTE FOR RUN COMMAND
3216 C07810 03020 SAVER CALL BYTE
3219 FED0 03030 CP 0D0H
321B 20A1 03040 JR NZ,SYNERR
3210 C07810 03050 CALL BYTE
3220 FE8E 03060 CP 8EH
3222 C2F433 03070 JP NZ,MACH
03080 ; CHECK FOR QUOTATION MARK DELIMITER
3225 C07810 03090 CALL BYTE
322B FE22 03100 CP 022H
322A 2092 03110 JR NZ,SYNERR
03120 ; CHECK TO SEE THAT NAME IS IN PLACE
322C C07810 03130 CALL BYTE
322F CAAD24 03140 JP Z,24A0H
03150 ; SAVE BASIC POINTERS IN STACK
3232 E5 03160 PUSH HL
3233 F5 03170 PUSH AF
3234 05 03180 PUSH OE
3235 C5 03190 PUSH BC
03200 ; DEFINE TAPE DRIVE 0, TURN ON RECORDER
03210 ; THEN WRITE LEADER AND SYNC BYTE
03220 ; AND WRITE MACHINE PROGRAM CODE 55H
3236 AF 03230 XOR A
3237 C01202 03240 CALL 0212H
323A C0B702 03250 CALL 02B7H
3230 3E55 03260 LO A,55H
323F C06402 03270 CALL WRT8YT
03280 ; WRITE PROGRAM NAME TO TAPE
3242 0606 03290 LD B,06
3244 2B 03300 DEC HL
3245 C07810 03310 NAME5 CALL BYTE
3248 FE22 03320 CP 22H
324A 2807 03330 JR Z,NEXT8T
324C C06402 03340 CALL WRT8YT
324F 10F4 03350 DJNZ NAME5
3251 1B07 03360 JR OUMP
03370 ; FILL OUT WITH 20H (ASCII BLANKS) IF NECESSARY
3253 3E20 03380 NEXT8T LD A,2CH
3255 C06402 03390 CALL WRT8YT
325B 10F9 03400 DJNZ NEXT8T
03410 ; OUMP FIRST TWO PAGES (40DD TO 41FF) TO TAPE
325A 21DD40 03420 OUMP LO HL,5ETPT5
3250 C0C032 03430 CALL OUTSEQ
3260 C0C032 03440 CALL OUTSEQ
03450 ; OUMP REST OF POINTERS (4200 TO 42E9) TO TAPE
3263 06E9 03460 LD B,0E9H
3265 C0C032 03470 CALL OUTSEQ
03480 ; FIND END OF PROGRAM VARIABLES AND ARRAYS
3268 E058F040 03490 LO DE,(VAREND)
03500 ; DUMP FIRST SEGMENT OF PROGRAM TO TAPE
326C 7B 03510 LO A,E
326D 95 03520 SUB L
326E 47 03530 LD B,A
326F C0C032 03540 CALL OUTSEQ
03550 ; OUMP PROGRAM TO TAPE PAGE BY PAGE
3272 7C 03560 NXTPGE LO A,H
3273 30 03570 DEC A
3274 BA 03580 CP 0
3275 28D5 03590 JR Z,FINSH1
3277 C0C032 03600 CALL OUTSEQ
327A 1BF6 03610 JR NXTPGE
03620 ; FIND BEGINNING OF STRING STORAGE AREA
327C 2AEB40 03630 FINSH1 LO HL,(STACKR)
03640 ; FIND TOP OF AVAILABLE MEMORY
327F E058B140 03650 LD OE,(MEMTOP)
03660 ; OUMP FIRST SEGMENT OF STRING-TO-MEMORY END
3283 13 03670 INC DE
3284 7B 03680 LD A,E
3285 95 03690 SUB L
3286 47 03700 LD B,A
3287 C0C032 03710 CALL OUTSEQ
03720 ; OUMP REMAINDER OF MEMORY PAGE BY PAGE
328A 7C 03730 NXT8CH LO A,H
328B 3D 03740 DEC A
328C 8A 03750 CP 0
328D 2B05 03760 JR Z,KEEPIT
328F C0C032 03770 CALL OUTSEQ
3292 1BF6 03780 JR NXT8CH
03790 ; OUMP KEEPIT CONTROL BYTES
3294 2A8140 03800 KEEPIT LO HL,(4081H)
3297 0610 03810 LO B,10H

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34B7 C3CC06 0754C PREP2 JP READY
07550 ; THIS IS THE STEPPER ROUTINE
34BA 07560 JMPPOS EQU $
34BA F5 07570 STEPPR PUSH AF
34BB C5 07580 PUSH BC
07590 ; WAIT FOR SHIFT TO BE PRESSED
34BC 3A803B 07600 LUPER LD A,(3B80H)
34BF A7 07610 ANO A
34CD 2BFA 07620 JR Z,LUPER
07630 ; LOAD DELAY VALUE INTO BC
34C2 4F 07640 LO C,A
34C3 F04606 07650 LO B,(IY+6)
07660 ; CALL DELAY IN ROM
34C6 C06000 07670 CALL 0060H
34C9 C1 07680 POP BC
34CA F1 07690 POP AF
07700 ; BACK TO REST OF TEST SEQUENCE
34CB C37631 07710 JP BEGIN
07720 ; BEGIN MEMORY RESET SEQUENCE
07730 ; CHECK FOR QUOTE MARK DELIMITER
34CE C07B10 07740 MEMSET CALL BYTE
34D1 07750 CP 22H
34D3 20CE 07760 JR NZ,SYN3
07770 ; CONVERT *MEM OPERAND TO HEX
34D5 C0D133 07780 CALL XX99
07790 ; CHECK FOR >4400H MEMORY ADDRESSES
07800 ; GO TO OM ERROR IF NOT ENOUGH
34DB 210D44 07810 LO HL,4400H
34DB AF 07820 XOR A
34DC E052 07830 SBC HL,DE
34DE 027A19 07840 JP NC,197AH
07850 ; TEST FOR MEMORY RESET <FBBBH
07860 ; OM ERROR IF TOO MUCH
34E1 05 07870 PUSH DE
34E2 E1 07880 POP HL
34E3 0607 07890 LO B,7
34E5 23 07900 INC HL
34E6 1UFD 07910 QJNZ $-1
34E8 7E 07920 LO A,(HL)
34E9 47 07930 LO B,A
34EA 2F 07940 CPL
34EB 77 07950 LO (HL),A
34EC BE 07960 CP (HL)
34ED C27A19 07970 JP NZ,197AH
34FD 70 07980 LO (HL),B
07990 ; PUT NEW MEMORY SIZE INTO PLACE
34F1 E053B140 09000 LO {40B1H},DE
0B010 ; TRANSFER KEEPIT DATA ROW
0B020 ; AND PUT NEW ADDRESS IN IY
34F5 F0E5 0B030 PUSH IY
34F7 E1 0B040 POP HL
34FB 13 0B050 INC DE
34F9 05 0B060 PUSH DE
34FA F0E1 0B070 POP IY
34FC 010700 0B080 LO BC,7
34FF E0B0 0B090 LOIR
3501 F0E5 0B100 PUSH IY
3503 E1 0B110 POP HL
0B120 ; MAKE FRE(AS) MEM SIZE - 50 DECIMAL
3504 0632 0B130 LO B,32H
3506 1B 0B140 STRING DEC OE
3507 10FD 0B150 QJNZ STRING
0B160 ; PUT NEW STRING POINTER IN PLACE
3509 E053A040 0B170 LO {40A0H},DE
3500 C36C30 0B180 JP GRIN5
0B190 ; THIS IS LOWER CASE DETERMINATION
3510 F5 0B200 LOWER PUSH AF
3511 3A1940 0B210 LO A,(SHIFTR)
3514 FE01 0B220 CP 1
3516 2604 0B230 JR Z,LOWER1
3518 F1 0B240 POP AF
3519 C35B04 0B250 JP 045BH
351C F1 0B260 LOWER1 POP AF
3510 006E03 0B270 LD L,(IX+3)
3520 006604 0B280 LD H,(IX+4)
3523 0A9A04 0B290 JP C,049AH
3526 007E05 0B300 LO A,(IX+5)
3529 B7 0B310 OR A
352A 2B01 0B320 JR Z,GETCHR
352C 77 0B330 LD (HL),A
352D 79 0B340 GETCHR LO A,C
352E FE20 0B350 CP 20H
3530 0A0605 0B360 JP C,0506H
3533 FE80 0B370 CP B0H
3535 02A604 0B380 JP NC,04A6H
3538 C37004 0B390 JP 0470H
353B 2A 0B400 INTR01 OEFM '*FIX = SET DEBOUNCE/BEEP/AUTOREPEAT'
353E 0A 0B410 OEFB OAH
353F 2A 0B420 OEFM '*KILL = RESET TO NORMAL KEYBOARD'
357F 0A 0B430 OEFB OAH
3580 2A 0B440 OEFM '*MEM = RESET MEMORY SIZE (NOTE 1)'
35A1 00 0B450 OEFB 00H
35A2 00 0B460 OEFB 00H
35A3 2A 0B470 INTR02 OEFM '*PRINT = LOWERCASE KEYBOARD/DISPLAY ON'
35CB 0A 0B480 OEFB OAH
35CA 2A 0B490 OEFM '*OUT = LOWERCASE KEYBOARD/DISPLAY OFF'
35EF 0A 0B500 OEFB OAH
35FD 2A 0B510 OEFM '*NEW = RESTORE PROGRAM VICTIM OF NEW'
3614 0A 0B520 OEFB OAH
3615 2A 0B530 OEFM '*OPEN = HEX/ASCII/GRAPHICS MONITOR (NOTE 1)'
3640 00 0B540 OEFB 00H
3641 00 0B550 OEFB 00H
3642 2A 0B560 INTR03 OEFM '*SAVE/RUN = SAVE RUNNING PROGRAM (NOTE 2)'
366B 0A 0B570 OEFB OAH
366C 2A 0B580 OEFM '*SAVE/OPEN = MEMORY BLOCK SAVE (NOTE 3)'
3693 0A 0B590 OEFB OAH
3694 2A 0B600 OEFM '*STEPXX = SINGLE STEPPER ON, XX = DELAY'
36BB 0A 0B610 OEFB OAH
36BC 2A 0B620 OEFM '*STEP = SINGLE STEPPER OFF'
36C6 00 0B630 OEFB 00H
36C7 00 0B640 OEFB 00
36C8 4E 0B650 INTR04 OEFM 'NOTE 1. REQUIRES 4-CHARACTER HEX VALUE IN QUOTES.'
370A 0A 0B660 OEFB OAH
370B 4E 0B670 OEFM 'NOTE 2. REQUIRES 6-CHARACTER NAME IN QUOTES.'
3738 0A 0B680 OEFB OAH
3739 4E 0B690 OEFM 'NOTE 3. SAME AS ABOVE PLUS HEX START,
3773 00 0B700 OEFB 00H
3774 00 0B710 OEFB 00H
303B 0B720 ENO START
00000 TOTAL ERRORS
1B637 TEXT AREA BYTES LEFT
AAAA 334F 05030 05000
AAAB 3355 05070 05040
AAAC 335F 05130 05090
AAD 336E 05230 05190
AAAE 337B 05280 05430
AAAF 33BF 05410 05380
ABCO 34A6 07430 07380
AROUND 300C 01400 01360
ASCHEX 33A7 05590 05460 05810 05990 06080
BACKUP 3003 01340 01390
BBBA 3335 04830 04890
BEEEEP 3132 01870
BEEPER 3163 02120 02200
BEGIN 3176 02260 07710
BIPOFF 31F0 02850 02580
BLEEEP 314F 02000 01890 01920
BREEK 3367 05180 05140
BYTE 1078 00390 02330 02350 02400 03020 03050 03090 03130
03310 04360 05940 06210 06330 06450 06520
06620 07180 07360 07740
COEF 32A1 03860 03890
CLRMEM 309E 00980 01000
CNTRL 30F8 01550 01490
CONTO2 3318 04630 04740
CONTNT 32F5 04440 04950 05550
DECA 30C6 01230 01170
DELAY 33CA 05880 05560
DUMP 325A 03420 03360
DUMP99 3422 06450 06380
EDITOR 3345 04970 05340 05400 05570
EXECU 344C 06700 06640
FGHIJ 3213 03000
FINSH1 327C 03630 03590
FOUND 30CB 01270 01090
FREQUY 316C 02180 02180
GETCHR 3520 0B340 0B320
GOAWAY 3131 01860 01520 01540 01570 01650 01670 01770
GRIN5 306C 00660 0B180
HEXASC 33B0 05660 04500 04580 04680 05760
IIII 3494 07290 07210
INBEEP 31F9 02890 02540
INKEYS 4099 00260 01100
INTR01 353B 0B400 02650
INTR02 35A3 0B470 02670
INTR03 3642 0B560 02690
INTR04 360B 0B650 02710
JMPPOS 34BA 07560 07470 07480
KBPFIX 30B5 00790 02890
KEEPIT 32B4 03800 03760
KEYPRS 30B0 00B30 00950
KPLACE 401A 00290 01050 01130 01150 01240
KYSCAN 03FB 00270
LLLLS 33C2 05810 05490 06030 06120
LOWCAS 310B 02750 02560
LOWER 3510 0B200 02770
LOWER1 351C 0B260 0B230
LP99 3465 06980 07030
LUPER 34BC 07600 07620
MACH 33F4 06180 03070
MEMSET 34CE 07740 02520
MENTOP 40B1 00360 03650
MENU 31C1 02650 02620
NAME99 3400 06330 06370

```

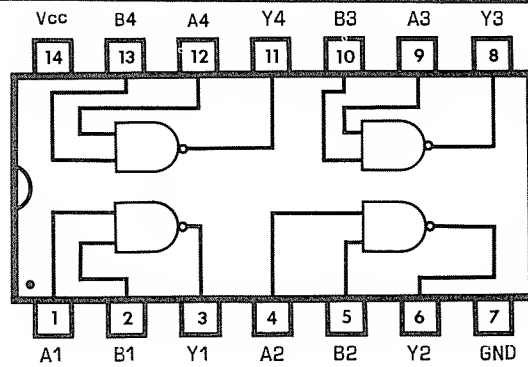
```

NAMES 3245 03310 03350
NEXT99 3342 04950 04420
NEXT8T 3253 03380 03330 03400
NOTROY 317F 02320 02290
NXT899 3418 06410 06350 06430
NXT8CH 328A 03730 03780
NXTPGE 3272 03560 03610
OKSTAR 3180 02400 02370
OPENER 32E4 04360 02480
OUTSEQ 32C0 04080 03430 03440 03470 03540 03600 03710 03770
03820 06990
PORTFF 4030 00280 02070
PREOUT 328E 04070 03870
PREP2 3487 07540 07130 07270
READY 06CC 00350 00770 02910 04040 05210 07540
READYX 31FE 02910 02730 02790 02830 02870
RECHK 3093 00890 01120 01220
RENEW 3201 02930 02460
RESTRT 4004 00310 00590 00690
RRRRS 3388 05720 04470 04550 04650
RSLVII 3076 00730 03000
SAVE 3193 02430
SAVER 3216 03020 02440
SETPTS 4000 00370 03420
SHIFTR 4019 00300 00710 01930 01950 02760 08210
SKPSTP 3176 02250 00630 07230 07240
SSSS 3303 05940 05960
STACKR 40E8 00330 03630
START 3039 00440 08720
STEPPR 34BA 07570
STGEN0 40A0 00320
STNORD 339F 05550 05020 05060 05120 05170
STPSET 3480 07160 02500
STRING 3506 08140 08150
STROKE 30AA 01080 00860
SYN2 33F0 06230 06190 06470
SYN3 34A3 07410 07760
SYNERR 318E 02630 02410 03040 03110 04380 06230 07410
TABLET 3121 01780 01720
TMWSTE 30C1 01200 01200
TTTT 3439 06590 06540
UPPPER 31E8 02810 02600
VARENO 40F0 00340 03490
VIDEO 3C00 00400 03840
WRIT99 3453 06740 06680
WR8TYT 0264 00380 03270 03340 03390 03920 03950 03970 04090
04120 04160 04220 04280 04330 06290 06360
06420 07060 07090 07110
WRTPGE 3205 04240 04300
XX99 3301 05930 04400 06490 06590 06700 07780
Z0429H 30F0 01580 01450
Z0435H 3109 01640 01620
Z0430H 3111 01680 01590
Z0443H 3117 01720 01700

```

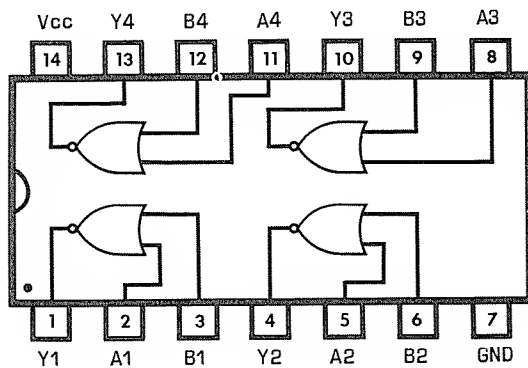
KEEPIT 3.2 is a 2K utility program created to extend the capabilities of Level II BASIC. It was originally sold by The Alternate Source in a RAM-based format (version 2.1), and has appeared variously from Personal Micro Computers, The Peripheral People, and Computer Accessory Technology. It is still available from C.A.T., and (together with the Memory Sidecar) from MSB Electronics, Drawer 766, Barre, Vermont 05641. However, since it contains many of the software drivers and other routines presented in The Custom TRS-80, and represents a complete implementation of the custom interpreter (Chapter 3), it is offered here as a completely revised version in source code format. Its current origin is 3000H, for use with the Memory Sidecar project presented in Chapter 8.

74LS00



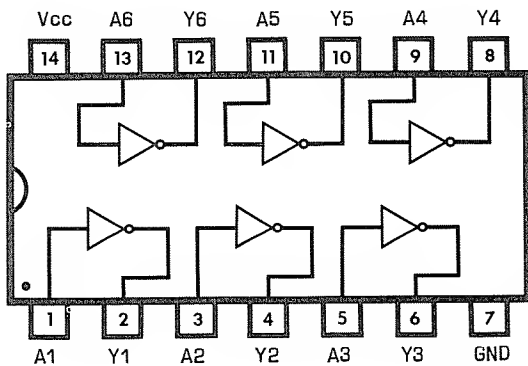
$$Y = \overline{AB}$$

74LS02



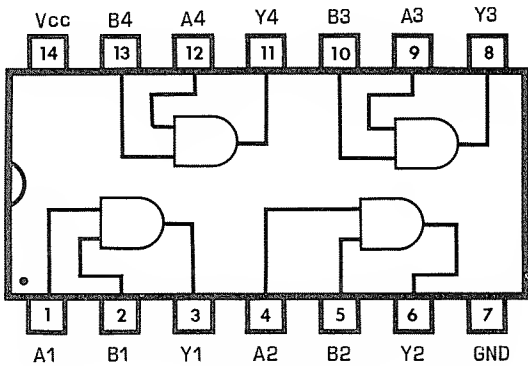
$$Y = \overline{A+B}$$

74LS04

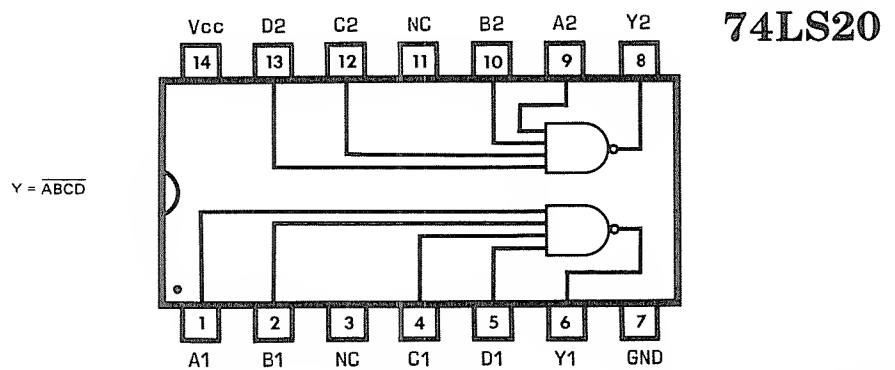
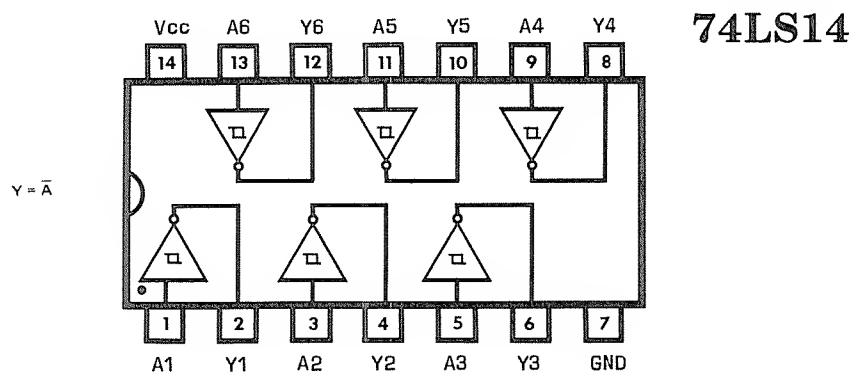
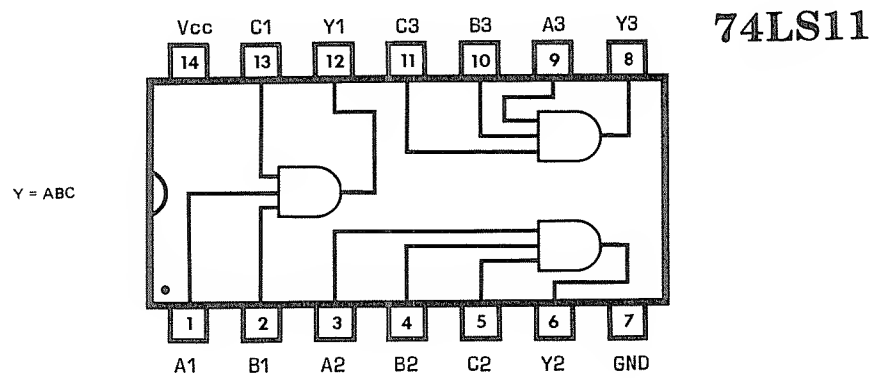
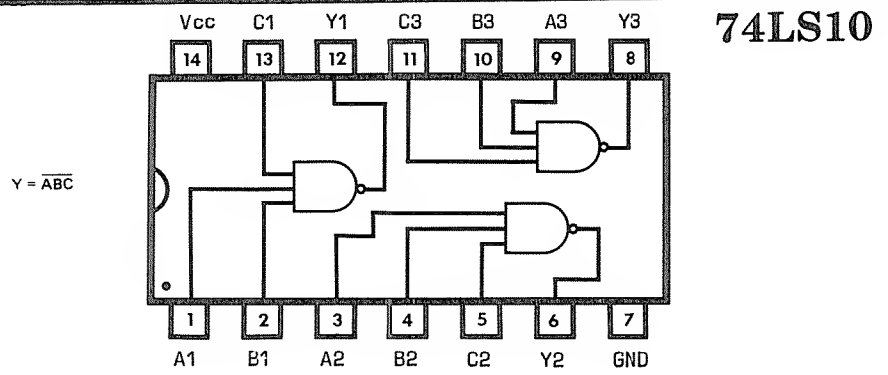


$$Y = \overline{A}$$

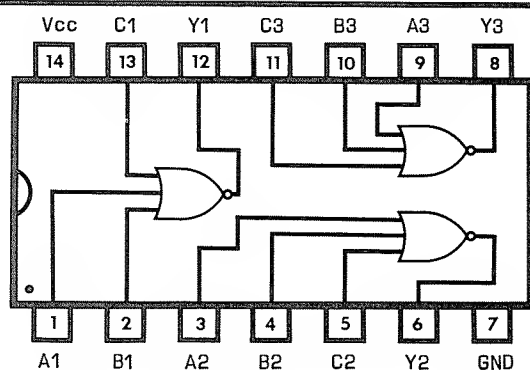
74LS08



$$Y = AB$$

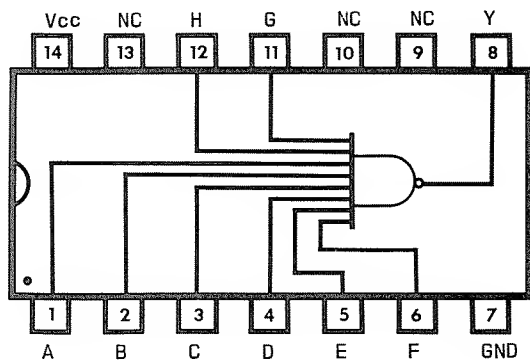


74LS27



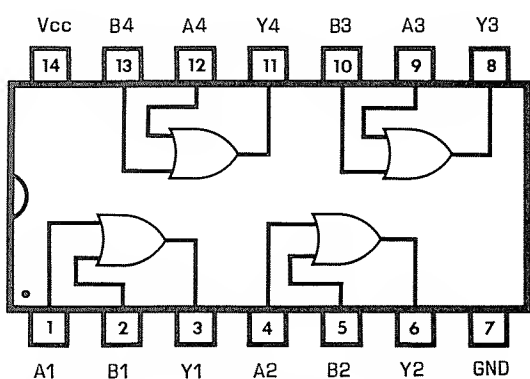
$$Y = A + B + C$$

74LS30



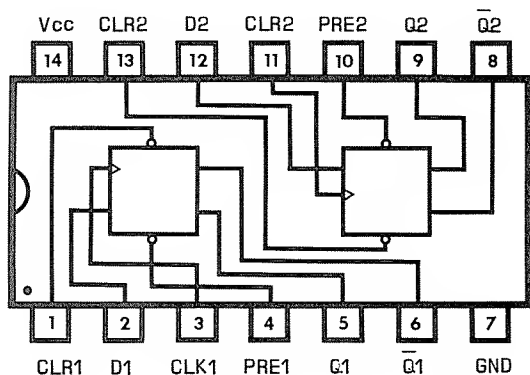
$$Y = \overline{ABCDEFGH}$$

74LS32



$$Y = A + B$$

74LS74

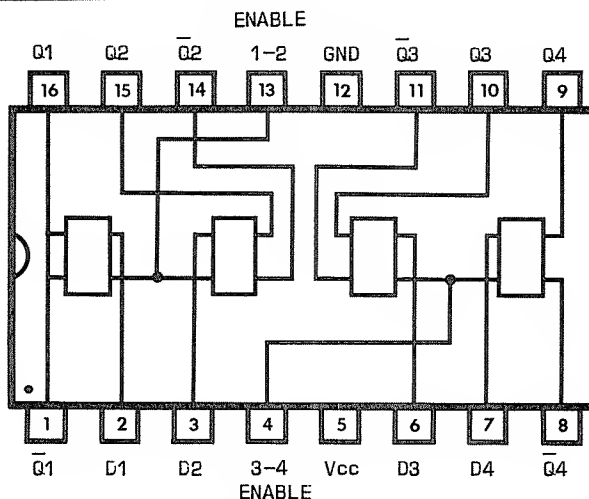


INPUTS				OUTPUTS	
PR	CLR	CLK	D	Q	\bar{Q}
L	H	X	X	H	L
H	L	X	X	L	H
L	L	X	X	H*	H*
H	H	↑	H	H	L
H	H	↑	L	L	H
H	H	L	X	00	00

74LS75

INPUTS		OUTPUTS	
Q	G	Q	\bar{Q}
L	H	L	H
H	H	H	L
X	L	Q_0	\bar{Q}_0

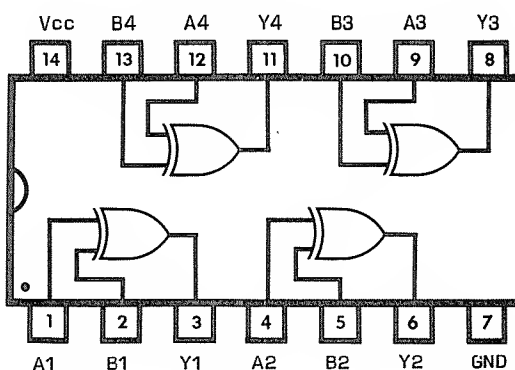
H = High Level, L = Low Level, X = Don't Care
 Q_0 = The Level of Q Before the High-to-Low Transition of G



74LS86

INPUTS		OUTPUT
A	B	Y
L	L	L
L	H	H
H	L	H
H	H	L

$$Y = A \oplus B = \bar{A}B + A\bar{B}$$



LS90
BCD COUNT SEQUENCE
(See Note A)

COUNT	OUTPUT			
	Q_D	Q_C	Q_B	Q_A
0	L	L	L	L
1	L	L	L	H
2	L	L	H	L
3	L	L	H	H
4	L	H	L	L
5	L	H	L	H
6	L	H	H	L
7	L	H	H	H
8	H	L	L	L
9	H	L	L	H

LS90 BI-QUINARY (5-2)
(See Note B)

COUNT	OUTPUT			
	Q_A	Q_D	Q_C	Q_B
0	L	L	L	L
1	L	L	L	H
2	L	L	H	L
3	L	L	H	H
4	L	H	L	L
5	H	L	L	L
6	H	L	L	H
7	H	L	H	L
8	H	L	H	H
9	H	H	L	L

LS90
RESET/COUNT TRUTH TABLE

RESET INPUTS				OUTPUT			
R0(1)	R0(2)	R9(1)	R9(2)	Q_D	Q_C	Q_B	Q_A
H	H	L	X	L	L	L	L
H	H	X	L	L	L	L	L
X	X	H	H	H	L	L	H
X	L	X	L	COUNT			
L	X	L	X	COUNT			
L	X	X	L	COUNT			
X	L	L	X	COUNT			

Notes:

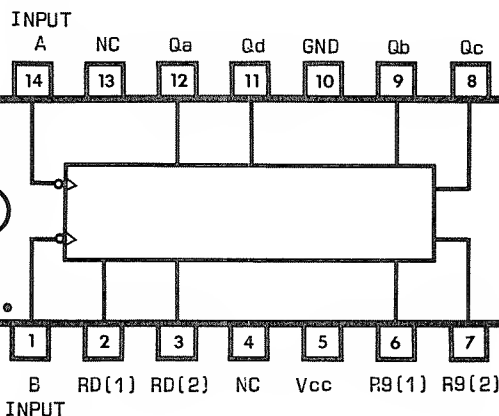
- (A) Output Q_A is connected to input B for BCD count.
- (B) Output Q_D is connected to input A for bi-quinary count.
- (C) Output Q_A is connected to input B.
- (D) H = High Level, L = Low Level, X = Don't Care.

LS92
COUNT SEQUENCE
(See Note C)

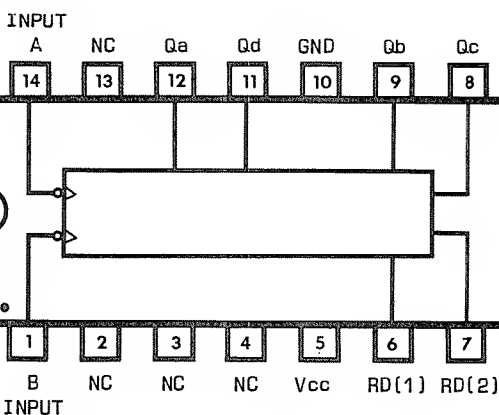
COUNT	OUTPUT			
	Q_D	Q_C	Q_B	Q_A
0	L	L	L	L
1	L	L	L	H
2	L	L	H	L
3	L	L	H	H
4	L	H	L	L
5	L	H	L	H
6	H	L	L	L
7	H	L	L	H
8	H	L	H	L
9	H	L	H	H
10	H	H	L	L
11	H	H	L	H

LS92
RESET/COUNT TRUTH TABLE

RESET INPUTS		OUTPUT			
R0(1)	R0(2)	Q_D	Q_C	Q_B	Q_A
H	H	L	L	L	L
L	X	COUNT			
X	L	COUNT			



74LS90



74LS92

CDUNT SEQUENCE (See Note C)				
CDUNT	DUTPUT			
	Q _O	Q _C	Q _B	Q _A
0	L	L	L	L
1	L	L	L	H
2	L	L	H	L
3	L	L	H	H
4	L	H	L	L
5	L	H	L	H
6	L	H	H	L
7	L	H	H	H
8	H	L	L	L
9	H	L	L	H
10	H	L	H	L
11	H	L	H	H
12	H	H	L	L
13	H	H	L	H
14	H	H	H	L
15	H	H	H	H

LS93
Notes:

- (A) Output D_A is connected to input B for BCD count.
- (B) Output Q_D is connected to input A for bi-quinary count.
- (C) Output Q_A is connected to input B
- (D) H = High Level, L = Low Level, X = Don't Care.

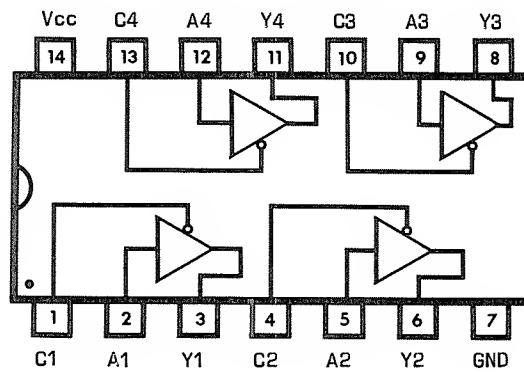
[illegible][illegible]

287

74LS125

INPUTS		OUTPUT
A	C	Y
H	L	H
L	L	L
X	H	Hi-Z

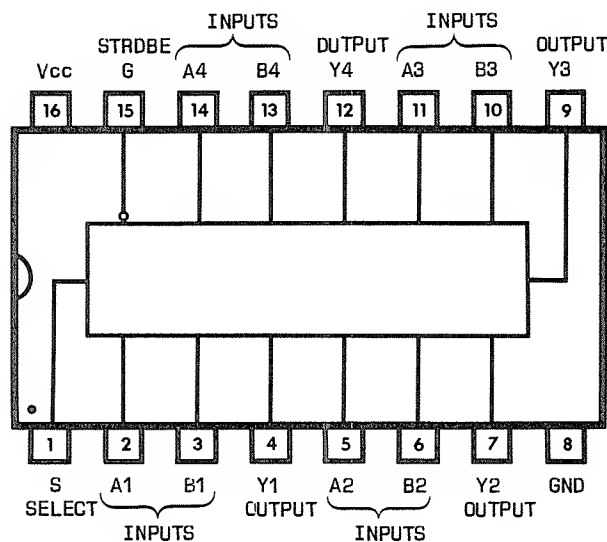
Y = A



74LS154

INPUTS				OUTPUT Y	
STROBE	SELECT	A	B	157, L157A LS157, S157	LS158 S158
H	X	X	X	L	H
L	L	L	X	L	H
L	L	H	X	H	L
L	H	X	L	L	H
L	H	X	H	H	L

H = High Level, L = Low Level, X = Don't Care



74LS157

INPUTS						INTERNAL OUTPUTS		OUTPUT
CLEAR	SHIFT/LOAD	CLOCK INHIBIT	CLOCK	SERIAL	PARALLEL A...H	Q _A	Q _B	Q _H
L	X	X	X	X	X	L	L	L
H	X	L	L	X	X	Q _{A0}	Q _{B0}	Q _{H0}
H	L	L	↑	X	a...h	a	b	h
H	H	L	↑	H	X	H	Q _{An}	Q _{Gn}
H	H	L	↑	L	X	L	Q _{An}	Q _{Gn}
H	X	H	↑	X	X	Q _{A0}	Q _{B0}	Q _{H0}

H = High Level (steady state), L = Low Level (steady state)

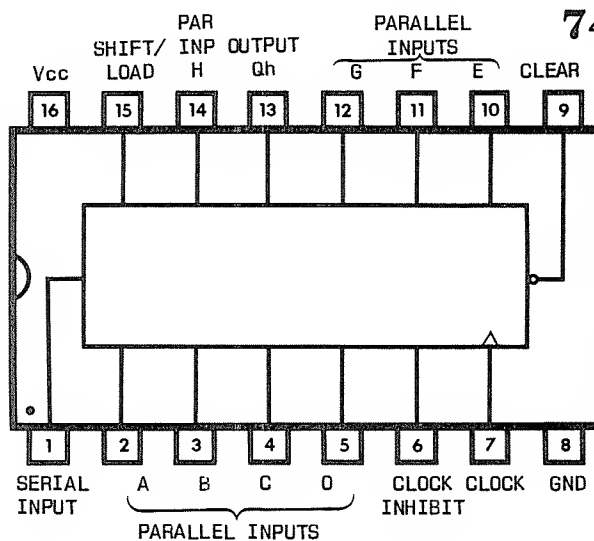
X = Don't Care (any input, including transitions)

↑ = Transition from low to high level

a...h = The level of steady-state input at inputs A through H, respectively.

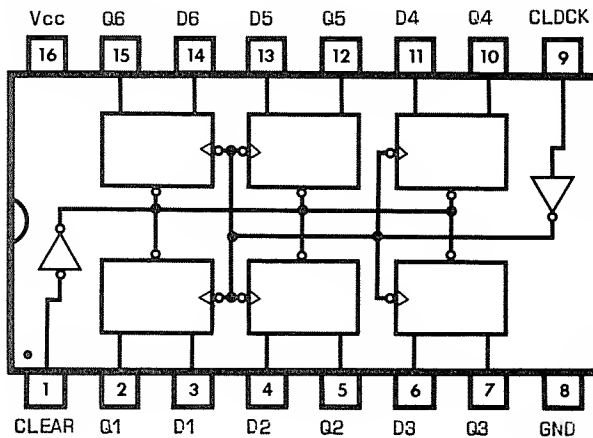
Q_{A0}, Q_{B0}, Q_{H0} = The level of Q_A, Q_B or Q_H, respectively, before the indicated steady-state input conditions were established.

Q_{An}, Q_{Gn} = The level of Q_A or Q_G, respectively, before the most recent ↑ transition of the clock



74LS166

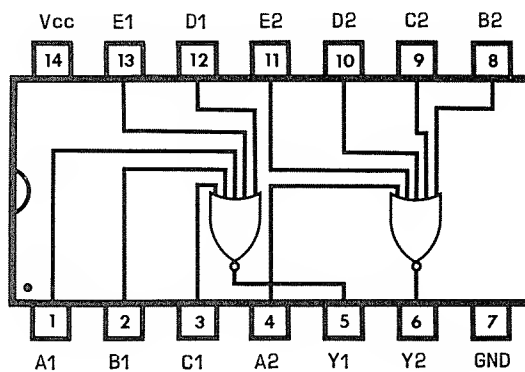
74LS174



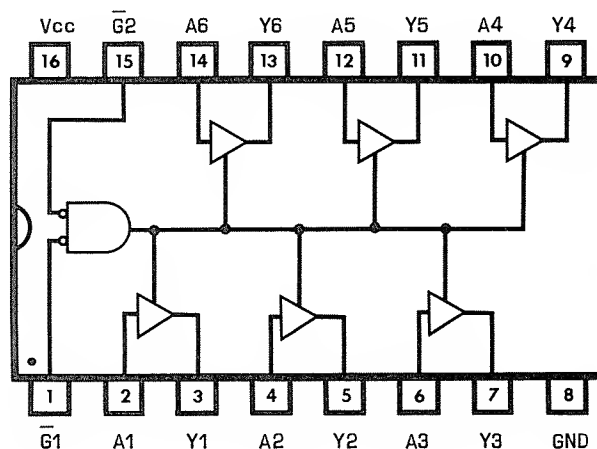
INPUTS			OUTPUTS	
CLEAR	CLOCK	O	Q	\bar{Q}
L	X	X	L	H
H	↑	H	H	L
H	↑	L	L	H
H	L	X	Q_0	\bar{Q}_0

H = High Level (steady state)
 L = Low Level (steady state)
 X = Don't Care
 ↑ = Transition from low to high level
 Q_0 = The level of Q before the indicated steady-state input conditions were established.
 † = 175, LS175, and S175 only

74LS260

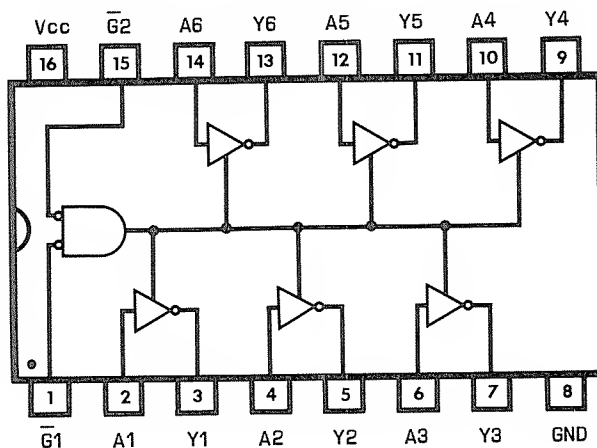


$$Y = A+B+C+D+E$$



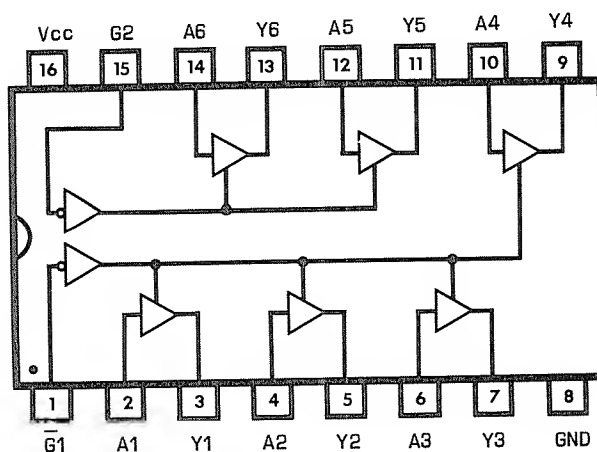
INPUTS			OUTPUT
$\bar{G}1$	$\bar{G}2$	A	Y
H	X	X	Z
X	H	X	Z
L	L	H	H
L	L	L	L

INPUTS			OUTPUT
$\bar{G}1$	$\bar{G}2$	A	Y
H	X	X	Z
X	H	X	Z
L	L	H	L
L	L	L	H



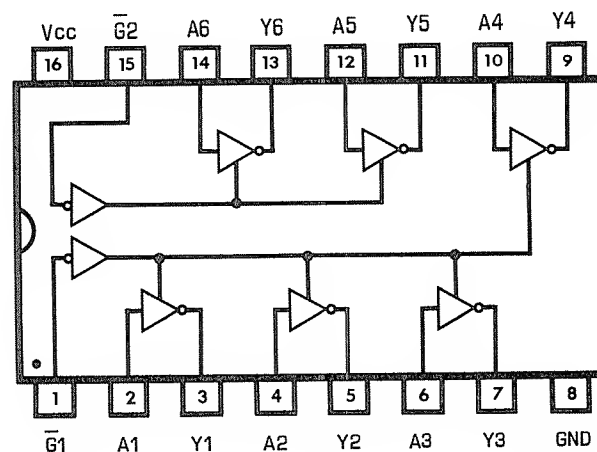
74LS366

INPUTS		OUTPUT
\bar{G}	A	Y
H	X	Z
L	H	H
L	L	L



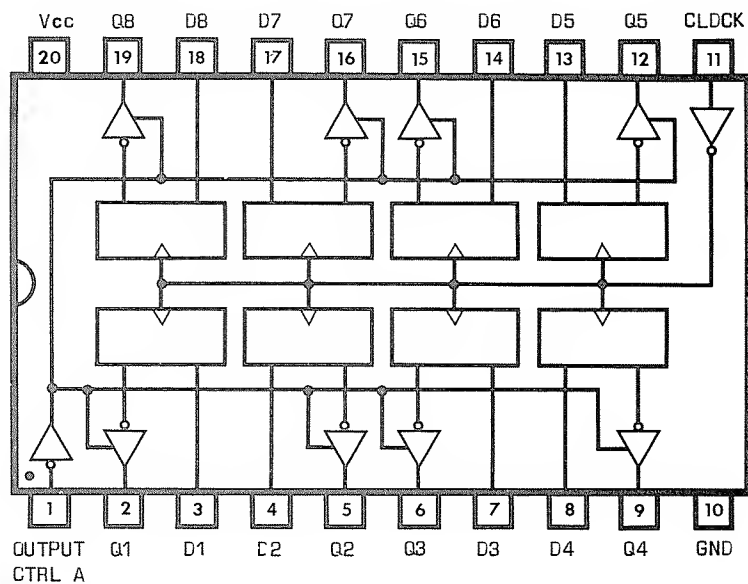
74LS367

INPUTS		OUTPUT
\bar{G}	A	Y
H	X	Z
L	H	L
L	L	H



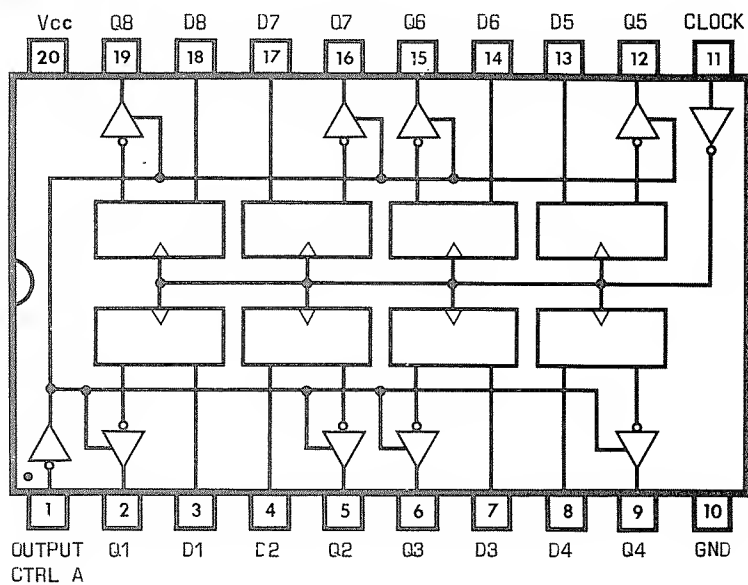
74LS368

74LS373



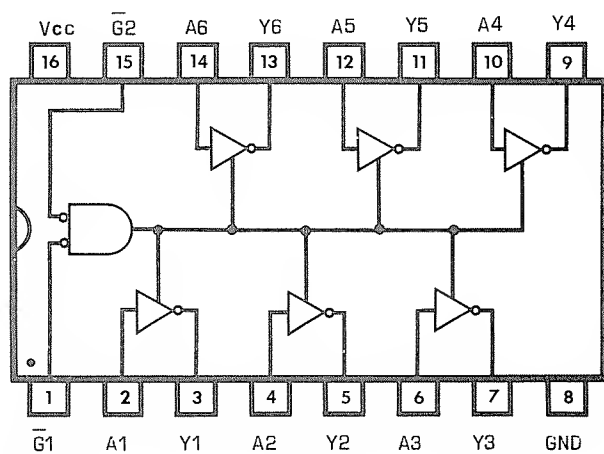
OUTPUT CONTROL	CLOCK	O	OUTPUT
L	H	H	H
L	H	L	L
L	L	X	00
H	X	X	Z

74LS374



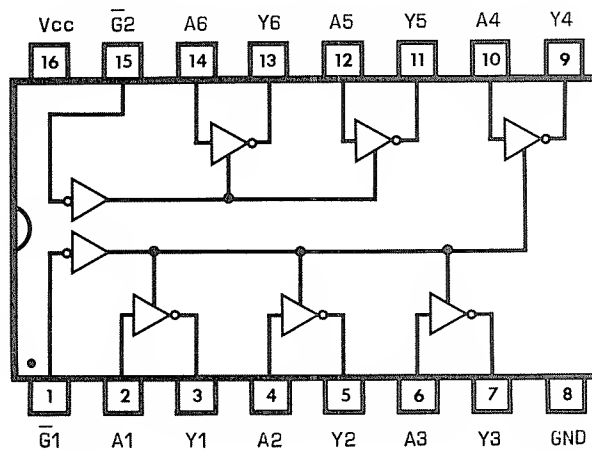
OUTPUT CONTROL	CLOCK	O	OUTPUT
L	↑	H	H
L	↑	L	L
L	L	X	00
H	X	X	Z

80C96



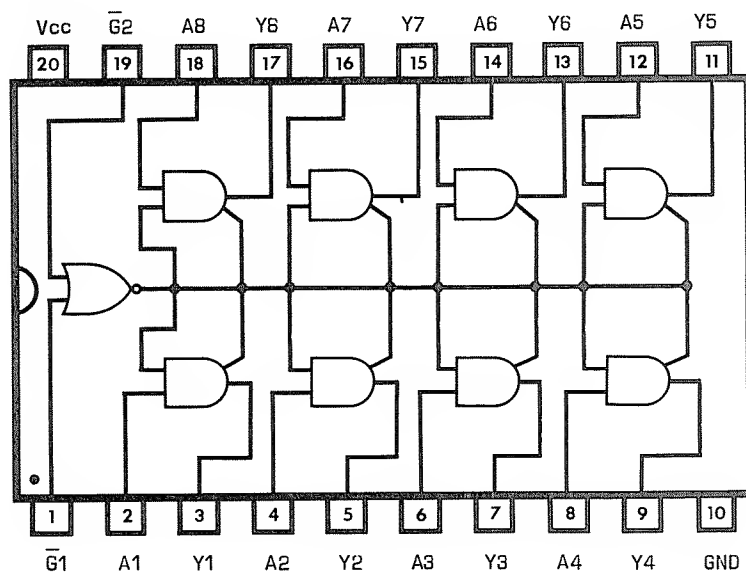
INPUTS			OUTPUT
G1	G2	A	Y
H	X	X	Hi-Z
X	H	X	Hi-Z
L	L	H	L
L	L	L	H

INPUTS		OUTPUT
\bar{G}	A	Y
H	X	Hi-Z
L	H	L
L	L	H



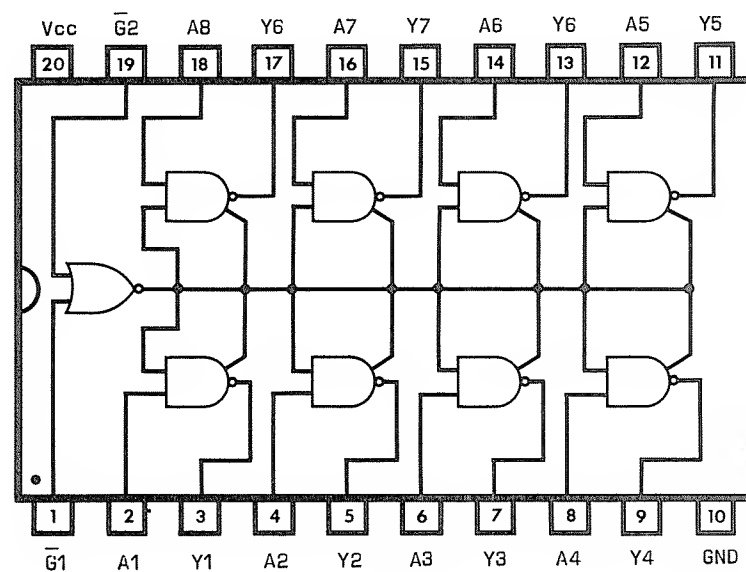
80C98

INPUTS			OUTPUT
$\bar{G}1$	$\bar{G}2$	A	Y
H	X	X	Z
X	H	X	Z
L	L	H	H
L	L	L	L

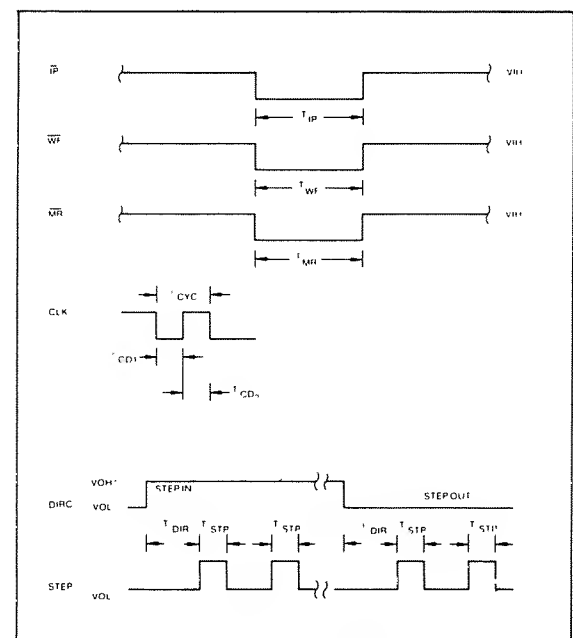
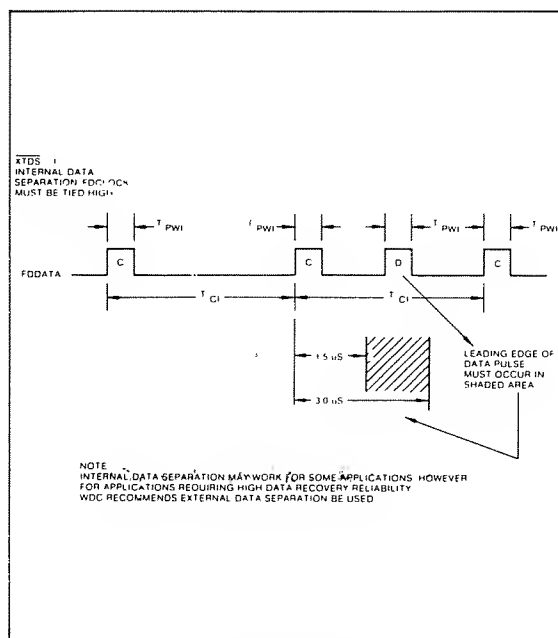
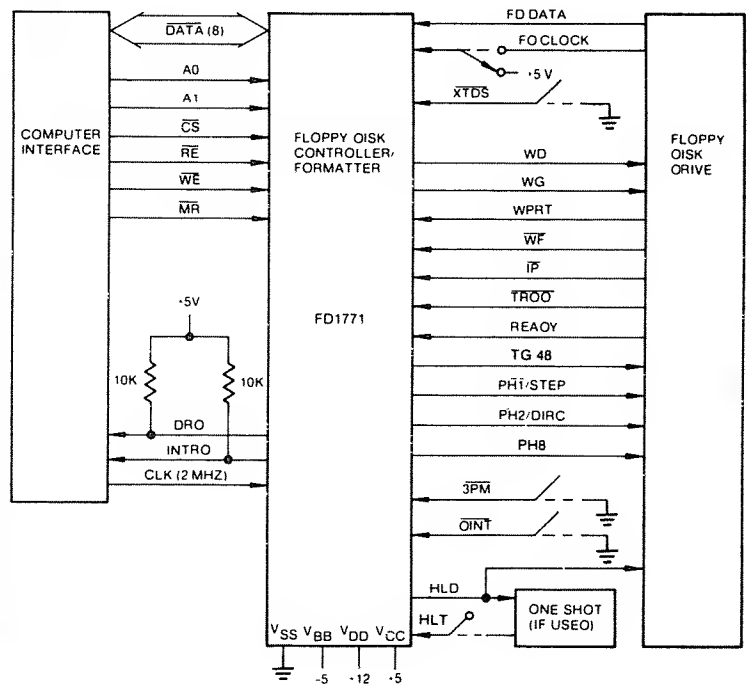


81LS95

INPUTS			OUTPUT
$\bar{G}1$	$\bar{G}2$	A	Y
H	X	X	Z
X	H	X	Z
L	L	H	L
L	L	L	H



81LS96



PIN OUTS

Pin No.	Pin Name	Symbol	Function																				
1 19	Power Supplies MASTER RESET	V_{BB}/NC \overline{MR}	-5V A logic low on this input resets the device and loads "03" into the command register. The Not Ready (Status bit 7) is reset during \overline{MR} ACTIVE. When \overline{MR} is brought to a logic high, a Restore Command is executed, regardless of the state of the Ready signal from the drive.																				
20		V_{SS}	Ground																				
21		V_{CC}	+5V																				
40		V_{DD}	+12V																				
Computer Interface																							
2	WRITE ENABLE	\overline{WE}	A logic low on this input gates data on the DAL into the selected register when \overline{CS} is low.																				
3	CHIP SELECT	\overline{CS}	A logic low on this input selects the chip and enables computer communication with the device.																				
4	READ ENABLE	\overline{RE}	A logic low on this input controls the placement of data from a selected register on the DAL when \overline{CS} is low.																				
5, 6	REGISTER SELECT LINES	A_0, A_1	These inputs select the register to receive/transfer data on the DAL lines under \overline{RE} and \overline{WE} control: <table> <tr> <td>A_1</td> <td>A_0</td> <td>\overline{RE}</td> <td>\overline{WE}</td> </tr> <tr> <td>0</td> <td>0</td> <td>Status Register</td> <td>Command Register</td> </tr> <tr> <td>0</td> <td>1</td> <td>Track Register</td> <td>Track Register</td> </tr> <tr> <td>1</td> <td>0</td> <td>Sector Register</td> <td>Sector Register</td> </tr> <tr> <td>1</td> <td>1</td> <td>Data Register</td> <td>Data Register</td> </tr> </table>	A_1	A_0	\overline{RE}	\overline{WE}	0	0	Status Register	Command Register	0	1	Track Register	Track Register	1	0	Sector Register	Sector Register	1	1	Data Register	Data Register
A_1	A_0	\overline{RE}	\overline{WE}																				
0	0	Status Register	Command Register																				
0	1	Track Register	Track Register																				
1	0	Sector Register	Sector Register																				
1	1	Data Register	Data Register																				
7-14	DATA ACCESS LINES	$DAL0-DAL7$	Eight bit inverted bidirectional bus used for transfer of data, control, and status. This bus is a receiver enabled by \overline{WE} or a transmitter enabled by \overline{RE} .																				
24	CLOCK	CLK	This input requires a free-running 2 MHz \pm 1% square wave clock for internal timing reference.																				
38	DATA REQUEST	DRQ	This open drain output indicates that the DR contains assembled data in Read operations, or the DR is empty in Write operations. This signal is reset when serviced by the computer through reading or loading the DR in Read or Write operation, respectively. Use 10K pull-up resistor to +5.																				
39	INTERRUPT REQUEST	INTRQ	This open drain output is set at the completion or termination of any operation and is reset when a new command is loaded into the command register. Use 10K pull-up resistor to +5.																				
Floppy Disk Interface:																							
15	Phase 1/Step	$\overline{PH1}/STEP$	If the $\overline{3PM}$ input is a logic low the three-phase motor control is selected and $\overline{PH1}$, $\overline{PH2}$, and $\overline{PH3}$ outputs																				

Pin No.	Pin Name	Symbol	Function
16	Phase 2/Direction	$\overline{\text{PH2/DIRC}}$	form a one active low signal out of three. $\overline{\text{PH1}}$ is active low after $\overline{\text{MR}}$. If the $\overline{\text{3PM}}$ input is a logic high the step and direction motor control is selected. The step output contains a 4 usec high signal for each step and the direction output is active high when stepping in; active low when stepping out.
17	Phase 3	PH3	
18	3-Phase Motor Select	$\overline{\text{3PM}}$	
22	$\overline{\text{TEST}}$	$\overline{\text{TEST}}$	This input is used for testing purposes only and should be tied to +5V or left open by the user.
23	HEAD LOAD TIMING	HLT	The HLT input is sampled after 10 ms. When a logic high is sampled on the HLT input the head is assumed to be engaged.
25	$\overline{\text{EXTERNAL DATA SEPARATION}}$	$\overline{\text{XTDS}}$	A logic low on this input selects external data separation. A logic high or open selects the internal data separator.
26	FLOPPY DISK CLOCK (External Separation)	FDCLOCK	This input receives the externally separated clock when $\overline{\text{XTDS}} = 0$. If $\overline{\text{XTDS}} = 1$, this input should be tied to a logic high.
27	FLOPPY DISK DATA	FDDATA	This input receives the raw read disk data if $\overline{\text{XTDS}} = 1$, or the externally separated data if $\overline{\text{XTDS}} = 0$.
28	HEAD LOAD	HLD	The HLD output controls the loading of the Read-Write head against the media.
29	Track Greater than 43	TG43	This output informs the drive that the Read-Write head is positioned between tracks 44-76. This output is valid only during Read and Write commands.
30	WRITE GATE	WG	This output is made valid when writing is to be performed on the diskette.
31	WRITE DATA	WD	This output contains both clock and data bits of 500 ns duration.
32	Ready	READY	This input indicates disk readiness and is sampled for a logic high before Read or Write commands are performed. If Ready is low, the Read or Write operation is not performed and an interrupt is generated. A Seek operation is performed regardless of the state of Ready. The Ready input appears in inverted format as Status Register bit 7.
33	$\overline{\text{WRITE FAULT}}$	$\overline{\text{WF}}$	This input detects wiring faults indications from the drive. When $\text{WG}=1$ and $\overline{\text{WF}}$ goes low, the current Write command is terminated and the Write Fault status bit is set. The $\overline{\text{WF}}$ input should be made inactive (high) when WG becomes inactive.
34	$\overline{\text{TRACK 00}}$	$\overline{\text{TR00}}$	This input informs the FD1771 that the Read-Write head is positioned over Track 00 when a logic low.
35	$\overline{\text{INDEX PULSE}}$	$\overline{\text{IP}}$	Input, when low for a minimum of 10 usec, informs the FD1771 when an index mark is encountered on the diskette.
36	$\overline{\text{WRITE PROTECT}}$	$\overline{\text{WPRT}}$	This input is sampled whenever a Write command is received. A logic low terminates the command and sets the Write Protect status bit.
37	$\overline{\text{DISK INITIALIZATION}}$	$\overline{\text{DINT}}$	The input is sampled whenever a Write Track command is received. If $\overline{\text{DINT}} = 0$, the operation is terminated and the Write Protect status bit is set.

COMMAND DESCRIPTION

The FD1771 will accept and execute eleven commands. Command words should only be loaded in the Command Register when the Busy status bit is off (status bit 0). The one exception is the Force Interrupt command. Whenever a command is being executed, the Busy status bit is set. When a command is completed, an interrupt is generated and the Busy status bit is reset. The Status Register indicates whether the completed command encountered an error or was fault-free. For ease of discussion, commands are divided into four types. Commands and types are summarized in Table 2.

TYPE 1 COMMANDS

The Type 1 Commands include the RESTORE, SEEK, STEP, STEP-IN, and STEP-OUT commands. Each of the Type 1 Commands contain a rate field (r_0r_1), which determines the stepping motor rate as defined in Table 1, page 4.

The Type 1 Commands contain a head load flag (h) which determines if the head is to be loaded at the

Table 2. COMMAND SUMMARY

TYPE	COMMAND	BITS							
		7	6	5	4	3	2	1	0
I	Restore	0	0	0	0	h	V	r_1	r_0
I	Seek	0	0	0	1	h	V	r_1	r_0
I	Step	0	0	1	u	h	V	r_1	r_0
I	Step In	0	1	0	u	h	V	r_1	r_0
I	Step Out	0	1	1	u	h	V	r_1	r_0
II	Read Command	1	0	0	0	1	1	0	0
II	Write Command	1	0	1	0	1	1	0	a_0
III	Read Address	1	1	0	0	0	E	0	0
III	Read Track	1	1	1	0	0	1	0	\bar{s}
III	Write Track	1	1	1	1	0	1	0	0
IV	Force Interrupt	1	1	0	1	l_3	l_2	l_1	l_4

Note: Bits shown in TRUE form.

Table 3. FLAG SUMMARY

TYPE I
<u>h = Head Load flag (Bit 3)</u> h = 1, Load head at beginning h = 0, Do not load head at beginning
<u>V = Verify flag (Bit 2)</u> V = 1, Verify on last track V = 0, No verify
<u>r_1r_0 = Stepping motor rate (Bits 1-0)</u> Refer to Table 1 for rate summary
<u>u = Update flag (Bit 4)</u> u = 1, Update Track register u = 0, No update

Table 4. FLAG SUMMARY

TYPE II
<u>m = Multiple Record flag (Bit 4)</u> m=0, Single Record m=1, Multiple Records
<u>b = Block length flag (Bit 3)</u> b=1, IBM format (128 to 1024 bytes) b=0, Non-IBM format (16 to 4096 bytes)
<u>a_1a_0 = Data Address Mark (Bits 1-0)</u> a_1a_0 = 00, FB (Data Mark) a_1a_0 = 01, FA (User defined) a_1a_0 = 10, F9 (User defined) a_1a_0 = 11, F8 (Deleted Data Mark)

Table 5. FLAG SUMMARY

TYPE III
<u>s = Synchronize flag (Bit 0)</u> \bar{s} =0, Synchronize to AM \bar{s} =1, Do Not Synchronize to AM
TYPE IV
<u>l_i = Interrupt Condition flags (Bits 3-0)</u> l_0 =1, Not Ready to Ready Transition l_1 =1, Ready to Not Ready Transition l_2 =1, Index Pulse l_3 =1, Immediate interrupt
<u>E = Enable HLD and 10 msec Delay</u> E=1, Enable HLD, HLT and 10 msec Delay E=0, Head is assumed Engaged and there is no 10 msec Delay

beginning of the command. If h=1, the head is loaded at the beginning of the command (HLD output is made active). If h=0, HLD is deactivated. Once the head is loaded, the head will remain engaged until the FD1771 receives a command that specifically disengages the head. If the FD1771 does not receive any commands after two revolutions of the disk, the head will be automatically disengaged (HLD made inactive). The Head Load Timing Input is sampled after a 10 ms delay, when reading or writing on the disk is to occur.

The Type 1 Commands also contain a verification (V) flag which determines if a verification operation is to take place on the destination track. If V=1, a verification is performed; if V=0, no verification is performed.

During verification, the head is loaded and after an internal 10 ms delay, the HLT input is sampled. When

ELECTRICAL CHARACTERISTICS**Maximum Ratings**

V _{DD} with respect to V _{BB} (Ground)	+20 to -0.3V
Max Voltage to any input with respect to V _{BB}	+20 to -0.3V
Operating Temperature	0°C to 70°C
Storage Temperature	-55°C to +125°C

OPERATING CHARACTERISTICS (DC)

T _A = 0°C to 70°C, V _{DD} = +12.0V ± .6V,
V _{BB} = -5.0 ± .5V, V _{SS} = 0V, V _{CC} = +5V ± .25V
I _{DD} = 10 ma Nominal, I _{CC} = 30 ma Nominal,
I _{BB} = 0.4 μa Nominal

Symbol	Characteristic	Min.	Typ.	Max.	Units	Conditions
I _{LI}	Input Leakage			10	μA	V _{IN} = V _{DD}
I _{LO}	Output Leakage			10	μA	V _{OUT} = V _{DD}
V _{IH}	Input High Voltage	2.6			V	
V _{IL}	Input Low Voltage (All Inputs)			0.8	V	
V _{OH}	Output High Voltage	2.8			V	I _O = -100 μA
V _{OL}	Output Low Voltage			0.45**	V	I _O = 1.0 mA

**Write Gate V_{OL} ≤ 0.5V.

TIMING CHARACTERISTICS

T _A = 0°C to 70°C, V _{DD} = +12V ± .6V,
V _{BB} = -5V ± .25V, V _{SS} = 0V, V _{CC} = +5V ± .25V

NOTE: Timings are given for 2 MHz Clock. For those timings noted, values will double when chip is operated at 1 MHz. Use 1 MHz when using mini-floppy.

Read Operations

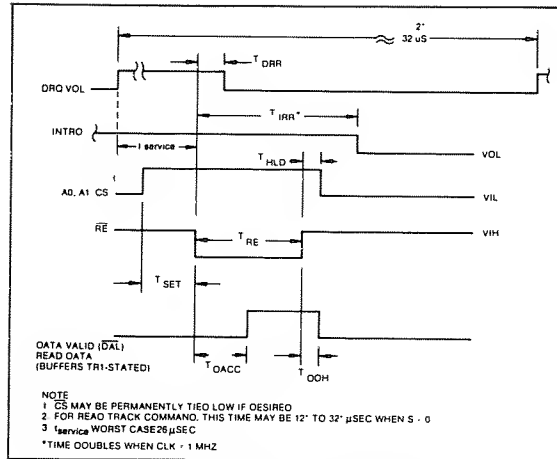
Symbol	Characteristic	Min.	Typ.	Max.	Units	Conditions
TSET	Setup ADDR and CS to \overline{RE}	100			nsec	
THLD	Hold ADDR and CS from \overline{RE}	10			nsec	
TRE	\overline{RE} Pulse Width	500			nsec	C _L = 25 pf
TDRR	DRQ Reset from \overline{RE}			500	nsec	
TIRR	INTRQ Reset from \overline{RE}			3000	nsec	
TDACC	Data Access from \overline{RE}			450	nsec	C _L = 25 pf
TDOH	Data Hold from \overline{RE}	50		150	nsec	C _L = 25 pf

Write Operations

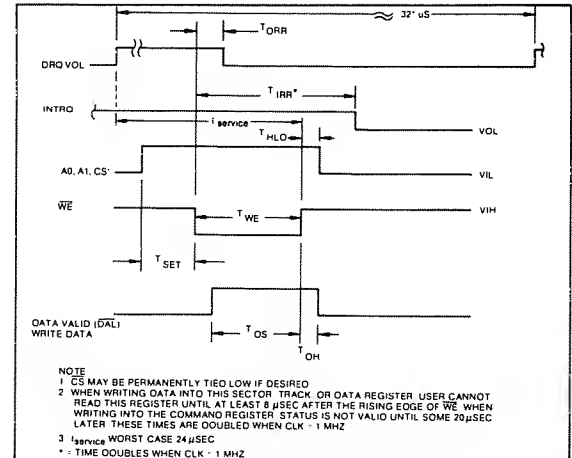
Symbol	Characteristic	Min.	Typ.	Max.	Units	Conditions
TSET	Setup ADDR and CS to \overline{WE}	100			nsec	
THLD	Hold ADDR and CS from \overline{WE}	10			nsec	
TWE	\overline{WE} Pulse Width	350			nsec	
TDRR	DRQ Reset from \overline{WE}			500	nsec	
TIRR	INTRQ Reset from \overline{WE}			3000	nsec	
TDS	Data Setup to \overline{WE}	250			nsec	
TDH	Data Hold from \overline{WE}	150			nsec	See Note

External Data Separation (XTDS = 0)

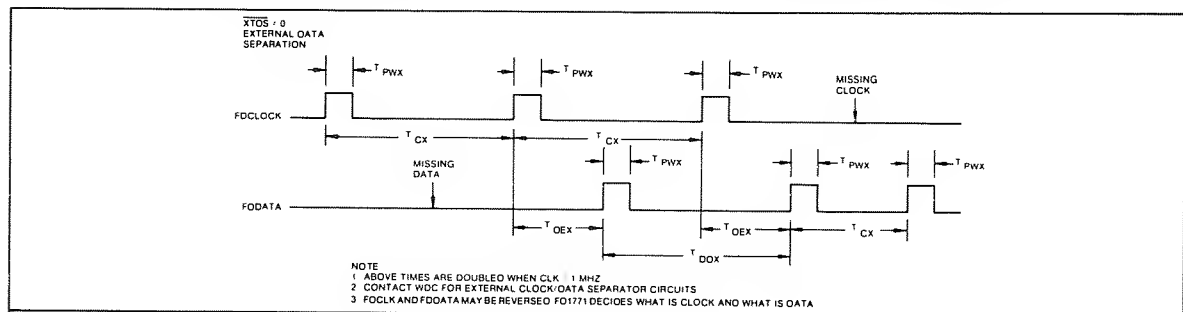
Symbol	Characteristic	Min.	Typ.	Max.	Units	Conditions
TPWX	Pulse Width Read Data & Read Clock	150		350	nsec	
TCX	Clock Cycle External	2500			nsec	
TDEX	Data to Clock	500			nsec	
TDDX	Data to Data Cycle	2500			nsec	



READ ENABLE TIMING



WRITE ENABLE TIMING



READ TIMING (XTDS = 0)

Internal Data Separation (XTDS = 1)

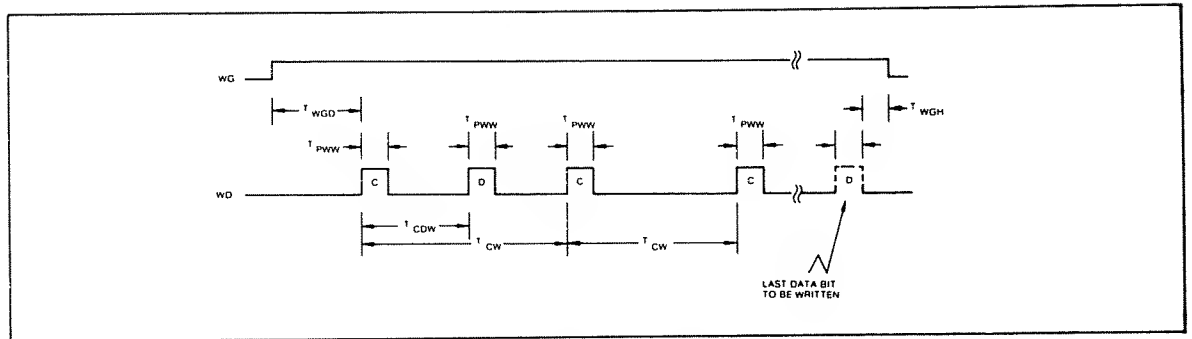
Symbol	Characteristic	Min.	Typ.	Max.	Units	Conditions
TPWI	Pulse Width Data and Clock	150		1000	nsec	
TCI	Clock Cycle Internal	3500		5000	nsec	

Write Data Timing

Symbol	Characteristic	Min.	Typ.	Max.	Units	Conditions
TWGD	Write Gate to Data		1200		nsec	300 nsec ± CLK tolerance
TPWW	Pulse Width Write Data	500		600	nsec	
TCDW	Clock to Data		2000		nsec	± CLK tolerance
TCW	Clock Cycle Write		4000		nsec	± CLK tolerance
TWGH	Write Gate Hold to Data	0		100	nsec	


Miscellaneous Timing

Symbol	Characteristic	Min.	Typ.	Max.	Units	Conditions
TCD ₁	Clock Duty	175			nsec	2 MHz ± 1% See Note These times doubled when CLK = 1 MHz
TCD ₂	Clock Duty	210			nsec	
TSTP	Step Pulse Output	3800		4200	nsec	
TDIR	Direct Setup to Step	24			usec	
TMR	Master Reset Pulse Width	10			usec	
TIP	Index Pulse Width	10			usec	
TWF	Write Fault Pulse Width	10			usec	

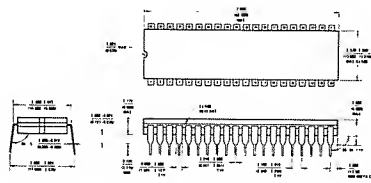


WRITE DATA TIMING

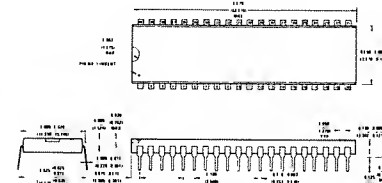
Mode Definition Summary Table

Port Bits	Mode 0		Mode 1		Mode 2
	IN	OUT	IN	OUT	Group A Only
PA ₀	IN	OUT	IN	OUT	Bidirectional  Bidirectional
PA ₁	IN	OUT	IN	OUT	
PA ₂	IN	OUT	IN	OUT	
PA ₃	IN	OUT	IN	OUT	
PA ₄	IN	OUT	IN	OUT	
PA ₅	IN	OUT	IN	OUT	
PA ₆	IN	OUT	IN	OUT	
PA ₇	IN	OUT	IN	OUT	
PB ₀	IN	OUT	IN	OUT	(Mode 0 or Mode 1 only)
PB ₁	IN	OUT	IN	OUT	
PB ₂	IN	OUT	IN	OUT	
PB ₃	IN	OUT	IN	OUT	
PB ₄	IN	OUT	IN	OUT	
PB ₅	IN	OUT	IN	OUT	
PB ₆	IN	OUT	IN	OUT	
PB ₇	IN	OUT	IN	OUT	
PC ₀	IN	OUT	INTR _B	INTR _B	I/O
PC ₁	IN	OUT	IBF _B	OBF _B	I/O
PC ₂	IN	OUT	STB _B	ACK _B	I/O
PC ₃	IN	OUT	INTR _A	INTR _A	INTR _A
PC ₄	IN	OUT	STB _A	I/O	STB _A
PC ₅	IN	OUT	IBF _A	I/O	IBF _A
PC ₆	IN	OUT	I/O	ACK _A	ACK _A
PC ₇	IN	OUT	I/O	OBF _A	OBF _A

Physical Dimensions



Ceramic Dual-in-Line Package [Cerdip (J)]
Order Number INS8255J



Plastic Dual-in-Line Package (N)
Order Number INS8255N



National Semiconductor Corporation
2000 Semiconductor Drive
Santa Clara, CA 95051
Tel: (408) 737-5000
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National Semiconductor GmbH
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NS International Inc., Japan
Mayake Building
1-9 Yotsuya Shinjuku-ku 160
Tokyo, Japan
Tel: (03) 335-3711
TWX: 330-2015 NSCJ-J

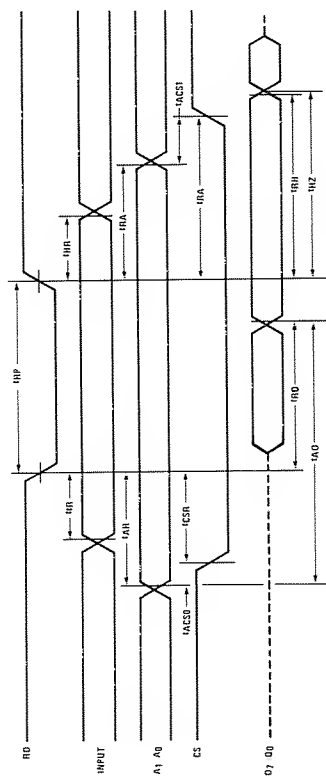
National Semiconductor (Hong Kong) Ltd.
8th Floor
Cheung Kong Electronic Bldg
4 Hing Yip Street
Kwun Tong
Kowloon, Hong Kong
Tel: 3-899235
Telex: 43668 NSEHK HK
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NS Electronics do Brasil
Avda. Brigadier Faria Lima 844
11 Andar Conjunto 1104
Jardim Paulistano
São Paulo, Brasil
Telex: 1121008 Cabine São Paulo

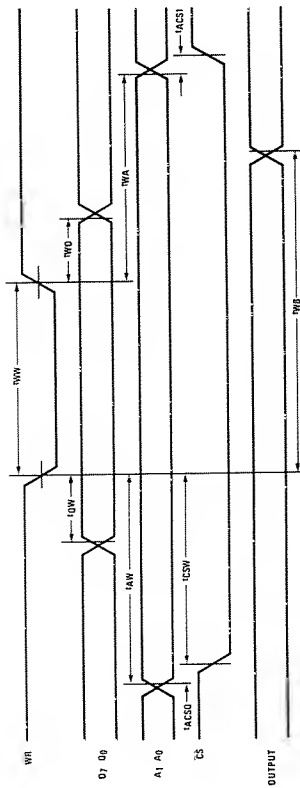
NS Electronics Pty. Ltd.
Cnr. Studd Rd & Main Highway
Glenwater Victoria 3153
Australia
Tel: 03-726-6333
Telex: 32096

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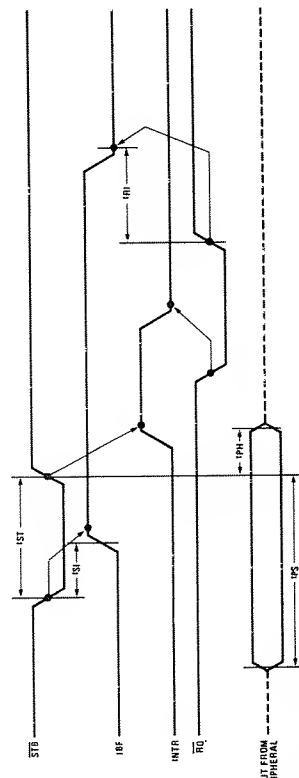
Timing Waveforms



Mode 0 (Basic Input)



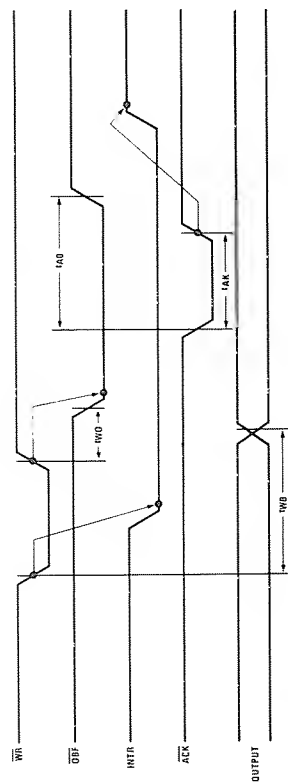
Mode 0 (Basic Output)



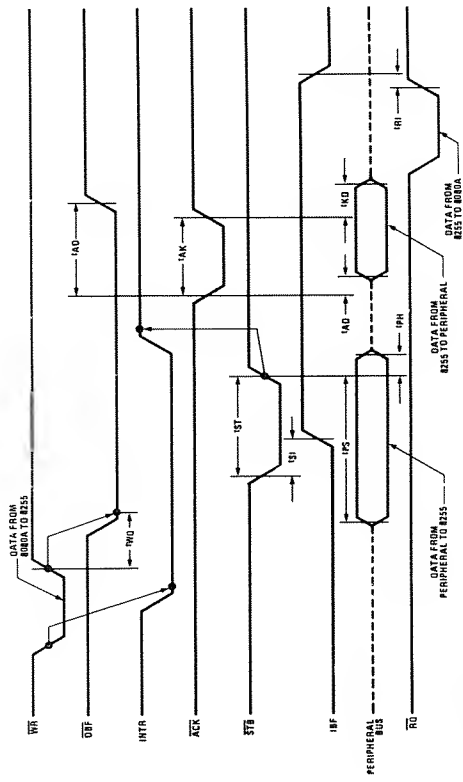
Mode 1 (Strobed Input)

3

Timing Waveforms (cont'd.)



Mode 1 (Strobed Output)



Mode 2 (Bidirectional)

4

INPUT/OUTPUT SIGNALS

Data (D7-D0) Bus, Pins 27-34: This bus comprises eight TRI-STATE® input/output lines. The bus provides bi-directional communication between the INS8255 and the INS8080A. Data is routed to or from the internal data bus buffer upon execution of an OUT or IN Instruction, respectively, by the INS8080A. In addition, control words and status information are transferred through the data bus buffer.

Port A (PA7-PA0), Pins 37-40, 1-4: This 8-bit input/output port forms one 8-bit data output latch/buffer and/or one 8-bit data input latch.

NOTE

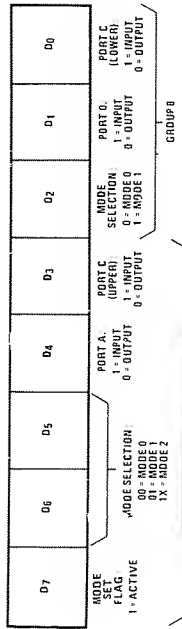
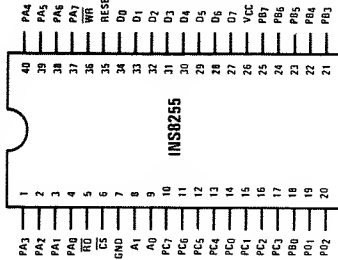
The system software uses a Mode Definition Control Word (see figure) as the second byte of OUT Instruction(s) to program the functional configuration of Ports A through C. Whenever the mode is changed, all output registers (and status flip-flops) are reset.

Port B (PB7-PB0), Pins 18-25: This 8-bit input/output port forms one 8-bit data output latch/buffer or one 8-bit data input buffer.

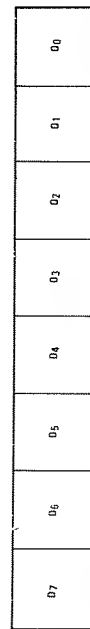
Port C (PC7-PC0), Pins 10-17: This 8-bit input/output port forms one 8-bit data output latch/buffer or one 8-bit data input buffer. The port can be split into two 4-bit ports under the mode control. Each of these 4-bit ports contains a 4-bit latch that may be used for the control and status signals, in conjunction with Ports A

and 8. The system software includes a Bit Set/Reset Control Word (see figure) for setting or resetting any of the eight bits of Port C. When Port C is being used as a status/control for Port A or 8, the Port C bits can be set or reset by using the Bit Set/Reset Control Word as the second byte of OUT Instruction(s).

Pin Configuration



Mode Definition Control Word Format



Bit Set/Reset Control Word Format (Port C Only)

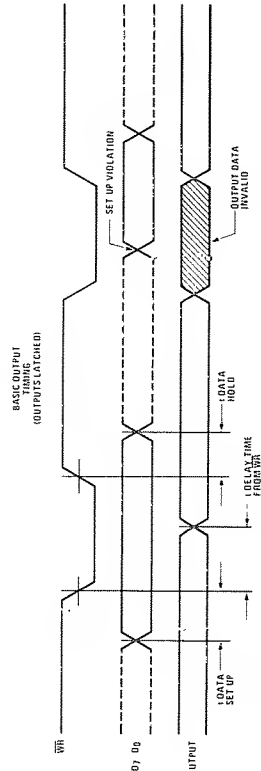
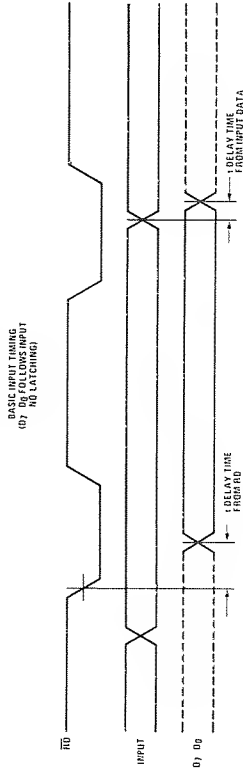
Operating Modes

Mode 0 (Basic Input/Output)

In this mode, simple input and output operations for each of the three ports are provided. No "handshaking" is required; data is simply written to or read from a specified port.

Mode 0 Port Definition Chart

No.	Control Word Bits								Group A			Group B		
	D7	D6	D5	D4	D3	D2	D1	D0	Port A	Port C (Upper)	Port B	Port C (Lower)	Port B	Port C (Lower)
0	1	0	0	0	0	0	0	0	OUTPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT
1	1	0	0	0	0	0	0	1	OUTPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT
2	1	0	0	0	0	0	1	0	OUTPUT	OUTPUT	INPUT	INPUT	OUTPUT	OUTPUT
3	1	0	0	0	0	0	1	1	OUTPUT	OUTPUT	INPUT	INPUT	INPUT	INPUT
4	1	0	0	0	1	0	0	0	OUTPUT	INPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT
5	1	0	0	0	1	0	0	1	OUTPUT	INPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT
6	1	0	0	0	1	0	1	0	OUTPUT	INPUT	INPUT	INPUT	INPUT	INPUT
7	1	0	0	0	1	0	1	1	OUTPUT	INPUT	INPUT	INPUT	INPUT	INPUT
8	1	0	0	1	0	0	0	0	INPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT
9	1	0	0	1	0	0	0	1	INPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT
10	1	0	0	1	0	0	1	0	INPUT	OUTPUT	INPUT	INPUT	INPUT	INPUT
11	1	0	0	1	0	0	1	1	INPUT	OUTPUT	INPUT	INPUT	INPUT	INPUT
12	1	0	0	1	1	0	0	0	INPUT	INPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT
13	1	0	0	1	1	0	0	1	INPUT	INPUT	OUTPUT	OUTPUT	OUTPUT	OUTPUT
14	1	0	0	1	1	0	1	0	INPUT	INPUT	INPUT	INPUT	INPUT	INPUT
15	1	0	0	1	1	0	1	1	INPUT	INPUT	INPUT	INPUT	INPUT	INPUT

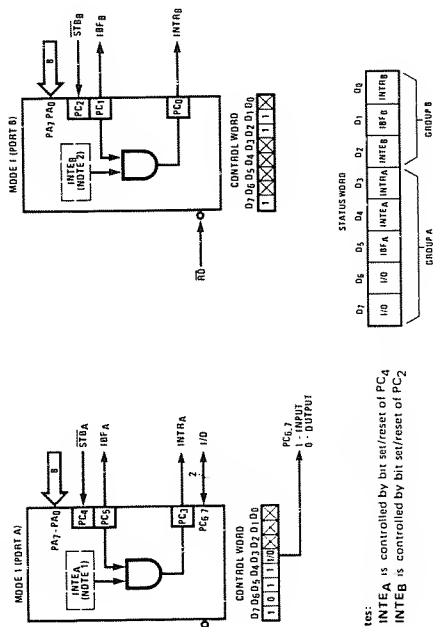


Mode 0 Timing

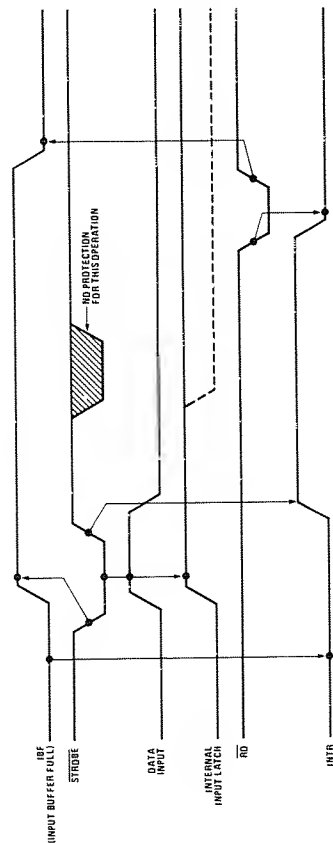
Operating Modes (cont'd.)

Mode 1 (Strobed Input/Output)

In this mode, a means for transferring input/output data to or from a specified port in conjunction with strobes or "handshaking" signals is provided. Port A and Port B use the lines on Port C to generate or accept these "handshaking" signals in Mode 1. The programmer can read the contents of Port C to test or verify the status of each peripheral device. Since no special instruction is provided in the INS8080A microcomputer system to read the Port C status information, a normal read operation must be executed to perform this function.

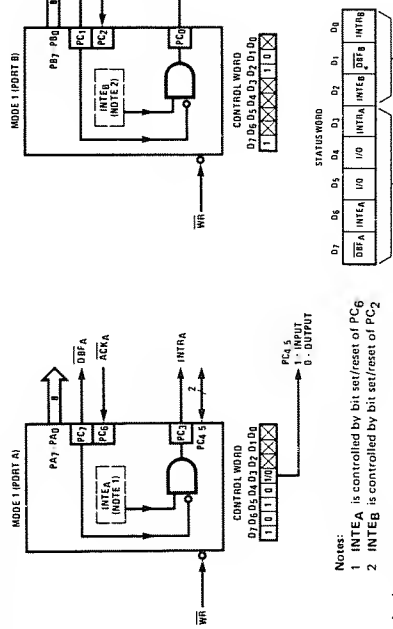


Mode 1 Input

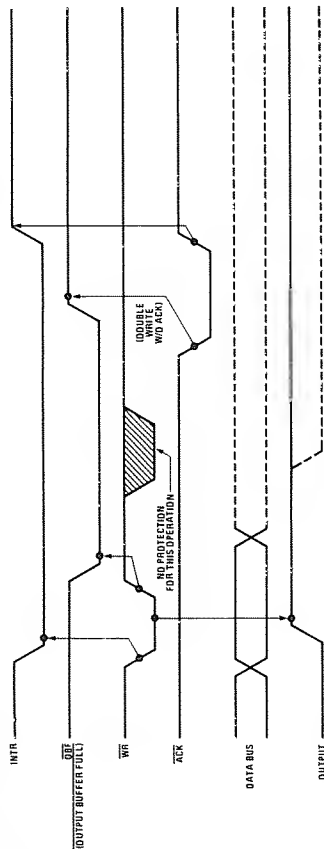


Mode 1 Input Timing

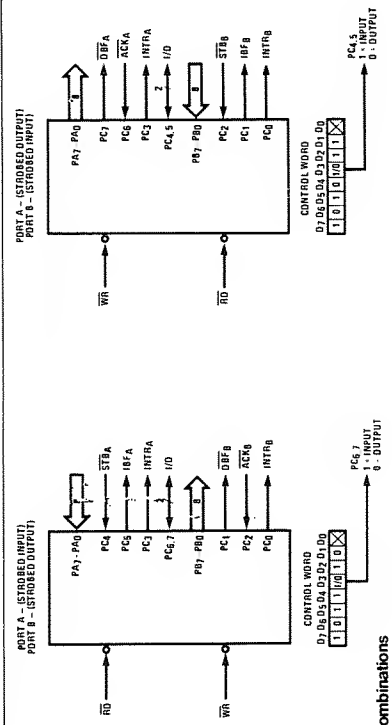
Operating Modes (cont'd.)



Mode 1 Output



Mode 1 Output Timing

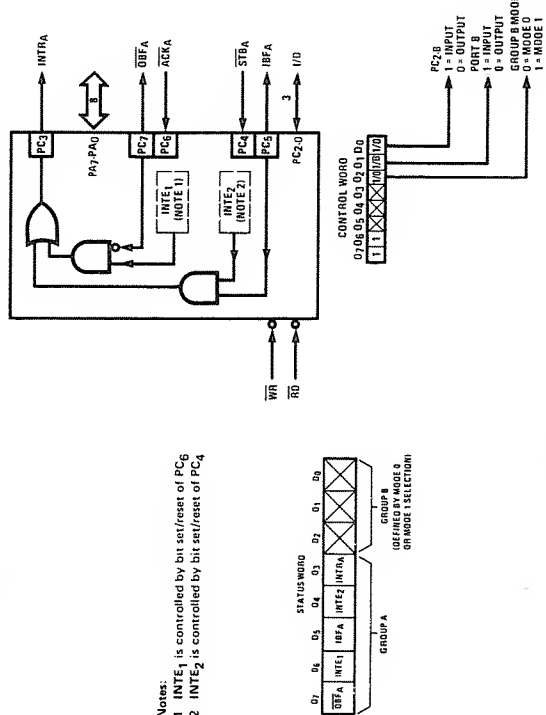


Mode 1 Combinations

Operating Modes (cont'd.)

Mode 2 (Strobed Bidirectional Bus Input/Output)

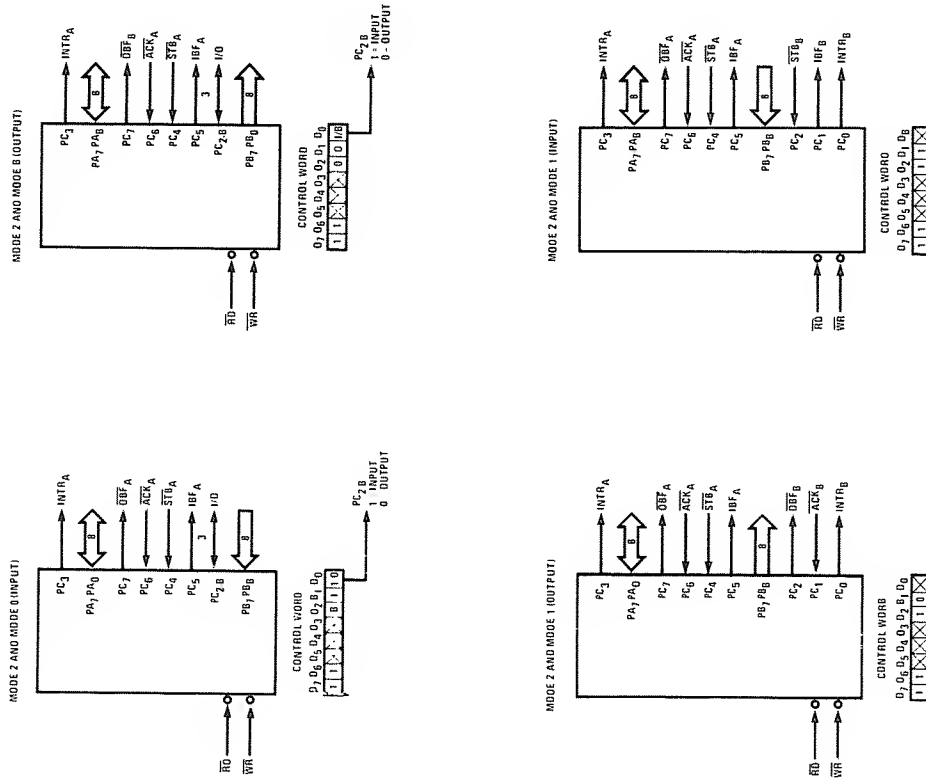
This mode enables communication with a peripheral device or structure on a single 8-bit bus for both transmitting and receiving data (bidirectional bus input/output). "Handshaking" signals are provided to maintain proper bus flow discipline in a manner similar to Mode 1. In addition, interrupt generation and enable/disable functions are available in Mode 2.



- Notes:
- 1. INTA is controlled by bit sel/reset of PC6
 - 2. INT1 is controlled by bit sel/reset of PC4

Mode 2

Operating Modes (cont'd.)



Mode 2 Combinations

Z80[®]-CPU Z80A-CPU



Product Specification MARCH 1978

The Zilog Z80 product line is a complete set of microcomputer components, development systems and support software. The Z80 microcomputer component set includes all of the circuits necessary to build high-performance microcomputer systems with virtually no other logic and a minimum number of low cost standard memory elements.

The Z80 and Z80A CPU's are third generation single chip microprocessors with unrivaled computational power. This increased computational power results in higher system through-put and more efficient memory utilization when compared to second generation microprocessors. In addition, the Z80 and Z80A CPU's are very easy to implement into a system because of their single voltage requirement plus all output signals are fully decoded and timed to control standard memory or peripheral circuits. The circuit is implemented using an N-channel, ion implanted, silicon gate MOS process.

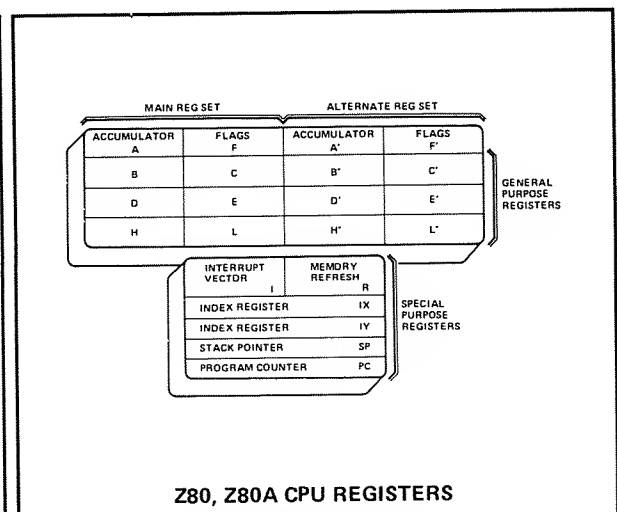
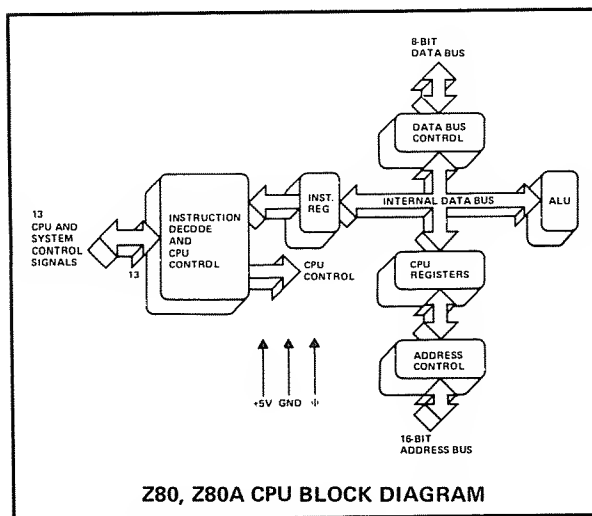
Figure 1 is a block diagram of the CPU, Figure 2 details the internal register configuration which contains 208 bits of Read/Write memory that are accessible to the programmer. The registers include two sets of six general purpose registers that may be used individually as 8-bit registers or as 16-bit register pairs. There are also two sets of accumulator and flag registers. The programmer has access to either set of main or alternate registers through a group of exchange instructions. This alternate set allows foreground/background mode of operation or may be reserved for very fast Interrupt response. Each CPU also contains a 16-bit stack pointer which permits simple implementation of

multiple level interrupts, unlimited subroutine nesting and simplification of many types of data handling.

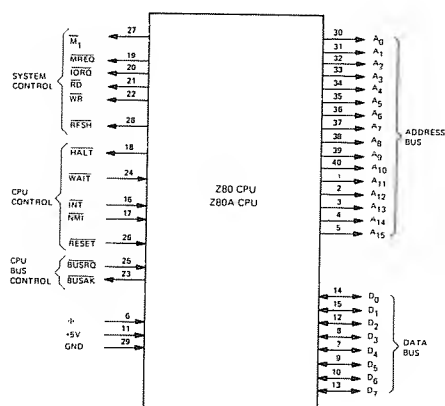
The two 16-bit index registers allow tabular data manipulation and easy implementation of relocatable code. The Refresh register provides for automatic, totally transparent refresh of external dynamic memories. The I register is used in a powerful interrupt response mode to form the upper 8 bits of a pointer to a interrupt service address table, while the interrupting device supplies the lower 8 bits of the pointer. An indirect call is then made to this service address.

FEATURES

- Single chip, N-channel Silicon Gate CPU.
- 158 instructions—includes all 78 of the 8080A instructions with total software compatibility. New instructions include 4-, 8- and 16-bit operations with more useful addressing modes such as indexed, bit and relative.
- 17 internal registers.
- Three modes of fast interrupt response plus a non-maskable interrupt.
- Directly interfaces standard speed static or dynamic memories with virtually no external logic.
- 1.0 μ s instruction execution speed.
- Single 5 VDC supply and single-phase 5 volt Clock.
- Out-performs any other single chip microcomputer in 4-, 8-, or 16-bit applications.
- All pins TTL Compatible
- Built-in dynamic RAM refresh circuitry.



Z80, Z80A-CPU Pin Description



Z80, Z80A CPU PIN CONFIGURATION

A₀-A₁₅
(Address Bus) Tri-state output, active high. A₀-A₁₅ constitute a 16-bit address bus. The address bus provides the address for memory (up to 64K bytes) data exchanges and for I/O device data exchanges.

D₀-D₇
(Data Bus) Tri-state input/output, active high. D₀-D₇ constitute an 8-bit bidirectional data bus. The data bus is used for data exchanges with memory and I/O devices.

M₁
(Machine Cycle one) Output, active low. $\overline{M_1}$ indicates that the current machine cycle is the OP code fetch cycle of an instruction execution.

MREQ
(Memory Request) Tri-state output, active low. The memory request signal indicates that the address bus holds a valid address for a memory read or memory write operation.

IORQ
(Input/Output Request) Tri-state output, active low. The \overline{IORQ} signal indicates that the lower half of the address bus holds a valid I/O address for a I/O read or write operation. An \overline{IORQ} signal is also generated when an interrupt is being acknowledged to indicate that an interrupt response vector can be placed on the data bus.

RD
(Memory Read) Tri-state output, active low. \overline{RD} indicates that the CPU wants to read data from memory or an I/O device. The addressed I/O device or memory should use this signal to gate data onto the CPU data bus.

WR
(Memory Write) Tri-state output, active low. \overline{WR} indicates that the CPU data bus holds valid data to be stored in the addressed memory or I/O device.

RFSH
(Refresh) Output, active low. \overline{RFSH} indicates that the lower 7 bits of the address bus contain a refresh address for dynamic memories and the current MREQ signal should be used to do a refresh read to all dynamic memories.

HALT
(Halt state) Output, active low. \overline{HALT} indicates that the CPU has executed a HALT software instruction and is awaiting either a non-maskable or a maskable interrupt (with the mask enabled) before operation can resume. While halted, the CPU executes NOP's to maintain memory refresh activity.

WAIT
(Wait) Input, active low. \overline{WAIT} indicates to the Z-80 CPU that the addressed memory or I/O devices are not ready for a data transfer. The CPU continues to enter wait states for as long as this signal is active.

INT
(Interrupt Request) Input, active low. The Interrupt Request signal is generated by I/O devices. A request will be honored at the end of the current instruction if the internal software controlled interrupt enable flip-flop (IFF) is enabled.

NMI
(Non Maskable Interrupt) Input, active low. The non-maskable interrupt request line has a higher priority than \overline{INT} and is always recognized at the end of the current instruction, independent of the status of the interrupt enable flip-flop. NMI automatically forces the Z-80 CPU to restart to location 0066H.

RESET Input, active low. \overline{RESET} initializes the CPU as follows: reset interrupt enable flip-flop, clear PC and registers I and R and set interrupt to 8080A mode. During reset time, the address and data bus go to a high impedance state and all control output signals go to the inactive state.

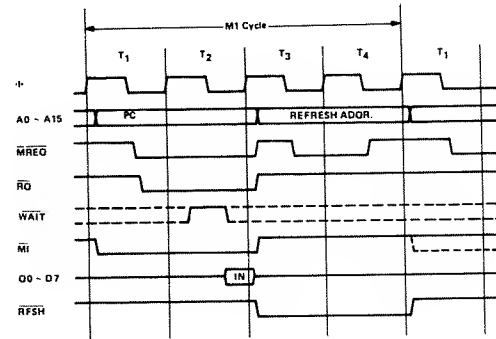
BUSRQ
(Bus Request) Input, active low. The bus request signal has a higher priority than \overline{NMI} and is always recognized at the end of the current machine cycle and is used to request the CPU address bus, data bus and tri-state output control signals to go to a high impedance state so that other devices can control these busses.

BUSAK
(Bus Acknowledge) Output, active low. Bus acknowledge is used to indicate to the requesting device that the CPU address bus, data bus and tri-state control bus signals have been set to their high impedance state and the external device can now control these signals.

Timing Waveforms

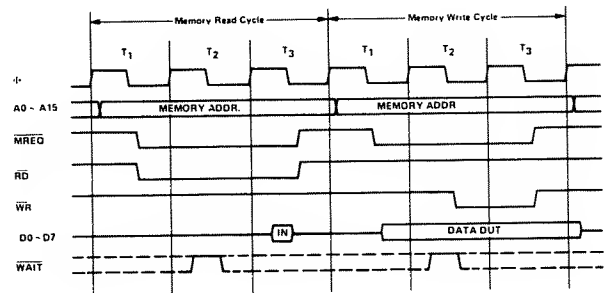
INSTRUCTION OP CODE FETCH

The program counter content (PC) is placed on the address bus immediately at the start of the cycle. One half clock time later \overline{MREQ} goes active. The falling edge of \overline{MREQ} can be used directly as a chip enable to dynamic memories. \overline{RD} when active indicates that the memory data should be enabled onto the CPU data bus. The CPU samples data with the rising edge of the clock state T_3 . Clock states T_3 and T_4 of a fetch cycle are used to refresh dynamic memories while the CPU is internally decoding and executing the instruction. The refresh control signal \overline{RFSH} indicates that a refresh read of all dynamic memories should be accomplished.



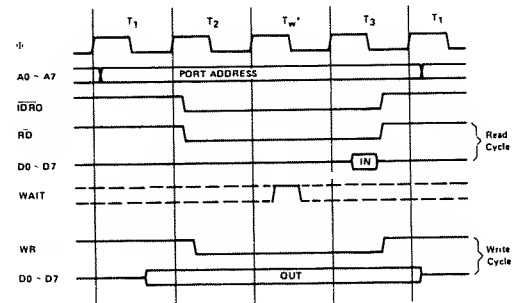
MEMORY READ OR WRITE CYCLES

Illustrated here is the timing of memory read or write cycles other than an OP code fetch (M_1 cycle). The \overline{MREQ} and \overline{RD} signals are used exactly as in the fetch cycle. In the case of a memory write cycle, the \overline{MREQ} also becomes active when the address bus is stable so that it can be used directly as a chip enable for dynamic memories. The \overline{WR} line is active when data on the data bus is stable so that it can be used directly as a R/W pulse to virtually any type of semiconductor memory.



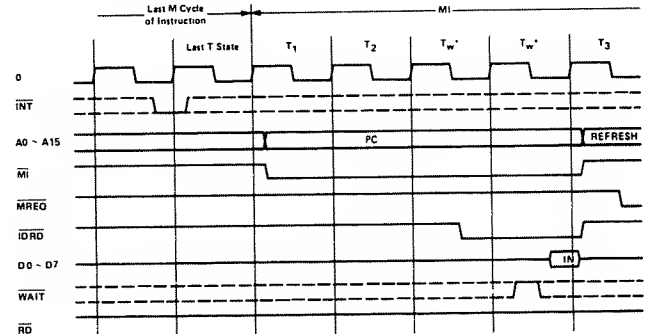
INPUT OR OUTPUT CYCLES

Illustrated here is the timing for an I/O read or I/O write operation. Notice that during I/O operations a single wait state is automatically inserted (T_w^*). The reason for this is that during I/O operations this extra state allows sufficient time for an I/O port to decode its address and activate the WAIT line if a wait is required.



INTERRUPT REQUEST/ACKNOWLEDGE CYCLE

The interrupt signal is sampled by the CPU with the rising edge of the last clock at the end of any instruction. When an interrupt is accepted, a special M_1 cycle is generated. During this M_1 cycle, the \overline{IORQ} signal becomes active (instead of \overline{MREQ}) to indicate that the interrupting device can place an 8-bit vector on the data bus. Two wait states (T_w^*) are automatically added to this cycle so that a ripple priority interrupt scheme, such as the one used in the Z80 peripheral controllers, can be easily implemented.



A.C. Characteristics

Z80-CPU

$T_A = 0^\circ\text{C}$ to 70°C , $V_{CC} = +5V \pm 5\%$, Unless Otherwise Noted.

Signal	Symbol	Parameter	Min	Max	Unit	Test Condition
Φ	t_c	Clock Period	.4	11.21	μsec	
	$t_w(\Phi H)$	Clock Pulse Width, Clock High	180	[E]	nsec	
	$t_w(\Phi L)$	Clock Pulse Width, Clock Low	180	2000	nsec	
	$t_{r,f}$	Clock Rise and Fall Time		30	nsec	
A_{0-15}	$t_D(AD)$	Address Output Delay		145	nsec	$C_L = 50\text{pF}$
	$t_F(AD)$	Delay to Float		110	nsec	
	t_{acm}	Address Stable Prior to \overline{MREQ} (Memory Cycle)	111		nsec	
	t_{aci}	Address Stable Prior to \overline{IORQ} , \overline{RD} or \overline{WR} (I/O Cycle)	121		nsec	
	t_{ca}	Address Stable From \overline{RD} , \overline{WR} , \overline{IORQ} or \overline{MREQ}	131		nsec	
	t_{caf}	Address Stable From \overline{RD} or \overline{WR} During Float	141		nsec	
D_{0-7}	$t_D(D)$	Data Output Delay		230	nsec	$C_L = 50\text{pF}$
	$t_F(D)$	Delay to Float During Write Cycle		90	nsec	
	$t_S(\Phi D)$	Data Setup Time to Rising Edge of Clock During M1 Cycle	50		nsec	
	$t_{SF}(\Phi D)$	Data Setup Time to Falling Edge of Clock During M2 to M5	60		nsec	
	t_{dem}	Data Stable Prior to \overline{WR} (Memory Cycle)	151		nsec	
	t_{dei}	Data Stable Prior to \overline{WR} (I/O Cycle)	161		nsec	
	t_{cdf}	Data Stable From \overline{WR}	171		nsec	
	t_H	Any Hold Time for Setup Time	0		nsec	
\overline{MREQ}	$t_{DL\Phi}(\overline{MR})$	\overline{MREQ} Delay From Falling Edge of Clock, \overline{MREQ} Low		100	nsec	$C_L = 50\text{pF}$
	$t_{DH\Phi}(\overline{MR})$	\overline{MREQ} Delay From Rising Edge of Clock, \overline{MREQ} High		100	nsec	
	$t_w(\overline{MRL})$	\overline{MREQ} Delay From Falling Edge of Clock, \overline{MREQ} High Pulse Width, \overline{MREQ} Low	[8]		nsec	
	$t_w(\overline{MRH})$	\overline{MREQ} Delay From Rising Edge of Clock, \overline{MREQ} High Pulse Width, \overline{MREQ} High	[9]		nsec	
\overline{IORQ}	$t_{DL\Phi}(\overline{IR})$	\overline{IORQ} Delay From Rising Edge of Clock, \overline{IORQ} Low		90	nsec	$C_L = 50\text{pF}$
	$t_{DH\Phi}(\overline{IR})$	\overline{IORQ} Delay From Falling Edge of Clock, \overline{IORQ} Low		110	nsec	
	$t_{DL\Phi}(\overline{IR})$	\overline{IORQ} Delay From Rising Edge of Clock, \overline{IORQ} High		100	nsec	
	$t_{DH\Phi}(\overline{IR})$	\overline{IORQ} Delay From Falling Edge of Clock, \overline{IORQ} High		110	nsec	
\overline{RD}	$t_{DL\Phi}(\overline{RD})$	\overline{RD} Delay From Rising Edge of Clock, \overline{RD} Low		100	nsec	$C_L = 50\text{pF}$
	$t_{DL\Phi}(\overline{RD})$	\overline{RD} Delay From Falling Edge of Clock, \overline{RD} Low		130	nsec	
	$t_{DH\Phi}(\overline{RD})$	\overline{RD} Delay From Rising Edge of Clock, \overline{RD} High		100	nsec	
	$t_{DH\Phi}(\overline{RD})$	\overline{RD} Delay From Falling Edge of Clock, \overline{RD} High		110	nsec	
\overline{WR}	$t_{DL\Phi}(\overline{WR})$	\overline{WR} Delay From Rising Edge of Clock, \overline{WR} Low		80	nsec	$C_L = 50\text{pF}$
	$t_{DL\Phi}(\overline{WR})$	\overline{WR} Delay From Falling Edge of Clock, \overline{WR} Low		90	nsec	
	$t_{DH\Phi}(\overline{WR})$	\overline{WR} Delay From Rising Edge of Clock, \overline{WR} High		100	nsec	
	$t_w(\overline{WRL})$	\overline{WR} Delay From Falling Edge of Clock, \overline{WR} High Pulse Width, \overline{WR} Low	[10]		nsec	
$\overline{M1}$	$t_{DL}(\overline{M1})$	$\overline{M1}$ Delay From Rising Edge of Clock, $\overline{M1}$ Low		130	nsec	$C_L = 50\text{pF}$
	$t_{DH}(\overline{M1})$	$\overline{M1}$ Delay From Rising Edge of Clock, $\overline{M1}$ High		130	nsec	
\overline{RFSH}	$t_{DL}(\overline{RF})$	\overline{RFSH} Delay From Rising Edge of Clock, \overline{RFSH} Low		180	nsec	$C_L = 50\text{pF}$
	$t_{DH}(\overline{RF})$	\overline{RFSH} Delay From Rising Edge of Clock, \overline{RFSH} High		150	nsec	
\overline{WAIT}	$t_s(\overline{WT})$	\overline{WAIT} Setup Time to Falling Edge of Clock	70		nsec	
\overline{HALT}	$t_D(\overline{HT})$	\overline{HALT} Delay Time From Falling Edge of Clock		300	nsec	$C_L = 50\text{pF}$
\overline{INT}	$t_s(\overline{IT})$	\overline{INT} Setup Time to Rising Edge of Clock	80		nsec	
\overline{NMI}	$t_w(\overline{NML})$	Pulse Width, \overline{NMI} Low	80		nsec	
\overline{BUSRQ}	$t_s(\overline{BQ})$	\overline{BUSRQ} Setup Time to Rising Edge of Clock	80		nsec	
$\overline{BUSA\overline{K}}$	$t_{DL}(\overline{BA})$	$\overline{BUSA\overline{K}}$ Delay From Rising Edge of Clock, $\overline{BUSA\overline{K}}$ Low		120	nsec	$C_L = 50\text{pF}$
	$t_{DH}(\overline{BA})$	$\overline{BUSA\overline{K}}$ Delay From Falling Edge of Clock, $\overline{BUSA\overline{K}}$ High		110	nsec	
\overline{RESET}	$t_s(\overline{RS})$	\overline{RESET} Setup Time to Rising Edge of Clock	90		nsec	
	$t_F(C)$	Delay to Float (\overline{MREQ} , \overline{IORQ} , \overline{RD} and \overline{WR})		100	nsec	
	t_{mr}	$\overline{M1}$ Stable Prior to \overline{IORQ} (Interrupt Ack.)	[11]		nsec	

$$[12] \quad t_c = t_w(\Phi H) + t_w(\Phi L) + t_r + t_f$$

$$[1] \quad t_{acm} = t_w(\Phi H) + t_r - 75$$

$$[2] \quad t_{aci} = t_c - 80$$

$$[3] \quad t_{ca} = t_w(\Phi L) + t_r - 40$$

$$[4] \quad t_{caf} = t_w(\Phi L) + t_r - 60$$

$$[5] \quad t_{dem} = t_c - 210$$

$$[6] \quad t_{dei} = t_w(\Phi L) + t_r - 210$$

$$[7] \quad t_{cdf} = t_w(\Phi L) + t_r - 80$$

$$[8] \quad t_w(\overline{MRL}) = t_c - 40$$

$$[9] \quad t_w(\overline{MRH}) = t_w(\Phi H) + t_r - 30$$

$$[10] \quad t_w(\overline{WRL}) = t_c - 40$$

$$[11] \quad t_{mr} = 2t_c + t_w(\Phi H) + t_r - 80$$

NOTES:

A Data should be enabled onto the CPU data bus when \overline{RD} is active. During interrupt acknowledge data should be enabled when $\overline{M1}$ and \overline{IORQ} are both active.
 B All control signals are internally synchronized, so they may be totally asynchronous with respect to the clock.

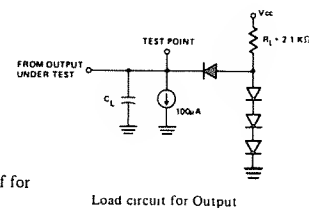
C The \overline{RESET} signal must be active for a minimum of 3 clock cycles.

D Output Delay vs. Loaded Capacitance

$T_A = 70^\circ\text{C}$ $V_{CC} = +5V \pm 5\%$

Add 10nsec delay for each 50pf increase in load up to a maximum of 200pf for the data bus & 100pf for address & control lines

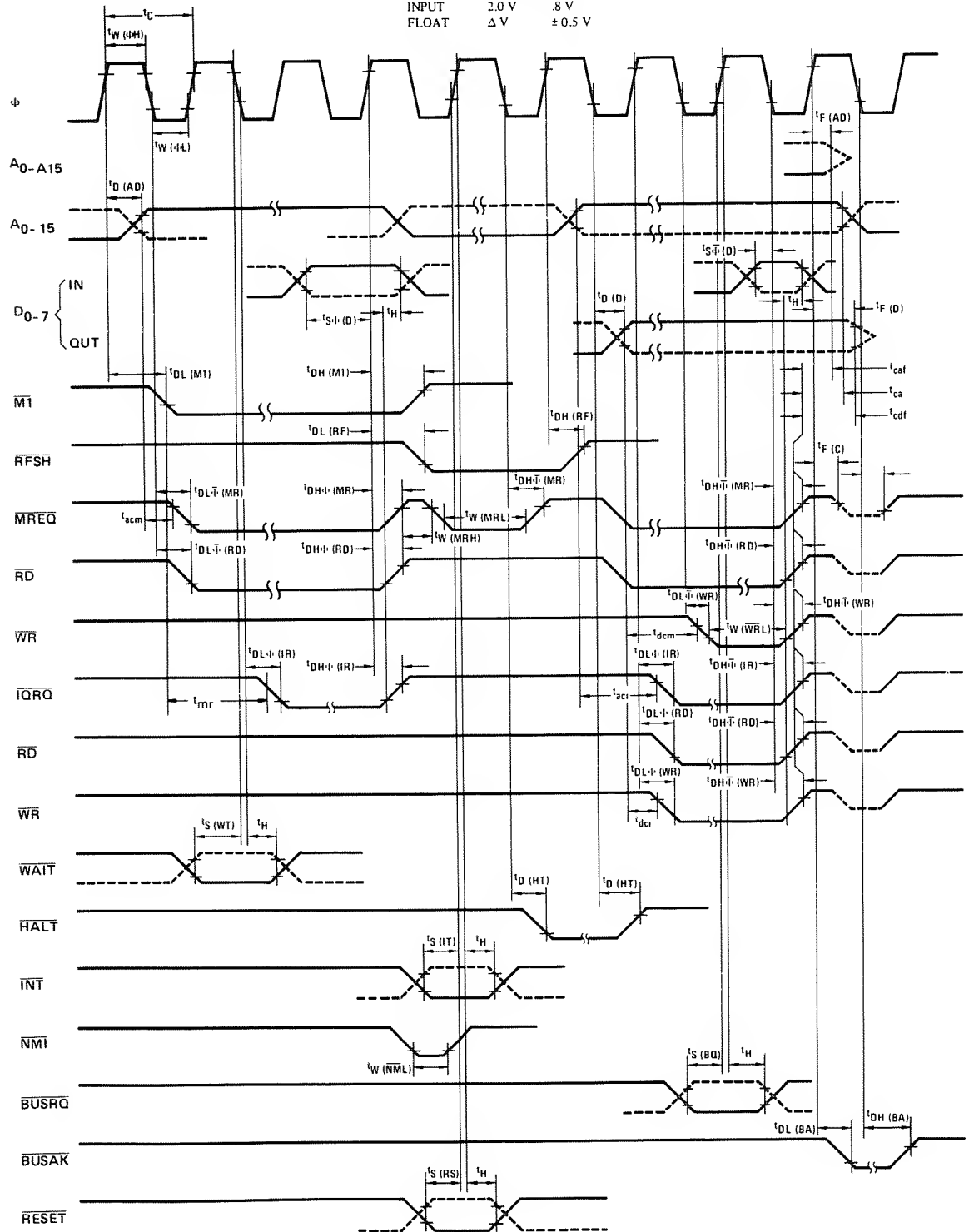
E Although static by design, testing guarantees $t_w(\Phi H)$ of 200 μsec maximum



A.C. Timing Diagram

Timing measurements are made at the following voltages, unless otherwise specified:

	"1"	"0"
CLOCK	$V_{CC} - .6V$.45V
OUTPUT	2.0 V	.8 V
INPUT	2.0 V	.8 V
FLOAT	ΔV	$\pm 0.5 V$



Absolute Maximum Ratings

Temperature Under Bias	Specified operating range.
Storage Temperature	-65°C to +150°C
Voltage On Any Pin with Respect to Ground	-0.3V to +7V
Power Dissipation	1.5W

*Comment

Stresses above those listed under "Absolute Maximum Rating" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other condition above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

Note: For Z80-CPU all AC and DC characteristics remain the same for the military grade parts except I_{CC} .

$$I_{CC} = 200 \text{ mA}$$

Z80-CPU D.C. Characteristics

$T_A = 0^\circ\text{C}$ to 70°C , $V_{CC} = 5V \pm 5\%$ unless otherwise specified

Symbol	Parameter	Min.	Typ.	Max.	Unit	Test Condition
V_{ILC}	Clock Input Low Voltage	-0.3		0.45	V	
V_{IHC}	Clock Input High Voltage	$V_{CC} - 0.6$		$V_{CC} + 0.3$	V	
V_{IL}	Input Low Voltage	-0.3		0.8	V	
V_{IH}	Input High Voltage	2.0		V_{CC}	V	
V_{OL}	Output Low Voltage			0.4	V	$I_{OL} = 1.8 \text{ mA}$
V_{OH}	Output High Voltage	2.4			V	$I_{OH} = -250 \mu\text{A}$
I_{CC}	Power Supply Current			150	mA	
i_{LI}	Input Leakage Current			10	μA	$V_{IN} = 0 \text{ to } V_{CC}$
I_{LOH}	Tri-State Output Leakage Current in Float			10	μA	$V_{OUT} = 2.4 \text{ to } V_{CC}$
I_{LOL}	Tri-State Output Leakage Current in Float			-10	μA	$V_{OUT} = 0.4 \text{ V}$
I_{LD}	Data Bus Leakage Current in Input Mode			± 10	μA	$0 \leq V_{IN} \leq V_{CC}$

Capacitance

$T_A = 25^\circ\text{C}$, $f = 1 \text{ MHz}$,

unmeasured pins returned to ground

Symbol	Parameter	Max.	Unit
C_Φ	Clock Capacitance	35	pF
C_{IN}	Input Capacitance	5	pF
C_{OUT}	Output Capacitance	10	pF

Z80-CPU

Ordering Information

C — Ceramic

P — Plastic

S — Standard $5V \pm 5\%$ 0° to 70°C

E — Extended $5V \pm 5\%$ -40° to 85°C

M — Military $5V \pm 10\%$ -55° to 125°C

Z80A-CPU D.C. Characteristics

$T_A = 0^\circ\text{C}$ to 70°C , $V_{CC} = 5V \pm 5\%$ unless otherwise specified

Symbol	Parameter	Min.	Typ.	Max.	Unit	Test Condition
V_{ILC}	Clock Input Low Voltage	-0.3		0.45	V	
V_{IHC}	Clock Input High Voltage	$V_{CC} - 0.6$		$V_{CC} + 0.3$	V	
V_{IL}	Input Low Voltage	-0.3		0.8	V	
V_{IH}	Input High Voltage	2.0		V_{CC}	V	
V_{OL}	Output Low Voltage			0.4	V	$I_{OL} = 1.8 \text{ mA}$
V_{OH}	Output High Voltage	2.4			V	$I_{OH} = -250 \mu\text{A}$
I_{CC}	Power Supply Current		90	200	mA	
i_{LI}	Input Leakage Current			10	μA	$V_{IN} = 0 \text{ to } V_{CC}$
I_{LOH}	Tri-State Output Leakage Current in Float			10	μA	$V_{OUT} = 2.4 \text{ to } V_{CC}$
I_{LOL}	Tri-State Output Leakage Current in Float			-10	μA	$V_{OUT} = 0.4 \text{ V}$
I_{LD}	Data Bus Leakage Current in Input Mode			± 10	μA	$0 \leq V_{IN} \leq V_{CC}$

Capacitance

$T_A = 25^\circ\text{C}$, $f = 1 \text{ MHz}$,

unmeasured pins returned to ground

Symbol	Parameter	Max.	Unit
C_Φ	Clock Capacitance	35	pF
C_{IN}	Input Capacitance	5	pF
C_{OUT}	Output Capacitance	10	pF

Z80A-CPU

Ordering Information

C — Ceramic

P — Plastic

S — Standard $5V \pm 5\%$ 0° to 70°C

A.C. Characteristics

Z80A-CPU

$T_A = 0^\circ\text{C}$ to 70°C , $V_{CC} = +5V \pm 5\%$, Unless Otherwise Noted.

Signal	Symbol	Parameter	Min	Max	Unit	Test Condition
ϕ	t_c	Clock Period	.25	[12]	μsec	
	$t_{w(\phi H)}$	Clock Pulse Width, Clock High	110	[E]	nsec	
	$t_{w(\phi L)}$	Clock Pulse Width, Clock Low	110	2000	nsec	
	$t_{r, f}$	Clock Rise and Fall Time		.30	nsec	
A_{0-15}	$t_D(AD)$	Address Output Delay		110	nsec	
	$t_F(AD)$	Delay to Float		90	nsec	
	t_{acm}	Address Stable Prior to \overline{MREQ} (Memory Cycle)	[11]		nsec	
	t_{aci}	Address Stable Prior to \overline{IORQ} , \overline{RD} or \overline{WR} (I/O Cycle)	[2]		nsec	
	t_{ca}	Address Stable from \overline{RD} , \overline{WR} , \overline{IORQ} or \overline{MREQ}	[3]		nsec	
	t_{caf}	Address Stable From \overline{RD} or \overline{WR} During Float	[4]		nsec	
D_{0-7}	$t_D(D)$	Data Output Delay		150	nsec	
	$t_F(D)$	Delay to Float During Write Cycle		90	nsec	
	$t_{SF}(D)$	Data Setup Time to Rising Edge of Clock During M1 Cycle	35		nsec	
	$t_{SF}(D)$	Data Setup Time to Falling Edge of Clock During M2 to M5	50		nsec	
	t_{dcn}	Data Stable Prior to \overline{WR} (Memory Cycle)	[5]		nsec	
	t_{dci}	Data Stable Prior to \overline{WR} (I/O Cycle)	[6]		nsec	
	t_{cdf}	Data Stable From \overline{WR}	[7]		nsec	
	t_H	Any Hold Time for Setup Time		0	nsec	
\overline{MREQ}	$t_{DL\phi}(\overline{MR})$	\overline{MREQ} Delay From Falling Edge of Clock, \overline{MREQ} Low		85	nsec	
	$t_{DH\phi}(\overline{MR})$	\overline{MREQ} Delay From Rising Edge of Clock, \overline{MREQ} High		85	nsec	
	$t_{w}(\overline{MRL})$	Pulse Width, \overline{MREQ} Low	[8]		nsec	
	$t_{w}(\overline{MRH})$	Pulse Width, \overline{MREQ} High	[9]		nsec	
\overline{IORQ}	$t_{DL\phi}(\overline{IR})$	\overline{IORQ} Delay From Rising Edge of Clock, \overline{IORQ} Low		75	nsec	
	$t_{DL\phi}(\overline{IR})$	\overline{IORQ} Delay From Falling Edge of Clock, \overline{IORQ} Low		85	nsec	
	$t_{DH\phi}(\overline{IR})$	\overline{IORQ} Delay From Rising Edge of Clock, \overline{IORQ} High		85	nsec	
	$t_{DH\phi}(\overline{IR})$	\overline{IORQ} Delay From Falling Edge of Clock, \overline{IORQ} High		85	nsec	
\overline{RD}	$t_{DL\phi}(\overline{RD})$	\overline{RD} Delay From Rising Edge of Clock, \overline{RD} Low		85	nsec	
	$t_{DL\phi}(\overline{RD})$	\overline{RD} Delay From Falling Edge of Clock, \overline{RD} Low		95	nsec	
	$t_{DH\phi}(\overline{RD})$	\overline{RD} Delay From Rising Edge of Clock, \overline{RD} High		85	nsec	
	$t_{DH\phi}(\overline{RD})$	\overline{RD} Delay From Falling Edge of Clock, \overline{RD} High		85	nsec	
\overline{WR}	$t_{DL\phi}(\overline{WR})$	\overline{WR} Delay From Rising Edge of Clock, \overline{WR} Low		65	nsec	
	$t_{DL\phi}(\overline{WR})$	\overline{WR} Delay From Falling Edge of Clock, \overline{WR} Low		80	nsec	
	$t_{DH\phi}(\overline{WR})$	\overline{WR} Delay From Rising Edge of Clock, \overline{WR} High		80	nsec	
	$t_w(\overline{WRL})$	Pulse Width, \overline{WR} Low	[10]		nsec	
$\overline{M1}$	$t_{DL}(\overline{M1})$	$\overline{M1}$ Delay From Rising Edge of Clock, $\overline{M1}$ Low		100	nsec	
	$t_{DH}(\overline{M1})$	$\overline{M1}$ Delay From Rising Edge of Clock, $\overline{M1}$ High		100	nsec	
\overline{RFSH}	$t_{DL}(\overline{RF})$	\overline{RFSH} Delay From Rising Edge of Clock, \overline{RFSH} Low		130	nsec	
	$t_{DH}(\overline{RF})$	\overline{RFSH} Delay From Rising Edge of Clock, \overline{RFSH} High		120	nsec	
\overline{WAIT}	$t_s(\overline{WT})$	\overline{WAIT} Setup Time to Falling Edge of Clock	70		nsec	
\overline{HALT}	$t_D(\overline{HT})$	\overline{HALT} Delay Time From Falling Edge of Clock		300	nsec	
\overline{INT}	$t_s(\overline{IT})$	\overline{INT} Setup Time to Rising Edge of Clock	80		nsec	
\overline{NMI}	$t_w(\overline{NML})$	Pulse Width, \overline{NMI} Low	80		nsec	
\overline{BUSRQ}	$t_s(\overline{BQ})$	\overline{BUSRQ} Setup Time to Rising Edge of Clock	50		nsec	
$\overline{BUSA\overline{K}}$	$t_{DL}(\overline{BA})$	$\overline{BUSA\overline{K}}$ Delay From Rising Edge of Clock, $\overline{BUSA\overline{K}}$ Low		100	nsec	
	$t_{DH}(\overline{BA})$	$\overline{BUSA\overline{K}}$ Delay From Falling Edge of Clock, $\overline{BUSA\overline{K}}$ High		100	nsec	
\overline{RESET}	$t_s(\overline{RS})$	\overline{RESET} Setup Time to Rising Edge of Clock	60		nsec	
	$t_F(C)$	Delay to Float (\overline{MREQ} , \overline{IORQ} , \overline{RD} and \overline{WR})		80	nsec	
	t_{mr}	$\overline{M1}$ Stable Prior to \overline{IORQ} (Interrupt Ack.)	[11]		nsec	

$$[12] \quad t_c = t_{w(\phi H)} + t_{w(\phi L)} + t_r + t_f$$

$$[1] \quad t_{acm} = t_{w(\phi H)} + t_f - 65$$

$$[2] \quad t_{aci} = t_c - 70$$

$$[3] \quad t_{ca} = t_{w(\phi L)} + t_f - 50$$

$$[4] \quad t_{caf} = t_{w(\phi L)} + t_f - 45$$

$$[5] \quad t_{dcn} = t_c - 170$$

$$[6] \quad t_{dci} = t_{w(\phi L)} + t_f - 170$$

$$[7] \quad t_{cdf} = t_{w(\phi L)} + t_f - 70$$

$$[8] \quad t_{w}(\overline{MRL}) = t_c - 30$$

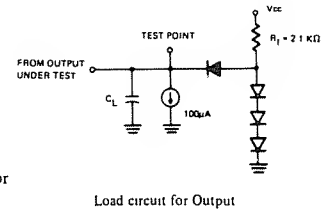
$$[9] \quad t_{w}(\overline{MRH}) = t_{w(\phi H)} + t_f - 20$$

$$[10] \quad t_{w}(\overline{WRL}) = t_c - 30$$

$$[11] \quad t_{mr} = 2t_c + t_{w(\phi H)} + t_f - 65$$

NOTES:

- Data should be enabled onto the CPU data bus when \overline{RD} is active. During interrupt acknowledge data should be enabled when $\overline{M1}$ and \overline{IORQ} are both active.
- All control signals are internally synchronized, so they may be totally asynchronous with respect to the clock.
- The \overline{RESET} signal must be active for a minimum of 3 clock cycles.
- Output Delay vs. Loaded Capacitance
 $T_A = 70^\circ\text{C}$ $V_{CC} = +5V \pm 5\%$
 Add 10nsec delay for each 50pf increase in load up to maximum of 200pf for data bus and 100pf for address & control lines.
- Although static by design, testing guarantees $t_{w(\phi H)}$ of 200 μsec maximum



NOTES

I GLOSSARY

ACCESS The operation of seeking, reading or writing data on a storage unit (in this case, the diskette).

ACCESS TIME The time that elapses between any instruction being given to access some data and that data becoming available for use.

ACTIVE RECORDS TABLE (ART) A table of binary values in which the relative position of a single value determines the status of a record with the same relative position; i.e., the Nth binary number determines the status of the Nth record. **EXAMPLE:** If the 8th binary number in the table is a zero, then the 8th record is inactive. Conversely, if the 8th binary number in the table is a one, then the 8th record is active.

ADDRESS An identification (number, name, or label) for a location in which data is stored.

ALGORITHM A computational procedure.

ALPHANUMERIC (CHARACTERS) A generic term for numeric digits, alphabetic characters, punctuation characters and special characters.

ALPHANUMERIC STRING A group of characters which may include digits, alphabetic characters, punctuation characters and special characters, and may include spaces. (NOTE: a space is a 'character' to the computer, as it must generate a code for spaces as well as symbols.)

ASCII Abbreviation for American Standard Code for Information Interchange. Pronounced: Ass-KEY. Usually refers to a standard method of encoding letter, numeral, symbol and special function characters, as used by the computer industry.

ASSEMBLY LANGUAGE A machine level language for programming, such as Radio Shack's "EDITOR/ASSEMBLER" which uses Z-80 processor mnemonics and automatically 'assembles' machine readable code from the mnemonics.

BASE A quantity of characters for use in each of the digital positions of a numbering system.

BASE 2 The 'BINARY' numbering system consisting of more than one symbol, representing a sum, in which the individual quantity represented by each figure is based on a multiple of 2.

BASE 10 The 'DECIMAL' numbering system - consisting of more than one symbol, representing a sum, in which the individual quantity represented by each symbol is based on a multiple of 10.

BASE 16 The 'HEXADECIMAL' numbering system - consisting of more than one symbol representing a sum, in which the individual quantity represented by each symbol is based on a multiple of 16.

BINARY See 'BASE 2'

BIT A single 'BINARY' digit whose value is 'zero' or 'one'.

BOOLEAN A form of algebra applied to binary numbers which is similar in form to ordinary algebra. It is especially useful for logical analysis of binary numbers as used in computers.

'BOOT' - BOOTSTRAP A machine language program file that is put onto every diskette by the 'FORMAT' routine. This routine is invoked when reset or power-on occurs. It automatically loads the necessary programs (SYS0/SYS) to cause the computer to respond to the DOS commands; i.e., the machine is 'BOOTSTRAPPED' or 'BOOTED' into operation.

BUFFER A small area of memory used for the temporary storage of data to be processed.

BUFFER TRACK A track on a diskette used for the temporary storage of data or program material during a recovery process.

BUG A software fault that results in the malfunction of a program. May also refer to hardware malfunctions.

BYTE Eight 'BITS'. A 'BYTE' may represent any numerical value between '0' and '255'.

CLOBBERED A slang term referring to the non-operation of software, hardware, computer device, or storage media (such as disk) usually as the result of a program or hardware error.

COMMAND FILE A file consisting of a list of commands, to be executed in sequence.

CONTIGUOUS Adjacent or adjoining.

CONTROL CODE In programming, instructions which determine conditional jumps are often referred to as control instructions and the time sequence of execution of instructions is called the flow of control.

CRC ERROR Cyclic Redundancy Check. A means of checking for errors by using redundant information used primarily to check disk I/O on the TRS-80.

DATA BASE A collection of interrelated data stored together with controlled redundancy to serve one or more applications. The data are stored so that they are independent of programs which use the data. A common and controlled approach is used in adding new data and in modifying and retrieving existing data within a data base. A system is said to contain a collection of data-based information if they are disjoint in structure.

DATA BASE MANAGEMENT SYSTEM The collection of software required for using a data base.

DATA ELEMENT Synonymous with 'DATA ITEM' or 'FIELD'

DATA TYPE The form in which data is stored; i.e., integer, single precision, double precision, 'alphanumeric' character strings or 'strings'.

DEC Initials for Directory Entry Code.

DECIMAL See 'BASE 10'.

DIRECT ACCESS Retrieval or storage of data by a reference to its location on a disk, rather than relative to the previously retrieved or stored data.

DIRECT STATEMENT (IN FILE) A program statement that exists in the disk file that is not assigned a line number.

DIRECTORY A table giving the relationships between items of data. Sometimes a table or an index giving the addresses of data.

DISPLACEMENT A specified number of sectors, at the top or beginning of the file, in which the 'bookkeeping' and file parameters are stored for later use by the various program modules.

DISTRIBUTED FREE SPACE Space left empty at intervals in a data layout to permit the possible insertion of new data.

DOUBLE PRECISION A positive or negative numeric value, 16 digits in length, not including a decimal point (EXAMPLE: 99999999999999.99).

DUMP To transfer all or part of the contents of one section of computer memory or disk into another section, or to some other computer device.

DYNAMIC STORAGE ALLOCATION The allocation of storage space by a procedure based on the instantaneous or actual demand for storage space by that procedure, rather than allocating storage space to a procedure based on its anticipated or predicted demand.

EATEN (DIRECTORY/DISK) Slang term. See 'CLOBBED'.

EMBEDDED POINTERS Pointers in the data records rather than in a directory.

ENTITY Something about which data is recorded.

EOF Initials for 'END OF FILE'. It is common practice to say that the EOF is record number nn or that the EOF is byte 15 of sector 12. Hence, it is a convenient term to use in describing the location of the last record or last byte in a file.

EXTENT A contiguous area of data storage.

FILE A collection of related records treated as a unit; The word file is used in the general sense to mean any collection of informational items similar to one another in purpose, form and content.

FILE PARAMETERS The data that describes or defines the structure of the file.

FILESPEC A file specification and may include the 'FILE NAME', 'FILE NAME EXTENSION', 'PASSWORD', and 'DISK DRIVE' specification.

FIELD See 'DATA ITEM'.

FLAKY Slang term -Alludes to less than acceptable performance.

FILE AREA The physical location of the file, on the disk, or in memory.

'FPDE' Initials for File Primary Directory Entry; a file's entry and file area pointers in the disk directory.

'FXDE' Initials for File Extended Directory Entry; a file's entry and file area pointers, in the case of an overflow in the 'FPDE'.

GAT Initials for Granule Allocation Table; A table from which available file areas are assigned to file entries.

GRANULE Unit of 5 sectors. On the TRS-80 disk operating system, a 'granule' is the basic unit of disk storage allocation. The diskette 'DIRECTORY' file keeps track of free and assigned disk space in terms of 'granules'.

HASH CODE A code number generated and used as a direct addressing technique in which the key is converted to a pseudo-random number from which the required address is derived.

HEADER RECORD A record containing common, constant or identifying information for a group of records which follow.

HEXADECIMAL See 'BASE 16'

HIT Initials for Hash Index Table; an addressing technique in which a disk file is referenced by a code number in a table, and the position of that code in the table relates to the file entry in the directory.

INDEX A table used to determine the location of a record.

INDIRECT ADDRESSING Any method of specifying or locating a storage location whereby the key (of itself or through calculation) does not represent an address. For example, locating an address through indices.

INSTRING (INSTRING SEARCH) Refers to the capability of locating a substring of characters that may exist in another character string. An example would be: Substring = "THE" String = "NOW IS THE TIME". An INSTRING routine would locate the substring and return its starting position within that string. In this example, it would return a value of eight.

INTEGER A natural or whole number. In the TRS-80, integer values may not exceed the range of +32767 to -32768.

INVERTED FILE A file structure which permits fast spontaneous searching for previous unspecified information. Independent lists or indices are maintained in records' keys which are accessible according to the values of specific fields.

INVERTED LIST A list organized by a secondary key --- not a primary key.

IPL Initials for Initial Program Loader; a program usually executed upon pressing of the 'RESET' button.

KEY A data item used to identify or locate a record or other data grouping.

LABEL A set of symbols used to identify or describe an item, record, message or file. Occasionally, it may be the same as the address in storage.

LEAST SIGNIFICANT BYTE The significant byte contributing the smallest quantity to the value of a numeral.

LIST An ordered set of data items. A 'chain'.

LOAD MODULE A program developed for loading into storage and being executed when control is passed to the program.

LOCK-OUT (TRACKS) Unusable tracks, on the disk, that are not accessible because of damage or by user option.

LOGICAL An adjective describing the form of data organization, hardware or system that is perceived by an application program, programmer, or user; it may be different than the real (PHYSICAL) form.

LOGICAL DATA-BASE DESCRIPTION A schema. A description of the overall data-base structure, as perceived for the users, which is employed by the data base management software.

LOGICAL FILE A file as perceived by an application program; it may be in a completely different form from that in which it is stored on the storage units.

LOGICAL OPERATOR A mathematical symbol that represents a mathematical process to be performed on an associated operand. Such operators are 'AND', 'OR', 'NOT', 'AND NOT' and 'OR NOT'.

LOGICAL RECORD A record or data item as perceived by an application program; it may be in a completely different form from that in which it is stored on the storage units.

LSB See LEAST SIGNIFICANT BYTE.

MACHINE LANGUAGE Direct machine readable code.

MAINTENANCE OF A FILE (1) The addition, deletion, changing or updating of records in the database. (2) Periodic reorganization of a file to better accommodate items that have been added.

MONITOR A program that may supervise the operation of another program for operation or debugging or other purposes.

MOST SIGNIFICANT BYTE The significant byte contributing the greatest quantity to the value of a numeral.

MSB See MOST SIGNIFICANT BYTE.

MULTIPLE-KEY RETRIEVAL Retrieval which requires searches of data based on the values of several key fields (some or all of which are secondary keys).

NULL An absence of information as contrasted with zero or blank for the presence of no information.

NYBBLE The four right most or left most binary digits of a byte.

ON-LINE An on-line system is one in which the input data enter the computer directly from their point of origin, and/or output data are transmitted directly to where they are used. The intermediate stages such as writing tape, loading disks or off-line printing are avoided.

ON-LINE STORAGE Storage devices and especially the storage media which they contain under the direct control of a computing system, not off-line or in a volume library.

OPEN RECORDS TABLE (ORT) A table of binary values in which the relative position of a single value determines the status of a record with the same relative position; i.e., the Nth binary number determines the status of the Nth record. EXAMPLE: If the 8th binary number in the table is a zero, then the 8th record is open. Conversely, if the 8th binary number in the table is a one, then the 8th record is on file.

OPERATING SYSTEM Software which enables a computer to supervise its own operations, automatically calling in programs, routines, language and data as needed for continuous throughput of different types of jobs.

PARITY Parity relates to the maintenance of a sameness of level or count, i.e., keeping the same number of binary ones in a computer word and thus be able to perform a check based on an even or odd number for all words under examination.

PHYSICAL An adjective, contrasted with logical, which refers to the form in which data or systems exist in reality. Data is often converted by software from the form in which it is physically stored to a form in which a user or programmer perceives it.

PHYSICAL DATA BASE A data base in the form in which it is stored on the storage media, including pointers or other means of interconnecting it. Multiple logical data bases may be derived from one or more physical data bases.

PHYSICAL RECORD A collection of bits that are physically recorded on the storage medium and which are read or written by one machine input/output instruction.

POINTER The address of a record (or other data groupings) contained in another record so that a program may access the former record when it has retrieved the latter record. The address can be absolute, relative or symbolic, hence, the pointer is referred to as absolute, relative or symbolic.

PRIMARY ENTRY The main entry made to the directory. Also see 'FPDE'.

RANDOM ACCESS To obtain data directly from any storage location regardless of its position, with respect to the previously referenced information. Also called 'DIRECT ACCESS'.

RANDOM ACCESS STORAGE A storage technique in which the time required to obtain information is independent of the location of the information most recently obtained.

READ To accept or copy information or data from input devices or a memory register; i.e., to read out, to read in.

RECORD A group of related fields of information treated as a unit by an application program.

RELATIONAL OPERATOR A mathematical symbol that represents a mathematical process to perform a comparison describing the relationship between two values (< less than....> greater than... = equal.... <> not equal... and combinations thereof (see TRS-80 LEVEL II manual, Section 1, Page 5). On the TRS-80, relational comparisons may be made on string values as well as numerical values.

RELATIVE (as pertains to position) An address or position that is referenced to a point of origin; i.e. X+20 is a specific position, 20 places from the reference point. If the reference point was at 50, then the absolute position would be at 70 (50+20=70). Also, 50 (since it is the starting reference point) is at relative position 0.

SCHEMA A map of the overall logical structure of a database.

SEARCH To examine a series of items for any that have a desired property or properties.

SECONDARY INDEX An index composed of secondary keys rather than primary keys.

SECTOR The smallest addressable portion of storage on a diskette (a unit of 256 bytes on a TRS-80 diskette).

SEEK To position the access mechanism of a direct-access storage device at a specified location.

SEQUENTIAL ACCESS Access in which records must be read serially or sequentially one after the other; i.e., ASCII files, tape.

SINGLE PRECISION A positive or negative numerical value of 6 digits in length, not including a decimal point (EXAMPLE: 99999.9).

SORT To arrange a file or data in a sequence by a specified key (may be alphabetic or numeric and in descending or ascending order).

SOURCE CODE The text from which code that may be executed is derived.

SYSTEM FILE A program used by the operating system to manage the executing program and/or the computer's resources.

SUB-STRINGS SUB-STRING SEARCH See INSTRING

TABLE A collection of data suitable for quick reference, each item being uniquely identified either by a label or its relative position.

TALLY To add or subtract a digit from a quantity.

TOKEN A one byte code representing a larger word consisting of 2 or more characters.

TRACK The circular recording surface traversed by a read/write head on the disk. On the TRS-80 a track contains 10 sectors (2 granules).

TRANSACTION An input record applied to an established file. The input record describes some "event" that will either cause a new file record to be generated, an existing record to be changed or an existing record to be deleted.

TRANSPARENT Complexities that are hidden from the programmers or users (made transparent to them) by the software.

VECTOR A line representing the properties of magnitude and direction. Since such a 'line' can be described in mathematical terms, a mathematical description (expressed in numbers, of course) of a given 'direction' and 'magnitude' is referred to as a "vector".

VERIFY To check a data transfer or transcription.

WORKING STORAGE A portion of storage, usually computer main memory, reserved for the temporary results of operations.

WRITE To record information on a storage device.

ZAP To change a byte or bytes of data in memory or on diskette by using a software utility program.

ZEROETH Zeroeth is to '0' as first is to '1'; in computer terms the first position of anything is usually described as the 'zeroeth' and the next position is the 'first' and so on.

NOTES

PARTS LIST
SIMPLE H08BYIST INTERFACE

IC	TYPE	+5 VOLTS	GROUND
Z1	81LS95	20	10
Z2	81LS96	20	10
Z3	81LS95	20	10
Z4	81LS95	20	10
Z5	74LS75	5	12
Z6	74LS75	5	12
Z7	74LS90	14	7
Z8	74LS02	14	7

PARTS LIST
ONE-SECOND INTERRUPT REAL-TIME CLOCK

IC	TYPE	+5 VOLTS	GROUND
Z1	7805	3 (OUTPUT)	2
Z2	74LS14	14	7
Z3	74LS90	5	10
Z4	74LS92	5	10
Z5	74LS74	14	7
Z6	74LS75	5	12

PART	VALUE	NOTES
C1	4700 mf, 16 volts	
C2	220 mf, 16 volts	
C3	100 nf (0.1 mf)	
C4	100 nf (0.1 mf)	

PARTS LIST
MSM5832 REAL-TIME CLOCK/CALENDAR

IC	TYPE	+5 VOLTS	GROUND
Z1	74LS90	14	7
Z2	74LS260	14	7
Z3	INS8255	26	7
Z4	MSM5832	1	13

PART	VALUE	NOTES
C1	20 pf	
C2	20 pf	
C3	100 nf (0.1 mf)	
C4	100 nf (0.1 mf)	
C5	100 nf (0.1 mf)	
C6	100 nf (0.1 mf)	
R1-R12	10k (10,000 ohms)	

Battery backup:

PART	VALUE	NOTES
C6	470 mf, 16 volts	
C7	10 mf, 16 volts	
C8	100 nf (0.1 mf)	
C9	10 mf, 16 volts	
D1	Bridge rectifier, 1A 50V	
D2	1N914 or equivalent	
Q1	PNP transistor, Vce=0.1 volt	
Q2	NPN transistor	
R13	47k (47,000 ohms)	
R14	10k (10,000 ohms)	
R15	10k (10,000 ohms)	
R16	100R (100 ohms), 0.5 W	
T1	6V3 (6.3 volt), 1A transformer	
Z5	7805 5-volt regulator	

PARTS LIST
QUAD SOUND, 800PS AND 8LEEPS GENERATOR

IC	TYPE	+5 VOLTS	GROUND
Z1	74LS374	20	10
Z2	74LS125	14	7
Z3	74LS30	14	7
Z4	74LS30	14	7
Z5	7LS02	14	7

PART	VALUE	NOTES
R1-R5	1k (10,000 ohms)	

PARTS LIST
MEMORY SIOECAR ROM AND RAM ADDITION

IC	TYPE	+5 VOLTS	GROUND
Z1	74LS30	14	7
Z2	74LS02	14	7
Z3	74LS00	14	7
Z4	74LS00	14	7
Z5	74LS125	14	7
Z6,7,8,9	2114-AN4L	18	9
Z10	2716	24	12

PART	VALUE	NOTES
R1-R16	1kO (10,000 ohms)	

PARTS LIST
MUSIC SYNTHESIZER INTERFACE 80AR0

IC	TYPE	+5 VOLTS	GROUND
Z1	81LS95	20	10
Z2	74LS154	24	12
Z3	74LS00	14	7
Z4	74LS04	14	7
Z5	INS8255	26	7
Z6	MC1408L8 OR 0AC0808	13	2
Z7	MC1408L8 OR 0AC0808	13	2
Z8	74LS123	14	7
Z9	74LS123	14	7
Z10	LM324	4	11

Note that Z6 and Z7 also require -15 volts on pin number 8 and +7.5 volt reference at pin 14 (through a variable resistor).

PART	VALUE	NOTES
C1,2	20 pf	
C3,4,5,6	100 nf (0.1 mf) (may be changed)	
R1,3	4k7 (4700 ohms)	
R2,4	1kO (1000 ohms) (may be changed)	
R5,6,7,8	1MO (1,000,000 ohms)	

PARTS LIST
BANK-SELECT ROM/RAM ADDITION

IC	TYPE	+5 VOLTS	GROUND
Z1	81LS95	20	10
Z2	81LS95	20	10
Z3	81LS95	20	10
Z4	81LS95	20	10
Z5-11	2716 OR 4118	24	12
Z12	74LS20	14	7
Z13	74LS260	14	7
Z14	74LS02	14	7
Z15	74LS75		
Z16	74154	24	12
Z17	74LS04	14	7
Z18-26	2716 OR 4118	24	12

PARTS LIST
8-TRACK MASS STORAGE SYSTEM

IC	TYPE	+5 VOLTS	+12 VOLTS	GROUND
Z1	70C98/60C98	.	.	.
Z2	74LS04	14		7
Z3	74LS125	14		7
Z4	LF353		.	.
Z5	LM339		.	.
Z6	70C98/80C98	.	.	.
Z7	LF353		.	.
Z8	74LS00	14		7
Z9	75452	8		4
Z10	70C96/80C96	.	.	.
Z11	75452	8		4
Z12	74LS373	20		10
Z13	74LS30	14		7
Z14	74LS04	14		7

PART	VALUE	NOTES
C1	100 pf	
C2	220 pf	
C3	220 nf (0.22 mf)	
C4	10 mf, 16 volts	
C5	100 pf	
C6	220 pf	
C7	220 nf (0.22 mf)	
C8	10 mf, 16 volts	
C9,10,11	100 nf (0.1 mf)	
C12	470 mf, 35 volts	
D1,2,3,4,5	1N914 or equivalent	
K1,2,3	Miniature SPST 5-volt relay	
R1,2,14,17, 18,30,38,39, 40,41	10k (10,000 ohms)	
R3,4,12,19, 20,28	470R (470 ohms)	
R5,6,9,11,16, 22,25,27,32	2k2 (2200 ohms)	
R7,8,23,24	220k (220,000 ohms)	
R10,26	100k (100,000 ohms)	
R13,29	22k (22,000 ohms)	
R15,31	220R (220 ohms)	
R33,34,35,36, 37,42	1k0 (1000 ohms)	
R43	100R (100 ohms)	
R44	75R (75 ohms)	May need adjustment
S1	Cartridge-in-place leaf switch	
S2	Track change light switch	
S3	Split-pole fail sensor switch	

PARTS LIST
MICRO FRONT PANEL MONITOR

IC	TYPE	+5 VOLTS	GROUND
Z1	74LS373	20	10
Z2	74LS373	20	10
Z3	74LS373	20	10
Z4	74LS20	14	7

PART	VALUE	NOTES
DIS1,2,3	10-segment bar LED	
LED1-24	Subminiature LED	
R1-24	270R (270 ohms)	
R25,26,27,28	1k0 (1000 ohms)	
S1	4-position DIP switch	

PARTS LIST
4K DYNAMIC RAM ADDITION

IC	TYPE	+5 VOLTS	GROUND
Z1	81LS95	20	10
Z2	74LS157	16	8
Z3	81LS95	20	10
Z4	74LS157	16	8
Z5	81LS95	20	10
Z6-13	MK4114 4K RAMS	9	16
Z14	74LS20	14	7
Z15	74LS86	14	7
Z16	81LS95	20	10

Note that Z6 through Z13 also require -5 volts on pin number 1 and +12 volts on pin number 8.

PART	VALUE	NOTES
R1,2,3	1k0 (1000 ohms)	
S1	4-position DIP switch	

PARTS LIST
HIGH-SPEED, REVERSE VIDEO, UPPER/LOWER CASE,
INDIVIDUAL REVERSE VIDEO M0US

IC	TYPE	+5 VOLTS	GROUND
F45 (ZMEM)	2102-4L	10	9
P25 (ZBITS)	74LS10	14	7
P6 (ZFAZE)	74LS04	14	7
P27 (ZMODE)	74LS368	20	10
P24 (ZMUXX)	74LS86	14	7
P26 (ZFLOP)	74LS74	14	7
P53 (ZPORT)	74LS02	14	7
P44 (ZFAST)	74LS367	20	10

PC VERSION ONLY:

PART	VALUE	NOTES
C1	330 pf	
C2	33 nf (.033 mf)	
R1	10k (10,000 ohms)	
VCR1	100k (100,000 ohms)	Variable

Miscellaneous:

One 16-pin wire-wrap integrated circuit socket for piggybacking the PC board version onto Z45.
No socket is needed for the hard-wired version.

PARTS LIST
HIGH-RESOLUTION GRAPHICS BOARD
POWER SUPPLY

PART	VALUE	NOTES
C1	4700 mf, 35 volts	
C2	100 mf, 25 volts	
C3	100 nf (0.1 mf)	
C4	10,000 mf, 35 volts	
C5	100 mf, 16 volts	
C6	100 nf (0.1 mf)	
C7	470 mf, 16 volts	
C8	100 mf, 16 volts	
C9	100 nf (0.1 mf)	
R1	220 ohms, 1/2 watt	

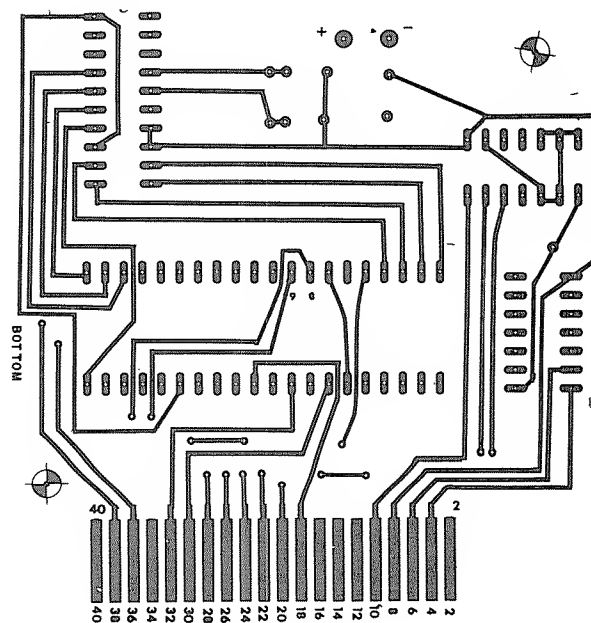
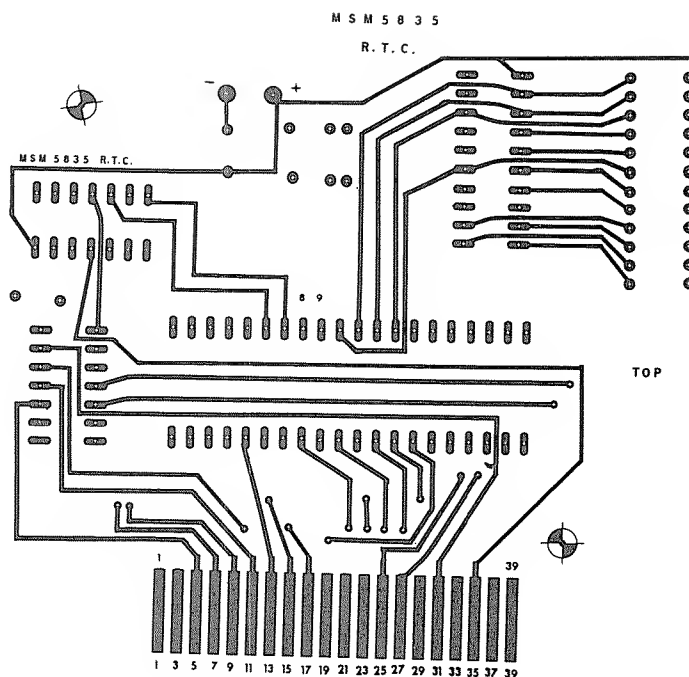
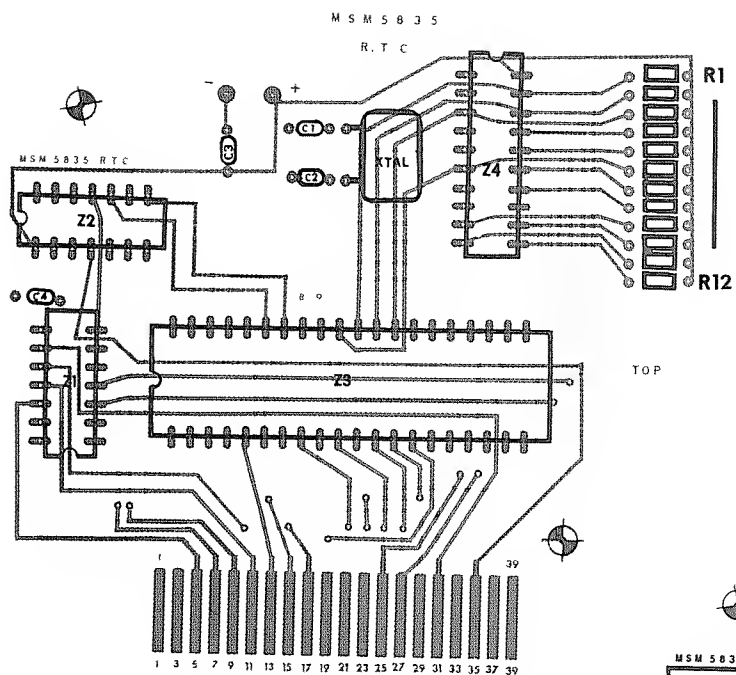
PARTS LIST
HIGH-RESOLUTION GRAPHICS BOARD

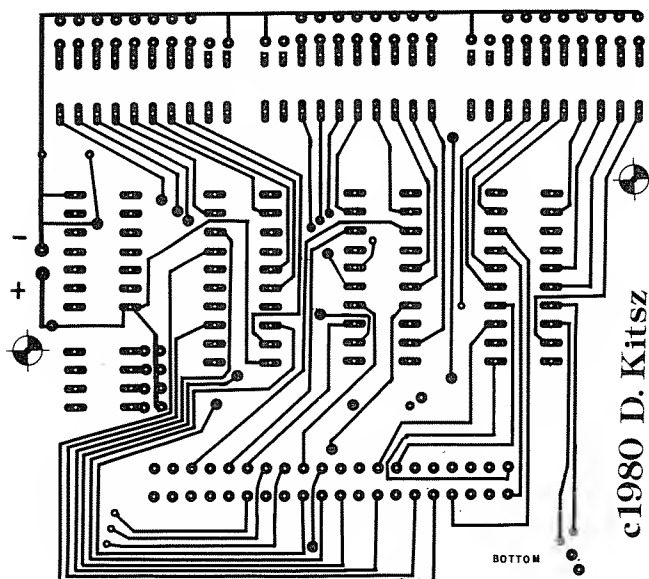
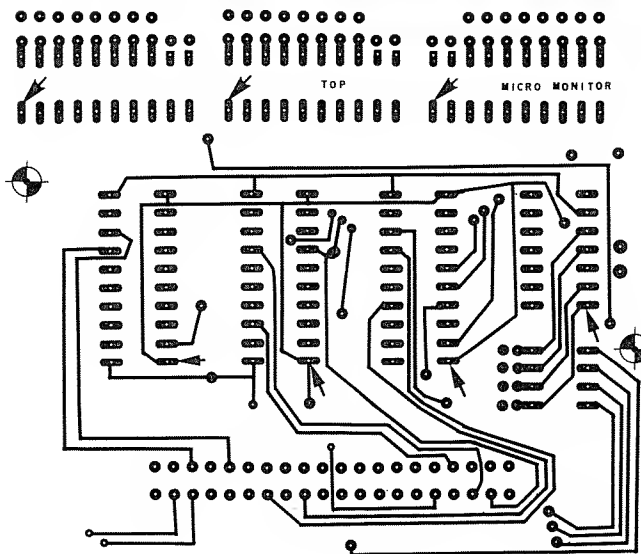
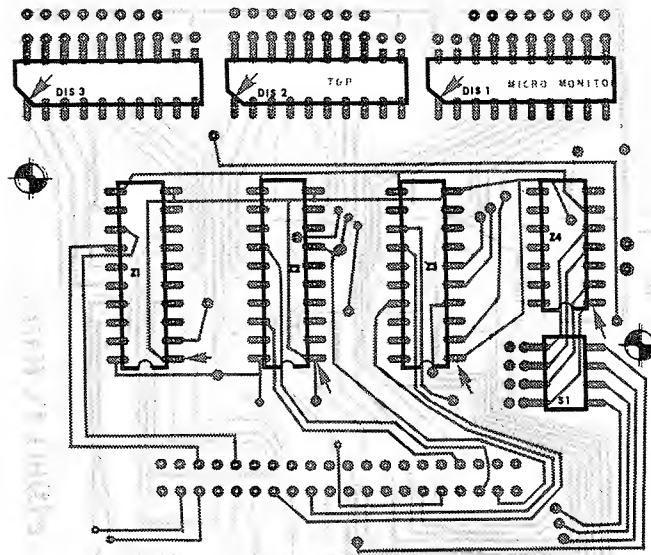
IC	TYPE	+5 VOLTS	GROUND
Z1	74LS04	14	7
Z2	74LS92	5	10
Z3	74LS74	14	7
Z4	74LS93	5	10
Z5	74LS93	5	10
Z6	74LS93	5	10
Z7	74LS93	5	10
Z8	74LS11	14	7
Z9	74LS02	14	7
Z10	74LS157	16	8
Z11	74LS157	16	8
Z12	74LS157	16	8
Z13	74LS157	16	8
Z14	74LS74	14	7
Z15	74LS00	14	7
Z16	74LS174	16	8
Z17	74LS166	16	8
Z18	74LS157	16	8
Z19	74LS157	16	8
Z20	74LS157	16	8
Z21	MK4116	9	16
Z22	MK4116	9	16
Z23	MK4116	9	16
Z24	MK4116	9	16
Z25	MK4116	9	16
Z26	MK4116	9	16
Z27	74C04	14	7
Z28	74C04	14	7
Z29	74C00	14	7
Z30	75452	8	4
Z31	74LS00	14	7

Note that Z21 through Z26 also need -5 volts on pin number 1, and +12 volts on pin number 8.

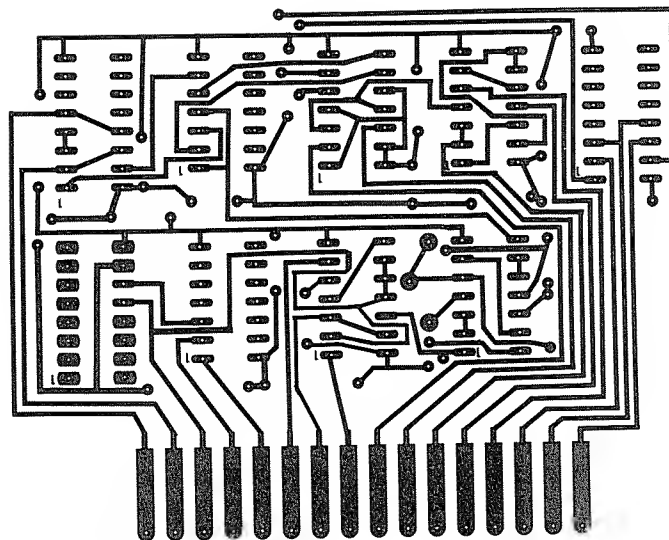
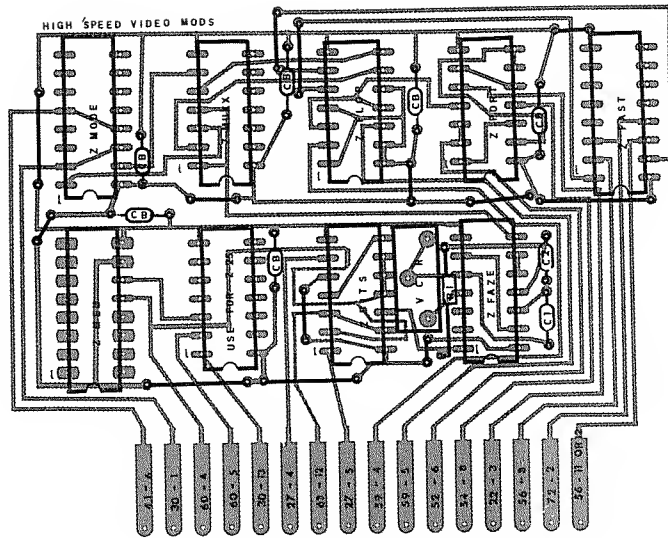
PART	VALUE	NOTES
C1*	47 pf	Miniature variable
C2	10 nf (.01 mf)	
C3	10 mf, 16v	
C4	330 pf	10% or better
C5	47 nf (.047 mf)	10% or better
C6	750 pf	10% or better
C7	22 nf (.022 mf)	10% or better
Q1	2N3904	NPN switching
Q2	2N3904	NPN switching
R1	1k0 (1000 ohms)	
R2	1k0 (1000 ohms)	
R3	1k0 (1000 ohms)	
R4*	510R (510 ohms)	
R5*	510R (510 ohms)	
R6	1k8 (1800 ohms)	
R7	47R (47 ohms)	
R8	270R (270 ohms)	
R9	120R (120 ohms)	
R10	330R (330 ohms)	
R11	75R (75 ohms)	
R12	10k (10,000 ohms)	
R13	10k (10,000 ohms)	
VCR1	100k (100,000 ohms)	Variable
VCR2	100k (100,000 ohms)	Variable
X1*	10.6445 MHz	

*Not needed if computer's internal clock is used.

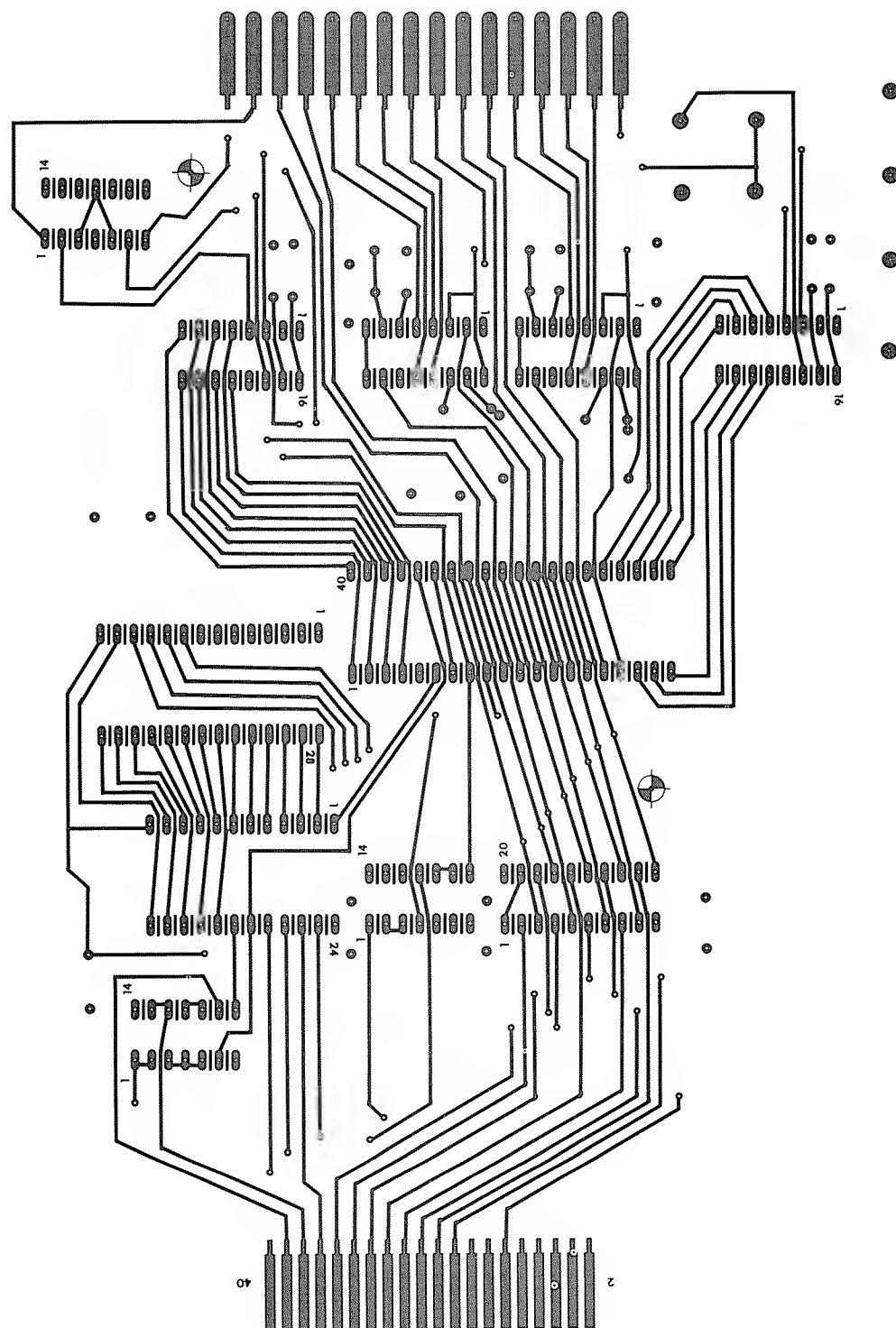


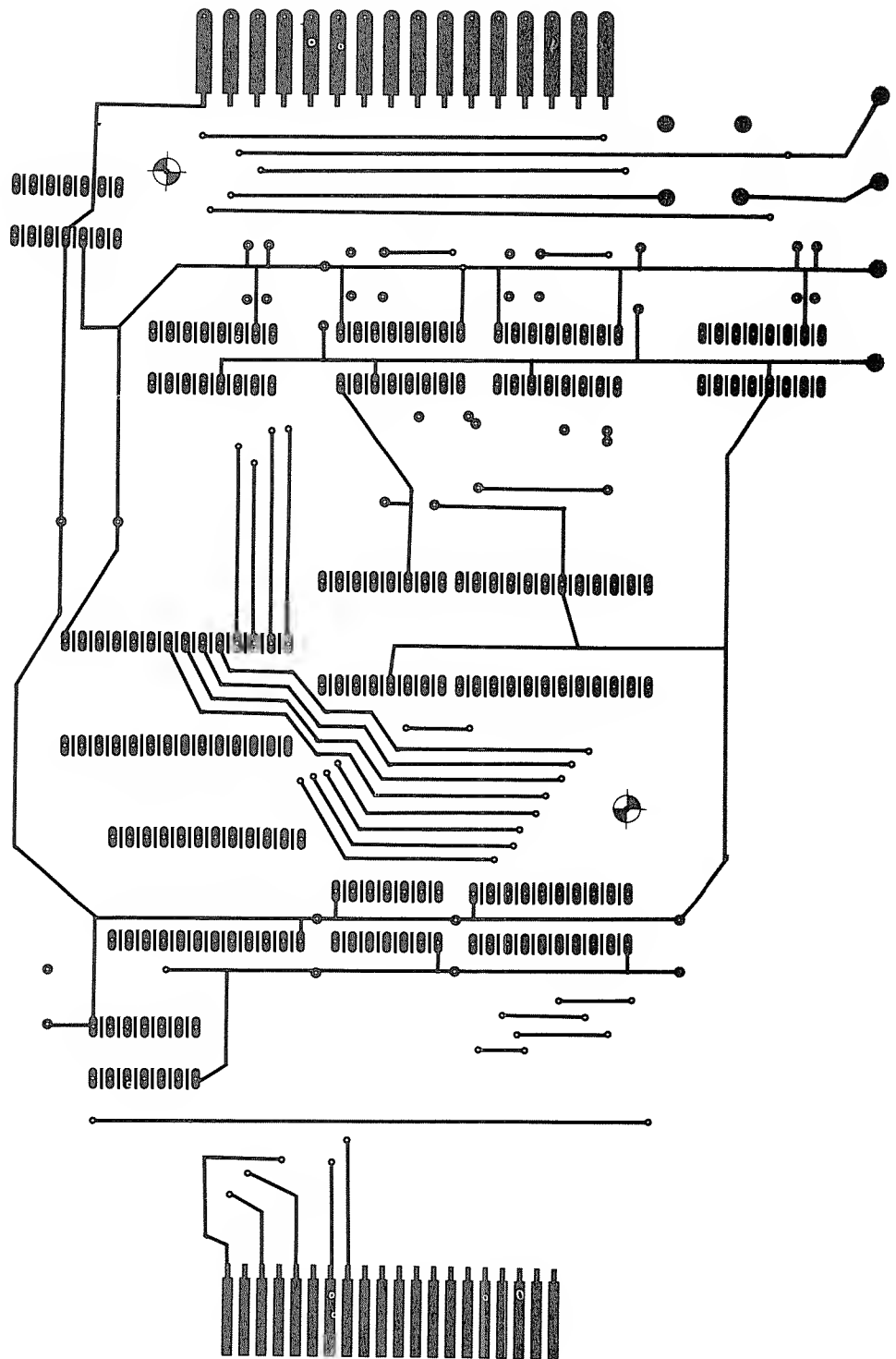


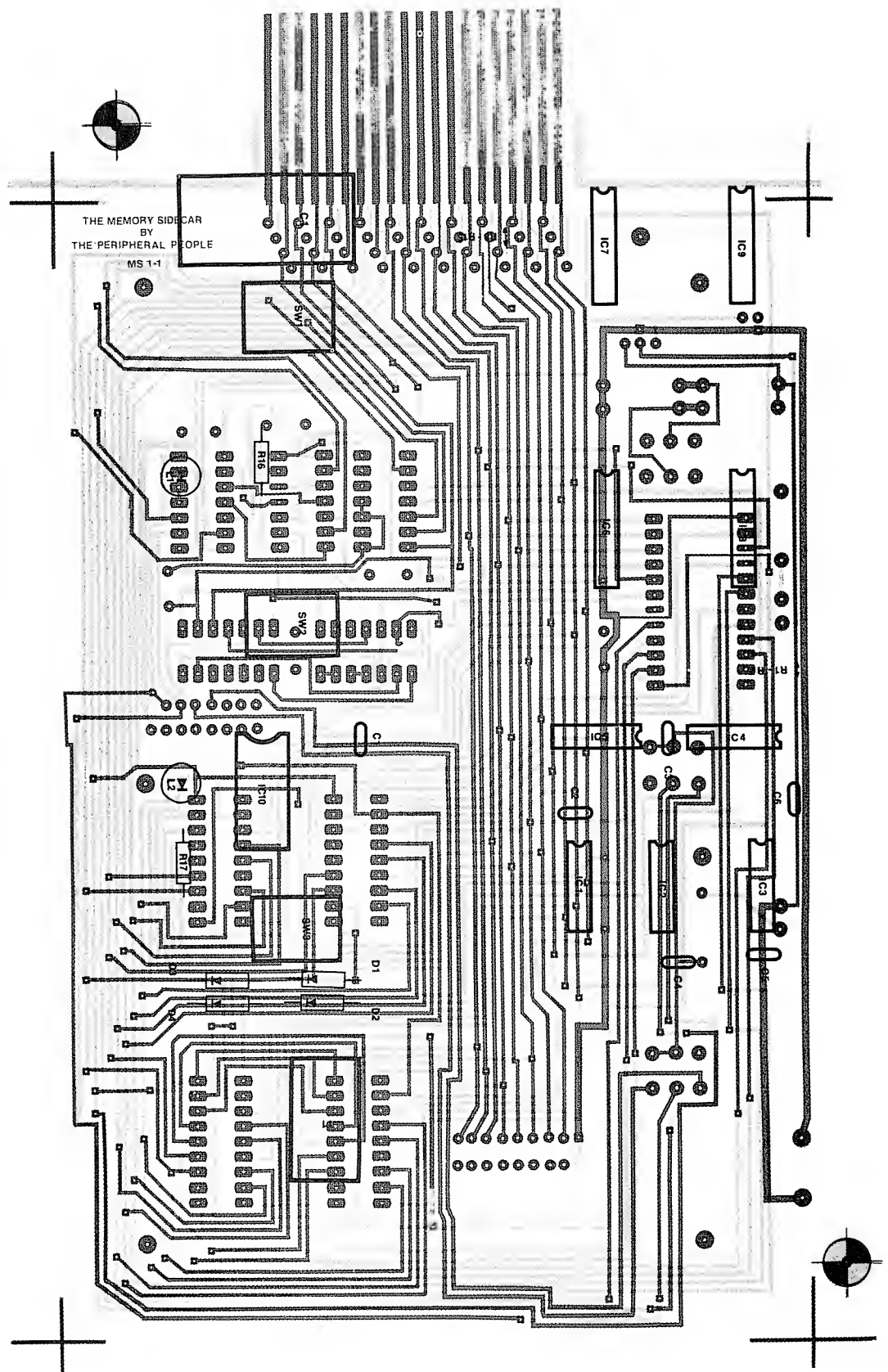
c1980 D. Kitz

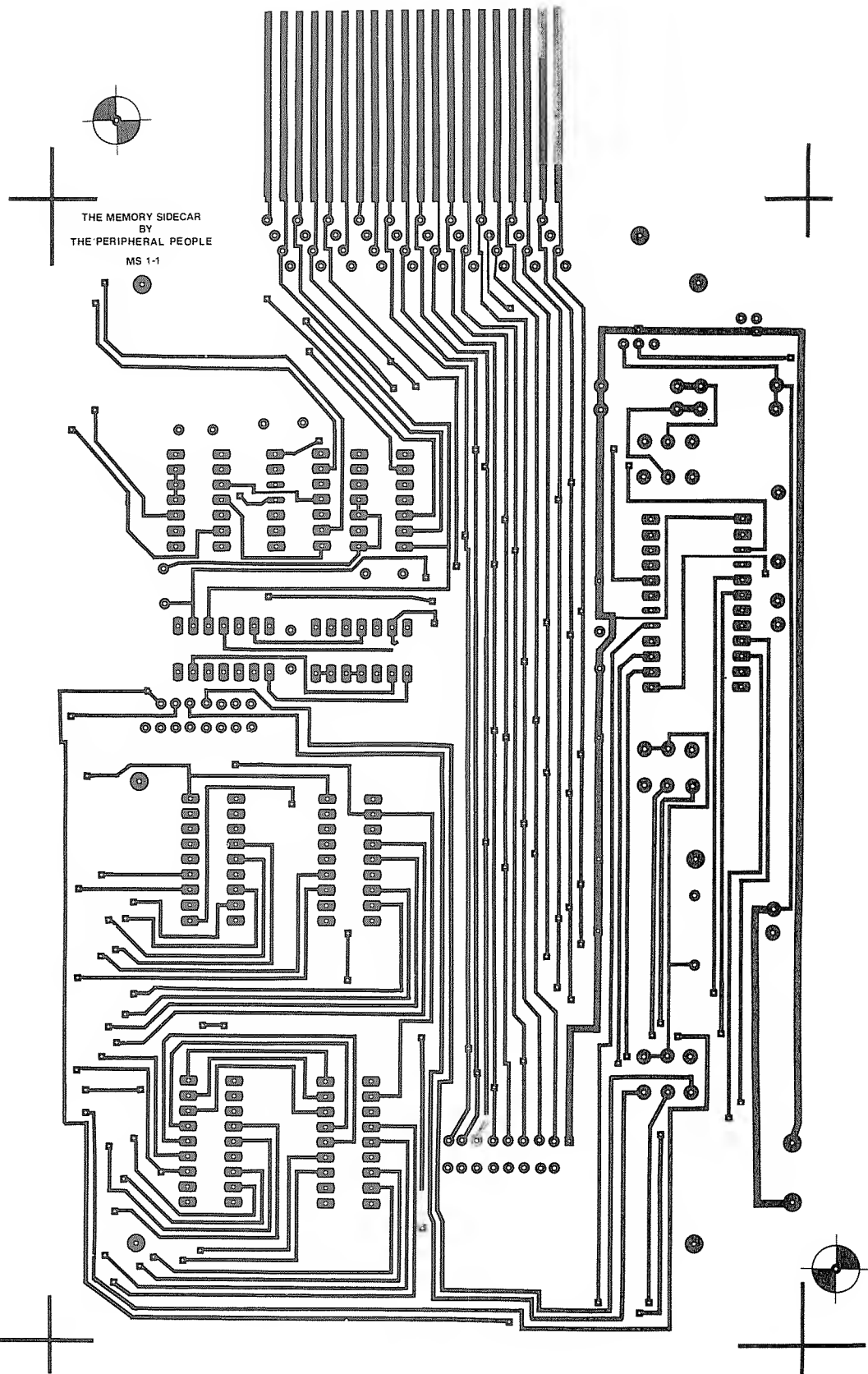


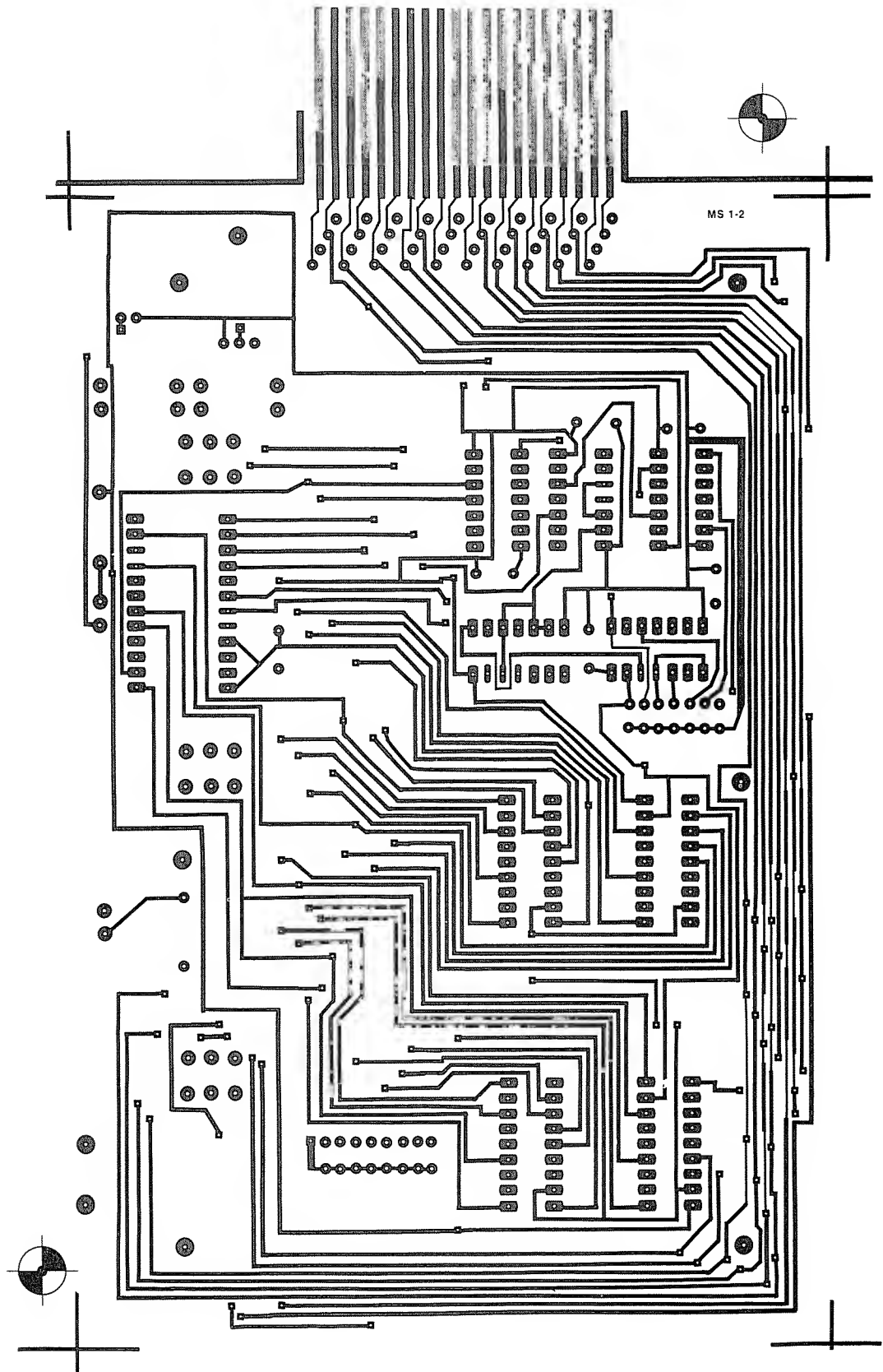


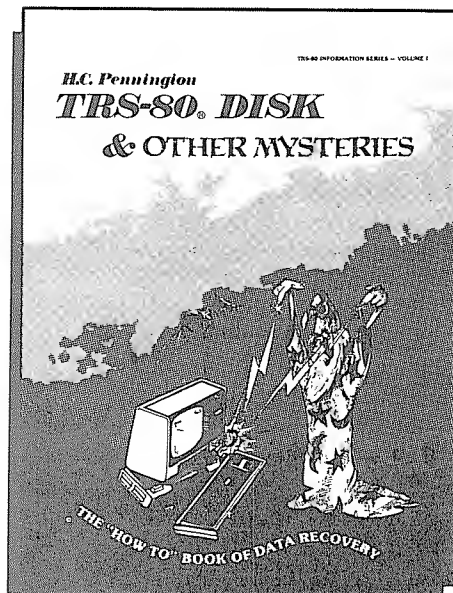












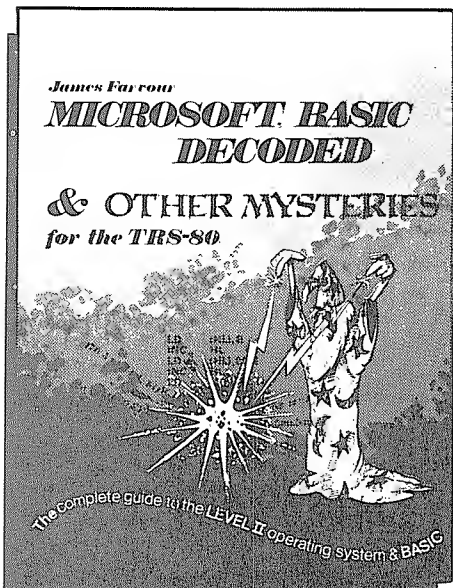
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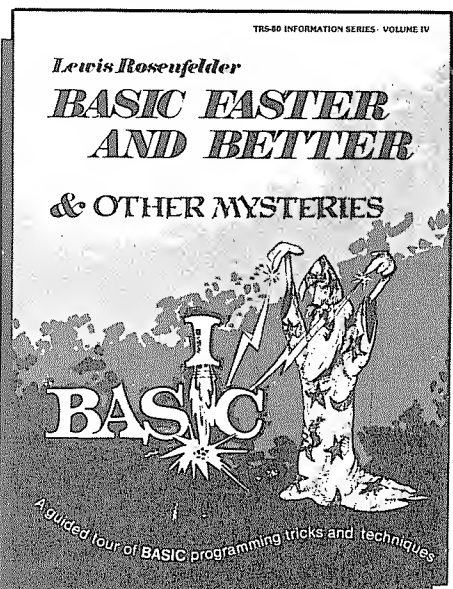
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